Extending propositional separation logic for robustness properties

F·R·I·E·N·D·S of separation logic

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What we will see

An extension of propositional separation logic that

- can express some interesting properties for program verification,
- is `PSpace`-complete,
- has very weak extensions that are `Tower`-hard.

A modal logic on trees that

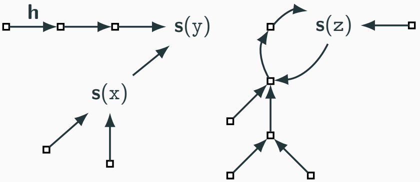
- is `Tower`-complete,
- it is very easily captured by logics that were independently found to be `Tower`-complete.

Memory states

Separation Logic is interpreted over **memory states** $(\mathbf{s}, \mathbf{h})$ where:

- **store**, $\mathbf{s} : \mathtt{VAR} \to \mathtt{LOC}$
- **heap**, $\mathbf{h} : \mathtt{LOC} \to_{\mathbf{fin}} \mathtt{LOC}$

where $\mathtt{VAR} = \{\mathrm{x}, \mathrm{y}, \mathrm{z}, \dots\}$ set of (program) variables,
$\mathtt{LOC}$ set of locations. $\mathtt{VAR}$ and $\mathtt{LOC}$ are countably infinite sets.



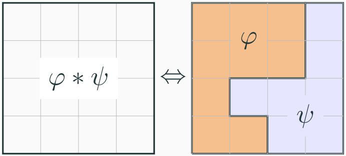
here, $\mathbf{h}(\mathbf{s}(\mathrm{x})) = \mathbf{s}(\mathrm{y})$

- Disjoint heaps: $\mathrm{dom}(\mathbf{h}_1) \cap \mathrm{dom}(\mathbf{h}_2) = \emptyset$
- Union of disjoint heaps $(\mathbf{h}_1 + \mathbf{h}_2)$: union of partial functions.

Propositional Separation Logic $\mathtt{SL}(*, \mathrel{-\!\!*})$

$$\varphi := \neg\varphi \mid \varphi_1 \wedge \varphi_2 \mid \mathtt{emp} \mid \mathtt{x} = \mathtt{y} \mid \mathtt{x} \hookrightarrow \mathtt{y} \mid \varphi_1 * \varphi_2 \mid \varphi_1 \mathrel{-\!\!*} \varphi_2$$

$$(\mathbf{s}, \mathbf{h}) \models \varphi * \psi$$



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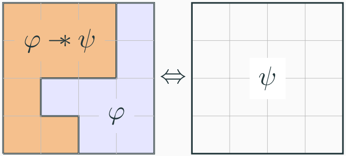


Note: the satisfiability problem SAT($\mathtt{SL}(*, \mathrel{-\!\!*})$) is $\mathtt{PSpace}$-complete.

From where it started

Theorem (Demri, Lozes, M. – 2018, Fossacs)

$\mathtt{SL}(*, -\!\!*)$ *enriched with* $\mathtt{reach}(\mathtt{x}, \mathtt{y}) = 2$ *and* $\mathtt{reach}(\mathtt{x}, \mathtt{y}) = 3$ *is undecidable.*

- reduction from $\mathtt{SL}(\forall, -\!\!*)$ (Brochenin et al.'12)
- $\mathtt{SL}(*, -\!\!*) + \mathtt{reach}(\mathtt{x}, \mathtt{y}) = 2$ is $\mathtt{PSpace}$-complete (Demri et al.'14)

Robustness Properties (Jansen, et al. – ESOP'17)

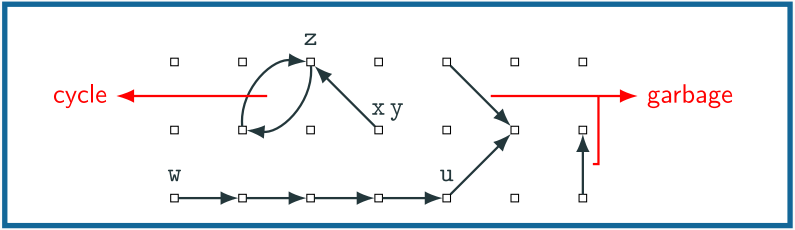
- φ comply with the **acyclicity** property iff every model of φ is acyclic.
- φ comply with the **garbage freedom** property iff in every model $(\mathbf{s}, \mathbf{h}) \models \varphi$, for each $\ell \in \mathrm{dom}(\mathbf{h})$ there is $\mathtt{x} \in \mathsf{v}(\varphi)$ s.t. $\mathbf{s}(\mathtt{x})$ reaches ℓ.

Checking for robustness properties is `ExpTime`-complete for Symbolic Heaps with Inductive Predicates (IP).

Our Goal

Provide a similar result for **propositional** separation logic.

Robustness Properties (Jansen, et al. – ESOP'17)



ic.

ℓ.

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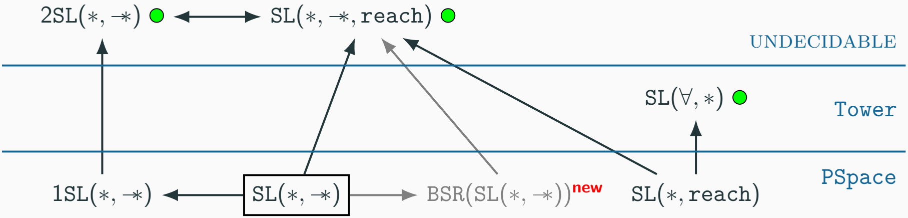
> **Our Goal**
>
> Provide a similar result for **propositional** separation logic.

Desiderata

We aim to an extension of propositional separation logic where

- satisfiability/entailment are decidable in PSpace (as $\mathtt{SL}(*, -\!\!*)$)
- robustness properties reduce to one of these classical problems

Known extensions



Let's start with reachability + 1 quantified variable

- $(\mathbf{s}, \mathbf{h}) \models \mathtt{reach}^{+}(\mathtt{x}, \mathtt{y}) \iff \mathbf{h}^{\mathbf{L}}(\mathbf{s}(\mathtt{x})) = \mathbf{s}(\mathtt{y})$ for some $\mathbf{L} \geq 1$
- $(\mathbf{s}, \mathbf{h}) \models \exists \mathtt{u}\ \varphi \iff$ there is $\ell \in \mathtt{LOC}$ s.t. $(\mathbf{s}[\mathtt{u} \leftarrow \ell], \mathbf{h}) \models \varphi$

It is only possible to quantify over the variable name $\mathtt{u}$.

Robustness properties reduce to entailment

- **Acyclicity**: $\varphi \models \neg \exists \mathtt{u}\ \mathtt{reach}^{+}(\mathtt{u}, \mathtt{u})$
- **Garbage freedom**: $\varphi \models \forall \mathtt{u}\ (\mathtt{alloc}(\mathtt{u}) \Rightarrow \bigvee_{\mathtt{x} \in \mathbf{fv}(\varphi)} \mathtt{reach}(\mathtt{x}, \mathtt{u}))$

where $\mathtt{u} \notin \mathbf{fv}(\varphi)$ and

- $\mathtt{alloc}(\mathtt{x}) \stackrel{\text{def}}{=} (\mathtt{x} \hookrightarrow \mathtt{x}) \twoheadrightarrow \perp$
- $\mathtt{reach}(\mathtt{x}, \mathtt{y}) \stackrel{\text{def}}{=} \mathtt{x} = \mathtt{y} \vee \mathtt{reach}^{+}(\mathtt{x}, \mathtt{y})$

Undecidability and Restrictions

Theorem (Demri, Lozes, M. – 2018, Fossacs)

$\mathtt{SL}(*, \mathrel{-\!\!*})$ *enriched with* $\mathtt{reach}(\mathtt{x}, \mathtt{y}) = 2$ *and* $\mathtt{reach}(\mathtt{x}, \mathtt{y}) = 3$ *is undecidable.*

$\implies$ SAT($1\mathtt{SL}(*, \mathrel{-\!\!*}, \mathtt{reach}^+)$) is undecidable.

We syntactically restrict the logic so that $\mathtt{reach}^+(\mathtt{x}, \mathtt{y})$ is s.t.

R1: it does not appear on the right side of its first $\mathrel{-\!\!*}$ ancestor (seeing the formula as a tree)

- $\varphi \mathrel{-\!\!*} (\psi * \mathtt{reach}^+(\mathtt{u}, \mathtt{u}))$ violates R1

R2: if $\mathtt{x} = \mathtt{u}$ then $\mathtt{y} = \mathtt{u}$ (syntactically)

- $\mathtt{reach}^+(\mathtt{u}, \mathtt{x})$ violates R2

Note: robustness properties are still expressible (formulae as before)!

Results

1. $\mathsf{SAT}(\mathtt{1SL}^{\mathtt{R2}}_{\mathtt{R1}}(*, -\!\!*, \mathtt{reach}^+))$ is $\mathtt{PSpace}$-complete
   - strictly subsumes $\mathtt{1SL}(*, -\!\!*)$ and $\mathtt{SL}(*, \mathtt{reach}^+)$.

2. $\mathsf{SAT}(\mathtt{1SL}_{\mathtt{R1}}(*, -\!\!*, \mathtt{reach}^+))$ is $\mathtt{Tower}$-hard.

Proof Techniques

(1) extend the *core formulae technique* used for $\mathtt{SL}(*, -\!\!*)$.

(2) reduction from "an auxiliary logic on trees".

Core formulae technique

(and a bit of $\mathtt{1SL}^{\mathtt{R2}}_{\mathtt{R1}}(*, \mathrel{-\!\!*}, \mathtt{reach}^{+})$)

First order theories: Gaifman Locality Theorem

Theorem (Gaifman – 1982, Herbrand Symposium)

*Every FO sentence is logically equivalent to a Boolean combination of **local formulae**.*

- application of Ehrenfeucht-Fraïssé games

Relation between models		EF-games		Semantics of logic
$\mathfrak{M} \leftrightarrow_n \mathfrak{M}'$	$\rightleftarrows$	Duplicator has a winning strategy	$\rightleftarrows$	$\mathfrak{M} \approx_n \mathfrak{M}'$
(partial iso. up to n)		(n round game)		(n nested quantifiers)

First order theories: Gaifman Locality Theorem

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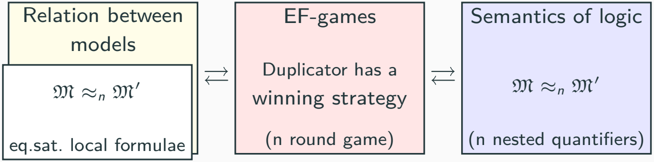

"Locality theorem" for $\mathtt{SL}(*, -\!\!*)$

Theorem (Lozes, 2004 – Space)

Every formula of $\mathtt{SL}(*, -\!\!*)$ is logically equivalent to a Boolean combination of ***core formulae****.*

From this theorem we can get:

- expressive power results
- complexity result (small model property)
- axiomatisation

When considering extensions of the logic, we need to derive new core formulae and reprove the theorem.
$\Longrightarrow$ It does not work (at all) for $\mathtt{1SL}^{\mathtt{R2}}_{\mathtt{R1}}(*, -\!\!*, \mathtt{reach}^+)$.

Core formulae for $\mathtt{SL}(*, -\!\!*)$

Fix $\mathtt{X} \subseteq \mathtt{VAR}$ and $\alpha \in \mathbb{N}^+$

$$\textbf{Core}(\mathtt{X}, \alpha) \stackrel{\text{def}}{=} \left\{ \begin{array}{cc} \mathtt{x} = \mathtt{y}, & \mathtt{x} \hookrightarrow \mathtt{y}, \\ \mathtt{alloc}(\mathtt{x}), & \mathtt{size} \geq \beta \end{array} \;\middle|\; \begin{array}{l} \beta \in [0, \alpha], \\ \mathtt{x}, \mathtt{y} \in \mathtt{X} \end{array} \right\}$$

where $(\mathbf{s}, \mathbf{h}) \models \mathtt{size} \geq \beta$ iff $\mathrm{card}(\mathrm{dom}(\mathbf{h})) \geq \beta$.

- indistinguishability $\boxed{\text{Relation}}$:

  $(\mathbf{s}, \mathbf{h}) \leftrightarrow^{\mathtt{X}}_{\alpha} (\mathbf{s}', \mathbf{h}')$ iff $\forall \varphi \in \textbf{Core}(\mathtt{X}, \alpha)$, $(\mathbf{s}, \mathbf{h}) \models \varphi$ iff $(\mathbf{s}', \mathbf{h}') \models \varphi$

- Both EF-game and winning strategy for Duplicator are hidden inside two (technical) elimination lemmas.

Core formulae: $*$ elimination lemma

Lemma

Suppose $(\boldsymbol{s}, \boldsymbol{h}) \leftrightarrow^{\mathtt{X}}_{\alpha} (\boldsymbol{s}', \boldsymbol{h}')$. *Then,*

for every $\alpha_1 + \alpha_2 = \alpha$ $(\alpha_1, \alpha_2 \in \mathbb{N}^+)$, *and every* $\boldsymbol{h}_1 + \boldsymbol{h}_2 = \boldsymbol{h}$, *(Spoiler)*

there are $\boldsymbol{h}'_1 + \boldsymbol{h}'_2 = \boldsymbol{h}'$ *such that* *(Duplicator)*

$$(\boldsymbol{s}, \boldsymbol{h}_1) \leftrightarrow^{\mathtt{X}}_{\alpha_1} (\boldsymbol{s}', \boldsymbol{h}'_1) \text{ and } (\boldsymbol{s}, \boldsymbol{h}_2) \leftrightarrow^{\mathtt{X}}_{\alpha_2} (\boldsymbol{s}', \boldsymbol{h}'_2).$$

- necessary to obtain a winning strategy for Duplicator

Core formulae: $*$ elimination lemma

Lemma

Suppose $(\boldsymbol{s}, \boldsymbol{h}) \leftrightarrow_{\alpha}^{\mathtt{X}} (\boldsymbol{s}', \boldsymbol{h}')$. *Then,*

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there are $\boldsymbol{h}_1' + \boldsymbol{h}_2' = \boldsymbol{h}'$ *such that* *(Duplicator)*

$$(\boldsymbol{s}, \boldsymbol{h}_1) \leftrightarrow_{\alpha_1}^{\mathtt{X}} (\boldsymbol{s}', \boldsymbol{h}_1') \text{ and } (\boldsymbol{s}, \boldsymbol{h}_2) \leftrightarrow_{\alpha_2}^{\mathtt{X}} (\boldsymbol{s}', \boldsymbol{h}_2').$$

- necessary to obtain a winning strategy for Duplicator

By Relation $\leftrightarrows$ EF-games $\leftrightarrows$ Semantics it leads to:

For every $\varphi \in \mathbf{Bool}(\mathbf{Core}(\mathtt{X}, \alpha_1))$ and $\psi \in \mathbf{Bool}(\mathbf{Core}(\mathtt{X}, \alpha_2))$ there is $\chi \in \mathbf{Bool}(\mathbf{Core}(\mathtt{X}, \alpha_1 + \alpha_2))$ such that

$$\varphi * \psi \iff \chi$$

Note: similar elimination lemma for $-\!\!*$.

Core formulae: after $*$ and $\mathrel{-\!\!*}$ elimination

Theorem

For every φ in $\mathtt{SL}(*, \mathrel{-\!\!*})$:

1. *there is en equivalent Boolean combination of core formulae.*
2. *for every $\alpha \geq |\varphi|$, $\mathtt{X} \supseteq \mathrm{v}(\varphi)$ and $(\boldsymbol{s}, \boldsymbol{h}) \leftrightarrow^{\mathtt{X}}_{\alpha} (\boldsymbol{s}', \boldsymbol{h}')$,*

$$(\boldsymbol{s}, \boldsymbol{h}) \models \varphi \text{ iff } (\boldsymbol{s}', \boldsymbol{h}') \models \varphi.$$

[2] allows to derive a small-model property which leads to a proof that SAT($\mathtt{SL}(*, \mathrel{-\!\!*})$) is in $\mathtt{PSpace}$.

$\mathtt{1SL}^{\mathtt{R2}}_{\mathtt{R1}}(*, \mathrel{-\!\!*}, \mathtt{reach}^+)$ is in $\mathtt{PSpace}$: Not so easy...

$$\pi \;\coloneqq \mathtt{x} = \mathtt{y} \;\mid\; \mathtt{x} \hookrightarrow \mathtt{y} \;\mid\; \mathtt{emp} \;\mid\; \underline{\mathcal{A} \mathrel{-\!\!*} \mathcal{C}\ (\mathtt{R1})}$$

$$\mathcal{C} \;\coloneqq \pi \;\mid\; \mathcal{C} \wedge \mathcal{C} \;\mid\; \neg\mathcal{C} \;\mid\; \exists \mathtt{u}\, \mathcal{C} \;\mid\; \mathcal{C} * \mathcal{C}$$

$$\mathcal{A} \;\coloneqq \pi \;\mid\; \underline{\mathtt{reach}^+(v_1, v_2)} \;\mid\; \mathcal{A} \wedge \mathcal{A} \;\mid\; \neg\mathcal{A} \;\mid\; \exists \mathtt{u}\, \mathcal{A} \;\mid\; \mathcal{A} * \mathcal{A}$$

where if $v_1 = \mathtt{u}$ then $v_2 = \mathtt{u}$ $(\mathtt{R2})$.

- Asymmetric $\mathcal{A} \mathrel{-\!\!*} \mathcal{C}$: design two sets of core formulae against
  - two $*$ and two $\exists$ elimination lemmas;
  - one $\mathrel{-\!\!*}$ elimination lemma that glues the two set of core formulae.

- instead of "$\mathtt{size} \geq \beta$ s.t. $\beta \in [1, \alpha]$", the βs of new core formulae are bounded by functions on α, e.g.

  $$\#\mathtt{loop}(\beta) \geq \gamma \qquad \gamma \in [1, \tfrac{1}{2}\alpha(\alpha + 3) - 1]$$

  bounds are found by solving a set of recurrence equations.

Core formulae: Example on a toy logic

$$\varphi := \neg\varphi \mid \varphi_1 \wedge \varphi_2 \mid \varphi_1 * \varphi_2 \mid \exists \mathtt{u}\ \varphi \mid \mathtt{alloc}(\mathtt{u}) \mid \mathtt{reach}^+(\mathtt{u},\mathtt{u})$$

Some formulae expressible in this logic:

- $\mathtt{size} \geq 0 \stackrel{\text{def}}{=} \top \qquad \mathtt{size} \geq \beta + 1 \stackrel{\text{def}}{=} \exists \mathtt{u}\ (\mathtt{alloc}(\mathtt{u}) * \mathtt{size} \geq \beta)$
- $\mathtt{reach}^+(\mathtt{u},\mathtt{u}){=}\beta$ iff there is a loop of size exactly β involving $\mathbf{s}(\mathtt{u})$.
- $\#\mathtt{loops}(\beta) \geq \gamma \stackrel{\text{def}}{=} \overbrace{\exists \mathtt{u}\, \mathtt{reach}^+(\mathtt{u},\mathtt{u}){=}\beta * \ldots * \exists \mathtt{u}\, \mathtt{reach}^+(\mathtt{u},\mathtt{u}){=}\beta}^{\gamma-1 \text{ times } *}$
- $\mathtt{rem} \geq \beta$ iff there are at least β memory cells not in a loop.

Designing Core Formulae

- Fix $\alpha \in \mathbb{N}^+$
- Let **Core**(α) be the **finite** set of predicates:

$$\left\{ \begin{array}{l|l} \texttt{rem} \geq \beta, & \\ \#\texttt{loops}(\beta) \geq \gamma, & \begin{array}{l} \beta \in [1, \mathcal{R}(\alpha)], \\ \gamma \in [1, \mathcal{L}(\alpha)] \end{array} \\ \#\texttt{loops}_{>\mathcal{R}(\alpha)} \geq \gamma, & \end{array} \right\}$$

for some functions $\mathcal{L}$ and $\mathcal{R}$ in $[\mathbb{N} \to \mathbb{N}]$.

$\#\texttt{loops}_{>\beta} \geq \gamma \;=\; \exists \mathtt{u}\, \mathtt{reach}^+(\mathtt{u},\mathtt{u}) {\geq} \beta + 1 * \ldots * \exists \mathtt{u}\, \mathtt{reach}^+(\mathtt{u},\mathtt{u}) {\geq} \beta + 1$

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  for some functions $\mathcal{L}$ and $\mathcal{R}$ in $[\mathbb{N} \to \mathbb{N}]$.

These formulae induce a partition on the heap:

- $\texttt{rem} \geq \beta$ speaks about memory cells not in a loop
- $\#\texttt{loops}(\beta) \geq \gamma$ speaks about locations in loops of size $\beta \in [1, \mathcal{R}(\alpha)]$
- $\#\texttt{loops}_{>\mathcal{R}(\alpha)} \geq \gamma$ speaks about locations in loops of size $> \mathcal{R}(\alpha)$.

$\#\texttt{loops}_{>\beta} \geq \gamma \ = \ \exists \texttt{u}\, \texttt{reach}^+(\texttt{u},\texttt{u}) \geq \beta + 1 * \ldots * \exists \texttt{u}\, \texttt{reach}^+(\texttt{u},\texttt{u}) \geq \beta + 1$

Find $\mathcal{R}$ and $\mathcal{L}$

Lemma

Suppose $(\boldsymbol{s}, \boldsymbol{h}) \leftrightarrow^{\mathtt{X}}_{\alpha} (\boldsymbol{s}', \boldsymbol{h}')$. *Then,*

for every $\alpha_1 + \alpha_2 = \alpha$ *(*$\alpha_1, \alpha_2 \in \mathbb{N}^+$*), and every* $\boldsymbol{h}_1 + \boldsymbol{h}_2 = \boldsymbol{h}$*,* *(Spoiler)*

...

- Test the core formulae against the $*$ elimination lemma.
- standard-ish way of doing things in EF-games.

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What happens to the locations corresponding to $\mathtt{rem} \geq \beta$,
when we split a heap?

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They correspond to $\mathtt{rem} \geq \beta$, also in the subheaps.

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...

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- st

$$\mathcal{R} \circlearrowleft +$$

$$\mathcal{R}(\alpha) \geq \max_{\substack{\alpha_1, \alpha_2 \in \mathbb{N}^+ \\ \alpha_1 + \alpha_2 = \alpha}} \big(\mathcal{R}(\alpha_1) + \mathcal{R}(\alpha_2)\big)$$

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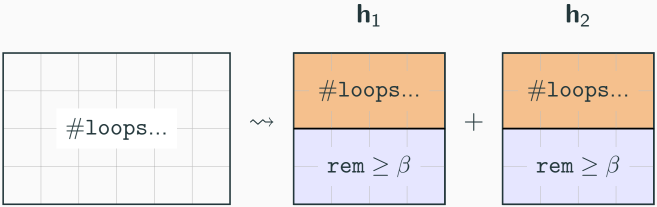
Find $\mathcal{R}$ and $\mathcal{L}$

For $\mathcal{L}$, roughly speaking...



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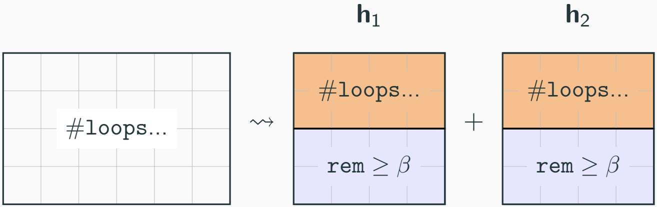


$$\mathcal{L}(\alpha) \geq \max_{\substack{\alpha_1,\alpha_2 \in \mathbb{N}^+ \\ \alpha_1+\alpha_2=\alpha}} \left(\mathcal{L}(\alpha_1) + \mathcal{L}(\alpha_2) + \mathcal{R}(\max(\alpha_1,\alpha_2))\right)$$

Find $\mathcal{R}$ and $\mathcal{L}$

We have the inequalities

$$\mathcal{R}(1) \geq 1 \qquad \mathcal{R}(\alpha) \geq \max_{\substack{\alpha_1,\alpha_2 \in \mathbb{N}^+ \\ \alpha_1+\alpha_2=\alpha}} (\mathcal{R}(\alpha_1) + \mathcal{R}(\alpha_2))$$

$$\mathcal{L}(1) \geq 1 \qquad \mathcal{L}(\alpha) \geq \max_{\substack{\alpha_1,\alpha_2 \in \mathbb{N}^+ \\ \alpha_1+\alpha_2=\alpha}} (\mathcal{L}(\alpha_1) + \mathcal{L}(\alpha_2) + \mathcal{R}(\alpha_1) + \mathcal{R}(\alpha_2))$$

Which admit $\mathcal{R}(\alpha) = \alpha$ and $\mathcal{L}(\alpha) = \frac{1}{2}\alpha(\alpha+1)$ as a solution.

To satisfy the $*$ elimination lemma, build $\leftrightarrow^{\mathtt{X}}_{\alpha}$ w.r.t.

$$\left\{ \begin{array}{l|l} \mathtt{rem} \geq \beta, & \beta \in [1, \alpha], \\ \#\mathtt{loops}(\beta) \geq \gamma, & \\ \#\mathtt{loops}_{>\alpha} \geq \gamma, & \gamma \in [1, \dfrac{1}{2}\alpha(\alpha+1)] \end{array} \right\}$$

(it is not a solution for the toy logic, we forgot the variable $\mathtt{u}$!)

First recap

$\mathtt{SL}(*, \mathrel{-\!\!*}, \mathtt{reach})$
undecidable

$\mathtt{1SL}_{\mathtt{R1}}(*, \mathrel{-\!\!*}, \mathtt{reach}^{+})$
unknown

$\mathtt{1SL}_{\mathtt{R1}}^{\mathtt{R2}}(*, \mathrel{-\!\!*}, \mathtt{reach}^{+})$
PSpace-complete

$\mathtt{1SL}(*, \mathrel{-\!\!*})$
PSpace-complete

$\mathtt{SL}(*, \mathtt{reach})$
PSpace-complete

- $\mathtt{1SL}_{\mathtt{R1}}^{\mathtt{R2}}(*, \mathrel{-\!\!*}, \mathtt{reach}^{+})$ strictly generalise other PSpace-complete extensions of propositional separation logic.
- It can be used to check for robustness properties.

ALT: An auxiliary logic on trees

(or, what happens if we allow $\mathtt{reach}^+(\mathtt{u}, \mathtt{x})$)

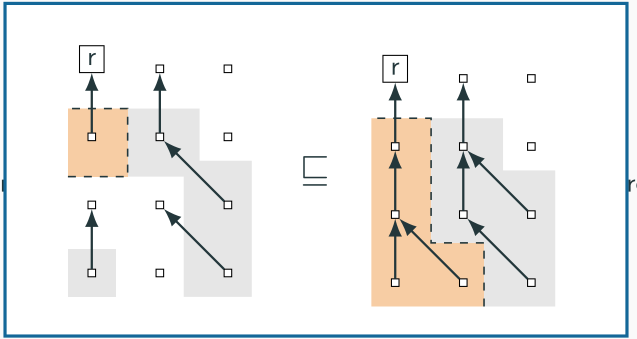
Auxiliary logic on trees (ALT)

$$\varphi := \varphi_1 \land \varphi_2 \;\mid\; \neg\varphi \;\mid\; \langle \mathrm{U} \rangle \varphi \;\mid\; \blacklozenge\, \varphi \;\mid\; \blacklozenge^{*} \varphi \;\mid\; \triangle \;\mid\; \oslash$$

- interpreted on *acyclic heaps* (finite forests, encoding parent relation)
- one *current node* $\mathsf{n} \in \mathtt{LOC}$, one fixed *target node* $\mathsf{r} \in \mathtt{LOC}$
- $\mathbf{h}, \mathsf{n} \models_{\mathsf{r}} \langle \mathrm{U} \rangle \varphi$ iff there is $\mathsf{n}' \in \mathtt{LOC}$ s.t. $\mathbf{h}, \mathsf{n}' \models_{\mathsf{r}} \varphi$
- $\mathbf{h}, \mathsf{n} \models_{\mathsf{r}} \triangle$ iff $\mathsf{n} \in \mathrm{dom}(\mathbf{h})$ and n reaches r in at least one step
- $\mathbf{h}, \mathsf{n} \models_{\mathsf{r}} \oslash$ iff $\mathsf{n} \in \mathrm{dom}(\mathbf{h})$ and n **does not** reach r in at least one step
- $\blacklozenge\, \varphi \equiv (\mathtt{size} = 1) * \varphi$, $\qquad \blacklozenge^{*} \varphi \equiv \top * \varphi$

We prove that $\mathrm{SAT}(\mathrm{ALT})$ is a $\mathtt{Tower}$-complete problem.

Auxiliary logic on trees (ALT)



- inter relation)
- one
- **h**, n
- **h**, n $\models_{\mathsf{r}} \triangle$ iff n $\in \mathrm{dom}(\mathbf{h})$ and n reaches r in at least one step
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- $\blacklozenge\, \varphi \equiv (\mathtt{size} = 1) * \varphi$, $\quad \blacklozenge^{*} \varphi \equiv \top * \varphi$

We prove that SAT(ALT) is a $\mathtt{Tower}$-complete problem.

What can ALT do?

Given a pointed model $(\mathbf{h}, \mathsf{n})$ and a target node r:

If we consider a portion of $\mathbf{h}$ with domain in $\{\mathsf{n}' \in \mathtt{LOC} \mid \mathbf{h}, \mathsf{n}' \models \oslash\}$, ALT **can only express** size bounds.

- Proof done with EF-games for ALT.

$$
\begin{aligned}
\mathtt{size}(\oslash) \geq 0 \quad &\stackrel{\text{def}}{=} \quad \top \\
\mathtt{size}(\oslash) \geq \beta{+}1 \quad &\stackrel{\text{def}}{=} \quad \langle \mathrm{U} \rangle \big(\oslash \wedge \blacklozenge(\neg \mathtt{alloc} \wedge \mathtt{size}(\oslash) \geq \beta)\big)
\end{aligned}
$$

where $\mathtt{alloc} \stackrel{\text{def}}{=} \oslash \vee \triangle$.

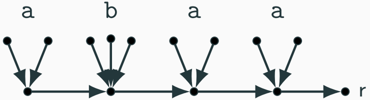
What can ALT do?

- If $\mathbf{h}, \mathsf{n} \models_{\mathsf{r}} \triangle$, ALT can check bounds on the number of descendants and children of n:

$$
\begin{aligned}
\#\mathtt{desc} \geq \beta &\stackrel{\text{def}}{=} \blacklozenge^{*}\big([\mathrm{U}]\neg\obslash \wedge \triangle \wedge \blacklozenge(\neg\mathtt{alloc} \wedge \mathtt{size}(\obslash) \geq \beta)\big)\\
\#\mathtt{child} \geq 0 &\stackrel{\text{def}}{=} \top\\
\#\mathtt{child} \geq \beta{+}1 &\stackrel{\text{def}}{=} \#\mathtt{desc} \geq \beta{+}1 \wedge \neg\blacklozenge^{\beta}(\triangle \wedge \neg\#\mathtt{desc} \geq 1)
\end{aligned}
$$

- Easy to encode words as acyclic memory states

$\boxed{\mathtt{abaa}}$ $\rightsquigarrow$



PITL **(Moszkowski'83)**

$$\varphi \coloneqq \mathtt{pt} \;\mid\; \mathtt{a} \;\mid\; \varphi_1 \,\mathbf{|}\, \varphi_2 \;\mid\; \neg\varphi \;\mid\; \varphi_1 \wedge \varphi_2$$

- interpreted on finite non-empty words over a finite alphabet Σ
- $\mathfrak{w} \models \mathtt{pt} \iff |\mathfrak{w}| = 1$
- $\mathfrak{w} \models \mathtt{a} \iff$ first letter of $\mathfrak{w}$ is $\mathtt{a} \in \Sigma$ (locality principle)
- $\mathfrak{w} \models \varphi_1 \,\mathbf{|}\, \varphi_2 \iff \mathfrak{w}[1:\mathtt{j}] \models \varphi_1$ and $\mathfrak{w}[\mathtt{j}:|\mathfrak{w}|] \models \varphi_2$
  for some $\mathtt{j} \in [1, |\mathfrak{w}|]$

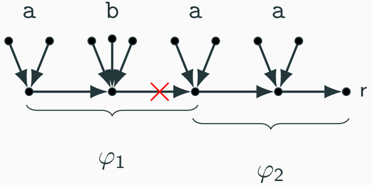
$$\boxed{\mathfrak{w}_1 \ldots \mathfrak{w}_{\mathtt{j}-1} \quad \mathfrak{w}_{\mathtt{j}} \quad \mathfrak{w}_{\mathtt{j}+1} \ldots \mathfrak{w}_{|\mathfrak{w}|}}$$

φ_1 φ_2

Note: $\mathrm{SAT}(\mathrm{PITL})$ is $\mathtt{Tower}$-complete.

Reducing PITL **to** ALT

- Set of models encoding words can be characterised in ALT
- However, difficult to translate $\varphi_1 \mathbin{|} \varphi_2$!



After the cut, left side does not reach r anymore.

$\implies$ nodes on the left side satisfy $\obslash$

$\implies$ We cannot express the satisfaction of φ_1.

PITL **to** ALT**: alternative semantics for** PITL

- $\boxed{\mathsf{a}}$ marked representation of $\mathsf{a} \in \Sigma$

$$\boxed{\mathfrak{w}_1 \dots \mathfrak{w}_{j-1}\ \mathfrak{w}_j\ \mathfrak{w}_{j+1} \dots \boxed{\mathfrak{w}_{|\mathfrak{w}|}}}$$

- $\varphi \,\boldsymbol{|}\, \psi$ on standard semantics:

$$\underbrace{\boxed{\mathfrak{w}_1 \dots \mathfrak{w}_{j-1}\ \mathfrak{w}_j}}_{\varphi_1} \qquad \underbrace{\boxed{\mathfrak{w}_j\ \mathfrak{w}_{j+1} \dots \mathfrak{w}_{|\mathfrak{w}|}}}_{\varphi_2}$$

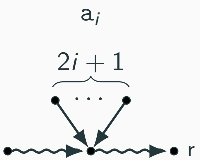
- $\varphi \,\boldsymbol{|}\, \psi$ on marked semantics

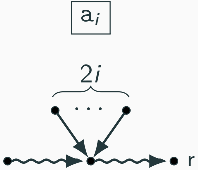
$$\boxed{\underbrace{\mathfrak{w}_1 \dots \mathfrak{w}_{j-1}\ \boxed{\mathfrak{w}_j}}_{\varphi_1}\ \mathfrak{w}_{j+1} \dots \boxed{\mathfrak{w}_{|\mathfrak{w}|}}} \qquad \underbrace{\boxed{\mathfrak{w}_j \mathfrak{w}_{j+1} \dots \boxed{\mathfrak{w}_{|\mathfrak{w}|}}}}_{\varphi_2}$$

- alternative semantics is equivalent to the original one.

ALT, marking an element

- Given an alphabet $\Sigma = \{\mathtt{a}_1, \ldots, \mathtt{a}_n\}$, $\mathtt{a}_i$ and $\boxed{\mathtt{a}_i}$ are encoded as





$\Longrightarrow$ marking a character $\sim$ removing a single child.

- SAT(PITL) can be reduced to SAT(ALT),
  (translated formula is in 2ExpSpace if Σ is coded in binary)

$\Longrightarrow$ ALT is Tower-complete (upper-bound from MSO).

Some logics that are $\texttt{Tower}$-hard

- It is easy to see that ALT is a fragment of $\texttt{1SL}_{\texttt{R1}}(*, \mathrel{-\!\!*}, \texttt{reach}^+)$:

  fix $\texttt{x} \in \texttt{VAR}$ to play the role of the target node r,

  $$\langle \mathrm{U} \rangle \varphi \equiv \exists \texttt{u}\ \varphi \qquad \triangle \equiv \texttt{reach}^+(\texttt{u}, \texttt{x}) \qquad \oslash \equiv \texttt{alloc}(\texttt{u}) \wedge \neg \triangle$$

  $+$ impose acyclic heaps: $\neg \exists \texttt{u}\ \texttt{reach}^+(\texttt{u}, \texttt{u})$.

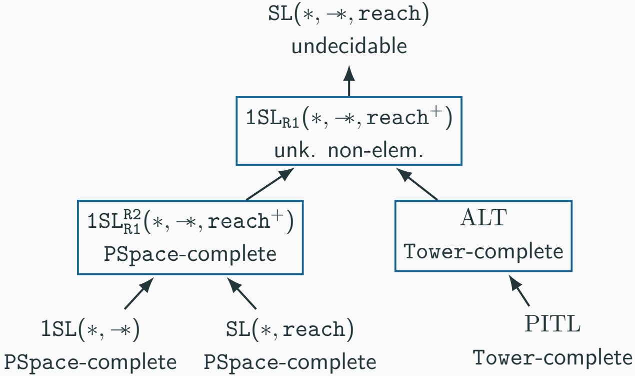
- ALT is a fragment of $\texttt{MSL}(*, \Diamond, \langle \mathrm{U} \rangle)$

- $\mathrm{ALT} \preceq_{\mathsf{SAT}} \texttt{MLH}(*, \Diamond, \langle \mathrm{U} \rangle)$ with modal depth 2.
  (then $*$, $\exists \texttt{u}$, $\texttt{alloc}(\texttt{u})$, $\texttt{alloc}^2(\texttt{u})$ is $\texttt{Tower}$-c.)

- $\mathrm{ALT} \preceq_{\mathsf{SAT}} \texttt{QCTL}(\texttt{U})$ without imbricated until operators $\texttt{U}$
  (or $\texttt{QCTL}(\texttt{EF})$ with 2 imbrication of $\texttt{EF}$)

  Note: in these results $*$ can always be replaced with $\blacklozenge$ and $\blacklozenge^*$.

Second Recap



- ALT improves the understanding of some Tower-complete logics.
- It seems to be an interesting tool to prove Tower-hardness.