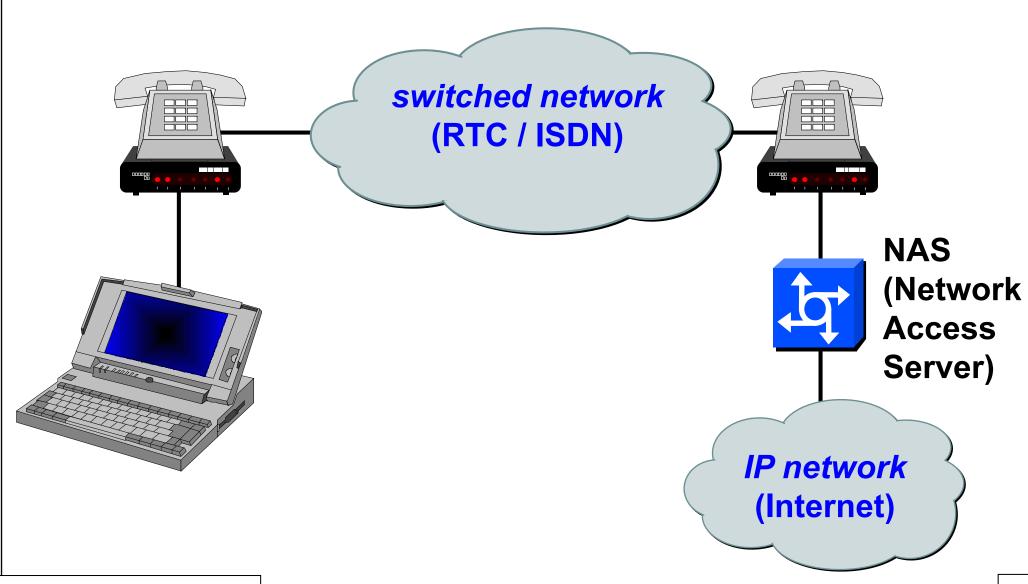
Security of IP networks

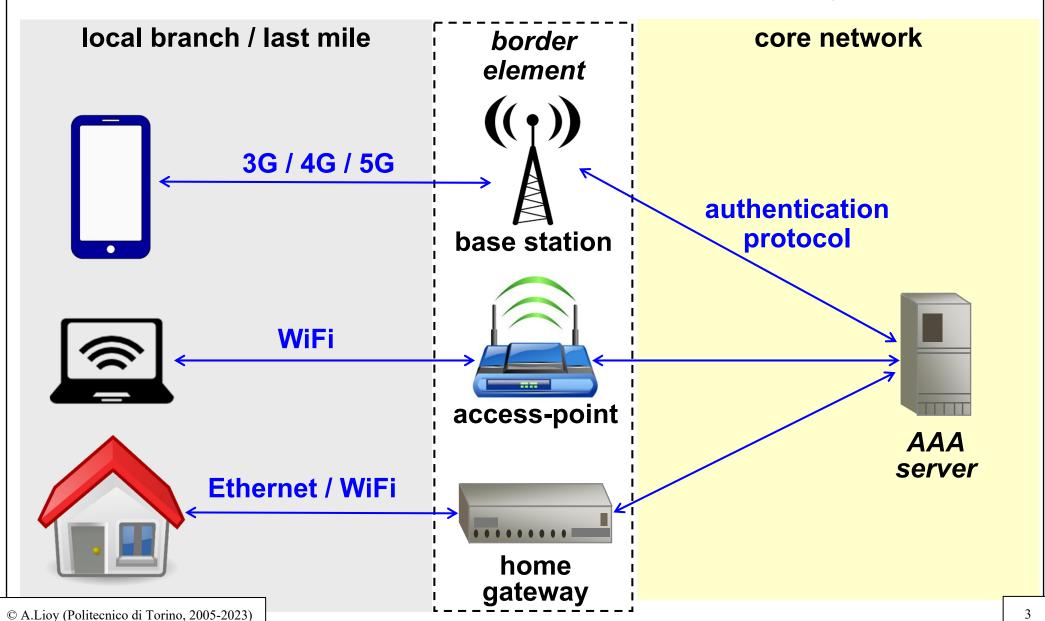
Antonio Lioy < lioy @ polito.it >

Politecnico di Torino Dip. Automatica e Informatica

Remote access via dial-up lines



Network access (modern way)



Authentication of PPP channels

- PPP (Point-to-Point Protocol) is able to encapsulate network packets (L3, e.g. IP) and carry them over a point-to-point link
 - physical (e.g. PSTN, ISDN)
 - virtual L2 (e.g. xDSL with PPPoE)
 - virtual L3 (e.g. L2TP over UDP/IP)
- activated in three sequential steps:
 - LCP (Link Control Protocol)
 - establishing, configuring, and testing the L2 connection
 - can negotiate also authN protocol and algorithm
 - authentication (optional; PAP, CHAP, or EAP)
 - L3 encapsulation via various NCPs (Network Control Protocols)
 - e.g. IPCP (IP Control Protocol) for IP packets

Authentication of network access

- protocols used not only for PPP
- PAP
 - Password Authentication Protocol
 - user password sent in clear
- CHAP
 - Challenge Handshake Authentication Protocol
 - symmetric challenge-response (based on user password)
- EAP
 - Extensible Authentication Protocol
 - it's an authentication framework
 - ... to use external techniques (e.g. challenge, OTP, TLS)

LCP Authentication-Protocol Configuration Option

this option contains:

- Type (8bit) = option type
- Length (8bit) = length of option in bytes
- Authentication Protocol (16bit) = protocol identifier
- [Algorithm (8bit)] = algorithm identifier (needed when a protocol supports various ones)

for PAP:

- Type=3, Length=4, Protocol=0xC023
- for CHAP:
 - Type=3, Length=5, Protocol=0xC223, Algorithm=5 (for MD5)

PAP

- Password Authentication Protocol
- RFC-1334 "PPP Authentication Protocols" (Oct 1992)
 - note: defines also initial version of CHAP
- user-id and password sent in clear between Peer and Authenticator
- authentication only once when the channel is created
- very dangerous!

PAP: 2-way handshake protocol

- (Peer > Authenticator) Authenticate-Request (code=1)
 - Code (8bit) + Identifier (8bit) + Length (16bit)
 - Peer-ID Length (8bit) + Peer-ID (0-255B)
 - Passwd-Length (8bit) + Password (0-255B)
- (Authenticator > Peer) Authenticate-Response (code=2, 3)
 - Code (8bit) + Identifier (8bit) + Length (16bit)
 - Msg-Length (8bit) + Message (0-255B)
 - code=2 (ACK), code=3 (NAK)
- Identifier needed to match Request and Response
- Authenticate-Request or -Response may be lost ... so, Authenticator MUST permit multiple requests

CHAP

- RFC-1994 "PPP Challenge Handshake Authentication Protocol (CHAP)" (Aug 1996)
- symmetric challenge (password-based)
 - initial challenge compulsory (at channel creation)
 - authentication request optionally repeated (with a different challenge) during transmission – decision taken by the NAS
 - challenge MUST be a nonce
- the Authenticators that support both PAP and CHAP must offer CHAP first

CHAP: 3-way handshake protocol

- (Authenticator > Peer) Challenge (code=1)
 - Code (8bit) + Identifier (8bit) + Length (16bit)
 - Challenge-Size (8bit) + Challenge-Value (0-255B)
- (Peer > Authenticator) Response (code=2)
 - Code (8bit) + Identifier (8bit) + Length (16bit)
 - Response-Size (8bit) + Response-Value (0-255B)
- (Authenticator > Peer) Result (code= 3 Success, 4 Failure)
 - Code (8bit) + Identifier (8bit) + Length (16bit)
- Response-Value = md5 (Identifier || pwd || Challenge-Value)
- Identifier needed to match Request and Response
- Challenge or Response may be lost ... so, Authenticator MUST resend Challenge if no Response (until retry limit)

MS-CHAP

MS-CHAPv1

- RFC 2433 "Microsoft PPP CHAP Extensions" (oct.1998)
- dropped by MS starting with Windows Vista
- MS-CHAPv2
 - RFC 2759 "Microsoft PPP CHAP Extensions, v2" (jan.2000)
 - dropped by MS starting with Win11 22H2
- LCP negotiates CHAP algorithm 0x80 (v1) or 0x81 (v2) for option 3 (Authentication Protocol)
- it's a MS-specific implementation of the CHAP concepts
 - yet supported by many other vendors (e.g. CISCO)

MS-CHAP: extensions over **CHAP**

- similar principles, different protocol
- common (v1 and v2):
 - authenticator-controlled password change
 - authenticator-controlled authentication retry
 - specific failure codes
- MS-CHAPv2 provides mutual authentication:
 - by piggybacking a peer challenge on the Response packet
 - plus an authenticator response on the Success packet
- each peer must know the plaintext password, or an MD4 hash of the password (thus not compatible with most password storage formats)

The MS-CHAPv2 protocol

```
[PEER]
                                                                [AUTHENTICATOR]
Hello >>
                                                   << (Server Challenge, 16 byte) SC
(Client Challenge, 16 byte) CC
(Challenge-Hash) H = sha1( SC || CC || username )[ 0...7 ]
(NT-Hash) K = md4( password )
(Challenge-Response) R = des(K[0...6],H) || des(K[7...13],H) || des(K[14...20],H)
R + H + username >>
                                R1
                                                 R2
                                                                    R3
                                         decrypt R and verify if the result matches H
                                     (NT-Hash-Hash) NHH = md4( md4( password ) )
                                                   (Digest) D = sha1( NHH || R || M1 )
                                   (Authentication-Response) A = sha1( D || H || M2 )
                                                                                << A
```

compute A' and verify if it matches A

M1 = "Magic server to client signing constant"
M2 = "Pad to make it do more than one iteration"

MS-CHAPv2: an attack

- we have a known ciphertext-plaintext pair, R and H
- we need to find the three keys, K[0-6] K[7-13] K[14-20]
- brute-force the pwd? too much time ...
- brute force the keys? 2^56 + 2^56 + 2^56 operations
- but K is just 128 bit (MD4 output) i.e. 16 B
 - so K[14-20] has got only TWO bytes of K[14-15] padded with 0
 - we need 2^16 to find K[14-20]
- to find K[0-6] and K[7-13] we perform only 2^56 operations and then compare the result with R1 and R2
- divide-and-conquer ~ 2^56 ops (<=23 hours with DES FPGA)</p>
- conclusion: MS-CHAPv2 should NEVER be used anymore
 - finally cancelled in Win11 22H2

EAP

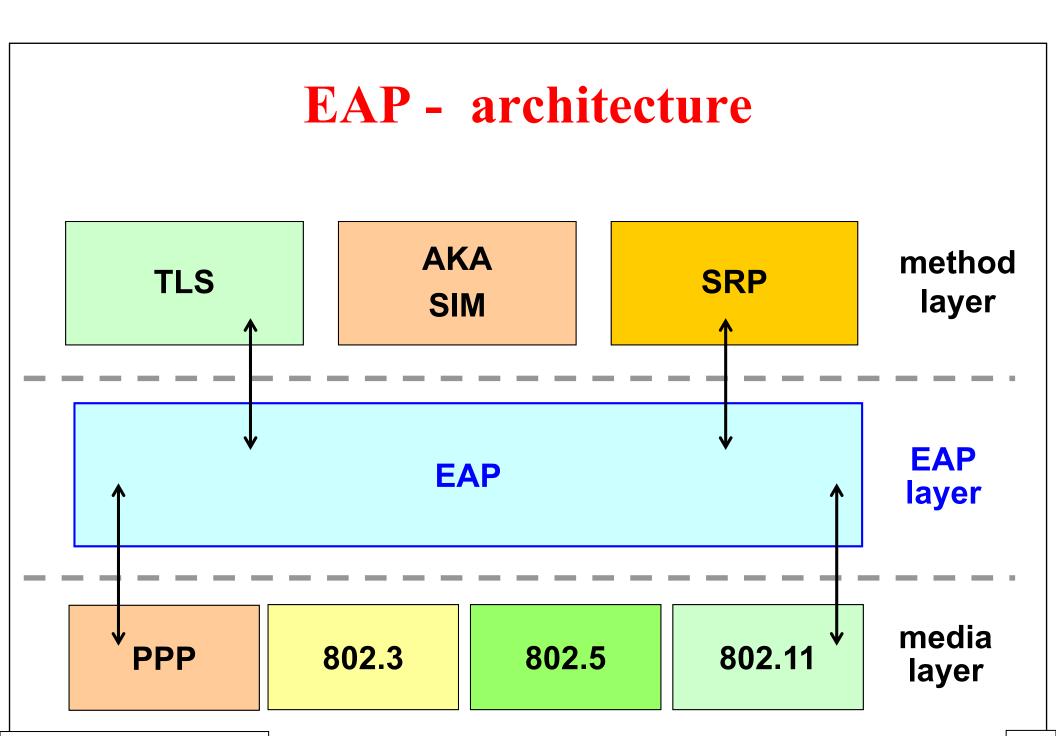
- RFC-3748 (extended by RFC-5247) "PPP Extensible Authentication Protocol (EAP)"
- a flexible L2 authentication framework
- authentication predefined mechanisms:
 - MD5-challenge (similar to CHAP)
 - OTP
 - generic token card
- other mechanisms may be added:
 - RFC-2716 "PPP EAP TLS authentication protocol"
 - RFC-3579 "RADIUS support for EAP"

EAP - encapsulation

- authentication data are transported via a specific EAP encapsulation protocol
 - because L3 packets are not yet available ...
- features of EAP encapsulation:
 - independent of IP
 - supports any link layer (e.g. PPP, 802, ...)
 - explicit ACK/NAK (no windowing)
 - assumes no reordering (PPP guarantees ordering, UDP and raw IP do not!)
 - retransmission (max 3-5 retransmissions)
 - no fragmentation (must be provided by EAP methods for a payload greater than the minimum EAP MTU)

EAP

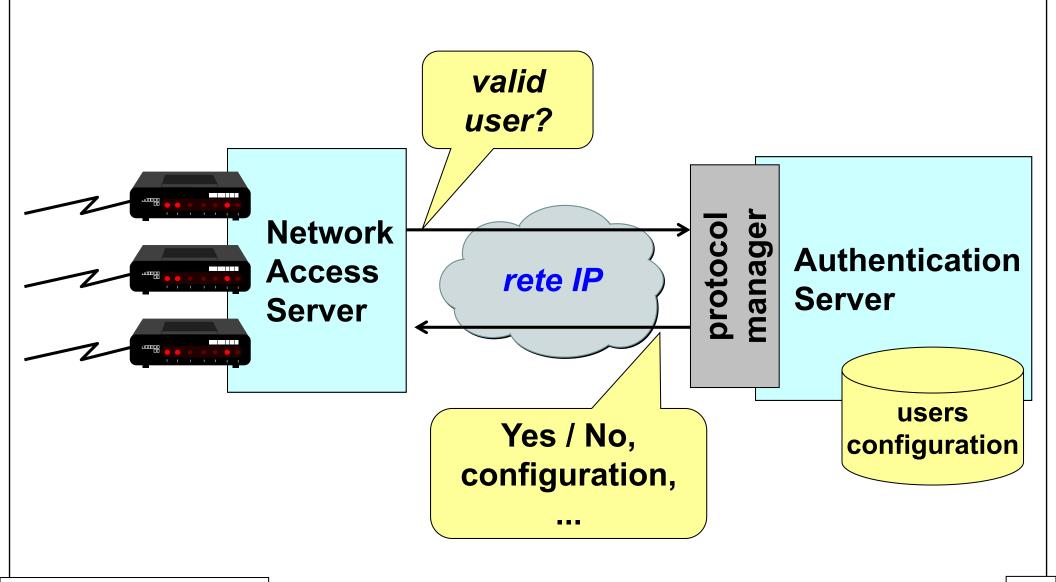
- the link is not assumed to be physically secure
 - EAP methods must provide security on their own
- some EAP methods:
 - EAP-TLS (RFC-5216) TLS mutual authentication
 - EAP-MD5 (RFC-3748) only EAP client authentication
 - EAP-TTLS = tunnelled TLS (to operate any authentication method protected by TLS, e.g. PAP, CHAP)
 - PEAP = TLS tunnel to protect an EAP method
 - EAP-SRP (Secure Remote Password)
 - GSS_API (includes Kerberos)
 - AKA-SIM (RFC-4186, RFC-4187)



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Authentication for network access



AAA

- the NAS manufacturers claim that security needs three functions:
 - Authentication entity's identity is authenticated based on credentials (e.g. password, OTP)
 - Authorization determining whether an entity is authorized to perform a given activity or gain access to resources/services
 - Accounting tracking network resource usage for audit support, capacity analysis or cost billing
- the AS performs exactly these three functions talking with one or more NAS via one or more protocols

Network authentication protocols

RADIUS

- the de-facto standard
- proxy towards other AS

DIAMETER

- evolution of RADIUS
- emphasis on roaming among different ISP
- takes care of security (IPsec, TLS)
- TACACS+ (TACACS, XTACACS)
 - originally technically better than RADIUS, achieved smaller acceptance because it was a proprietary solution (Cisco)

RADIUS

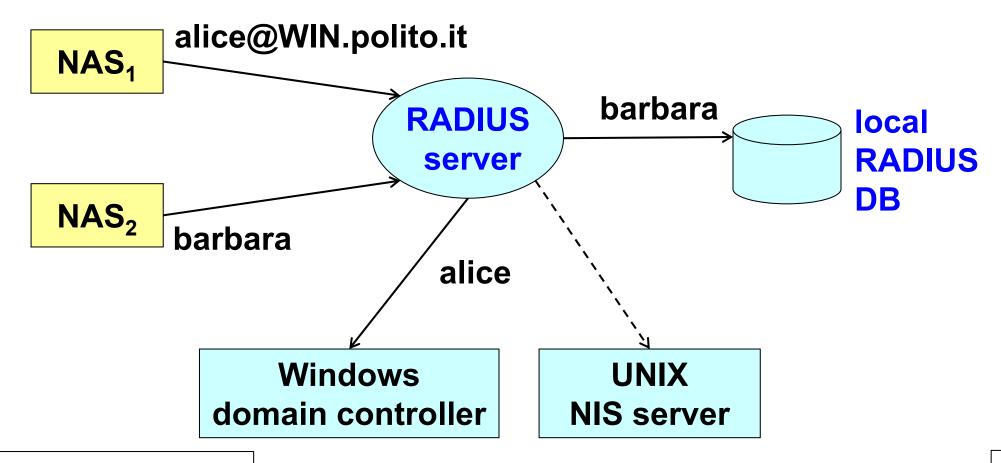
- Remote Authentication Dial-In User Service
- Livingston Technologies (1991) then IETF
- supports authentication, authorization and accounting to control network access:
 - physical ports (analogical, ISDN, IEEE 802)
 - virtual ports (tunnel, wireless access)
- centralized administration and accounting
- client-server schema between NAS and AS
 - port 1812/UDP (authentication) and 1813/UDP (accounting); unofficial ports: 1645 & 1646/UDP
 - timeout + retransmission
 - secondary server

RADIUS - RFC

- RFC-2865 (protocol)
- RFC-2866 (accounting)
- RFC-2867/2868 (tunnel accounting and attributes)
- RFC-2869 (extensions)
- RFC-3579 (RADIUS support for EAP)
- RFC-3580 (guidelines for 802.1X with RADIUS)

RADIUS proxy

the RADIUS server may act as a proxy towards other authentication servers



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Which security functionalities for Radius?

- sniffing NAS req (if contains pwd)
 - confidentiality, privacy
- fake AS resp (to block valid or allow invalid user)
- changing AS resp (Y > N or N > Y)
 - authN & integrity of AS resp
- replay of AS resp (if not properly tied to NAS req)
 - anti-replay of AS resp
- pwd enumeration (from fake NAS)
 - authN of NAS req
- DoS (many NAS req from fake NAS)
 - server scalability

RADIUS: data protection

- packet integrity and authentication via keyed-MD5:
 - key = shared-secret
 - client without key are ignored
- password transmitted "encrypted" with MD5 (after padding with NUL bytes to a multiple of 128 bit):

password

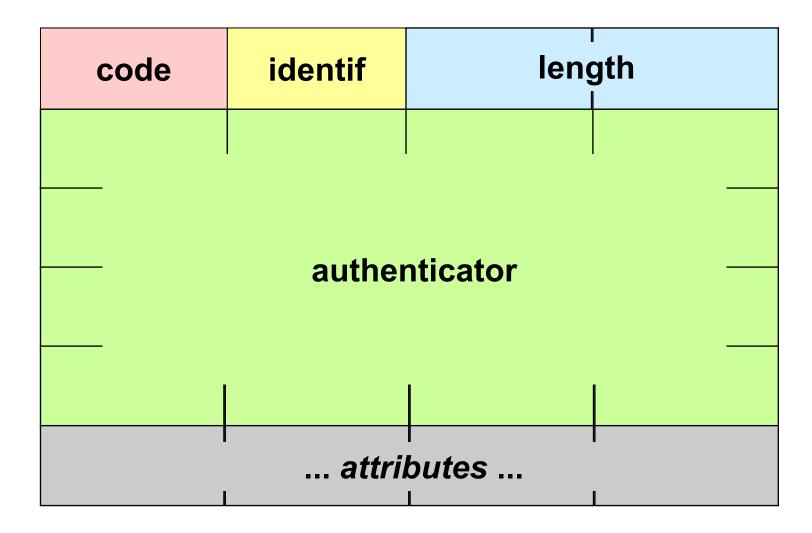
md5(key+authenticator)

RADIUS

- user authentication via PAP, CHAP, token-card and EAP
 - CISCO provides a free server for CryptoCard
 - others support SecurID
- attributes in TLV form, easily extensible without modification to installed base (by ignoring any unknown Type):

attribute type - length - value

RADIUS - format



RADIUS – packet types

ACCESS-REQUEST

contains access credentials (e.g. username + pwd)

ACCESS-REJECT

access is denied (e.g. due to bad username/pwd)

ACCESS-CHALLENGE

requests additional info from the user (e.g. a PIN, token code, secondary password)

ACCESS-ACCEPT (parameters):

- access is granted + network parameters are given
 - for SLIP/PPP: Framed-Protocol, Framed-IP-Address, Framed-IP-Netmask, MS-primary-DNS-server, MS-Secondary-DNS-server, ...
 - for terminal: host, port

RADIUS - authenticator

double purpose:

- server reply authentication and no replay
- masking the password
- in Access-Request:
 - it is named Request Authenticator
 - 16 byte randomly generated by the NAS
- in Access-Accept / Reject / Challenge
 - it is named Response Authenticator
 - it is computed via a keyed-digest:

md5(code | ID | length | RequestAuth | attributes | secret)

RADIUS - some attributes

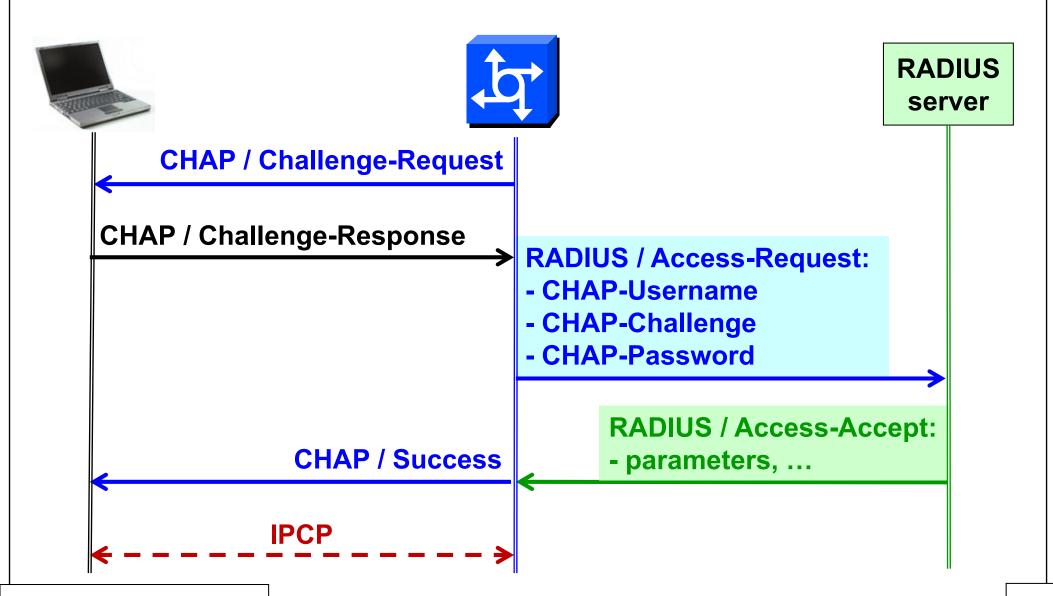
type length value

- type = 1 (User-Name)
 - value = text, network access identifier (NAI), DN
- type = 2 (User-Password)
 - value = password ⊕ md5 (key || RequestAuthent.)
- type = 3 (Chap-Password)
 - value = user CHAP response (128 bit)
- type = 60 (CHAP-Challenge)
 - value = challenge from the NAS to the user

NAI (Network Access Identifier)

- RFC-2486
- NAI = username [@ realm]
- all devices must support NAI up to 72 byte long
- the exact syntax for username and realm is in the RFC (note that only ASCII characters < 128 are allowed, but all of them are allowed)
- note that the username is the one used in the PPP authentication phase (does not necessarily match the application username)

Example - CHAP + RADIUS



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IEEE 802.1x

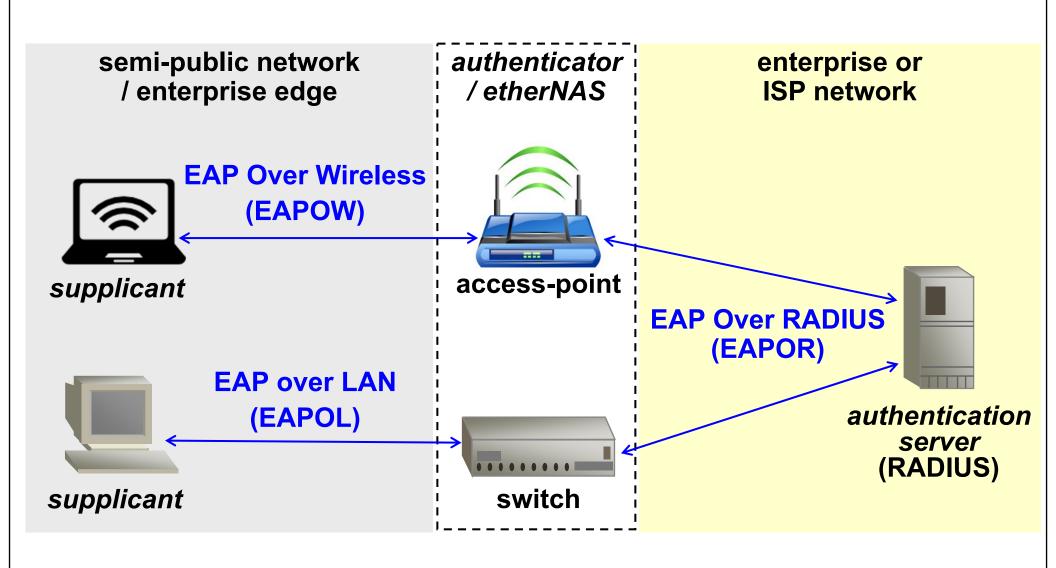
- Port-Based Network Access Control:
 - L2 authentication architecture
 - useful in a wired network to block access
 - absolutely needed in wireless networks
- first implementations (long ago):
 - Windows-XP and Cisco wireless access-points

http://standards.ieee.org/getieee802/download/802.1X-2001.pdf

IEEE 802.1x

- authentication and key-management framework:
 - may derive session keys for use in packet authentication, integrity and confidentiality
 - standard algorithms for key derivation (e.g. TLS, SRP, ...)
 - optional security services (authentication or authentication+encryption)

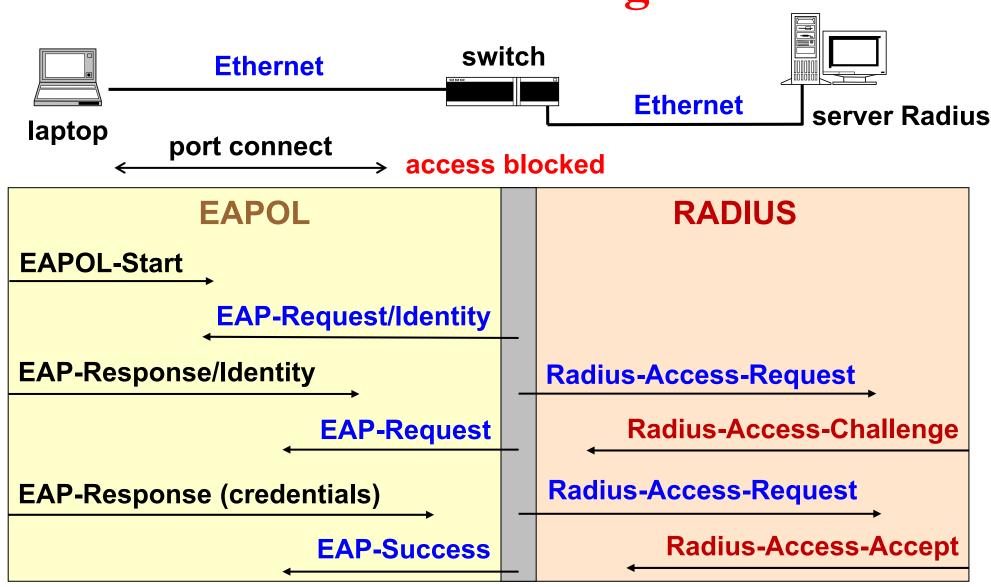
802.1x - architecture



802.1x - advantages

- exploits the application level for the actual implementation of the security mechanisms
 - direct dialogue between supplicant and AS
 - NIC and NAS operate as "pass-through device"
 - no change needed on NIC and NAS to implement new mechanisms
 - perfect integration in AAA





access allowed

eduroam

- WiFi access at research institutes (Italy, Europe, ...)
 - **(21/11/2021) 106 countries**
 - uses 802.1x + RADIUS federation

