Cryptography: Vigenère and Kasiski

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- ➤ Simple ciphers of either type were quite easily broken: substitution ciphers suffered from attacks using **frequency analysis**.
- ▶ The Vigenère cipher was developed and was thought to tackle these problems.

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- ▶ The alphabets form a square: the Vigenère square, or tabula recta.



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key	С	0	D	Е	S	С	0	D	Е	S	С	0	D	Ε	S	С	0	D	Е	S
plain	i	n	v	a	d	е	a	t	d	a	W	n	0	n	f	r	i	d	a	у
cipher	K	В	Y	E	V	G	0	W	Н	S	Y	В	R	R	X	T	W	G	Е	Q

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cipher	K	В	Y	Е	V	G	0	W	Н	S	Y	В	R	R	Х	T	W	G	Е	Q

Can also be thought of in terms of modular arithmetic:

keyword	2	14	3	4	18	2	14	3	4	18	2	14	3
plaintext	8	13	21	0	3	4	0	19	3	0	22	13	14
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► Encryption is achieved by performing the addition

$$c_i = m_i + k_i \mod 26$$
,

where m_i is the *i*th letter of the message and k_i is the *i*th letter of the keyword repeated over and over again.

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- Our simple frequency analysis will not work any more and this type of cipher was thought unbreakable right up until the end of World War 1 (1914-1918).
- ▶ A few people were involved in the cracking of this cipher, including quite a famous name in computing history.

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 - ► ANYVGYSTYNRPLWH...
 - These letters must have used the same row of the Vigenère square.

Breaking Vigenère

ZSHRSNAYEHVRHIUIZZQZXHWEFLXPOJFCXEFJAJMLSEURXXSVZXAGSEFYKCHYMXMLWJISKPRN MWUIWESATXQYQHDISEXCTRRTXSLIZPNCBRHVXPBKSEOILKFVMXXVHYMRFEBJMRWCSKMWFSFK MPTWVZESPRHYMXTWAVZFNWWVPXAIAJQPOIGRNSNXHYQMKZOIUSNWQFZGXVBJFLXCKVDILGFL FMGMGVPEGHGKGHBIRGQVAEDJMPFSGKMWGEFIAAECOJMQTRKZFLTQWTDSLGCGQQQBKVKEGKYHZ ZMLIHYQXKEBJUIGXQIQEMYFVEXAEHJIEKQQEPQNPHZBPRMBRPVHTCWIEMIFNUXAMBWURBXST AQIPOTQRVCAVZAXRHKAEGHTIASOIFKTMLKZFNITFCLFXAIWIXMMXZVMJYEWIEWXVSEQMGXVV UVTWGLDEGGSFRXAIWIQQIMFVAZXVARFXXVWKUWISGJUFEIHYMXMLSZZJNWCIUENRRVDXAIAZ OVHWQFBIWSHYQWTQSEASGIURHITXVFGKAXHFFLXSZUQVPSFCPWHJGGMGXEGJAYKGSJAJAYAR ZHTRUVDSKXVFGKAXCWFLXQCEXCMSRZEQBWGKTIBHSRAJEMTVGTHRHYQQTWWDBSLWWSXIHVWD BVHFOSXIBXWJOYKMCLEXHVSTMPEWCDQSYXVVYIGXOCTEUMHJAJMLCJQHXTOIFIWHOPEEMQCJ FXXVFVEXKMOCYIGJOEOMXHHYQVXQWXTXUICKTIKQSEGTHRARDWIIFYMTLMBWQVBSFKAXAIAJ OPOIGRZHKIOUKXHASCOSFIODUWLMCEMVRIBKQVIVWJQCXXOTDSLWHYQKNPTFRWIEQVYMGHGK TEMEFVFSHYFDURWWOJAYKWOIQXHXVFEIHJHYQFXEGKEXAEHGQVBWVZZXXPZVOXLZOJFEGHQF APTRRLZWRQDRFLXXWTDIZEFUQHMLWJQEKXVNUXAIBMUSNWSPQWTRRJXSPPMRZHLYFVXCWVSN FI.XMFGXFGWOXMMGWHI.F.

Assume we know the key length: k = 6. We'll look into how to figure this out later.

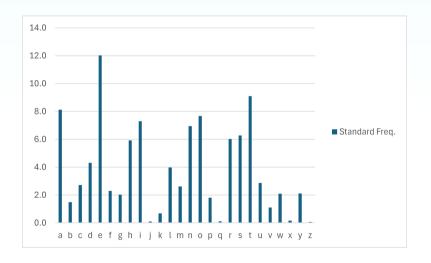
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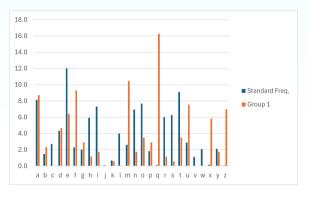
ZSHRSNAYEHVRHIUIZZQZXHWEFLXPOJFCXEFJAJMLSEURXXSVZXAGSEFYKCHYMXMLWJISKPRN MWUIWESATXQYQHDISEXCTRRTXSLIZPNCBRHVXPBKSEOILKFVMXXVHYMRFEBJMRWCSKMWFSFK MPTWVZESPRHYMXTWAVZFNWWVPXAIAJQPOIGRNSNXHYQMKZOIUSNWQFZGXVBJFLXCKVDILGFL FMGMGVPEGHGKGHBIRGQVAEDJMPFSGKMWGEFIAAECOJMQTRKZFLTQWTDSLGCGQQBKVKEGKYHZ ZMLIHYQXKEBJUIGXQIQEMYFVEXAEHJIEKQOEPQNPHZBPRMBRPVHTCWIEMIFNUXAMBWURBXST AQIPOTQRVCAVZAXRHKAEGHTIASOIFKTMLKZFNITFCLFXAIWIXMMXZVMJYEWIEWXVSEQMGXVV UVTWGLDEGGSFRXAIWIQQIMFVAZXVARFXXVWKUWISGJUFEIHYMXMLSZZJNWCIUENRRVDXAIAZ OVHWQFBIWSHYQWTQSEASGIURHITXVFGKAXHFFLXSZUQVPSFCPWHJGGMGXEGJAYKGSJAJAYAR ZHTRUVDSKXVFGKAXCWFLXQCEXCMSRZEQBWGKTIBHSRAJEMTVGTHRHYQQTWWDBSLWWSXIHVWD BVHFOSXIBXWJOYKMCLEXHVSTMPEWCDQSYXVVYIGXOCTEUMHJAJMLCJQHXTOIFIWHOPEEMQCJ FXXVFVEXKMOCYIGJOEOMXHHYQVXQWXTXUICKTIKQSEGTHRARDWIIFYMTLMBWQVBSFKAXAIAJ QPOIGRZHKIOUKXHASCOSFIODUWLMCEMVRIBKQVIVWJQCXXOTDSLWHYQKNPTFRWIEQVYMGHGK TEMEFVFSHYFDURWWOJAYKWOJQXHXVFEIHJHYQFXEGKEXAEHGQVBWVZZXXPZVOXLZOJFEGHQF APTRRLZWRQDRFLXXWTDIZEFUQHMLWJQEKXVNUXAIBMUSNWSPQWTRRJXSPPMRZHLYFVXCWVSN FI.XMFGXEGWOXMMGWHI.E.

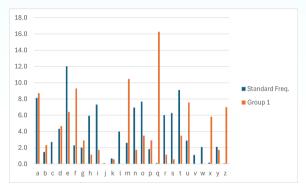
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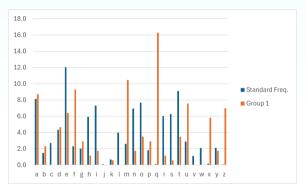
Patterns in the Standard Frequencies



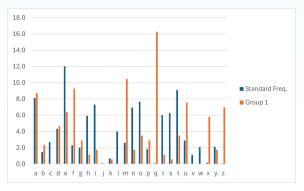




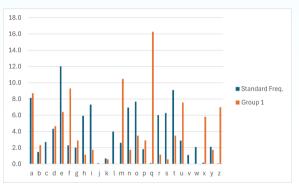
▶ Each group has been shifted by the same value: a Caesar cipher.



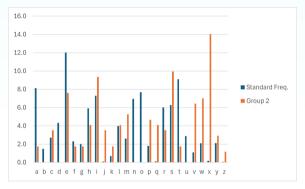
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- ▶ The charts will look the same, just shifted by a number of places.

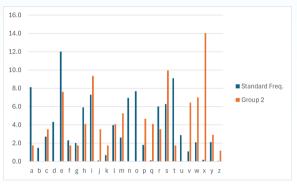


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- ▶ There is a gap of 4 letters from a to e, with e having the higher frequency.

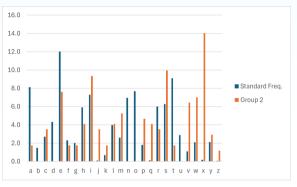


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- ▶ There is a gap of 4 letters from a to e, with e having the higher frequency.
- ▶ We can find this same gap, or other patterns, and deduce that the likely shift is a to m: shift value 12.

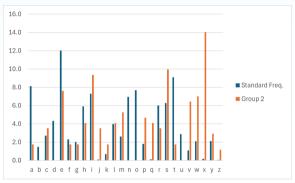




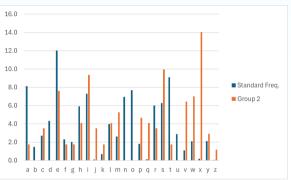
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- ▶ The charts will look the same, just shifted by a number of places.
- ► There looks to be a gap like before from s to x, but the gap is 5 so can't be the gap from a to e.
- Instead we can choose the gap e to i and deduce that the likely shift is a to e: shift value 4.

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- ▶ We make the following guesses:
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 - ► Group 2: shift 4, key letter e.

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 - ► Group 6: shift 17.

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 - ► Group 6: shift 17. key letter r.
- ▶ So the keyword appears to be meteor and we can then attempt to decrypt the text.

noonewouldhavebelievedinthelastyearsofthenineteenthcenturythatthisworldw asbeingwatchedkeenlyandcloselybyintelligencesgreaterthanmansandyetasmortalashi wnthatasmenbusiedthemselvesabouttheirvariousconcernstheywerescrutinisedandstud dperhapsalmostasnarrowlyasamanwithamicroscopemightscrutinisethetransientcreatusthatswarmandmultiplyinadropofwaterwithinfinitecomplacencymenwenttoandfroovert

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noone would have believed in the last years of the nineteenth century that this world was being watched keenly and closely by intelligences greater than mans and yet as mortal as his own that as men busied themselves about their various concerns they were scrutinised and studied perhaps almost as narrowly as a man with a microscope might scrutinise the transient creatures that swarm and multiply in a drop of water with infinite complacency men went to and fro over this globe about their little affairs serene in their assurance of their empire over matter it is possible that the infusoria under the microscope do the same noone gave a thought to the older worlds of space as sources of human danger or thought of them only to dismiss the idea of life upon them as impossible or improbable it is curious to recall some of the mental habits of those departed days at most terrestrial men fancied there might be other men upon mars perhaps inferior to themselves and ready to welcome a missionary enterprise yet across the gulf of space minds that are to our minds as ours are to those of the beasts that perish intellects vast and cool and unsympathetic regarded this earth with envious eyes and slowly and surely drew their plans against us

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- ▶ This gives repeated strings of letters in the ciphertext.
- This forms the basis of a Kasiski attack.
- Longer messages are more likely to yield to such an attack as there will be more repeated segments.

ZSHRSNAYEHVRHIUIZZQZXHWEFLXPOJFCXEFJAJMLSEURXXSVZXAGSEFYKCHYMXMLWJISKPRN MWUIWESATXQYQHDISEXCTRRTXSLIZPNCBRHVXPBKSEOILKFVMXXVHYMRFEBJMRWCSKMWFSFK MPTWVZESPRHYMXTWAVZFNWWVPXAIAJQPOIGRNSNXHYQMKZOIUSNWQFZGXVBJFLXCKVDILGFL FMGMGVPEGHGKGHBIRGQVAEDJMPFSGKMWGEFIAAECOJMQTRKZFLTQWTDSLGCGQQBKVKEGKYHZ ZMLIHYQXKEBJUIGXQIQEMYFVEXAEHJIEKQOEPQNPHZBPRMBRPVHTCWIEMIFNUXAMBWURBXST AQIPOTQRVCAVZAXRHKAEGHTIASOIFKTMLKZFNITFCLFXAIWIXMMXZVMJYEWIEWXVSEQMGXVV UVTWGLDEGGSFRXAIWIQQIMFVAZXVARFXXVWKUWISGJUFEIHYMXMLSZZJNWCIUENRRVDXAIAZ OVHWQFBIWSHYQWTQSEASGIURHITXVFGKAXHFFLXSZUQVPSFCPWHJGGMGXEGJAYKGSJAJAYAR ZHTRUVDSKXVFGKAXCWFLXQCEXCMSRZEQBWGKTIBHSRAJEMTVGTHRHYQQTWWDBSLWWSXIHVWD BVHFOSXIBXWJOYKMCLEXHVSTMPEWCDQSYXVVYIGXOCTEUMHJAJMLCJQHXTOIFIWHOPEEMQCJ FXXVFVEXKMOCYIGJOEOMXHHYQVXQWXTXUICKTIKQSEGTHRARDWIIFYMTLMBWQVBSFKAXAIAJ QPOIGRZHKIOUKXHASCOSFIODUWLMCEMVRIBKQVIVWJQCXXOTDSLWHYQKNPTFRWIEQVYMGHGK TEMEFVFSHYFDURWWOJAYKWOIQXHXVFEIHJHYQFXEGKEXAEHGQVBWVZZXXPZVOXLZOJFEGHQF APTRRLZWRQDRFLXXWTDIZEFUQHMLWJQEKXVNUXAIBMUSNWSPQWTRRJXSPPMRZHLYFVXCWVSN FI.XMFGXEGWOXMMGWHI.E.

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- ► The key length could be 2, 3, 4, 6, 8, 9, 12, 18, 24, 32, 36, 48, 72, 96, 133, or 288.
- ▶ We've narrowed it down, but that's still a lot of options.

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- We repeat this with other sequences.

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- ▶ We repeat this again with other sequences.

sequence	interval	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16
XAIA	288	Х	Χ	Х		Χ		Х	Χ			Χ				Х
FLX	60	Х	Χ	Х	Х	Χ				Χ		Χ			Χ	
HYQ	108	Х	Χ	Χ		Χ			Χ			Χ				
XVF	54	Х	Χ	Χ		Χ			Χ			Χ				
YMX	96	Х	Χ	Χ		Χ		Х				Χ				Χ
KAX	54	Х	Χ	Х		Χ			Χ			Χ				
FXXV	258	Χ	Χ			Χ										

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- ▶ If we analyse the factors in the table, the obvious candidates for key lengths are 2, 3, 4, and 6.
- ▶ Key lengths of 2, 3 or 4 are quite short, so we rule them out. (We may have to come back to them.) The key length is likely to be 6 and as we found out, the actual keyword was meteor.



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- German cryptographer and archaeologist Friedrick Kasiski attacked the problem in 1863.
- But Charles Babbage had cracked it in 1854. The British government wanted to keept it secret in the Crimean War.
- We know Babbage better as being the originator of the programmable computer, having designed the Analytical and Differences Engines.

Tutorials

In the tutorial this week we will:

- Create a spreadsheet to perform Vigenère encryption.
- Use a premade spreadsheet to perform a Kasiski-style attack on Vigenère encryption.