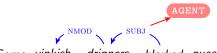
# L90: Overview of Natural Language Processing Lecture 9: Compositional Semantics

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syntax provides scaffolding for semantic composition

#### Lecture 9: Compositional Semantics

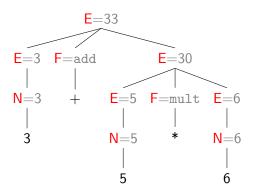
- 1. Being able to transform
- 2. Semantic composition
- 3. Graph-based meaning representations
- 4. Inference and RTE

many slides are from Ann Copestake Principle: Being Able to Transform

# Programming language interpreter

#### What is meaning of 3 + 5 \* 6?

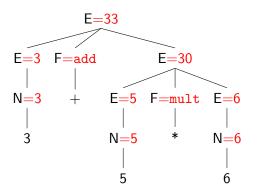
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- Now give a meaning to each node in the tree (bottom-up)



# Programming language interpreter

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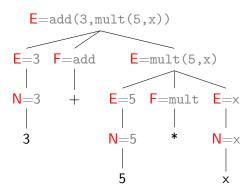
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### Interpreting in an environment

#### How about 3 + 5 \* x?

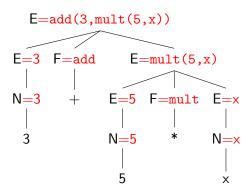
- Don't know x at compile time
- Meaning at a node is a piece of code, not a number



## Interpreting in an environment

#### How about 3 + 5 \* x?

- Don't know x at compile time
- Meaning at a node is a piece of code, not a number





There is in my opinion no important theoretical difference between natural languages and the artificial languages of logicians; indeed, I consider it possible to comprehend the syntax and semantics of both kinds of languages within a single natural and mathematically precise theory.

Richard Montague, 1930–1971

## What counts as understanding?

Charaterizing what we mean by *meaning* is a difficult philosophical issue.

a compiler is a translator (formal language to formal language)

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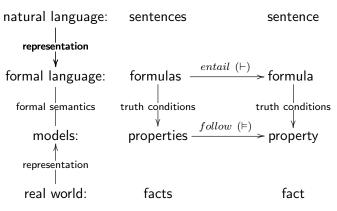
↓
Natural Language Understanding ▷ being able to translate

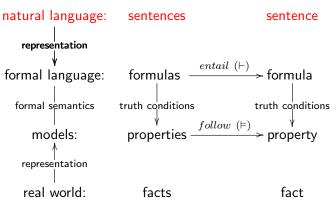
Natural language to natural language?

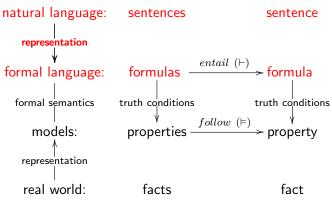
Reasonable. Sometimes requires deeper understanding

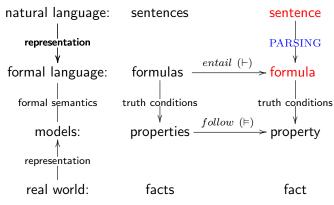
Natural language to formal language (defined by logic)?

Popular in NLP.



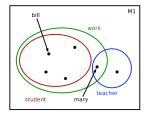


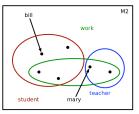




# Logic as a formal language

- [[every student works]] = true iff. student ⊆ work
- every student works  $\Rightarrow \forall x(\operatorname{stud}'(x) \to \operatorname{work}'(x))$





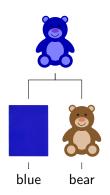
- Logic supports precise, consistent and controlled meaning representation via truth-conditional interpretation.
- Logic provides deduction systems to model inference processes, controlled through a formal entailment concept.
- Logic supports uniform modelling of the semantic composition process.

# Semantic Composition

#### The Principle of Compositionality

The meaning of an expression is a function of the meanings of its parts and of the way they are syntactically combined.

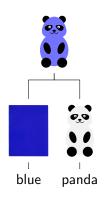
B. Partee



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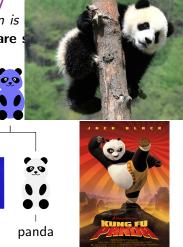


blue

The Principle of Compositionality

The meaning of an expression is

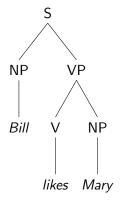
its parts and of the way they are

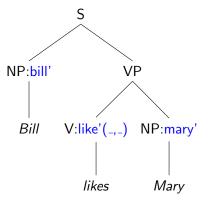


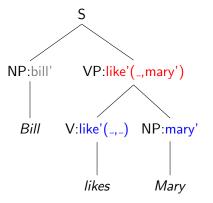
Syntactic parsing + Lexical interpretation

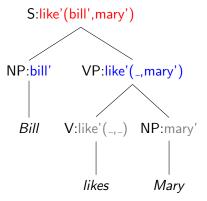
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Meaning representation







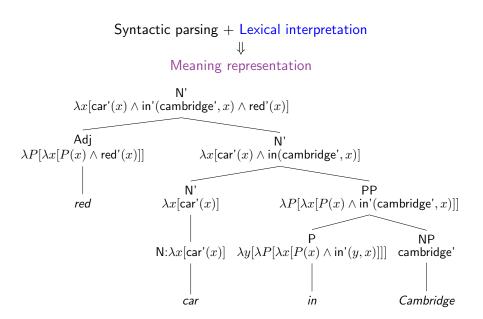


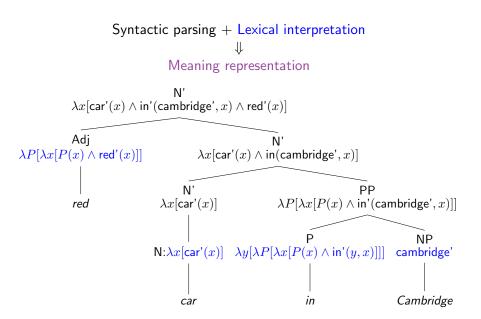
# Using $\lambda$ 's

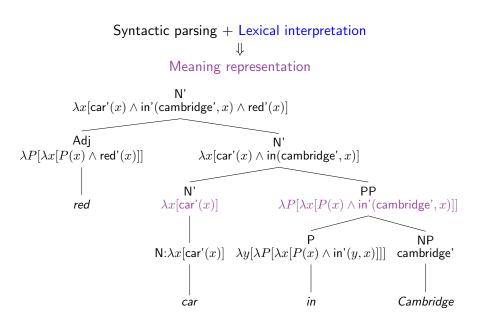
- Church defined an idealized programming language called the λ-calculus.
- A formal system in mathematical logic. A model of computation.
- λ-reduction:
  - $\lambda x[\text{sleep'}(x)](\text{john'})$  becomes sleep'(john')
  - $\lambda y[\lambda x.love'(x,y)](pizza')$  becomes  $\lambda x[love'(x,pizza')]$
  - $\lambda x[\mathsf{love'}(x,\mathsf{pizza'})](\mathsf{john'})$  becomes  $\mathsf{love'}(\mathsf{john'},\mathsf{pizza'})$

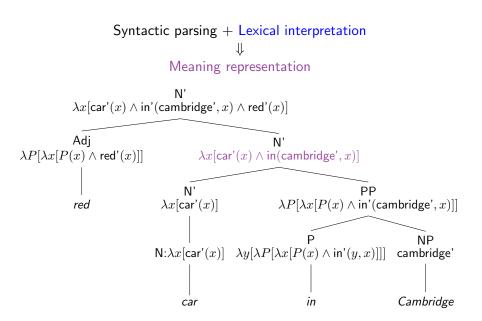
• 
$$f(5) = 25$$
  $\lambda x[x^2](5) = 25$ 

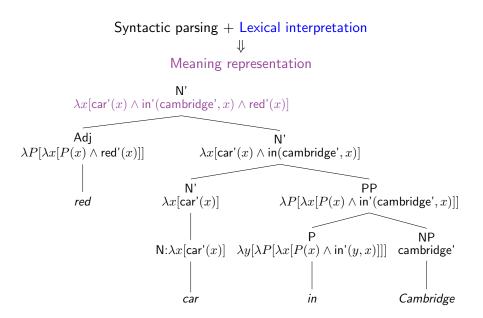
• 
$$g(x,y) = x^2 + y^2$$
 
$$\lambda x[\lambda y[x^2 + y^2]]$$











# Semantic composition rules are non-trivial

Ordinary pronouns contribute to the semantics:

- (1) a. It barked.
  - b.  $\exists x [\mathsf{bark'}(x) \land \mathsf{PRON}(x)]$

Pleonastic pronouns don't:

- (2) a. It rained.
  - b. rain'

Similar syntactic structures may have different meanings. Different syntactic structures may have the same meaning:

- (3) a. Kim seems to sleep.
  - b. It seems that Kim sleeps.

Differences in presentation but not in truth conditions.

# Beyond toy examples . . .

Use first order logic where possible (e.g., event variables, next slide).

However, First Order Predicate Calculus (FOPC) is sometimes inadequate: e.g., *most*, *may*, *believe*.

Quantifier scoping multiplies analyses:

- (4) a. Every cat chased some dog
  - b.  $\forall x[\mathsf{cat'}(x) \to \exists y[\mathsf{dog'}(y) \land \mathsf{chase'}(x,y)]]$
  - c.  $\exists y [\mathsf{dog'}(y) \land \forall x [\mathsf{cat'}(x) \to \mathsf{chase'}(x,y)]]$

Often no straightforward logical analysis e.g., Bare plurals such as *Ducks lay eggs*.

Non-compositional phrases (multiword expressions): e.g., *red tape* meaning bureaucracy.

#### Event variables

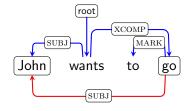
Allow first order treatment of adverbs and PPs modifying verbs by *reifying* the event.

- (5) a. Rover barked
  - b. bark'(r)
  - c.  $\exists e[\mathsf{bark'}(e,r)]$
- (6) a. Rover barked loudly
  - b.  $\exists e[\mathsf{bark'}(e,r) \land \mathsf{loud'}(e)]$

There was an event of Rover barking and that event was loud.

# Graph-Based Meaning Representations

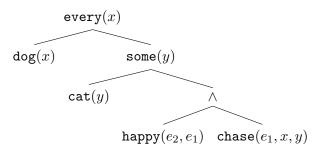
# Non-tree dependency structures



# Logical expression and semantic graph

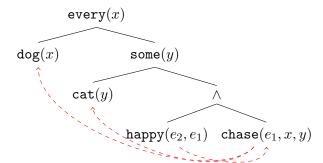
- Every dog chases some cats.
- $\exists y (\mathtt{cat}(y) \land \forall x (\mathtt{dog}(x) \to \mathtt{chase}(e, x, y)))$
- some(y, cat(y), every(x, dog(x), chase(e, x, y)))
- $\forall x (\operatorname{dog}(x) \to \exists y (\operatorname{cat}(y) \land \operatorname{chase}(e, x, y)))$
- every(x, dog(x), some(y, cat(y), chase(e, x, y)))

bracketing ⇒ tree

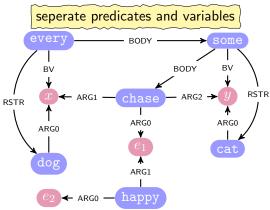


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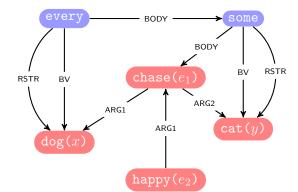
 $bracketing \Rightarrow tree$ 



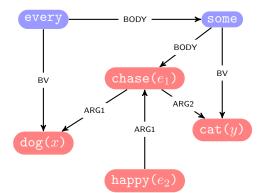
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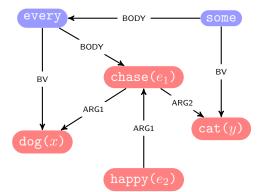
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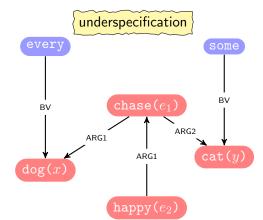
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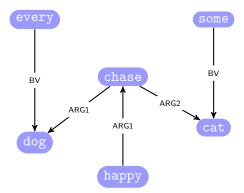
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# Inference and RTE

#### Natural language inference

**Inference on a knowledge base**: convert natural language expression to KB expression, valid inference according to KB.

- Precise
- Formally verifiable
- © Disambiguation using KB state
- © Limited domain, requires KB to be formally encodable

Language-based inference: does one utterance follow from another?

- Unlimited domain
- ☺/☺ Human judgement
- ⊗/ⓒ Approximate/imprecise

Both approaches may use logical form of utterance.

### Lexical meaning and meaning postulates

- Some inferences validated on logical representation directly, most require lexical meaning. What makes soup, soup?
- meaning postulates: e.g.,

$$\forall x [\mathit{bachelor'}(x) \rightarrow \mathit{man'}(x) \land \mathit{unmarried'}(x)]$$

usable with compositional semantics and theorem provers, e.g.

Problematic in general, OK for narrow domains or micro-worlds

# Recognising Textual Entailment (RTE) shared tasks

- T The girl was found in Drummondville earlier this month.
- *H* The girl was discovered in Drummondville.
  - Data: pairs of text (T) and hypothesis (H). H may or may not follow from T.
- *Task*: label true (if follows) or false (if doesn't follow), according to human judgements.

# RTE using logical forms

- T sentence has logical form T', H sentence has logical form H'
- If  $T' \Rightarrow H'$  conclude true, otherwise conclude false.
- The girl was found in Drummondville earlier this month.
- $T' \exists x, u, e[\mathsf{girl'}(x) \land \mathsf{find'}(e, u, x) \land \mathsf{in'}(e, \mathsf{drummondville'}) \land \mathsf{earlier\text{-}this\text{-}month'}(e)]$
- *H* The girl was discovered in Drummondville.
- $H' \ \exists x, u, e[\mathsf{girl'}(x) \land \mathsf{discover'}(e, u, x) \land \mathsf{in'}(e, \mathsf{drummondville'})]$
- $\mathsf{MP} \; \mathsf{find'}(x,y,z) \Rightarrow \mathsf{discover'}(x,y,z)$ 
  - So  $T' \Rightarrow H'$  and we conclude true

# More complex examples

- T Four Venezuelan firefighters who were traveling to a training course in Texas were killed when their sport utility vehicle drifted onto the shoulder of a highway and struck a parked truck.
- H Four firefighters were killed in a car accident.

Systems using logical inference are not robust to missing information: simpler techniques can be effective (partly because of choice of hypotheses in RTE).

# More examples

- T Clinton's book is not a big seller here.
- H Clinton's book is a big seller.
- T After the war the city was briefly occupied by the Allies and then was returned to the Dutch.
- H After the war, the city was returned to the Dutch.
- T Lyon is actually the gastronomic capital of France.
- *H* Lyon is the capital of France.

# An example from a linguist

- The Commissioner doesn't regret that the President failed to make him leave Athens before May 2.
- H The Commissioner was in Athens on May 2.

presupposition
negation
causation
event
semantic role
coreference
temporal expression

# Reading

- Ann's lecture notes.
- ACL tutorial on graph-based meaning representations https://github.com/cfmrp/tutorial