

Applied Security

DPA on AES (8)

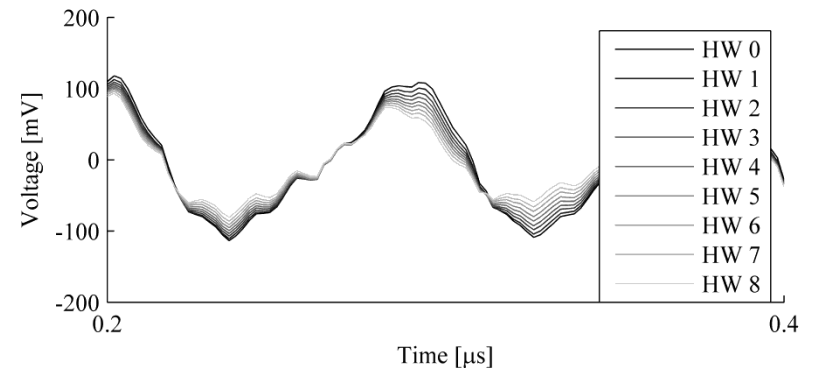
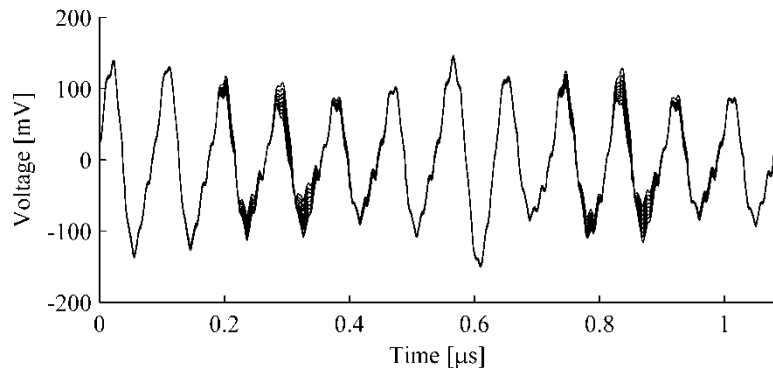
General Overview

- We go through a DPA attack on AES step by step and think about the impact of choices we have in each step
- This will allow you to complete part of the coursework

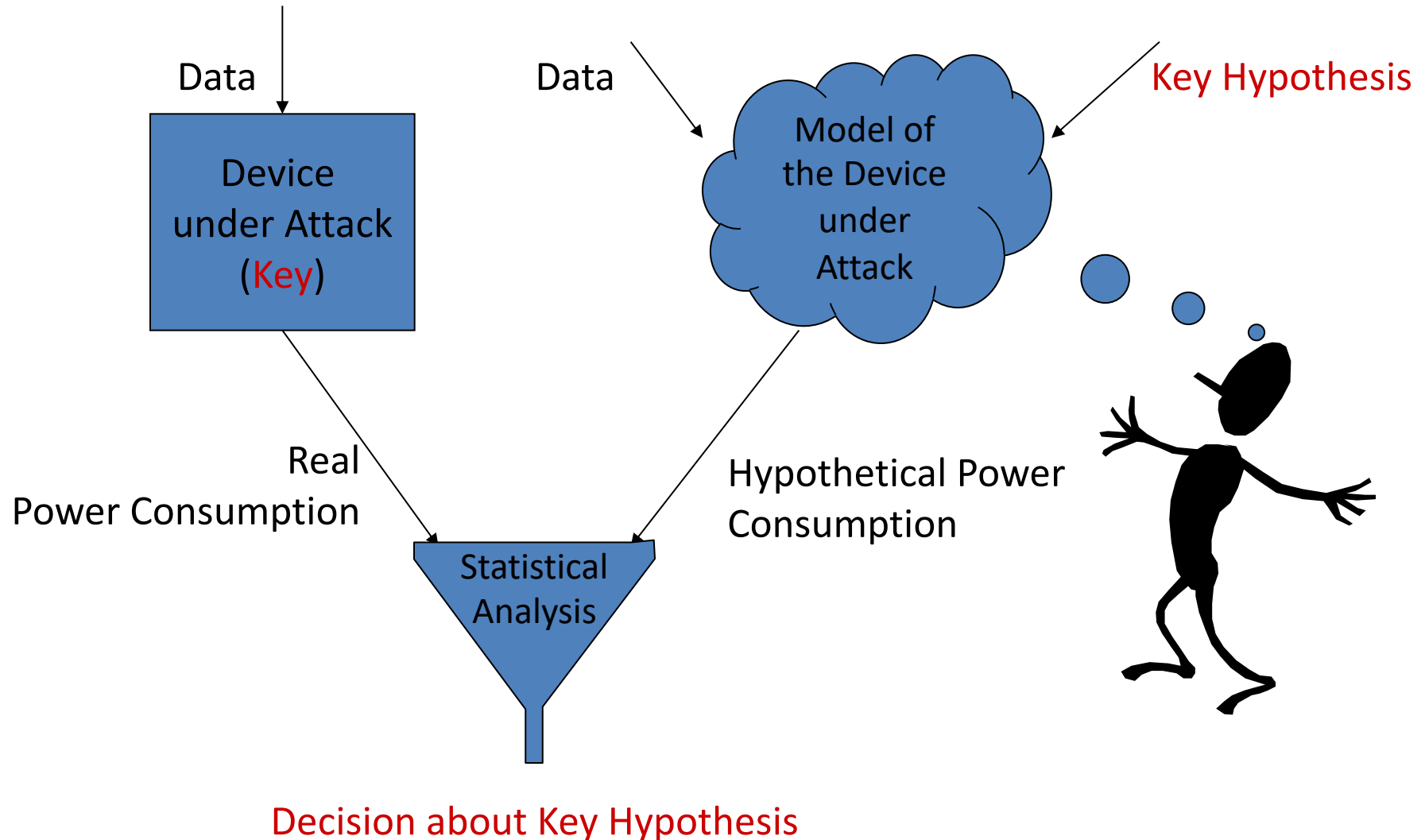
Beware: in order to follow the lectures you NEED to be familiar with various cryptographic algorithms and implementation techniques!

SPA vs DPA

- Typically SPA attacks exploit leakage that depends on the type of operation
- DPA attacks exploit the data dependent leakage:
 - Example below shows MOV instructions on an 8-bit: the shape is the same but the height of the curves in some of the clock cycles depends on the data

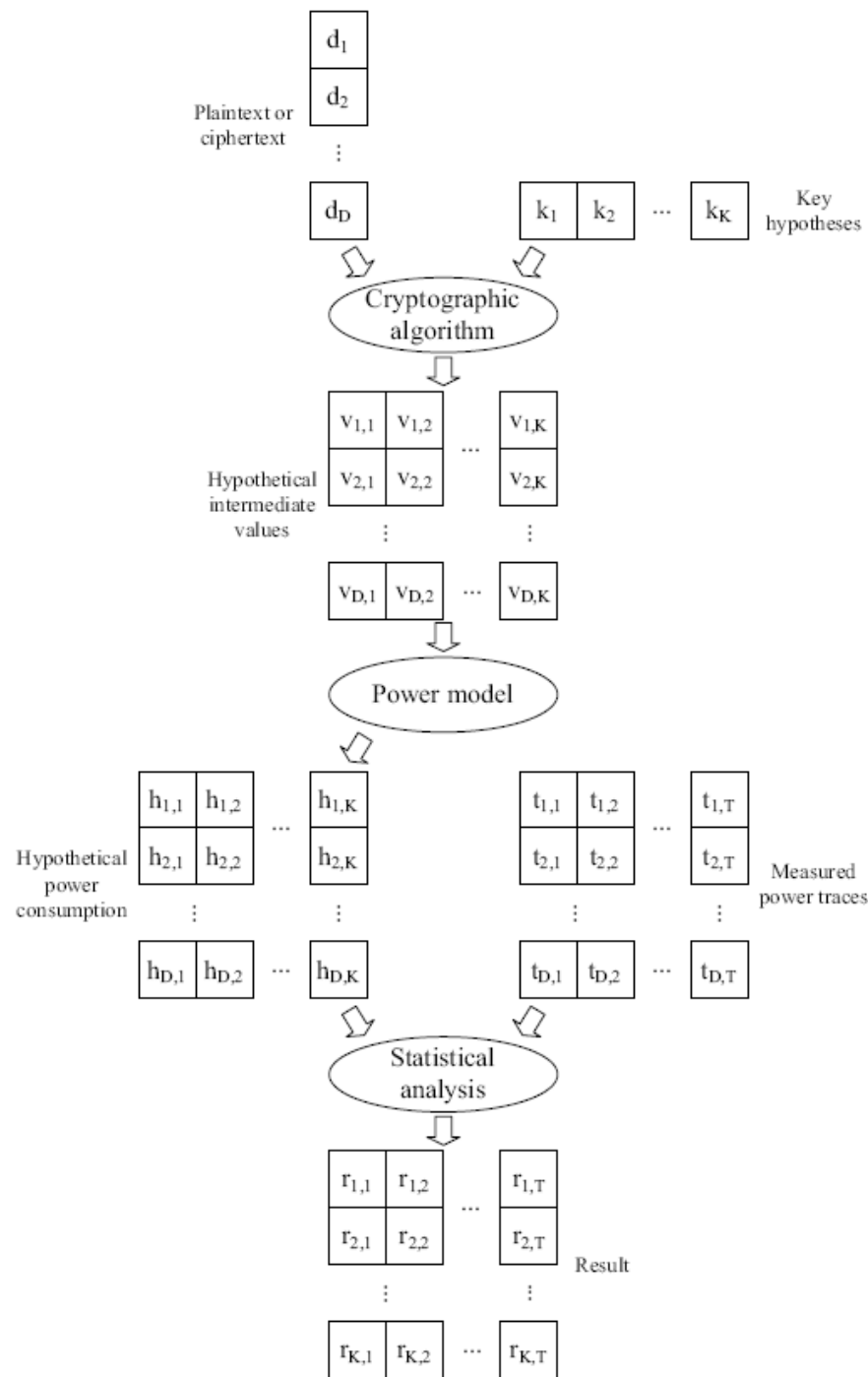


Principle of a differential power attack



5 Step Model of a DPA Attack

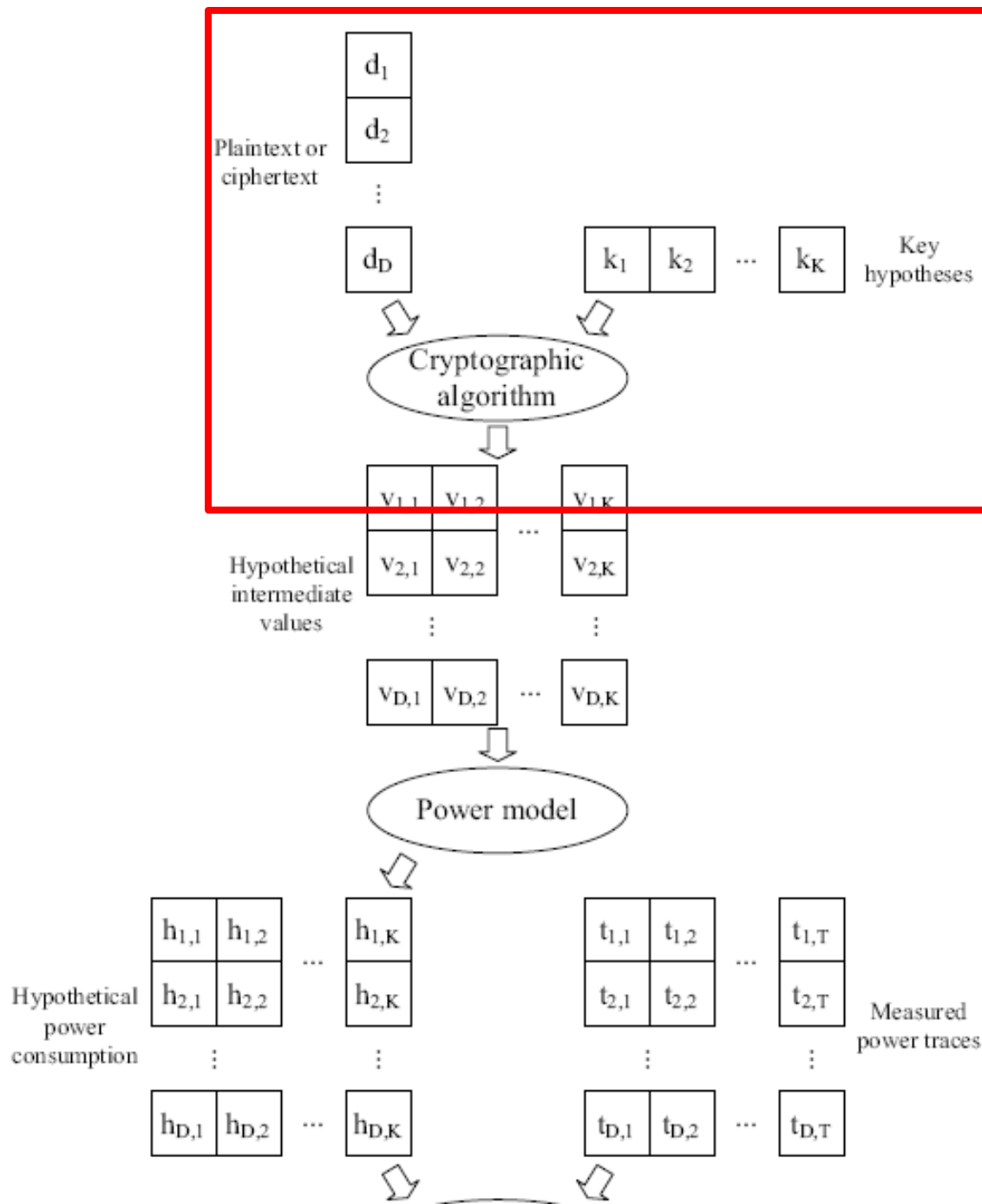
- 1) Selection of target intermediate variable
- 2) Measurement of the power consumption/EM
- 3) Calculation of the hypothetical intermediate values
- 4) Mapping of hypothetical values to predicted power consumption values
- 5) Statistical analysis using a distinguisher



5 Steps in detail

Step 1: Choosing an intermediate result of the executed cryptographic algorithm

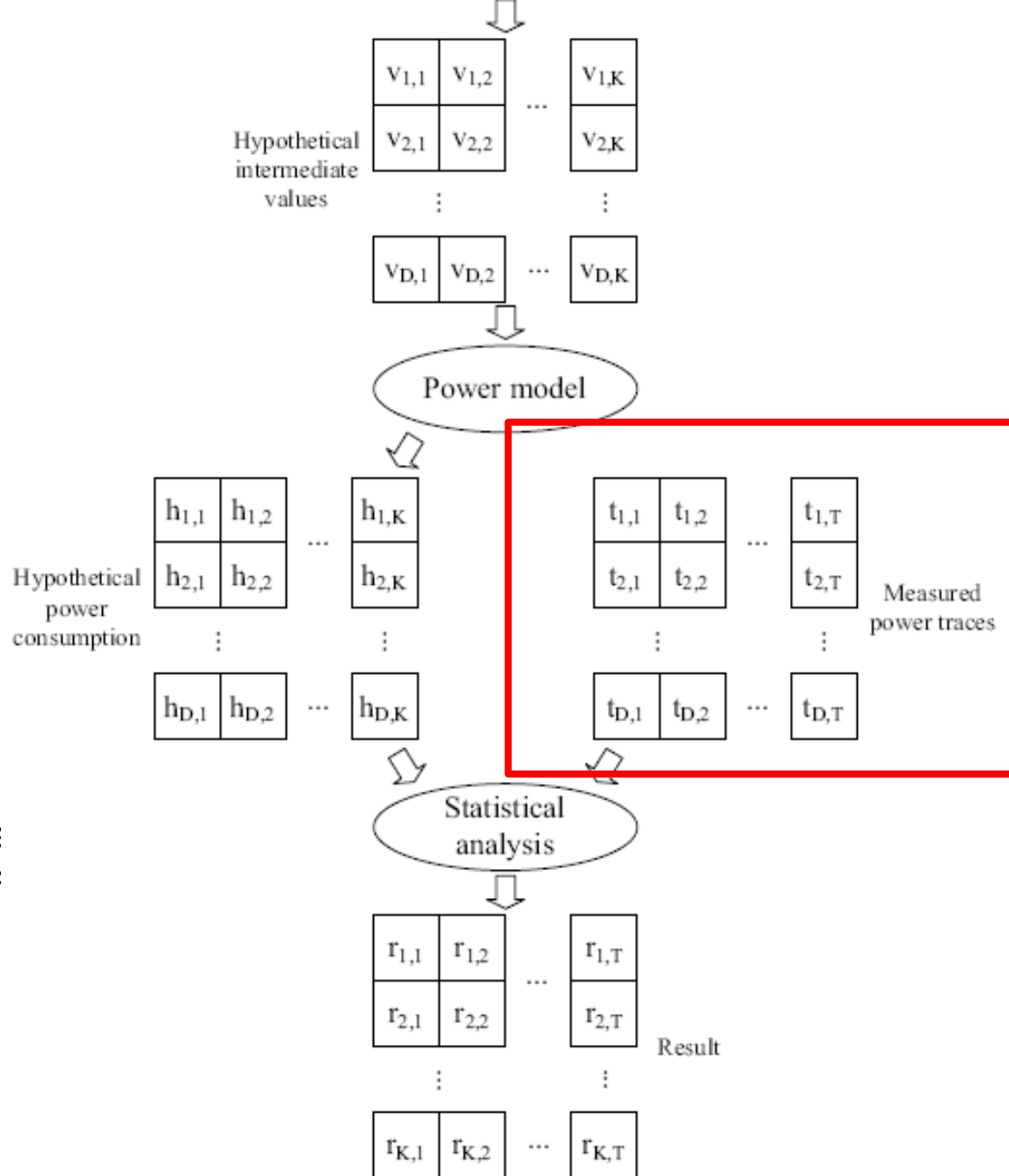
- Algorithm must be known
- Intermediate variable (the target) is function of key and data
- Good targets for DPA are functions which are highly non-linear as they facilitate distinguishability of the correct key hypothesis



5 Steps in detail

Step 2: Measuring the power consumption for D input values

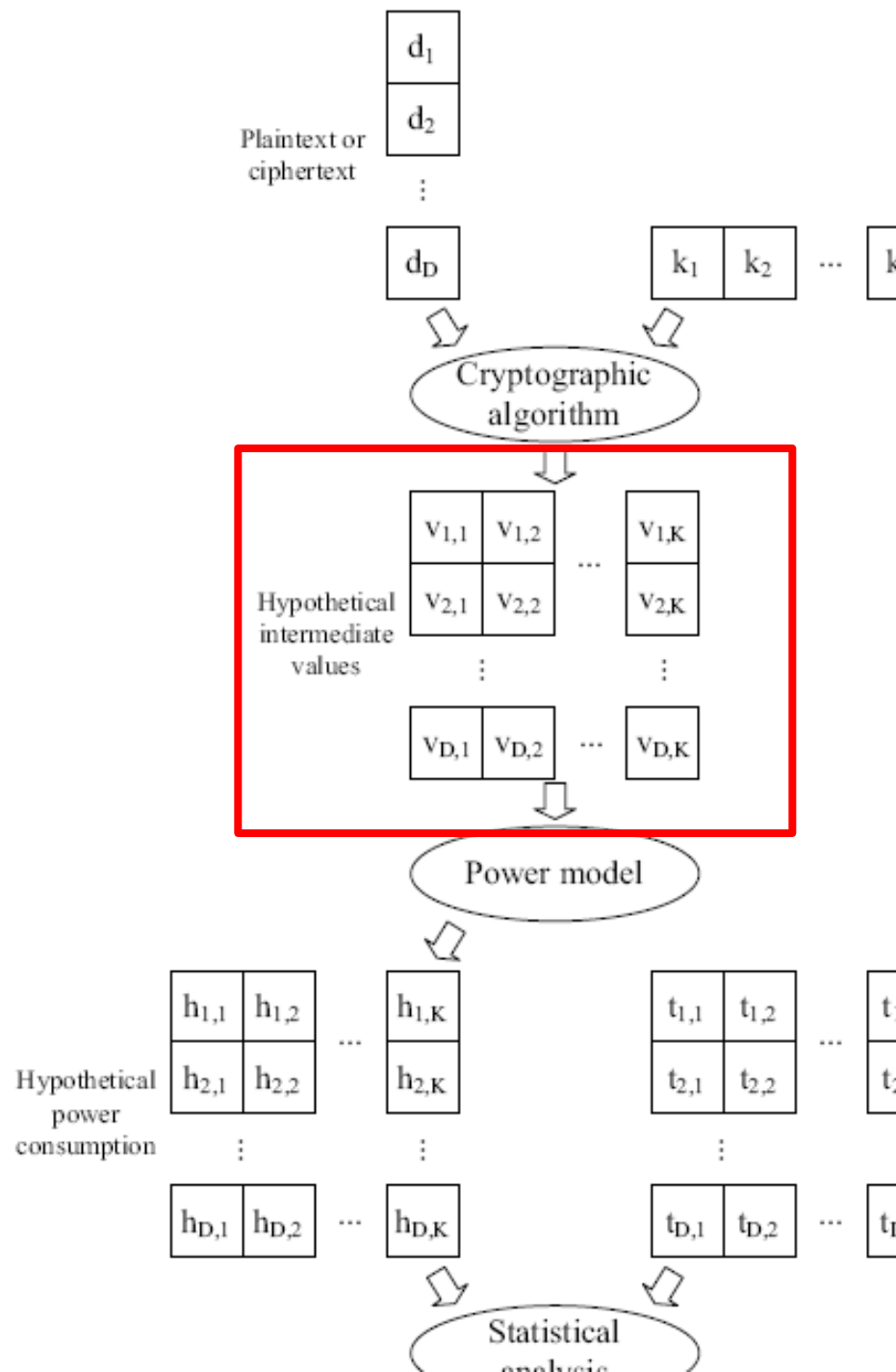
- If necessary process traces to reduce noise and strip away unnecessary information
- Result is a (D x T) matrix \mathbf{t}
 - For each input we store a power trace consisting of T data points



5 Steps in detail

Step 3: Calculating hypothetical intermediate values

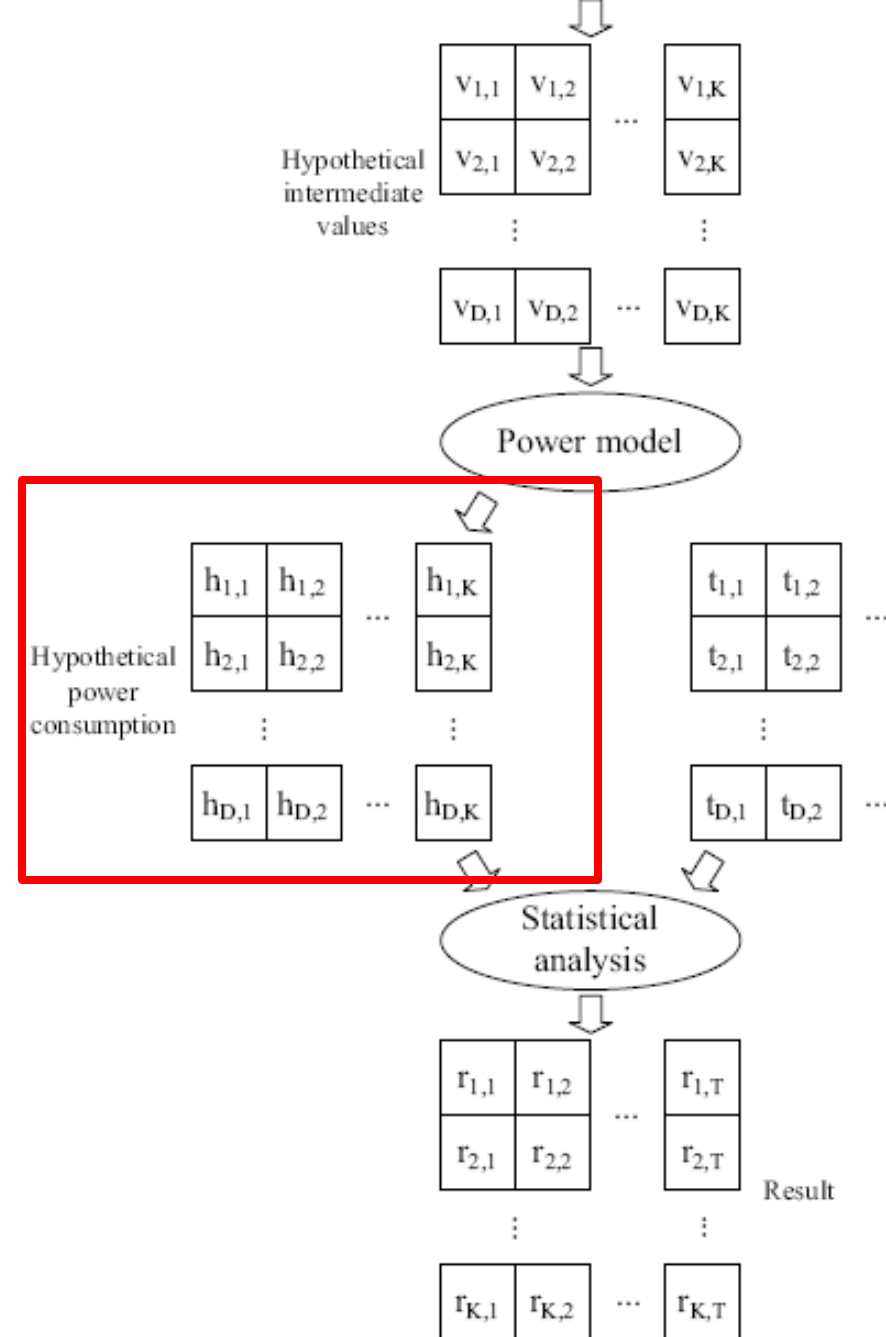
- For a K -bit guess we get 2^K key hypotheses:
 $\mathbf{k} = (k_1, \dots, k_{2^K})'$
- Together with the D inputs we can calculate a $(D \times 2^K)$ Matrix containing the predicted intermediate values $v_{i,j} = f(d_i, k_j)$
 - We write the column corresponding to the correct key as 'ck'



5 Steps in detail

Step 4: Calculating hypothetical power values

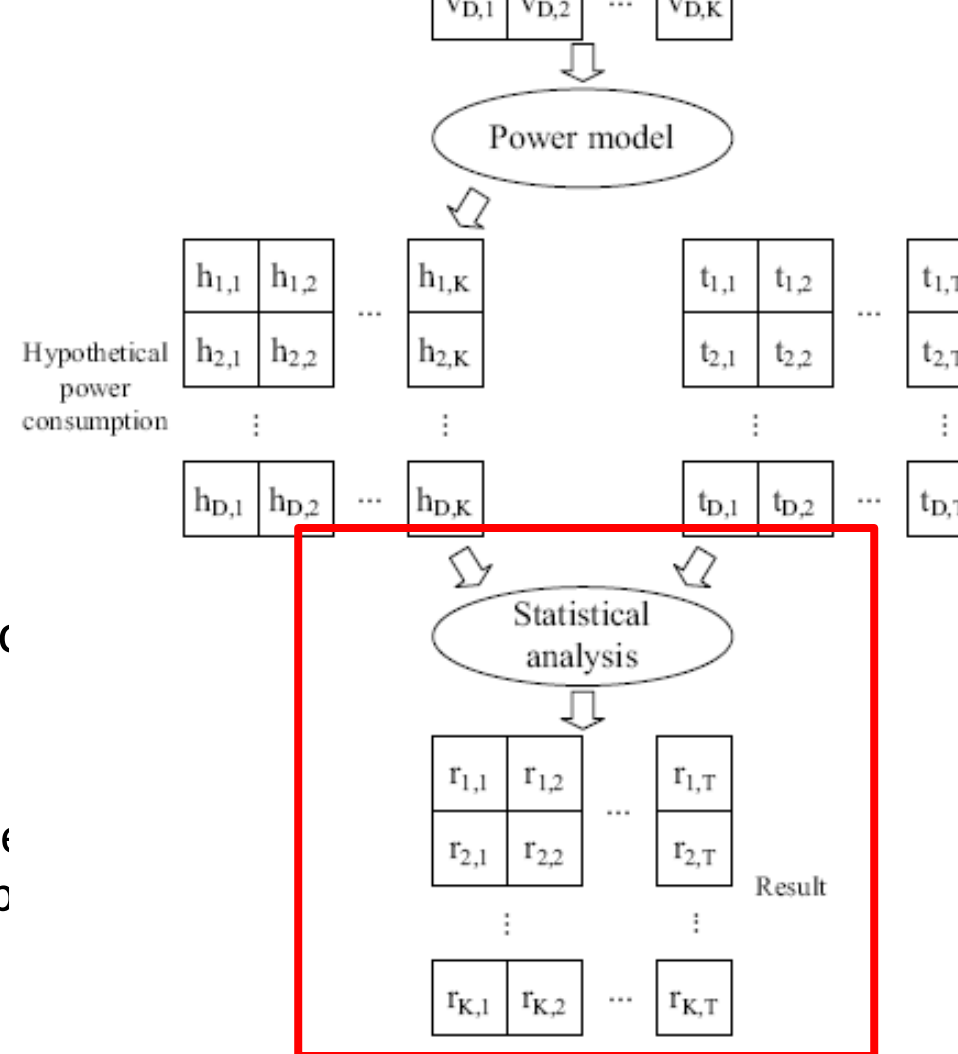
- We map the values in matrix V to matrix H using a power model
- Models requiring little knowledge:
 - Bit model
 - Zero-value model
- Models requiring moderate knowledge
 - Hamming Weight (HW)
 - Hamming Distance (HD)
- Models that require characterisation
 - Templates



5 Steps in detail

Step 5: statistical analysis

- Each column of **H** is compared with (column-wise)
 - Results in a $(K \times T)$ matrix
- The correct key hypothesis 'stands out'
 - How it stands out depends on the distinguisher
 - But the general principle is that the better the modelled traces resemble the real traces the more likely the associated key hypothesis is the correct key
 - We expect to see something 'distinguishable' in **R** at index (ck, ct) ($ct \dots$ time index when attacked intermediate variable is processed)

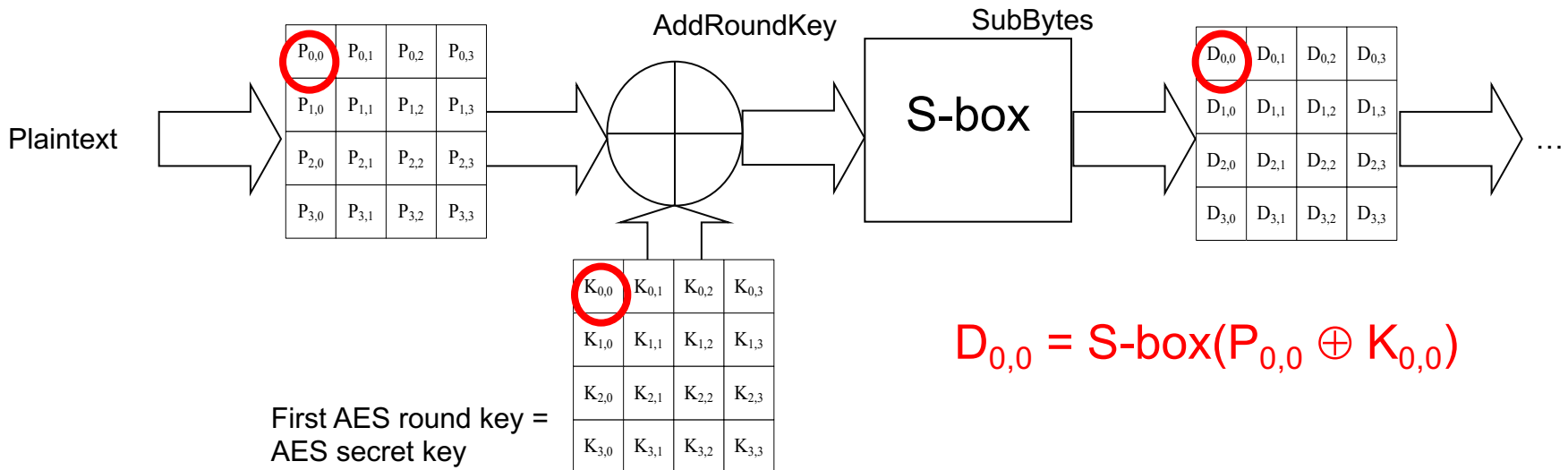


A DPA attack on an AES implementation

- We use traces that have been acquired using a simple microcontroller setup
 - Resemble traces one would get from a typical low-end DPA unprotected smart card
- AES has been implemented in software as described in the last lecture
 - ARK, SB, SR and the MC
- We assume the adversary has access to the plaintexts and aims to recover the key

Step 1: Intermediate variable

- **Choice:** output byte of the first AES S-box operation in round 1
 - Function of the first byte of the plaintext and the first byte of the key
- The first round of AES



Steps 2-4

- **Step 2:** Measurement sample:
1000 traces $\rightarrow \mathbf{T}$
- **Step 3:** $v_{i,j} = \text{S-box}(d_i \oplus k_j) \rightarrow \mathbf{V}$
 - size of \mathbf{V} : 1000x256
- **Step 4:** $h_{i,j} = \text{LSB}(v_{i,j}) \rightarrow \mathbf{H}$
 - size of \mathbf{H} : 1000x256

Steps 2-4

| Sample number i | Plaintext byte d_i | Key hypothesis $k_1 (j=1)$ | Hyp. intermediate value $v_{i,1} = \text{S-box}(d_i \oplus k_1)$ | Hyp. power consumption $h_{i,1} = \text{LSB}(v_{i,1})$ |
|----------------------|-------------------------|-------------------------------|---|---|
| 1 | 0D | 00 | D7 | 1 |
| 2 | 95 | 00 | 2A | 0 |
| 3 | 17 | 00 | F0 | 0 |
| 4 | C7 | 00 | C6 | 0 |
| 5 | 9B | 00 | 14 | 0 |
| 6 | 3B | 00 | E2 | 0 |
| 7 | 34 | 00 | 18 | 0 |

Steps 2-4

| Sample number i | Plaintext byte d_i | Key hypothesis k_2 ($j=2$) | Hyp. intermediate value $v_{i,2} = \text{S-box}(d_i \oplus k_2)$ | Hyp. power consumption $h_{i,2} = \text{LSB}(v_{i,2})$ |
|----------------------|-------------------------|-----------------------------------|---|---|
| 1 | 0D | 01 | FE | 0 |
| 2 | 95 | 01 | 22 | 0 |
| 3 | 17 | 01 | 47 | 1 |
| 4 | C7 | 01 | B4 | 0 |
| 5 | 9B | 01 | B8 | 0 |
| 6 | 3B | 01 | 80 | 0 |
| 7 | 34 | 01 | 96 | 0 |

Step 5

- **Step 5:** $R = D(H, T)$, we choose $D = \text{correlation}$

Key Hypotheses

Samples

| | | |
|---|---|---|
| 1 | 0 | 0 |
| 0 | 0 | 0 |
| 0 | 1 | 1 |
| 0 | 0 | 0 |

Hypothetical power consumption

Time

Samples

| | | | | |
|----|----|----|----|----|
| 10 | 5 | 4 | 45 | 45 |
| 67 | 56 | 45 | 4 | 8 |
| 37 | 54 | 12 | 45 | 5 |
| 27 | 12 | 69 | 8 | 2 |

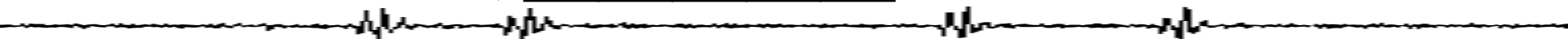
Measured power consumption

Time

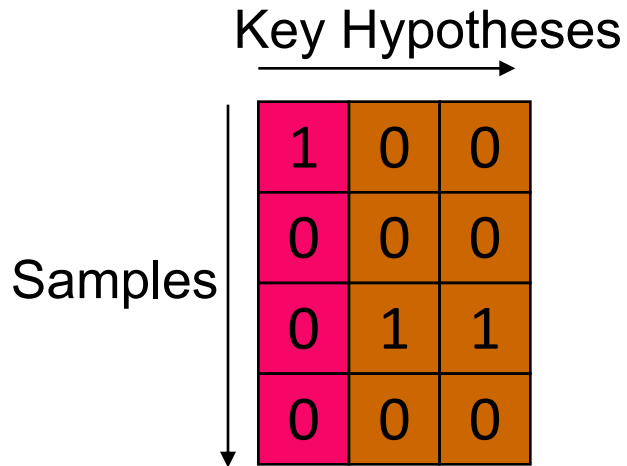
Key Hypotheses

| | | | | |
|--|--|--|--|--|
| | | | | |
| | | | | |
| | | | | |

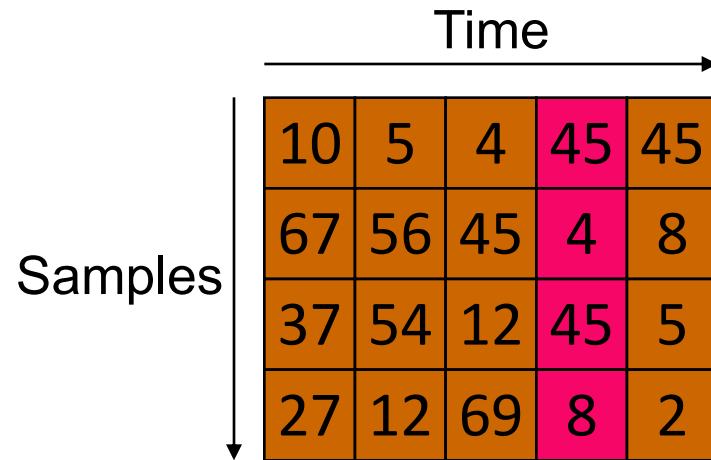
Correlation



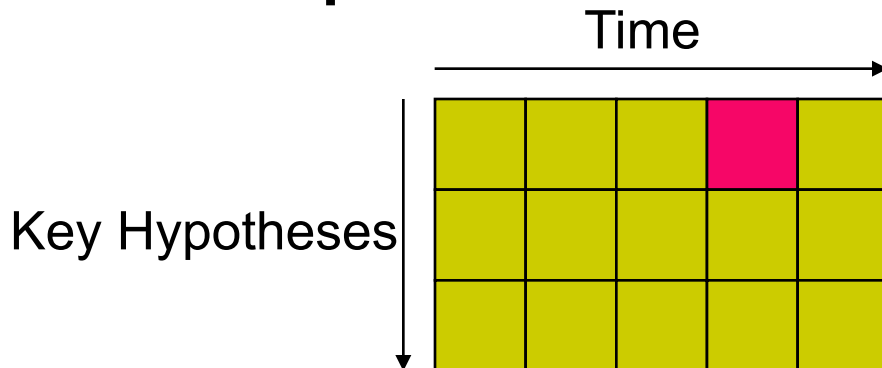
Step 5



Hypothetical power consumption

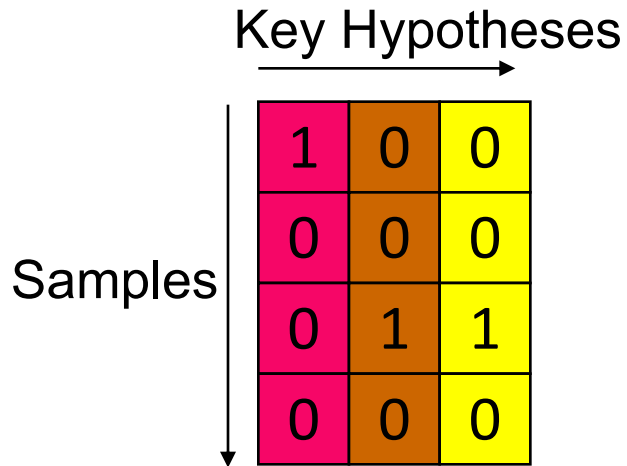


Measured power consumption

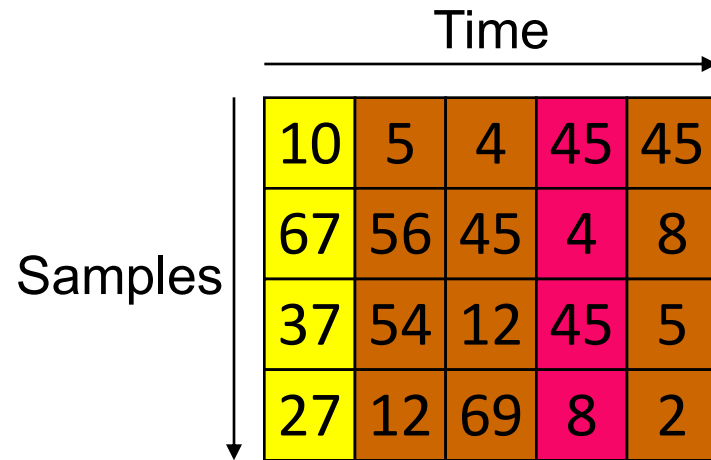


Correlation

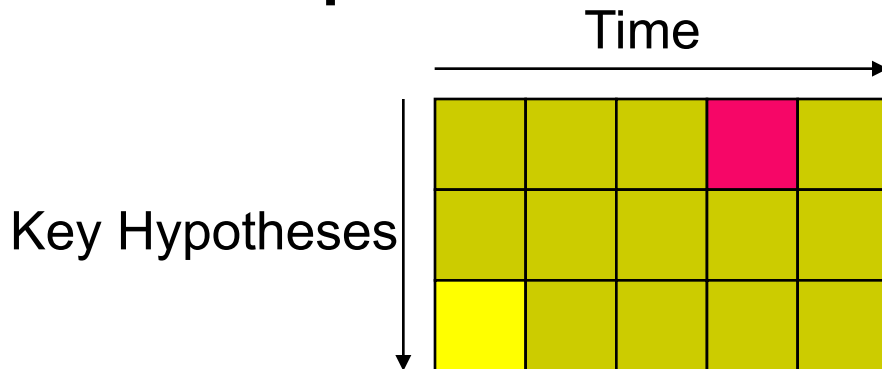
Step 5



Hypothetical power consumption



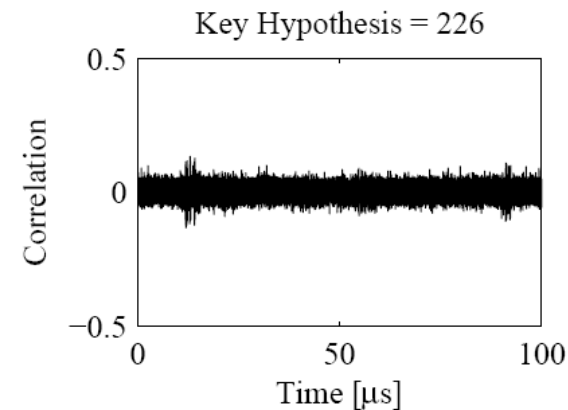
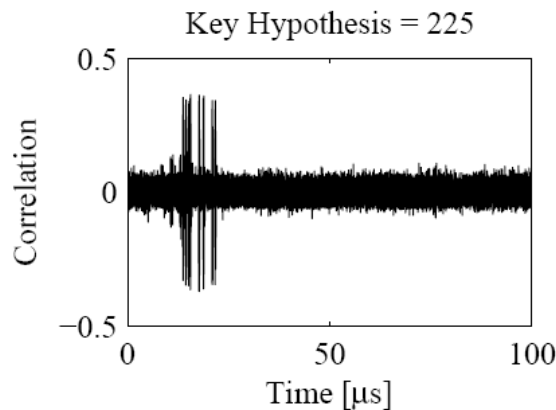
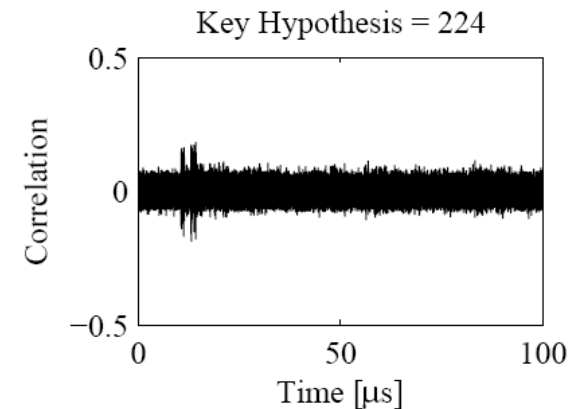
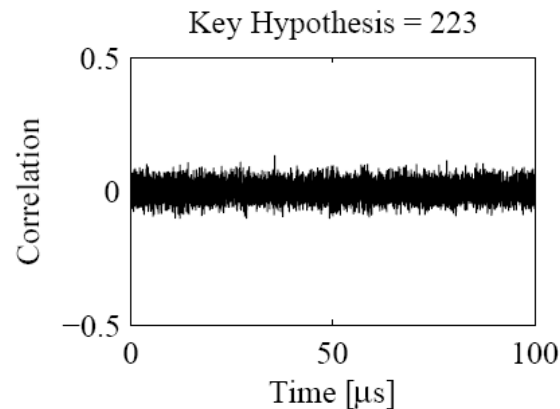
Measured power consumption



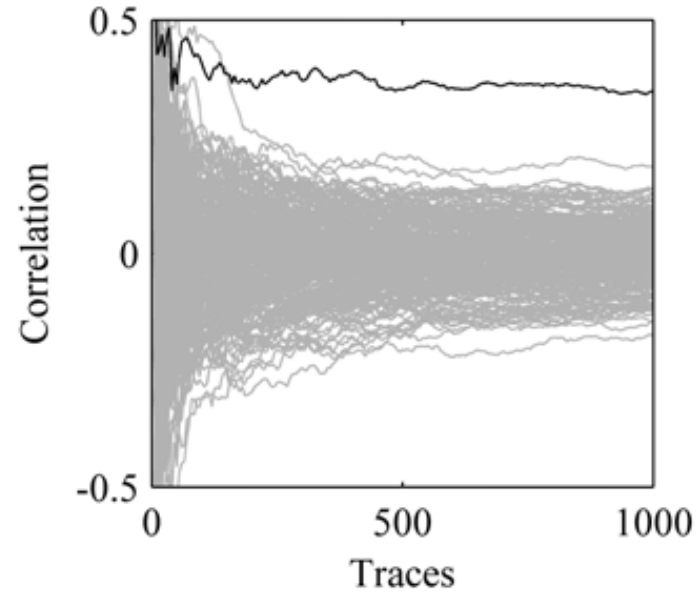
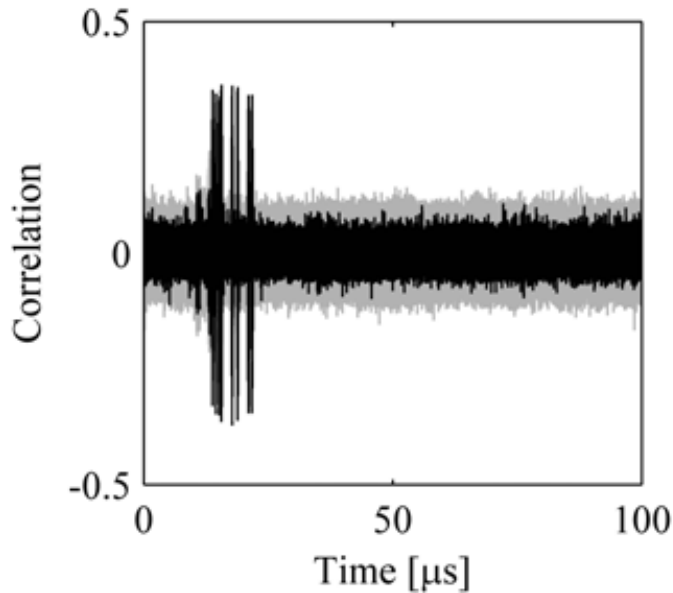
Correlation

Step 5: Result

- Inspect the resulting 256 traces
- We used correlation hence we expect the correct key hypothesis to lead to the highest correlation
- Key $ck = 226$

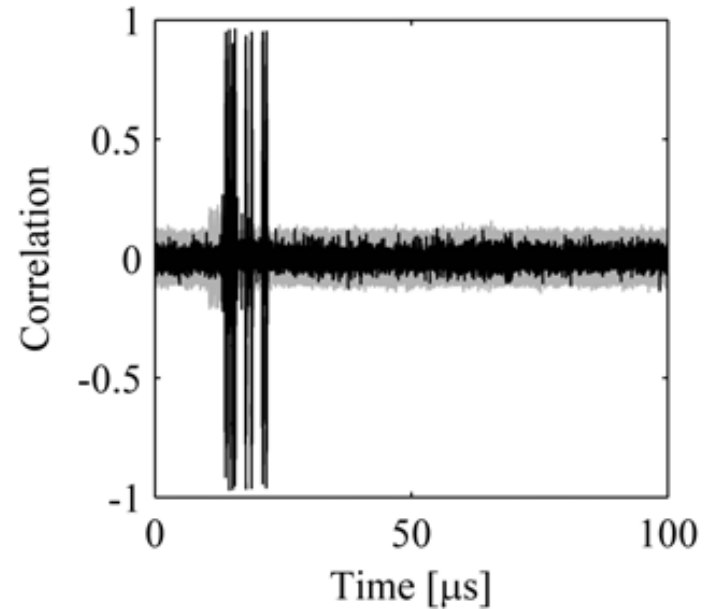
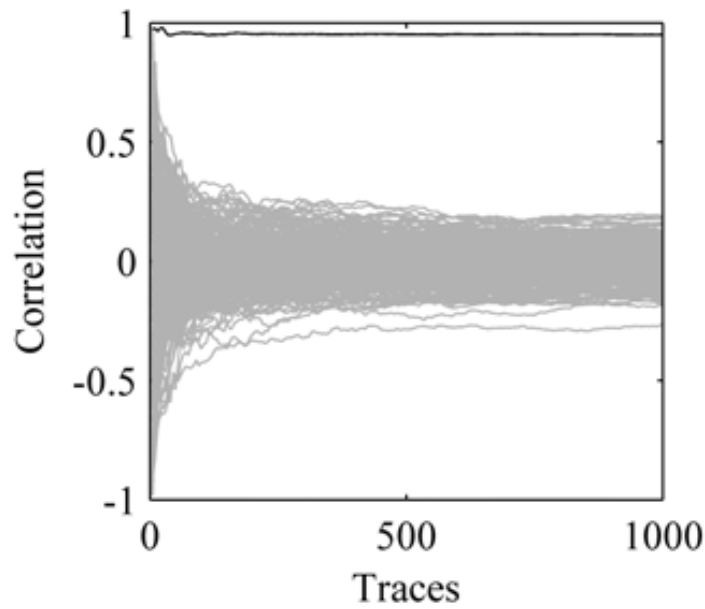


A closer look at the result



- The correct key hypothesis leads to peaks in the correlation trace that are significantly higher than the peaks that come from incorrect key hypotheses
- About 200 power traces suffice to determine the key

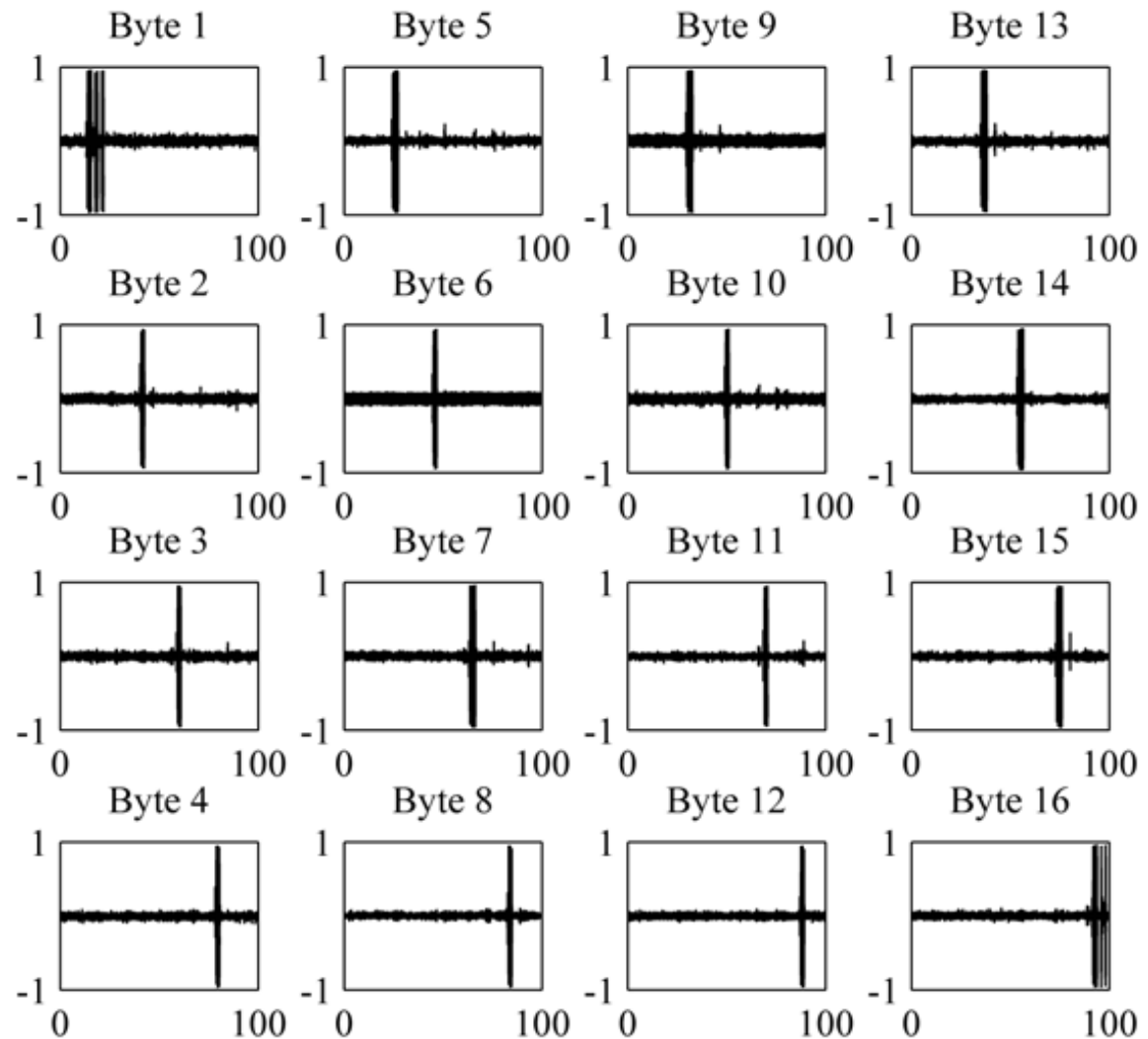
Improving the result



- Using a better power model improves the result of a DPA attack:
 - The correlation coefficient goes up to almost 1 if we compare the Hamming weight of the hypothetical intermediate values with the power traces

Peaks reveal information about implementation too

- High correlation coefficients indicate correct key byte
- Correlation coefficients in this example are almost maximal
- About 30 traces are sufficient to reliably determine the key
- Positions of DPA peaks reveal the points in time when attacked intermediate result is computed.



5 Step Model of a DPA Attack

- | | |
|---|---|
| 1) Selection of target intermediate variable | 1. Linear vs. Non-linear transformation on combination of input and key |
| 2) Measurement of the power consumption/EM | 2. Signal vs. Noise |
| 3) Calculation of the hypothetical intermediate values | 3. Computational resources depending on size of key hypothesis |
| 4) Mapping of hypothetical values to predicted power consumption values | 4. Quality of model |
| 5) Statistical analysis using a distinguisher | 5. Ability of distinguisher to use power model |

DPA Demo

- www.dpabook.org

A summary

- DPA attacks involve a number of steps which include choices that will influence the results
- A good power model is crucial for a correlation based attack to succeed (with few traces)
- Clearly the more distinguishable key hypotheses are the less traces are needed
- Note: you are now able to do the DPA assignment in your coursework!