

# A solution to “easycrackme” by “bexplode”

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## 1 Introduction

On the `crackmes.one` website, user “bexplode” published an easy, high quality,<sup>1</sup> “crack me” toy problem [1]. End-users are encouraged to work out a solution and submit their findings without the use of binary patching.<sup>2</sup> Here, the toy problem comes in the form of a binary executable where its respective solution comes in the form of a “key” that will cause the application to output a message affirming that a correct key has been worked out.

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Email: `hackermaneia@riseup.net`

`crackmes.one` Account: <https://crackmes.one/user/RelationalAlgebra>

GitHub Portfolio: <https://github.com/Alekseyyy>

<sup>1</sup>This “crack me” has a difficulty score of 1.4/6, and a quality score of 5/6.

<sup>2</sup>See the *crackmes.one* FAQ.

Before trying to find a solution, I should begin by setting up the problem. I proceeded by downloading the crack me ZIP archive and extracting the files.<sup>3</sup> The relevant file in the archive is `bxtumations_crackme.exe`. I ran the software executable, which gave me the prompt depicted in figure 1:

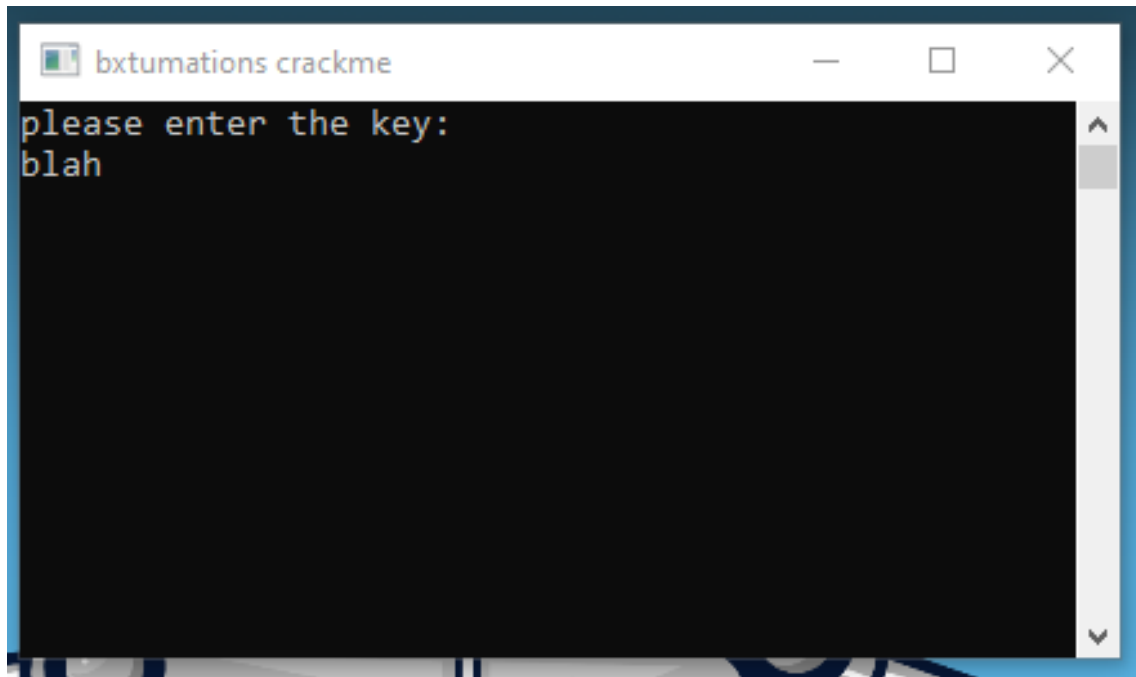


Figure 1: Prompt given when running `bxtumations_crackme.exe`

I typed in “blah” followed by the `enter`-key, and was presented with a message informing me that my solution was incorrect (figure 2). Hereforth, the problem defined is to work out a string that, when inputted into the application, will result in a message that is not “wrong key!!” and “try again :)”.

## 2 Solution: Static and Dynamic Analysis

I will go about working out the proper key with the methods of static and dynamic analysis.<sup>4</sup> This section will outline how I went about deriving the solution.

### 2.1 Initial Analysis of the Software System

The first step that I took is to load `bxtumations_crackme.exe` into a number of software reverse engineering tools, such as *Ghidra* [2], *IDA Pro*<sup>5</sup> [3], *Detect It Easy* [4], and *x64dbg* [5]. `crackmes.one` listed this binary as being written in C/C++ and written for Microsoft Windows. The *Detect It Easy* results (figure 3) confirms this.

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<sup>3</sup>The password for the ZIP archive is “crackmes.one”

<sup>4</sup>In a future writeup, I intend to discuss an alternative solution.

<sup>5</sup>Free Edition

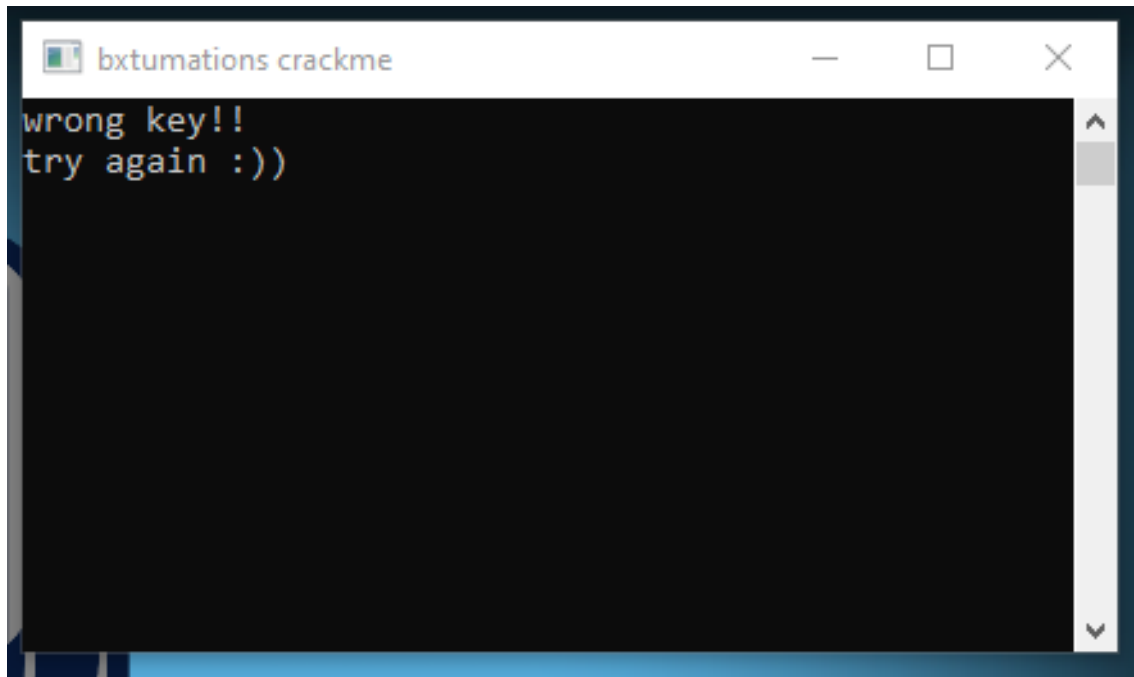


Figure 2: Error message stating that `blah` is not the correct key

When loading `bxtumations_crackme.exe` into *Ghidra*, I made sure that the format was set to “Portable Executable (PE)” and used all of *Ghidra*’s “analysis” options on it.<sup>6</sup> A cursory look at the number of calls that the executable makes tells me that trying to trace its execution flow by manually graphing it would take too long. So, I instead decided to try to identify the part of the executable where the key is worked out and stored in memory.

Using *Ghidra*’s export features, I had its decompiler save a high-level C/C++ representation of the binary executable onto disk.<sup>7</sup> I then used a standard text editor to search for the string `please enter the key` — where I discovered that the key is worked out in the `FUN_00401789` function.

I decided to give a cursory look at `FUN_00401789`, and the following is a snippet of the relevant source code:

```
undefined8 FUN_00401789(void){

    undefined8 uVar1;

    [... snip ...]

    while( true ) {
        FUN_004a0e90(&DAT_004a6860,local_c8);
        ppvVar3 = local_108;
        pplVar2 = local_c8;
```

---

<sup>6</sup>Including the beta/experimental analysers.

<sup>7</sup>This is done by going to **File > Export Program**, setting the “format” to C/C++, and specifying a location to save the decompile dump.

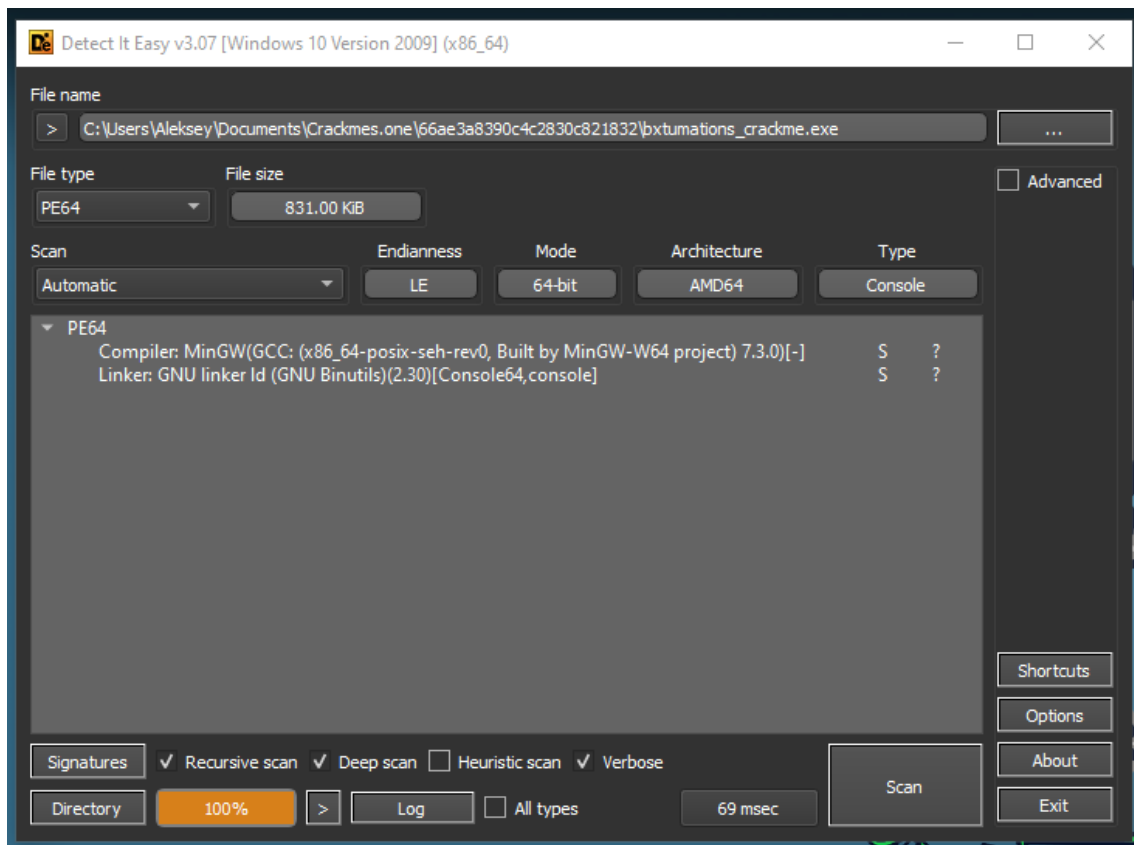


Figure 3: *Detect-It-Easy* Output

```

uVar1 = FUN_0049fbe0(pp1Var2,ppvVar3);
if ((char)uVar1 != '\0') break;
FUN_00401560();
FUN_00460b70();

[... snip ...]

}
DAT_004d1030 = 1;
FUN_0040157b(pp1Var2,ppvVar3,ppiVar4);
FUN_00491060(local_108);
FUN_00491060(local_e8);
FUN_00491060(local_c8);
return 0;
}

```

I removed a lot of what I see to be irrelevant information, but the reader may consult the “supplementary materials” and the Appendix<sup>8</sup> if they want the full source code. This relevant bits of the function works as follows:

<sup>8</sup>Under the section title “C/C++ Decompile Dumps”

1. First, a variable called `uVar1` of type “undefined8” is declared.<sup>9</sup> Other variables are declared, but I am not interested in them.
2. Next, other functions are called, and then the program goes to an infinite `while`-loop.
3. In this `while`-loop, the `uVar1` is set to the results of `FUN_0049fbe0(pp1Var2,ppvVar3)`, and then it is used in a conditional.
4. Regarding the conditional, the `uVar1` is casted to the `char` type — `(char)uVar1` — and then compared to the character `\0`.
  - (a) If `uVar1` is equal to `\0`, then the program will **break** out of the `while`-loop.
  - (b) Otherwise, the program will display a “wrong key” error message, and not break out of the `while`-loop.
5. Assuming that the program breaks out of the `while`-loop, a message will be printed onto the screen congratulating the end-user for working out the correct key.

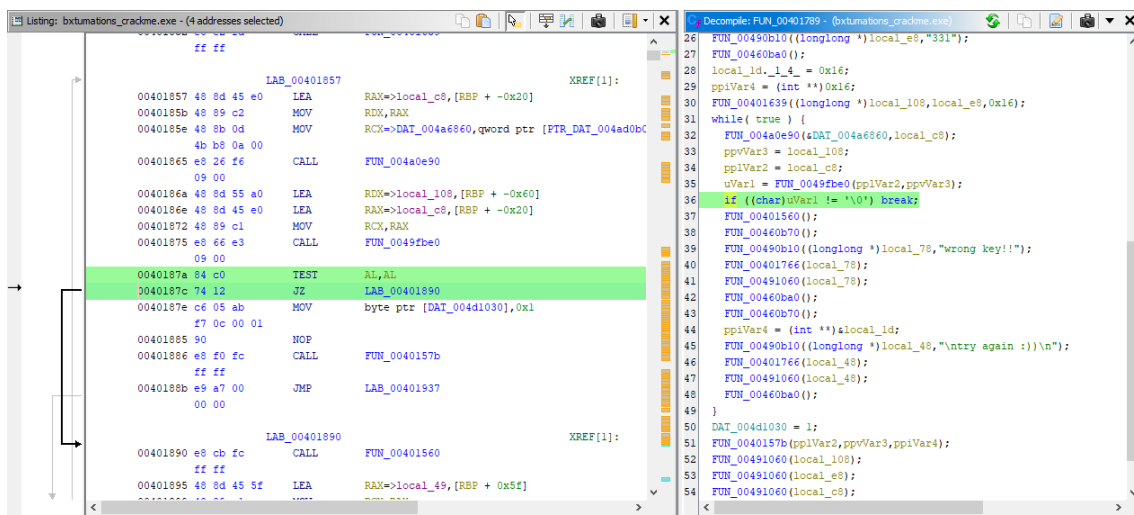


Figure 4: A comparison of `FUN_00401789`’s C/C++ representation to its disassembly

I focused on the disassembly of the conditional. On *Ghidra*, I compared the C/C++ decompile dump of the binary to its respective disassembly — which is depicted by figure 4 depicts this. In particular, the `if ((char)uVar1 != '\0') break;` corresponds to the following assembler instructions:

1. `TEST AL,AL`
2. `JZ LAB_00401890`

<sup>9</sup>I think that this might be an `unsigned int`, though I was told that it could be some 8-bit type.

Line 1 has an address value of 0040187a and line 2 has an address value of 0040187c. This will become relevant as I move on to dynamic analysis with the *x64dbg* debugger.

## 2.2 Breakpoints and Debugger Setup

The software binary is too complicated for me to understand with just the methods of manually charting the program’s execution flow with a directed graph. The aforementioned method would take too long, and if I assume that “time is of the essence,” then I should find a quicker way to work out a solution. This “quicker way” involves the *x64dbg* debugger: after loading the `bxtumations_crackme.exe` into it, I was presented with four panes showing states and information regarding the software binary pre-execution, and during execution — as depicted by figure 5.

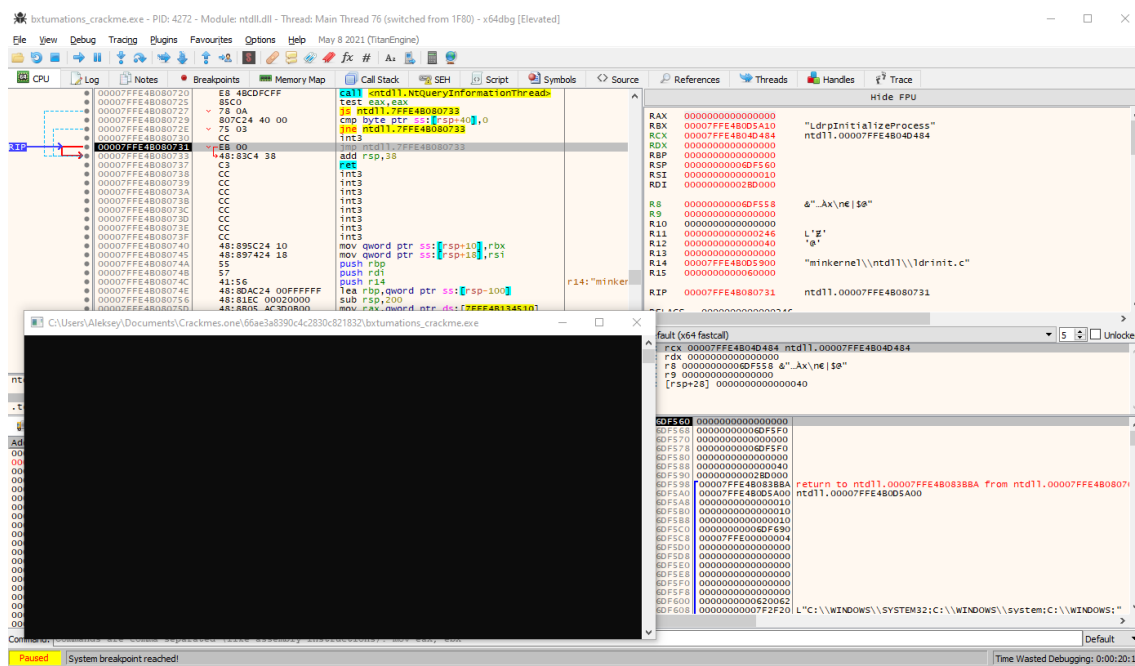


Figure 5: *x64dbg* view after loading `bxtumations_crackme.exe`

I proceeded by identifying the assembly instructions with the address values 0040187a and 0040187c, and proceeding to set breakpoints onto them — which is depicted by figure 6. I expect that, when I run the software under the debugger, it will run the necessary calculations and “self-decode” the key, and then load the key into memory, CPU registers, or other kinds of memory. I was right, and will discuss my findings in the next subsection.

## 2.3 Discovering the Key

After I have configured the breakpoints on the *x64dbg* debugger, I ran the executable, and it paused execution on the breakpoints. Figure 7 depicts what each of *x64dbg*’s four panes looked like after running until it reached the breakpoints.

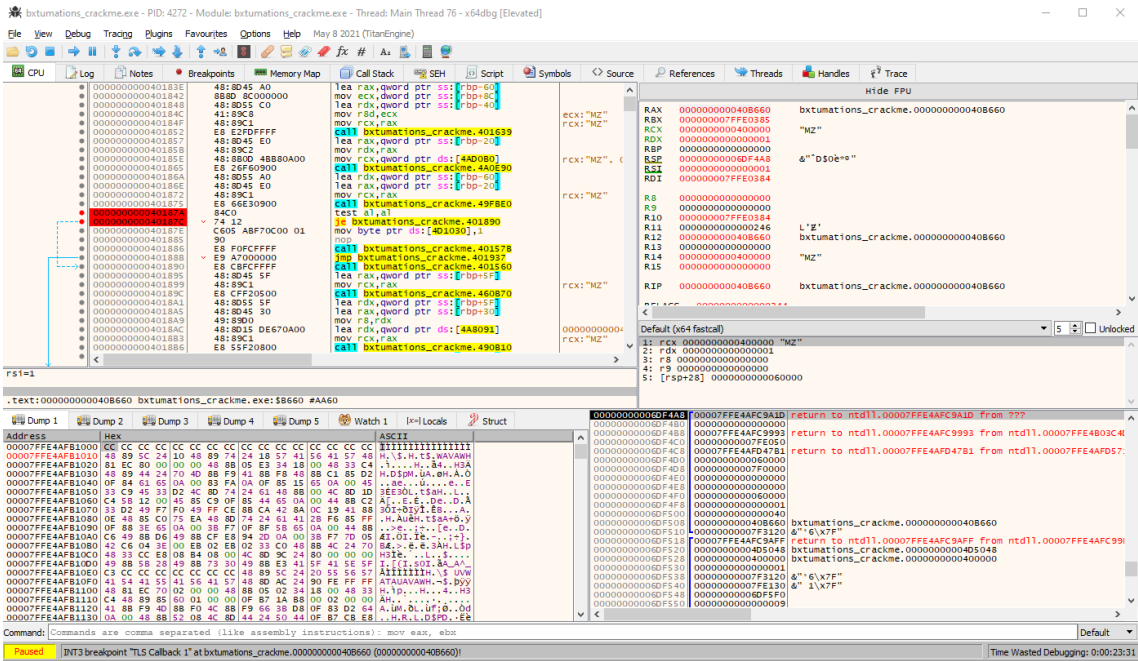


Figure 6: Setting breakpoints on 0040187a and 0040187c

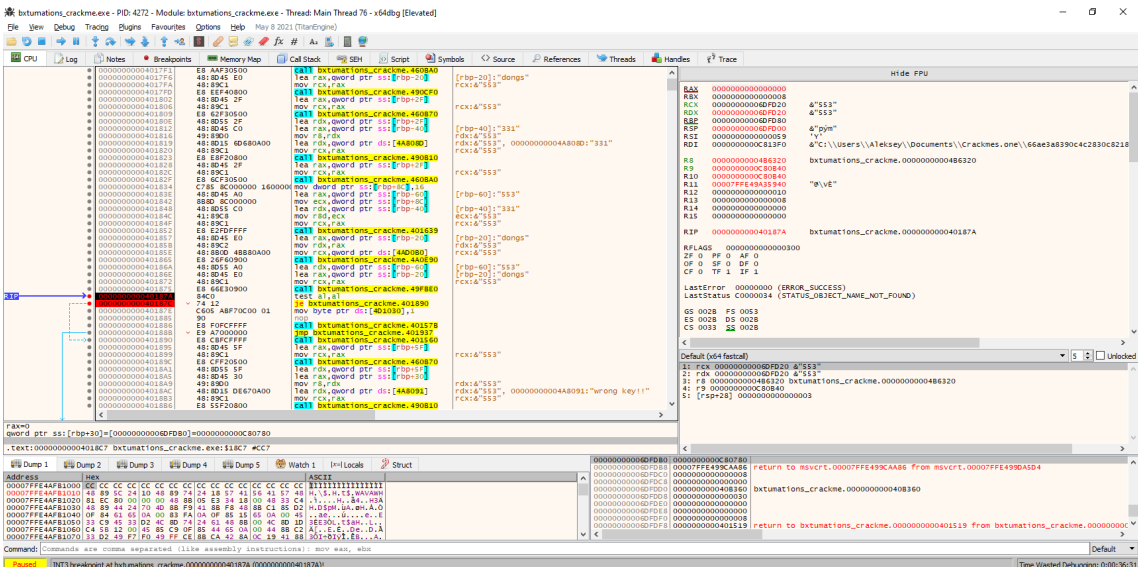


Figure 7: Execution state up until the set breakpoints (on 0040187a and 0040187c)

The following is some of what was reported on the top-right pane showing CPU register values:

```
RAX 0000000000000000
RBX 0000000000000000
RCX 0000000000000000 &"553"
RDX 0000000000000000 &"553"
RBP 0000000000000000
RSP 0000000000000000 &"pym"
```

[... snip ...]

Furthermore, the following was reported on the top-left pane showing the software binary's disassembly.

[... snip ...]

```
00401857    lea rax, qword ptr ss:[rbp-20] ; [rbp-20]:"dongs"
0040185B    mov rdx, rax ; rdx:&"553"
00401853    mov rcx, qword ptr ds:[4AD0B0] ; rcx:&"553"
00401865    call bxtumations_crackme.4A0E90
0040186A    lea rdx, qword ptr ss:[rbp-60] ; [rbp-50]:"553"
0040186E    lea rax, qword ptr ss:[rbp-20] ; [rbp-20]:"dongs"
00401872    mov rcx, rax ; rcx:"553"
00401875    call bxtumations_crackme.49FBE0
0040187A    test al, al
0040187C    je bxtumations_crackme.401890
```

[... snip ...]

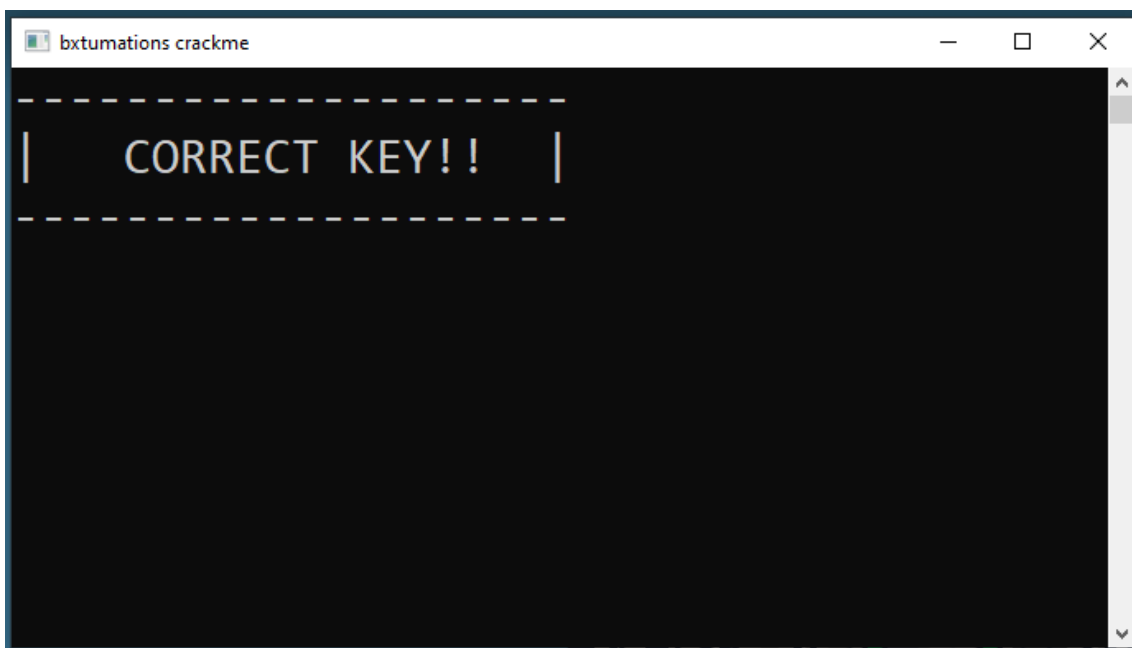


Figure 8: Message when 553 is entered as the key.

From an intuitive look at both the registers, stack values, and disassembler, I guessed that “553” is the correct key. I ran `bxtumations_crackme.exe` without a debugger, typed in “553”, and was presented with a message stating that I have supplied the “correct key” — as depicted by figure 8.<sup>10</sup>

---

<sup>10</sup>Unlike with the first two figures, I made the text bigger by changing the font settings in the command prompt.



While I initially worked out the key with a combination of intuition, experimentation, and luck, I do think that I should try to analyse *why* this is the correct key. Looking back at the CPU registers, I noticed that the number 553, represented as a string, has its address referenced into the RCX and RDX registers. I also noticed the following instructions<sup>11</sup> of the CPU assembler dump shown earlier:

[... snip ...]

```
05. 0040186A    lea rdx, qword ptr ss:[rbp-60] ; [rbp-50]:"553"
06. 0040186E    lea rax, qword ptr ss:[rbp-20] ; [rbp-20]:"dongs"
07. 00401872    mov rcx, rax ; rcx:"553"
08. 00401875    call bxtumations_crackme.49FBEO
09. 0040187A    test al, al
10. 0040187C    je bxtumations_crackme.401890
```

The fifth line<sup>12</sup> references “553”’s address in the RDX register, and then the sixth line<sup>13</sup> references the `dongs` string’s address into the RAX register. This RAX address is then copied in the RCX register,<sup>14</sup> and then a function is called.<sup>15</sup> I am not quite sure what that function does, but the `test` instruction is executed in the ninth line,<sup>16</sup> and the *jump if equal* instruction is executed on the tenth line.<sup>17</sup> This tells me that the string 553 is being compared to the input, and serves as a useful hint for guessing how it could affect the trajectory of the software binary’s execution path.

### 3 End Matter

By the procedure outlined in the previous section, I can confidently say that “553” is the correct key. The following are “takeaways” that I have learnt from doing this short project:

- Given that certain assumptions are made,<sup>18</sup> static analysis is useful for mapping out how a software binary’s system looks like.
- Regarding static analysis, a “easy-to-read” decompiler output can be compared to a more concrete assembler output to find clues when setting debugger breakpoints.

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<sup>11</sup>Lines 05-10, or addresses 0040186A to 0040187C.

<sup>12</sup>Ln. 05

<sup>13</sup>Ln. 06

<sup>14</sup>Ln. 07

<sup>15</sup>Ln. 08

<sup>16</sup>Ln. 09

<sup>17</sup>Ln. 10

<sup>18</sup>For example, in a criticism of a malware analysis report [6], Robert Graham states that reverse engineering tools make assumptions about a binary, and the reverse engineer or systems analyst must use intuition and good judgement for determining if assumptions are met.

- Software can be very complex,<sup>19</sup> even when it is the “easy” `crackmes.one` puzzle that is the subject of this writeup. Dynamic analysis done by a debugger can help reverse engineers “cut through the complexity” to solve a problem. In this case, I used a debugger to run the software binary and observe CPU registers and the software’s disassembly, which allowed me to work out the correct key.
- While this paper is not technically “original research,” it does share at most a few things in common with it: mainly that they both do not have a “fixed answer.”<sup>20</sup> Academic research, does not always give “clean” results, and sometimes guesswork and intuition is needed to interpret results and observation. But nonetheless, a well defined criteria for progress is needed for research to be successful.

### 3.1 Supplementary Materials

The project files can be accessed from the following GitHub repository:

<https://github.com/Alekseyyy/SNHU/tree/main/sundries/wargames/crackmes.one/writeups/66ae3a8390c4c2830c821832>

### 3.2 About the Author

At the time of this writing, I am a junior computer science undergraduate student, with a minor in mathematics and a concentration in data analysis. I currently do bug-bounty and responsible vulnerability disclosure. I also enjoy learning more about low-level aspects of computer hardware and assembler languages by solving toy-problems and through resources outside of school.<sup>21</sup>

### 3.3 Acknowledgements

I would like to take the time to acknowledge the helpful feedback from the following on earlier drafts of this paper:

- *Anonymous collaborator 1* for encouraging me to expand more on bulletpoints in the conclusion.
- *Anonymous collaborator 2* for informing me that an “undefined” data type might be an 8-bit data type (as opposed to specifically an `unsigned int`), for helping me better articulate how stack values get referenced in memory, and for helping me better articulate my points in general.

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<sup>19</sup>In a paper [7] by security researcher Greg Hoglund, he discusses how software that is made up of just simple rules can nonetheless produce complex behaviour, and how such complexity can lead to software bugs.

<sup>20</sup>I was informed by *Anonymous collaborators 2 and 3* that this writeup resembles more of a homework problem than a original research.

<sup>21</sup>Like by solving `crackmes.one` puzzles ;-)

- *Anonymous collaborator 3* for convincing me that the analogy between `crackmes.one` puzzles and original research has severe limitations.

I take full responsibility for any erratas in this writeup.

## A C/C++ Decompile Dumps

I included the following decompile dumps in this writeup and in the supplementary materials:

- `entry.c`: this is the entry point of the application, which executes the code blocks represented by C functions: `FUN_0040b450` and `FUN_00401180`.
- `FUN_0040b450.c`: this is one of the functions called by the entry point (I don't know exactly what it does).
- `FUN_00401180.c`: this is another one of the functions called by the entry point (I don't know exactly what it does).
- `FUN_00401789.c`: this is the function that I looked at when trying to work out what addresses to set breakpoints when debugging a the software binary. I documented a simplified version of it in the section where I looked at the decompile dump of the software binary.

### A.1 `entry.c`

```
void entry(undefined8 param_1,
           undefined8 param_2,undefined8 param_3)

{
    DAT_004d1610 = 0;
    FUN_0040b450();
    FUN_00401180(param_1,param_2,param_3);
    return;
}
```

### A.2 `FUN_0040b450.c`

```
void FUN_0040b450(void)

{
    _FILETIME _Var1;
    DWORD DVar2;
    DWORD DVar3;
    DWORD DVar4;
    _FILETIME local_38;
```

```

LARGE_INTEGER local_30;

local_38.dwLowDateTime = 0;
local_38.dwHighDateTime = 0;
if (DAT_004a7200 != 0x2b992ddfa232) {
    DAT_004a7210 = ~DAT_004a7200;
    return;
}
GetSystemTimeAsFileTime(&local_38);
_Var1 = local_38;
DVar2 = GetCurrentProcessId();
DVar3 = GetCurrentThreadId();
DVar4 = GetTickCount();
QueryPerformanceCounter(&local_30);
DAT_004a7200 = ((ulonglong)DVar4 ^
    (ulonglong)DVar3 ^ (ulonglong)DVar2 ^ (ulonglong)_Var1
    ^ local_30.QuadPart) & 0xffffffffffff;
if (DAT_004a7200 == 0x2b992ddfa232) {
    DAT_004a7210 = 0xffffd466d2205dcc;
    DAT_004a7200 = 0x2b992ddfa233;
}
else {
    DAT_004a7210 = ~DAT_004a7200;
}
return;
}

```

### A.3 FUN\_00401180.c

```

/* WARNING: Globals starting with '_'
    overlap smaller symbols at the same address */

ulonglong FUN_00401180(undefined8 param_1,
    undefined8 param_2,undefined8 param_3)

{
    int iVar1;
    void **ppvVar2;
    char cVar3;
    ulonglong uVar4;
    ulonglong uVar5;
    undefined8 *puVar6;
    int iVar7;
    char **ppcVar8;
    char *pcVar9;

```

```

undefined8 *puVar10;
size_t sVar11;
void *_Dst;
undefined8 *puVar12;
ulonglong uVar13;
longlong lVar14;
undefined8 uVar15;
undefined8 uVar16;
LPSTARTUPINFOA p_Var17;
undefined8 uVar18;
longlong unaff_GS_OFFSET;
bool bVar19;
undefined local_a8 [64];
ushort local_68;

p_Var17 = (LPSTARTUPINFOA)local_a8;
for (lVar14 = 0xd; lVar14 != 0; lVar14 = lVar14 + -1) {
    *(undefined8 *)p_Var17 = 0;
    p_Var17 = (LPSTARTUPINFOA)&p_Var17->lpReserved;
}
uVar13 = (ulonglong)DAT_004d1610;
if (DAT_004d1610 != 0) {
    GetStartupInfoA((LPSTARTUPINFOA)local_a8);
}
uVar4 = *(ulonglong *) (*(longlong *) (unaff_GS_OFFSET
    + 0x30) + 8);
while( true ) {
    LOCK();
    bVar19 = DAT_004d2420 == 0;
    DAT_004d2420 = DAT_004d2420 ^ (ulonglong)bVar19 *
        (DAT_004d2420 ^ uVar4);
    uVar5 = !bVar19 * DAT_004d2420;
    UNLOCK();
    if (uVar5 == 0) break;
    if (uVar4 == uVar5) {
        bVar19 = true;
        goto joined_r0x004011ff;
    }
    Sleep(1000);
}
bVar19 = false;
joined_r0x004011ff:
if (DAT_004d2428 == 1) {
    _amsg_exit(0x1f);
}
else if (DAT_004d2428 == 0) {

```

```

    DAT_004d2428 = 1;
    _initterm();
}
else {
    DAT_004d1004 = 1;
}
if (DAT_004d2428 == 1) {
    _initterm();
    DAT_004d2428 = 2;
}
if (!bVar19) {
    LOCK();
    DAT_004d2420 = 0;
    UNLOCK();
}
uVar18 = 0;
uVar16 = 2;
uVar15 = 0;
tls_callback_0(0,2);
FUN_0040ba50(uVar15,uVar16,uVar18,uVar13);
DAT_004d1640 = SetUnhandledExceptionFilter(
    (LPTOP_LEVEL_EXCEPTION_FILTER)&LAB_0040bfa0);
FUN_0040beb0();
FUN_004175f0(&LAB_00401000);
FUN_0040b850();
_DAT_004d2410 = &IMAGE_DOS_HEADER_00400000;
ppcVar8 = (char **)FUN_004176d0();
iVar7 = DAT_004d1020;
bVar19 = false;
pcVar9 = *ppcVar8;
if (pcVar9 != (char *)0x0) {
    do {
        cVar3 = *pcVar9;
        if (cVar3 < '!') {
            if ((cVar3 == '\0') || (!bVar19)) goto LAB_004012d0;
            bVar19 = true;
        }
        else if (cVar3 == '\"') {
            bVar19 = (bool)(bVar19 ^ 1);
        }
        pcVar9 = pcVar9 + 1;
    } while( true );
}
goto LAB_004012f7;
LAB_004012d0:
_DAT_004d2418 = pcVar9;

```

```

    if (cVar3 != '\0') {
        do {
            pcVar9 = pcVar9 + 1;
            _DAT_004d2418 = pcVar9;
            if (*pcVar9 == '\0') break;
        } while (*pcVar9 < '!');
    }
LAB_004012f7:
    if ((DAT_004d1610 != 0) &&
        (_DAT_004a4000 = 10, (local_a8[60] & 1) != 0)) {
        _DAT_004a4000 = (uint)local_68;
    }
    iVar1 = DAT_004d1020 + 1;
    puVar10 = (undefined8 *)malloc((longlong)iVar1 * 8);
    puVar6 = DAT_004d1018;
    puVar12 = puVar10;
    if (0 < iVar7) {
        lVar14 = 0;
        do {
            sVar11 = strlen(*(char **)((longlong)puVar6 + lVar14));
            _Dst = malloc(sVar11 + 1);
            *(void **)((longlong)puVar10 + lVar14) = _Dst;
            ppvVar2 = (void **)((longlong)puVar6 + lVar14);
            lVar14 = lVar14 + 8;
            memcpy(_Dst, *ppvVar2, sVar11 + 1);
        } while ((ulonglong)(iVar7 - 1) * 8 + 8 != lVar14);
        puVar12 = puVar10 + (longlong)iVar1 + -1;
    }
    *puVar12 = 0;
    DAT_004d1018 = puVar10;
    FUN_0040b410();
    *(undefined8 *)__initenv_exref = DAT_004d1010;
    uVar13 = FUN_00401789();
    DAT_004d100c = (uint)uVar13;
    if (DAT_004d1008 != 0) {
        if (DAT_004d1004 == 0) {
            _cexit();
            uVar13 = (ulonglong)DAT_004d100c;
        }
        return uVar13;
    }
    /* WARNING: Subroutine does not return */
    exit(DAT_004d100c);
}

```

## A.4 FUN\_00401789.c

```
undefined8 FUN_00401789(void)

{
    undefined8 uVar1;
    longlong **pplVar2;
    void **ppvVar3;
    int **ppiVar4;
    void *local_108 [4];
    void *local_e8 [4];
    longlong *local_c8 [4];
    void *local_a8 [6];
    void *local_78 [6];
    void *local_48 [5];
    undefined8 local_1d;

    FUN_0040b410();
    SetConsoleTitleA("bxtumations crackme");
    FUN_00460b70();
    FUN_00490b10((longlong *)local_a8,"please enter the key:\n");
    FUN_00401766(local_a8);
    FUN_00491060(local_a8);
    FUN_00460ba0();
    FUN_00490cf0((longlong *)local_c8);
    FUN_00460b70();
    FUN_00490b10((longlong *)local_e8,"331");
    FUN_00460ba0();
    local_1d._1_4_ = 0x16;
    ppiVar4 = (int **)0x16;
    FUN_00401639((longlong *)local_108,local_e8,0x16);
    while( true ) {
        FUN_004a0e90(&DAT_004a6860,local_c8);
        ppvVar3 = local_108;
        pplVar2 = local_c8;
        uVar1 = FUN_0049fbe0(pplVar2,ppvVar3);
        if ((char)uVar1 != '\0') break;
        FUN_00401560();
        FUN_00460b70();
        FUN_00490b10((longlong *)local_78,"wrong key!!");
        FUN_00401766(local_78);
        FUN_00491060(local_78);
        FUN_00460ba0();
        FUN_00460b70();
        ppiVar4 = (int *)&local_1d;
        FUN_00490b10((longlong *)local_48,"\ntry again :))\n");
        FUN_00401766(local_48);
    }
```



```

    FUN_00491060(local_48);
    FUN_00460ba0();
}
DAT_004d1030 = 1;
FUN_0040157b(pp1Var2,ppvVar3,ppiVar4);
FUN_00491060(local_108);
FUN_00491060(local_e8);
FUN_00491060(local_c8);
return 0;
}

```

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