## Assigment 10

#### 1(a)

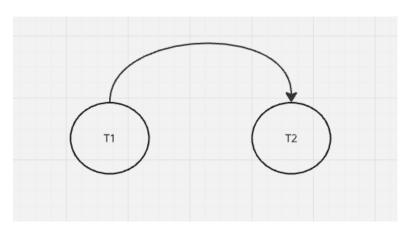
 $S_1 = \langle r_1(a), r_1(b), w_1(c), r_2(c), w_2(b), w_2(a) \rangle$  conflict pairs:

 $r_1(a) \nleftrightarrow w_2(a), T_1 \to T_2$ 

 $r_1(b) \leftrightarrow w_2(b), T_1 \to T_2$ 

 $w_1(c) \leftrightarrow r_2(c), T_1 \to T_2$ 

 $T_1 \to T_2$  not a circle, then conflict serializable



#### 1(b)

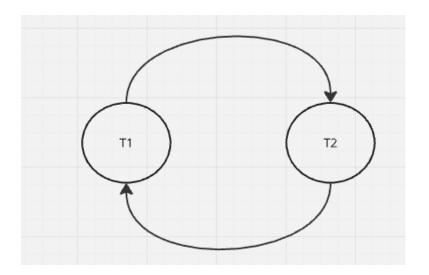
 $S_2 = \langle r_1(a), r_2(c), r_1(b), w_1(c), r_2(c), w_2(b), w_2(a) \rangle$  conflict pairs:

 $r_1(a) \nleftrightarrow w_2(a), T_1 \to T_2$ 

 $r_2(c) \nleftrightarrow w_1(c), T_2 \to T_1$ 

 $r_1(b) \leftrightarrow w_2(b), T_1 \to T_2$ 

 $T_1 \to T_2$  circle, then non-conflict serializable



1(c)

20

1(d)

1

1(e)

720

**2**(a)

 $S_b = \langle r_1(y), w_1(y), r_2(z), w_2(z), a, r_2(y), w_2(y) \rangle$ 

2(b)

Because it allows to read data only after it was committed

2(c). i.

Undo $(T_5)$ , Undo $(T_6)$ Redo $(T_1)$ , Redo $(T_2)$ , Redo $(T_3)$ , Redo $(T_4)$ , Redo $(T_7)$ Restart  $T_5, T_6$ 

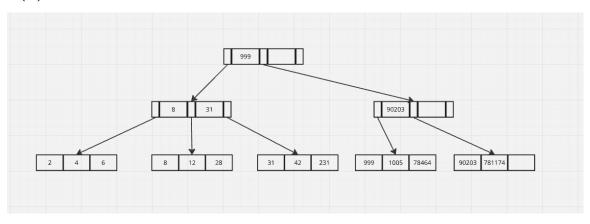
### 2(c). ii.

Undo $(T_5)$ , Undo $(T_6)$ Redo $(T_3)$ , Redo $(T_4)$ , Redo $(T_7)$ Restart  $T_5, T_6$ 

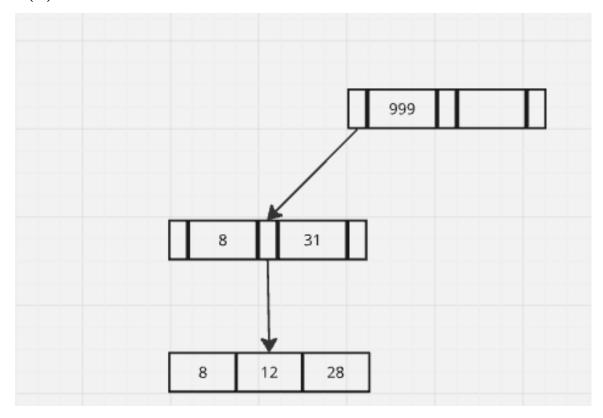
#### 2(d)

Checkpoints make amount of Redo's smaller. Log protocol is used to do Undo's backwards.

### 3(a)



# (b)



# **3(c)**

