## Solving Ax=5: Iterative methods

- · Recap : We have looked at direct methods
  - o Gaussian elimination
  - · LU decomposition -> A-1 -> x = A-15

In principle, exact methods.

- · Atternative approach: Iterative nethods.
  - · Iterate closer and closer to true solution (but not exact)
  - o Often faster, or with smaller mon footprint, than direct nethods when netrices are large.
- o Iterative sheme: Need to think about neg. for the iterative procedure to converge.
- · Methods:
  - facobite method (not to be confused with facobits votation method for  $A\overline{\times} = \lambda \overline{\times}$ )
  - Gauss-Seidel
  - Over-relaxation anethods

- · Checking convergence (when to stop iterations)
- o Typically, monitor the relative change in vector norm

$$\mathcal{E} = \left| \frac{\left| \overline{\mathbf{x}}^{(m+1)} \right|_{1} - \left| \overline{\mathbf{x}}^{(m)} \right|_{1}}{\left| \overline{\mathbf{x}}^{(m)} \right|_{1}} \right|$$

- o Intuitively: stop when x (m+1) is almost identical to x (m) i.e. when it doesn't change anymore.
- o Can use different nouns:

$$|\overline{X}|_{\ell} \equiv \left[\sum_{i=1}^{N} |\chi_{i}|^{\ell}\right]^{\frac{1}{\ell}}$$

$$|X|_1 = |X_1| + |X_2| + ... + |X_N|$$
 (Sum of abs. val.)  
 $|X|_2 = \sqrt{|X|_2 + |X_2|_2^2 + ... + |X_N|^2}$  (Eucl. Length)

$$0 = 0 : |\overline{X}|_{\infty} = \max_{i} (X_{i})$$

(Max element of x)

Intuition:  

$$0.9^{1000} + 1.2^{1000} + 1.19^{1000} \approx 1.2^{1000}$$
  
Formal:  
 $\lim_{l \to \infty} \left[ \chi_1^l + \chi_2^l + ... + \chi_N^l \right]^l = \left[ \max_{l \to \infty} (\chi_i)^l \right]^{l/2}$   
 $= \max_{l \to \infty} (\chi_i)$ 

## The Facosi method (for itemetive solution of Ax= 5)

Summary:

© Decompose A: 
$$A = L + D + U$$

o Iterate to find  $\overline{X}$ :  $\overline{X}^{(m+1)} = -D^{-7}(L+U)\overline{X}^{(m)} + D^{-7}\overline{b}$ 

o Start from some guess  $\overline{X}^{(0)}$ 

· Always converges if A is diagonally dominant, i.e.

$$|a_{ii}| > \sum_{j \neq i} |a_{ij}|$$
 for all vows i

· Rearrange A as sun L+D+U:

· Note: Unveloted to LV decomposition, which was a product A=LU

$$A \overline{x} = (L + D + U) \overline{x} = \overline{b}$$

$$D \overline{x} = -(L + U) \overline{x} + \overline{b}$$

$$\overline{x} = -D'(L + U) \overline{x} + D^{-7} \overline{b}$$

o Treat this as iterative equation, with 
$$\chi^{(m+1)}$$
 on the left-hand side, and  $\chi^{(n)}$  on the right-hand side.

get 
$$D^{-1}$$
 give  $D$  is diagonal:
$$D^{-1} = \operatorname{diag}\left(\frac{1}{d_{11}}, \frac{1}{d_{22}}, \dots, \frac{1}{d_{NN}}\right) = \begin{bmatrix} \frac{1}{d_{11}} & \frac{1}{d_{22}} & 0 \\ \vdots & \vdots & \vdots \\ \frac{1}{d_{NN}} & \vdots & \vdots \\ \frac{1}{d_{N$$

• 
$$4 + 4$$
 example :  $\overline{X} = \begin{bmatrix} x_1 \\ x_2 \\ x_3 \\ x_4 \end{bmatrix}$ 

$$- x_1^{(m+1)} = \left[ b_1 - a_{12} x_2^{(m)} - a_{13} x_2^{(m)} - a_{14} x_4^{(m)} \right] / a_{11}$$

$$= \left( b_{z} - \alpha_{z_{1}} x_{1}^{(m)} - \alpha_{z_{3}} x_{2}^{(m)} - \alpha_{z_{4}} x_{4}^{(m)} \right) / \alpha_{z_{2}}$$

$$= \left[ b_{4} - q_{41} x_{1}^{(n)} - q_{42} x_{2}^{(n)} - q_{43} x_{3}^{(n)} \right] / q_{44}$$

Tu general: 
$$x_i^{(m+1)} = \left[b_i - \sum_{\substack{j=1\\j\neq i}}^{N} a_{ij} \times_j^{(m)}\right] / a_{ii}$$

## Gauss-Seidel method

· Consider 4x4 example for the facobi method, but now use the (m+1) versions of already computed vector elements:

o We've doing forward subst.

· On matrix form we've rewriting 
$$A = (L + D + U) = L$$
 as

$$\mathbb{D}\,\overline{x}=-L\overline{x}-U\overline{x}-\overline{L}$$

and turning it into an iterative equation as

$$D \overline{X}^{(m+1)} = - L \overline{X}^{(m+1)} - U \overline{X}^{(m)} - \overline{L}$$

$$= -(L+D)^{-1} \cdot \sqrt{x^{(m)}} + (L+D)^{-1}$$
but this just looks confusing.

$$X_{i}^{(m+1)} = \begin{bmatrix} -\sum_{j=1}^{i-1} a_{ij} \times_{j}^{(m+1)} - \sum_{j=i+1}^{N} a_{ij} \times_{j}^{(m)} + b_{i} \end{bmatrix} / a_{ii} \begin{cases} -\sum_{j=1}^{N} a_{ij} \times_{j}^{(m+1)} - \sum_{j=i+1}^{N} a_{ij} \times_{j}^{(m)} + b_{i} \end{bmatrix} / a_{ii} \begin{cases} -\sum_{j=1}^{N} a_{ij} \times_{j}^{(m+1)} - \sum_{j=i+1}^{N} a_{ij} \times_{j}^{(m)} + b_{i} \end{bmatrix} / a_{ii} \end{cases}$$

## Successive over-relaxation (SOR)

- · Modified version of Gauss-Seidel, with better
- · But: has a free parameter w that must be chasen/turned for the specific problem.
- · Schematically:

- · For w=7 we get back Gauss-Seidel
- o Can be proven: For  $\omega \in (1,2]$  we have <u>better</u> convergence than G-S, but optimal choice is problem specific.
- o For w > ? , SOR fails. ("Wolbles" out of control ?)
- o Component formulation.

$$X_{i}^{(m+1)} = X_{i}^{(m)} + \frac{\omega}{\alpha_{ii}} \left[ -\sum_{j=1}^{i-1} a_{ij} X_{j}^{(m+1)} - \sum_{j=i+1}^{N} a_{ij} X_{j}^{(m)} + b_{i} - a_{ii} X_{i}^{(m)} \right]$$

i=1,2..., N (evaluated in this order)

$$\overline{X}^{(m+1)} = \left(\omega L + D\right)^{-1} \times \left[-\left(\omega U + (\omega - 1)D\right)\overline{X}^{(m)} + \omega \overline{b}\right]$$

- · Why are iterative methods often useful ?
  - Each iteration only requires matrix-vector multiplication, which has complexity  $O(N^1) \subset O(N^k) \subset O(N^2)$ If A is dense matrix
  - If the method requires M iterations for convergence, we have a total rost of

O(NkM) operations

which is typically much less than the typical  $O(N^2)$  cost of direct methods (assuming number of iterations M is smaller than matrix size N)

Second advantage: Can often get more accurate results
then direct (exact) methods, because
iterative methods are less
susceptible to round-off errors.