

Business & Computer Science

‘It is not from the benevolence of the butcher, the brewer, or the baker that we expect our dinner, but from their regard to their own interest.’

The Wealth of Nations, Books 1-3
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1

Course Introduction

1.1 Course Aims

Despite what the name might make you think, the course is *still* a practical one, it aims to introduce how computers work in a real world scenario, such as a company.

The Purpose of a Computer Through the lens of this course, the purpose of a software is to improve the company's efficiency and throughput, everything is evaluated through a framework of costs and profits, hence, even something as simple as buying a pen must be evaluated at multiple levels, from the cost of the pen itself, to the cost of the time spent by the employee to go to the store and buy it.

1.2 Course Structure

The course is usually organized in two parts, alternating on a weekly basis:

Day	Purpose
Wednesday	Lecture
Thursday	Seminars

Table 1.1: Course Structure

Naturally, the seminars are *included* in the course's materials and are hence expected

The exam is also split into multiple sections:

Part	Contents
Written	30 to 50 closed questions
Oral	Optional, to raise marks

Table 1.2: Exam Structure

The general information about the course is available at the [official course page](#).

1.3 Information Systems

As studied in the first unit of *Data Management & Analysis*, an information system is a set of compo-

nents that allows an organization to collect, process, store and distribute information, it is *not* necessarily something software-based, it can even be a file cabinet!

Definition 1. *An Information System is something that Manages the flow of information in an organization.*

Information Naturally, information is something *abstract* and *immaterial*, it is naturally difficult to manage, but it is possible to construct a set of rules, conventions and tools that allow us to represent it in a way that is *manageable*, this process is the *Flow*, and hence, as defined above, these systems are called *Information Systems*.

Computers When a computer is involved, the Information System is naturally labeled a Computer Information System, or *CIS* for short.

This chapter aims to give us the following capabilities

- Define what an Information System is
- Describe the history of Information Systems
- Describe the basic argument behind the article *Does IT matter?* by Nicholas Carr.

An information system is made up of *Three* main components:

- People: In charge of decision making and organization
- Technology: Hardware and Software that supports the business
- Processes: Collecting and storing information

In General This is a pretty broad definition, that allows a wide variety of software to be seen through the lens of an Information System, from network analyzers, to national healthcare systems, to online education platforms.

Acronyms Depending on the application, an Information System can be called by different acronyms, such as:

- **MRP:** *Manufacturing Resource Planning*
- **CIM:** *Computer Integrated Manufacturing*
- **SAP:** *Systems, Applications and Products*

1.3.1 Anthony's Triangle

Anthony's triangle is a diagram that categorizes the three purposes of an Information System:

1. Strategic (*Executive Information System*): for senior management decisions
2. Tactical (*Management Information System*): For middle management decisions
3. Operational (*Transaction Processing Systems*): For daily transactions of business

1.3.2 Information System Components

Several Components work together to add value to an organization, they are

1. Hardware: The physical components of the system
2. Software: The programs that run on the hardware, they can be:
 - Operating Systems: The programs that manage the hardware
 - Application Software: The programs that are used by the users
3. Data: The information that is stored in the system
4. People: The users of the systems, both producers and consumers of information
5. Processes: The procedures that are used to collect, process and store information¹

This is somewhat redundant to what we stated beforehand in definition 1, but it is important to restate it to emphasize the fact that the system is composed of both *People* and *Technology*.

1.3.3 Processes

One of the most important components of an Information System is the *Process*, it is the goal of the Computer Information System to optimize the processes of an organization, bringing about an increase in efficiency.

¹ Organizing something in *processes* that is, a series of well-defined steps, brings about a series of benefits in productivity.

Processes Formally Since processes are an advantageous way to organize work, it is important to spend some time to also give a formal definition of their nature.

Definition 2. *A process is a series of well-defined steps that are used to achieve a specific goal, it can be defined as a set of activities*

1.4 Does IT Matter?

Nicholas Carr, in *Harvard Business Review*, argues that Information Technology is not an *Investment*, but rather a commodity, so something that must be managed to optimize the company's profits by reducing its operational costs.

IT as a Marketing Tool It is also interesting to note how a company is perceived as better when it uses IT, and when this IT is of high quality, therefore IT can be seen as a *Marketing Tool* that can be used to attract customers.

1.5 Mainframes

A Mainframe is a class of computer that is usually used as the heart of an Information System, where everything is *centralized*, as opposed to a distributed systems, users, which in this architecture are defined *dumb*, are not allowed to access the system directly, but rather act as consumers of its services as allowed by the system's administrators and operating system.

Definition 3. *A Mainframe is an architecture in which a central computer with very high processing power is connected to a multitude of terminals through a star topology, where the central computer is the hub of the network.*

Even though *Mainframes* are not as popular as they used to be, they are still used in many field where they are unmatched in terms of performance.

IBM The current IBM's Operating System is z/OS, which is a *Monolithic* Operating System, it is being opened to different programming languages, such as **Java**, beforehand, however, it was only available in **COBOL**.

Batch Processing Mainframes are based on the *Batch Processing* paradigm, where a set of jobs are submitted to be executed in a *batch*, that is, a set of jobs that are executed in a single run, this is in contrast to *Real Time Processing*, where jobs are executed as soon as they are submitted.

This, provided that the company has a sufficient amount of jobs to be executed, allows for a continuous flow of work, and hence, a higher level of productivity.

2

Process Modeling

2.1 Modeling Business

In order to properly formalize the concepts that are going to be introduced through the rest of this course, we will introduce *mathematical models* that allow us to describe the ebbs and flow of human economy and companies.

Processes These models are called *Processes*, there are certain common classes of projects:

1. Service
2. Support
3. Management and Control
4. Physical
5. Information
6. Business

2.2 Process Descriptors

These processes are then described by some diagrams:

- Hierarchical
- State Diagrams (Automata)
- DFD – Data Flow Diagram
- Wide – Workflow on an Intelligent and Distributed database Environment
- Action Workflow
- Petri Nets

2.2.1 Hierarchical Process Model

Everything in a company can be described as a set of *Hierarchies*, that is, a tree of *Processes* that are organized in a *Top-Down* fashion, where the *Top* process is known as *Macroprocess*, the hierarchy goes as follows:

1. Macroprocess – Sales
2. Process – Sales Management
3. Phase – Order Processing
4. Activity – Shipment
5. Operation – Pricing, Packaging, etc. . .

Naturally, these hierarchy together form a *Forest*.

2.2.2 Data Flow Diagrams

Data Flow diagrams are what we also call *Flowcharts*, they are a graphical representation of the flow of the data through the company's Information Systems, they are not used often due to them being subject to *spaghettification*, that is, the diagrams become so complex that they are hard to understand.

These flowcharts are usually composed of a set of components that are used to represent the different parts of the system, they are:

- Processes – Circles
- Data Collections – Rectangles
- Interface – Bordered Rectangles
- Data Flows – Directed Arrows

Definition 4 (Data Flow). *A Data Flow represents any kind of flow in a system, the first component must be a process, the second can be either a process, a data collection or an interface, moreover, they can be either*

1. *Structured*
2. *TODO*

Data Dictionary Usually, to help with readability, we provide a *Data Dictionary* and a textual description of each process, to help the user understand what the process does.

2.2.3 WIDE

Wide Relies on three main components:

- Process Model
- Information Model
- Organization Model

Process Model Wide still relies on a *Process* model, it is somewhat similar to an *Activity Diagram* used to describe Object Oriented Code

Anti-Spaghetti Techniques The wide model introduces some more complex components with respect to the DFD model, such as forks and joins, to better describe the flow of the data through the system with fewer connections between the components.

Remark 1. *An analyst's role is to translate the business model in a way that is understandable in layman's terms, that is, when you are describing a process through a graph, strive to be clear, the important thing is that People Are Going to Read It.*

2.2.4 Petri Nets

Petri Nets are a formal model that is used to describe the flow of data through an Information Systems

Definition 5 (Petri Nets). *A petri net is a 3 – Tuple = (P, T, A) that forms a Bipartite Graph:*

$$G = (P \cup T, A), A \subseteq (P \times T) \cup (T \times P).$$

Where:

- *P is a set of places*
- *T is a set of transitions*
- *A is a set of arcs $A \subseteq (P \times T) \cup (T \times P)$*

There is a set of Initial Markings $M_0 \subseteq P$ that are the places that are initially marked, and we have that $M : P \rightarrow N$ is a function that maps each place to a non-negative integer, that is, the number of tokens in the place.

At the end we reach a set of Final Markings $M_f \subseteq P$ that are the places that are marked at the end of the process.

Essentially, it is a Finite State Automata in which a set of markings indicate the states that are 'Firable', once a transition happens, the markings are updated, and the process continues, so we are effectively running multiple Finite State Machines in parallel, each with the same topology.

Terminology Each transition $t \in T$ has an input set ${}^\circ t \subseteq P$ and an output set $t^\circ \subseteq P$ also called input and output places, the same notation also applies to places.

Petri Net Evolution An enabled transition can *Fire*, deleting a token in each input place and creating a token in each output place:

$$M_0 = (2, 1, 0, 0, 1)$$

$$M_0 \rightarrow M_1$$

$$M_1 = (1, 0, 1, 1, 1)$$

A *Firable Sequence* is a sequence of transition $\sigma = \langle t_1, t_2, \dots, t_n \rangle$ where

$$M_0 \xrightarrow{t_1} M_1 \xrightarrow{t_2} M_2 \dots \xrightarrow{t_n} M_n.$$

For which we can use the closure notation $M_0 \xrightarrow{\sigma} M_n$ to denote the evolution of the markings.

Further Definitions Now, given a net $P = (P, T, A, M_0)$ we have that:

- A *Potentially Firable Transition* $t \in T$ is such that

$$\exists \tau \in T^* \text{ s.t } \tau t \text{ Is Firable : } M_0^\tau \rightarrow^t.$$

- A *Potentially Firable Sequence* $\sigma \in T^*$ is such that there exists a prefix sequence $\tau \in T^*$ such that $\tau \sigma$ is firable ($M_0 \rightarrow^{\tau \sigma} M_n$)
- A *Reachable Marking* M is such that

$$\exists \sigma \text{ s.t } M_0 \rightarrow^\sigma M.$$

- $R(M_0)$ is the set of all reachable markings.
- P_r is the set of reachable places s.t

$$P_r = \{p \in P | \exists M \in R(M_0), M(p) > 0\}.$$

Petri Nets and Automata We have that Petri nets are a generalization of automata under particular conditions:

Definition 6 (Petri Nets and Automatas). *A Finite state machine is a petri net where, for each transition t both the input and the output places have cardinality 1:*

$$\forall t \in T, |{}^\circ t| = |t^\circ| = 1.$$

2.2.5 Workflow Nets

Workflow nets are a generalization of Petri Nets, in which we have two additional conditions:

1. A source s.t $|^{\circ}t| = 0$ for all $t \in T$ ¹
2. A sink s.t $|t^{\circ}| = 0$ for all $t \in T$

Hence for each node the node is *reachable* from the source and *can reach* the sink.

Workflow Nets as Models Workflow nets are more representative of real world systems, since they can model *producers* and *consumers* in a system, which are widespread in real world systems.

¹ AI generated math, please check (will fix once the slides are up).

3

Petri Nets – Cont.

3.1 Introduction

Petri nets are, as aforementionedly introduced, a *Graphical* and *Mathematical* modelling tool, that is, they are formalizable through math and representable pictorially.

Graphical Representation The graphical representation is a bipartite graph, where we have two kind of nodes:

- Places
- Transitions

The transitions are stylized as *black rectangles*, in order to distinguish them more easily, but they are easily interchangeable with the usual labeled circles that are used to represent states in a finite state machine.

Tokens are pictorially represented with dots, however, this is not that scalable, so we can also use mathematical notations and label a node with a number $n \in \mathbb{N}$ to represent n tokens in the place, be careful with this representation since it reduces the expressiveness of the model, especially towards non-mathematics oriented people.

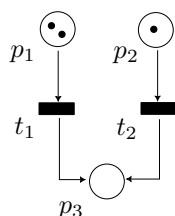


Figure 3.1: A Petri Net with 3 places and 2 transitions, we can see three tokens in p_1 and one token in p_2 .

Mathematical Representation Since petri nets are just a graph, they can be rigorously for-

malized, as in definition 5, so we omit the formal definition here to avoid redundancy.

Transition Enabling We can further complicate a Petri Net by introducing the concept of transition *enabling*: a transition is enabled if it has enough tokens in its input places to fire, and it is disabled otherwise, naturally, the firing of the transition removes a token from each input place and adds a token to each output place.

The system steps in time *discretely*, that is each event happens simultaneously in a single step in time.

Petri Net Non-Determinism The evolution of a petri net is not deterministic, that is, there is no guarantee that a transition will fire in a given step, and therefore we cannot predict the next state of the system from the current state.

Conflicts Whenever we have a set of transition T defined as such:

$$T := \{t_1, t_2, \dots, t_n\} \quad (3.1)$$

$$s.t \exists (t_i, t_j) \in T \quad (3.2)$$

$$s.t t_i \wedge t_j \text{ are both enabled} \quad (3.3)$$

where the firing of one of $t_i \vee t_j$ would lead to the disabling of the other, we have a *conflict*, these conflicts are not too dissimilar to what was explored in the Operating Systems course with the *deadlock* concept¹, these conflicts can be resolved by introducing *priority* to the transitions, that is, we can restructure the mathematical definition to allow for a hierarchy of transitions, or we can simply resolve the conflict by restructuring the model itself.

3.2 Reachability graphs

Reachability graphs are a way to represent the evolution of a Petri Net, they are a directed graph,

¹ Specifically the dining philosophers problem.

where each node is a marking of the Petri Net, and each edge represents a transition that can be fired from the current marking to the next marking.²

Reachability Graph Construction Reachability graphs are constructed by starting from the initial marking, and then, for each marking, we generate all the possible markings that can be reached from the current marking by firing a transition.

By properly identifying the frontier nodes, the generation of the reachability graph is performable in a finite amount of steps, even if the Petri Net is unbounded, this is polynomial time, and therefore the reachability graph can be generated in a finite amount of time.

Infinite Reproducibility A marking $M'' \geq M'$ is *Infinitely Reproducible* if:

$$M'' \geq M' \quad (3.4)$$

$$m_i'' \geq m_i' \quad \forall i \in (1, 2, \dots, n) \quad (3.5)$$

3.3 Petri Nets Extensions

Petri nets can be *extended* in order to model more complex systems, for example, we can have:

- Arc Multiplicity
- Inhibitor Arcs
- Priority Levels
- Enabling Functions (Guard Conditions)

Remark 2. The last three extensions destroy the infinitely reproducibility property of the Petri Net.

3.3.1 Arc Multiplicity

An *arc cardinality* may be associated with input and output arcs, whereby the enabling and firing rules change:

- Each input must contain at least as many tokens as the input arc's cardinality
- When the transition fires, the number of tokens removed from each input place is equal to the input arc's cardinality, as is the number of token added to each output place.

² This section is a placeholder, disregard it.

3.3.2 Inhibitor Arcs

Inhibitor arcs are arcs represented with a circle head, the transition can fire only if the inhibitor place does *not* contain any tokens.

They are used to implement constraints on the firing of a transition, in the restaurant example, we could think of a transition that represents the *delivery* of a dish, and we could have an inhibitor arc from the *delivery* transition to the *payment* transition, so that the dish can only be delivered once the payment has been made.

Inhibitor Arcs and Multithreading Once again, we also notice a stunning similarity with multithreading, where we can think of the inhibitor as some sort of synchronization primitive, such as a *mutex* or a *semaphore*.

Inhibitors and Transitions An inhibitor arc placed on a transition means that the transition can only fire if the inhibitor place does *not* contain any tokens, however, the cardinality of the output arc is not affected by the presence of the inhibitor arc, nor is the inhibitor affected by the cardinality.

3.3.3 Priority Levels

A Priority level can be attached to each Petri Net Transition, the standard execution rules are then modified to allow the highest priority to fire first in case of conflict, when two transitions have the same priority, that subset of transitions fires simultaneously.

3.3.4 Enabling Functions

An enabling function (or guard) is a boolean expression, composed with Petri Nets primitives, the enabling rules are modified so that, besides the standard conditions, the transition can only fire if the guard evaluates to *true*.

Hence we can define an enabling functions as:

$$\begin{aligned} f(tk) := & \\ & \times (P, T, A) \\ & \times (+, \cdot, \leq, \geq, \neq, \wedge, \vee, \neg, \text{true}, \text{false}) \\ & \times \{\text{true}, \text{false}\} \end{aligned}$$

3.3.5 Colored Petri Nets

We can extend petri nets with colors, to increase expressiveness:

- We say that C is a set of colors of cardinality $|C|$ and x is a color of C .
- Place p can contain tokens of any color $x \in C$
- Transition t can fire tokens of any color $x \in C$

3.3.6 Stochastic Petri Nets

Petri nets can be extended by clearly associating *time* with the firing of transitions, resulting in *Timed Petri Nets*.

A special case of timed Petri Nets is stochastic petri nets (SPN), where the firing times are considered to be random variables.

This can be useful to model the random natures of sales, that are based on a number of difficulty-to-predict factors, such as the customer's tastes, current season, advertisements, world events, etc. . .