MSP430X port - basic checkpointing design documentation

David Garriou, Jean-Luc Béchennec January 23, 2020

1 Goals and ideas

1.1 Consistency

Our main goal is to keep a system within a coherent state even after a recovery from power loss. In other words, the kernel execution must be atomic with respect to power loss.

1.1.1 Consistency of kernel space and user space

When the system recovers all of its code, user space and kernel space, must be consistent. OS_VAR section contains the kernel data structures. It means that we must not transfer OS_VAR section to FRAM otherwise the system would recover with two differents (not consistent) states, the last user space state and the latest kernel space state.

We have to save both kernel and user space variables at the same time (during a single checkpoint).

1.2 Big picture

Big picture of what to do:

- A periodic task, let's call it energy_task, will check for remaining energy in the super capacitor;
 - (a) user space execution;
 - (b) remaining energy consists in mesuring VCC voltage (on the terminals of the super capacitor);
 - (c) if the remaining energy is lower than a defined threshold, we must realize a checkpoint

2 Normal MSP430X startup sequence

The MSP430X startup sequence is as follow:

- 1. After a reset, the tpl_reset_handler executes (see tpl_startup.S file). It:
 - (a) stops the watchdog timer;
 - (b) disables the interrupts;
 - (c) sets up the stack;
 - (d) calls tpl_continue_reset_handler C function.
- 2. tpl_continue_reset_handler (see tpl_startup.c file) does:
 - (a) initialization of the .bss section to 0 (uninitialized variables) and initialization of the .data section by copying the initial values from FRAM (initialized variables)¹;
 - (b) clock setup by calling tpl_set_mcu_clock;
 - (c) initialization of the MPU;
 - (d) a call to main.
- 3. main is responsability of the user but usually it:
 - (a) initializes application level devices;
 - (b) calls StartOS.
- 4. StartOS:
 - (a) calls tpl_init_machine that calls;
 - i. tpl_init_machine_generic that calls:
 - A. tpl_init_mpu which is not implemented yet (and would be redundant with 2c).
 - ii. tpl_init_machine_specific that calls:
 - A. tpl_set_systick_timer.
 - (b) calls tpl_start_os that does a system call to tpl_start_os_service that.
 - i. calls tpl_init_os;
 - ii. calls tpl_enable_counters;
 - iii. calls StartupHook if any;
 - iv. calls tpl_start_scheduling.

when returning a task is schedule.

¹Why not use the DMA to do that, it would use less energy

3 Modified sequence to restore a checkpoint

Following items shall be modified when restarting from a checkpoint:

• item 2a should be replaced by a copy from the checkpoint data in FRAM to SRAM.

At this date, the .bss, .data and the constructor call to c++ global objects is done before the checkpoiting process. This is because some data should be initialized during the RestartOS() service call. Maybe a new section in the link script should be defined, not to initialize the SRAM twice in case of restart.

- item 2b should use the replaced by an init with the clock frequency when the checkpoint was done. This can be done by having a variable (in .data segment) to store the clock frequency so that it would restored as part of the checkpoint data.
- item 2d should be replaced by a call to a mandatory function used to initialize the devices for the application (UART for instance) and to a new service, let's call it RestartOS. RestartOS would:
 - 1. Call tpl_init_machine;
 - 2. Call tpl_restart_os that does a system call to tpl_restart_os_service that does a tpl_start. tpl_start moves the highest priority task from the ready list to the elected slot of tpl_kern. Conditions shall be NEED_SWITCH true and NEED_SAVE false.

When returning from the RestartOS service, the highest priority task is scheduled and the system continues execution.

A boolean variable stored in FRAM, let's call it tpl_checkpoint_available, shall be used to select, when true, the modified sequence instead of the normal one.

4 Checkpointing

A new service is necessary, let's call it Hibernate. When called, Hibernate, terminates the caller. It copies the SRAM to the FRAM (checkpoint), stops the Systick, stops application interrupts (a user function shall be provided for that), programs the RTC to emit an interrupt every x seconds, enables interrupts and goes in LPM3 (Lowest power mode that preserves the RAM).

The RTC ISR checks the voltage of the MCU. If it is above a threshold (RESUME_FROM_-HIBERNATE_THRESHOLD), it exists from LPM3. If this never happens, after a while, the MCU will power off. When in the future the MCU restart, what is described in section 3 applies.

If Hibernate resumes because the RTC ISR exited from LPM3 then it disables interrupt, starts application interrupts, starts the Systick, stop the RTC. It calls tpl_start to elect the highest priority task and returns.

In addition, a periodic basic task, energy_task checks (every 5 seconds? more?)

```
MB: should be defined in the .oil, as it depends on both supercapacitor value, and application through external peripherals

JLB: you are right
```

the voltage. If the voltage drops below a threshold (HIBERNATE_THRESHOLD), energy_task calls Hibernate (and terminates):

```
TASK(energy_task)
{
  uint16 voltage = readPowerVoltage();
  if (voltage < HIBERNATE_THRESHOLD) {
    Hibernate();
  }
  else {
    TerminateTask();
  }
}</pre>
```

HIBERNATE_THRESHOLD shall be chosen so that the worst voltage drop between 2 executions of energy_task plus the voltage drop due to checkpointing is lower than the threshold.

HIBERNATE_THRESHOLD shall be lower than RESUME_FROM_HIBERNATE_THRESHOLD.

4.1 OIL modifications

An OS attribute is added:

```
BOOLEAN CHECKPOINT = FALSE;
```

If TRUE, some complementary information can be given:

```
BOOLEAN CHECKPOINT = TRUE {
   HIBERNATE_THRESHOLD = 1900; //in mV
   ENERGY_TASK_PERIOD = 5000; //in ms
};
```

4.2 Non volatile data declaration

Non Volatile data is stored at the beginning of the FRAM. The MPU is configured at startup to let a section in R/W mode. However, limitations on the MPU hardware implementation require to use a memory area that is a multiple of 1024 bytes. The code section just follows, with R/X rights.

An OS variable in non volatile memory should be defined as:

```
#define OS_START_SEC_NON_VOLATILE_VAR_16BIT
#include "tpl_memmap.h"
VAR(uint16, OS_VAR) nvdata = 0;
#define OS_STOP_SEC_NON_VOLATILE_VAR_16BIT
```

```
#include "tpl_memmap.h"
```

for a task related non volatile variable the definition is (task is serial_TX here:

```
#define APP_Task_serial_TX_START_SEC_VAR_NON_VOLATILE_32BIT
#include "tpl_memmap.h"
VAR(uint32_t,AUTOMATIC) dataFRAM = 100;
#define APP_Task_serial_TX_STOP_SEC_VAR_NON_VOLATILE_32BIT
#include "tpl_memmap.h"
```

Note:

At this date, the non volatile data is initialized only when flashing the firmware. It should be updated so that the value can be updated on a non-restoring startup (no checkpoint).

4.3 DMA

Both checkpoint save and checkpoint restore are done using the DMA. The MSP430FR5994 has 6 DMA channels. Channel 0 is used. The block transfer mode is used. DMA is triggered by software. It copies the content of the SRAM which is used by the software (4kB at most, from 0x1C00 to 0x2BFF) to the FRAM. It would be safer to use double buffering because if a checkpoint cannot complete before the loss of power for any reason, the previous checkpoint will not be corrupted. An int variable in FRAM with value equal to 0, 1 or -1 is used to select the buffer. Let's call it tpl_checkpoint_buffer. If the value is -1, no checkpoint is saved. This is the initial value. If value is 0 a checkpoint is saved in buffer 0 and if value is 1, a checkpoint is saved in buffer 1. Saving a checkpoint would be as follow:

```
int buffer = (tpl_checkpoint_buffer + 1) & 1;
tpl_save_checkpoint(buffer);
tpl_checkpoint_buffer = buffer;
```

If power is lost during the tpl_save_checkpoint function, tpl_checkpoint_buffer is not updated and the checkpoint is not committed.

The content of tpl_checkpoint_buffer is used to select the normal startup described at section ?? or the startup from a checkpoint described at section 3.

```
#define OS_START_SEC_CONST_16BIT
#include "tpl_memmap.h"
static CONST(sint16, OS_CONST) NO_CHECKPOINT = -1;
#define OS_STOP_SEC_CONST_16BIT
#include "tpl_memmap.h"

#define OS_START_SEC_NON_VOLATILE_VAR_16BIT
#include "tpl_memmap.h"
VAR(sint16, OS_VAR) tpl_checkpoint_buffer = -1;
#define OS_STOP_SEC_NON_VOLATILE_VAR_16BIT
```

```
#include "tpl_memmap.h"

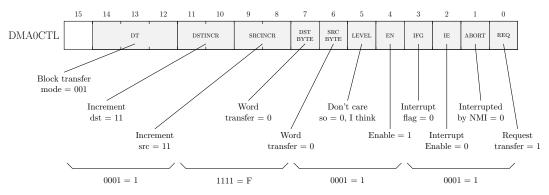
if (tpl_checkpoint_buffer == NO_CHECKPOINT) {
   /* normal startup */
}
else {
   /* startup from checkpoint */
}
```

The 2 buffers are located above the 64kB limit. The first one from 0x010000 to 0x010FFFh and the second one from 0x011000 to 0x011FFFh.

The SRAM starts at address 0x001C00 and ends at address 0x003BFF but we only use the first 4kB, from 0x001C00 to 0x002BFF.

The function to save the SRAM to a buffer using the DMA has to be written in assembly because writing to start address and end address registers shall be done using extended instructions when address is above the 64kB limit. See file framtr.S in examples/msp430x/small/msp430fr5994/launchpad/checkpoint2.

DMAOCTL programming explanation



5 New services

5.1 Service RestartOS

5.2 Service Hibernate

6 Peripherals

Specific code for peripherals re-initialization. TODO They have to be re-initialized before main call.