

$$\begin{array}{lcl} \max & z = & \frac{16}{5} - \frac{2}{5}x_3 - \frac{4}{5}x_4 \\ & x_1 = & \frac{7}{5} - \frac{1}{5}x_3 - \frac{3}{5}x_4 \\ & x_2 = & \frac{2}{5} - \frac{1}{5}x_3 + \frac{5}{5}x_4 \end{array}$$

$$\text{fnc}(\frac{1}{5})x_3 + \text{fnc}(-\frac{1}{5})x_4 \geq \text{fnc}(\frac{7}{5})$$

$$\frac{1}{5}x_3 + \frac{3}{5}x_4 \geq \frac{2}{5}$$

$$\frac{1}{5}x_3 + \frac{3}{5}x_4 - x_5 = \frac{2}{5} \rightarrow x_5 = -\frac{2}{5} + \frac{1}{5}x_3 + \frac{3}{5}x_4$$

$$(B \setminus \{x_5\}) \cup \{x_4\}$$

$$x_4 = \frac{5}{3} \left( x_5 + \frac{2}{5} - \frac{1}{5}x_3 \right) = \frac{2}{3} - \frac{1}{3}x_3 + \frac{5}{3}x_5$$

$$x_2 = \frac{7}{5} - \frac{1}{5}x_3 + \frac{2}{5} \left( \frac{2}{3} - \frac{1}{3}x_3 + \frac{5}{3}x_5 \right) = \frac{5}{3} - \frac{1}{3}x_3 + \frac{2}{3}x_5$$

$$x_1 = \frac{2}{5} - \frac{1}{5}x_3 - \frac{3}{5} \left( \frac{2}{3} - \frac{1}{3}x_3 + \frac{5}{3}x_5 \right) = 0 - x_5$$

$$z = \frac{16}{5} - \frac{2}{5}x_3 - \frac{4}{5} \left( \frac{2}{3} - \frac{1}{3}x_3 + \frac{5}{3}x_5 \right) = \frac{8}{3} - \frac{2}{15}x_3 - \frac{4}{3}x_5$$

$$\max z = \frac{8}{3} - \frac{2}{15}x_3 - \frac{4}{3}x_5$$

$$x_1 = 0 - x_5$$

$$x_2 = \frac{5}{3} - \frac{1}{3}x_3 + \frac{2}{3}x_5$$

$$x_4 = \frac{2}{3} - \frac{1}{3}x_3 + \frac{5}{3}x_5$$

Esercizio 2. (5 punti) Per il seguente problema — noto come 2-knapsack — dove  $p_i, w_i, v_i$  e  $b_1, b_2$  sono numeri interi positivi, proporre almeno due metodi di rilassamento, discutendone punti di forza e debolezze.

$$\begin{array}{l} \max z = \sum_{j=1}^n p_j x_j \\ \text{soggetto a } \sum_{j=1}^n w_j x_j \leq b_1 \quad \text{limite peso} \\ \sum_{j=1}^n v_j x_j \leq b_2 \quad \text{limite volume} \\ x_1, \dots, x_n \in \{0, 1\}. \end{array}$$

1) Rilasciamento continuo:

$$\max z = p^T x$$

$$w^T x \leq b_1 \quad (1)$$

$$v^T x \leq b_2 \quad (2)$$

$$x_1, \dots, x_n \in [0, 1]$$

→ efficiente: greedy  $O(n)$

→ Efficienza:

$$\frac{\text{GREEDY}(I, p, w, v, b_1, b_2)}{n}$$

$$v^T x \leq b_2 \quad (2)$$

$$x_1, \dots, x_n \in [0, 1]$$

$h=2$

i	1	2
$\uparrow$	1	1
$w$	1	2
$v$	1	2

GREEDY ( $I, \uparrow, w, v, b_1, b_2$ )

$$\bar{w} = w$$

$$\bar{v} = v$$

$$i = 1$$

while ( $\bar{w} - w_i \geq 0$  &  $\bar{v} - v_i \geq 0$ )

$$x_i = 1$$

$$\bar{w} = \bar{w} - w_i$$

$$\bar{v} = \bar{v} - v_i$$

$$i = i + 1$$

// o  $\bar{w} - w_i < 0$  o  $\bar{v} - v_i < 0$

$$f_1 = \infty, f_2 = \infty$$

if ( $\bar{w} - w_i < 0$ )

$$f_1 = \sum_{j=1}^i w_j^* / w_i$$

if ( $\bar{v} - v_i < 0$ )

$$f_2 = \sum_{j=1}^i v_j^* / v_i$$

if  $f_1 \leq f_2$ :

$$z = \sum_{j=1}^i n_j + f_1 \cdot n_i$$

else:

$$z = \sum_{j=1}^i n_j + f_2 \cdot n_i$$

2) Elm. uncod.:

~~$x_1$~~   ~~$x_2$~~  2 problemi dello zaino

→ si prende la soluzione minima (UB migliore)

i	1	2	3	4
$n$	10	8	7	5
$w$	2	1	1	4
$v$	4	3	6	7

$$b_1 = 5 = b_2$$

$$z^* = 25$$

$$z^* = 10$$

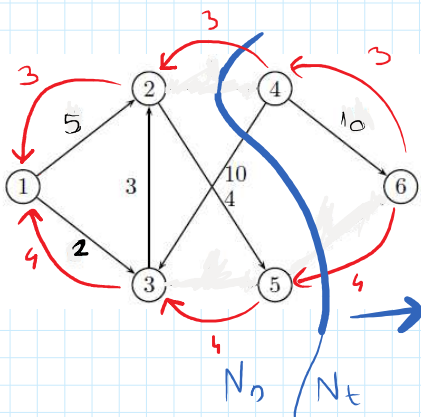
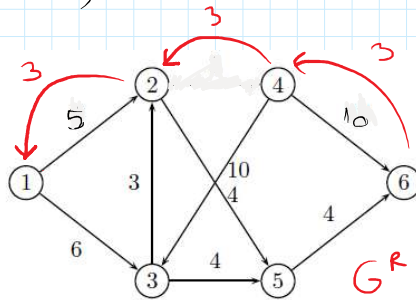
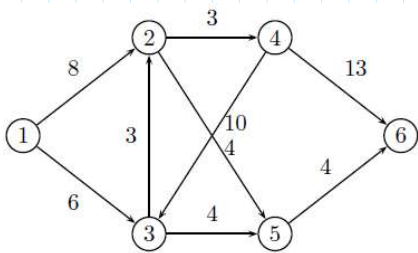
3) bil. lagrangiano:

3) Fil. Lagrangiano:

$$P(\lambda) = \max \left\{ \sum_{i=1}^m p_i x_i + \lambda \left( \sum_{i=1}^m w_i x_i - h_1 \right) \mid \sum_{i=1}^m v_i x_i \leq h_2 \right\}$$

- è più facile risolvere dato un  $\lambda$ , come  $1 < P < 4$
- si può trovare il  $\lambda$  migliore con il metodo del gradiente

### 3) Max Flow (Ford-Fulkerson)

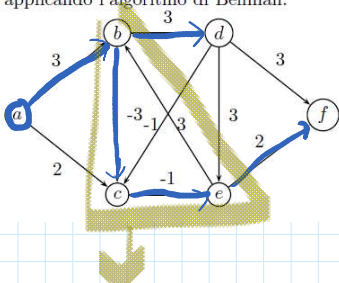


Flusso massimo

$$X = x_{1,2} + x_{1,3} = 3 + 4 = 7$$

taglia di capacità minima  
 $U(N_s, N_t) = 3 + 4 = 7 = X$   
 $\Rightarrow$  Max flow 7

Esercizio 4. (5 punti) Dato il grafo in figura trovare l'albero dei cammini minimi al nodo a applicando l'algoritmo di Bellman.



a	b	c	d	e	f
0	3 2	2 0	6 5	-1	9 1

CONDA: ~~a~~ ~~b~~ ~~c~~ ~~d~~ ~~e~~ ~~f~~ ~~a~~

ciclo di  
 controregressione!

Esercizio 5. (7 punti) Risolvere il seguente problema dello zaino con il metodo del branch and bound.

$$\begin{aligned} \max z &= 10x_1 + 8x_2 + 15x_3 + 7x_4 + 9x_5 \\ \text{sogetto a} \quad &4x_1 + 3x_2 + 7x_3 + 2x_4 + 5x_5 \leq 8 \\ &x_1, \dots, x_5 \in \{0, 1\} \end{aligned}$$

i	1	2	3	4	5
p	10	8	15	7	9
w	4	3	7	2	5
$\frac{p}{w}$	2,5	2,6	2,14	3,5	1,8

[i]	4	2	1	3	5
p	7	8	10	15	9
w	2	3	4	7	5

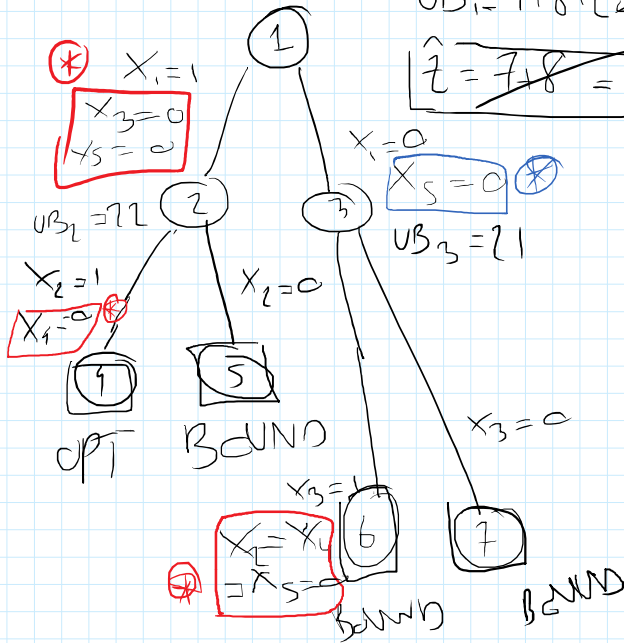
$b = 8$

$$UB = 7 + 8 + \left\lfloor \frac{3}{4} \cdot 10 \right\rfloor = 22$$

$$\hat{z} = 7 + 8 = 15$$

⊗  $x_2 = 0$  per DOMINANZA  
 $x_1 \rightarrow x_5$

⊗  $x_1 = 0$   
per VARIABLE FIXING



$$UB_2 = 10 + 7 + \left\lfloor \frac{2}{3} \cdot 8 \right\rfloor = 22$$

$$\hat{z} = 17$$

$$UB_3 = 7 + 8 + \left\lfloor \frac{3}{7} \cdot 15 \right\rfloor = 21$$

$$\hat{z} = 7 + 8 = 15$$

$$UB_4 = 18 = \hat{z}$$

$$UB_5 = 10 + 7 = 17 \text{ BOUND}$$

$$UB_6 = 15$$

$$UB_7 = 7 + 8 + \left\lfloor \frac{3}{5} \cdot 9 \right\rfloor = 18 \text{ BOUND}$$