

ASSIGNMENT 4 – COMP 252

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1. BROWSING THE SMALL ELEMENTS IN A RED-BLACK TREE.

2. GREEDY ALGORITHM. On a flat table, we have placed n disks of radii r_1, \dots, r_n , numbered from left to right. We push them together without creating overlap, as in the figure below. Give an $O(n)$ time algorithm to compute the size of the smallest axis-aligned rectangle that can hold the disks.

Solution:

Algorithm 1: Greedy circle packing**Input:** An ordered list $\Omega := \{1, 2, 3, \dots, n\}$ of n circles.**Output:** The minimum width of a rectangle that can hold the disks.

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// returns the radius of the circle a
1 Function Radius(circle a):
2   return radius of a;

// returns the distance between the centres of a and b if they are pushed
// together.
3 Function centre_dist(circle a, circle b):
4   return  $\sqrt{(\text{Radius}(a) + \text{Radius}(b))^2 - (\text{Radius}(a) - \text{Radius}(b))^2}$ ;

// largest subarray of  $\Omega$  ending with  $k$ , decreasing in radii.
5 decreasing_radii[]  $\leftarrow$  new array of circles;

// distance from the left side of the rectangle to circle at index
6 left_to_circle[]  $\leftarrow$  new array;
7 left_to_circle[0]  $\leftarrow$  0;
8 for  $i \leftarrow 1$  to  $n$  do
9   if  $i = 1$  then
10    left_to_circle[ $i$ ]  $\leftarrow$  Radius( $i$ );
11    decreasing_radii[ $i$ ]  $\leftarrow$   $i$ ;
12   else
13    max_d  $\leftarrow$  0;
14    // iterate through array of decreasing radii, starting with the
    // smallest, until larger radius is hit
15    for  $j \leftarrow \text{length}(\text{decreasing\_radii}) - 1$  to  $j \leftarrow 0$  do
16      // get the distance from the left side of the rectangle to the
      // centre of circle  $i$  when it is adjacent to circle at
      // decreasing_radii [ $j$ ]
17      max_d  $\leftarrow$  max{max_d, left_to_circle[decreasing_radii[ $j$ ]] +
      // centre_dist(decreasing_radii[ $j$ ],  $i$ )};
18      if Radius( $i$ )  $\geq$  Radius( $j$ ) then
19        // if we encounter a circle of larger radius than  $i$ , stop
        // checking distances, as there would clearly be an overlap.
20        exit for loop;
21    // max_d holds the maximum distance occurring when circle  $i$  is
    // adjacent to a circle in decreasing_radii. This
22    left_to_circle[ $i$ ]  $\leftarrow$  max_d;
    // update array of decreasing radii.
23    while Radius( $i$ )  $\geq$  Radius(decreasing_radii[ $-1$ ]) do
24      remove(decreasing_radii[ $-1$ ]);
25    decreasing_radii[length(decreasing_radii)]  $\leftarrow$   $i$ ;
26 return left_to_circle

```
