Operads, Equivariance, and Pseudo-Commutativity

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Slides: alex-corner.github.io/slides/corner-pssl110.pdf

Operads

$\mathsf{Data} \qquad n \in \mathbb{N} = \{0, 1, 2, \ldots\}$

- sets P(n) thought of as 'n-ary operations'
- functions 'composition' or 'multiplication'

$$\mu^P \colon P(n) \times P(k_1) \times \ldots \times P(k_n) \to P(k_1 + \ldots + k_n)$$

• identity element $e \in P(1)$

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Axioms

multiplication:

$$\mu(\mu(\rho; \mathbf{q}_{1}, \dots, \mathbf{q}_{n}); r_{11}, \dots, r_{1k_{1}}, \dots, r_{n1}, \dots, r_{nk_{n}}) = \mu(\rho; \mu(\mathbf{q}_{1}; r_{11}, \dots, r_{1k_{1}}), \dots, \mu(\mathbf{q}_{n}; r_{n1}, \dots, r_{nk_{n}}))$$

• identity:

$$\mu(p;e,\ldots,e)=p=\mu(e;p)$$

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Idea

• elements of P(n) are n-ary operations:

$$n \text{ inputs } \left\{ \begin{array}{c} \vdots \\ \vdots \\ \end{array} \right.$$

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Operads - Symmetric Actions

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Symmetric Operads

• for each $n \in \mathbb{N}$, an action:

$$P(n) \times \Sigma_n \to P(n)$$

satisfying:

$$\mu(p \cdot g; q_1, \ldots, q_n) = \mu(p; q_{g^{-1}(1)}, \ldots, q_{g^{-1}(n)}) \cdot \delta(g)$$

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Plain Operads

$$P(n) \times 1 \rightarrow P(n)$$

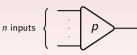
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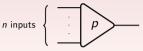




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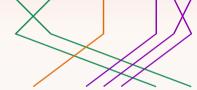
















$$\beta(p_1,\ldots,p_n)=\mu(e_n;p_1,\ldots,p_n)$$

Extension

$$\delta(p) = \mu(p; e_{k_1}, \dots, e_{k_n})$$

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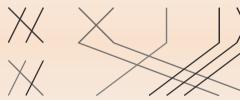
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Action Operads

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$$n \in \mathbb{N} = \{0, 1, 2, \ldots\}$$

- groups ∧(n)
- ullet plain operad structure $m(\mu,em)$
- group homomorphisms $\pi_n \colon \Lambda(n) \to \Sigma_n$

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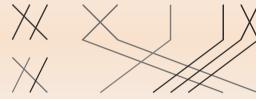
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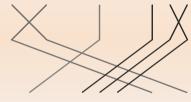














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- $\Lambda(n) = \Sigma_n$
- operad structure: β and δ , $e_1 = (1) \in \Sigma_1$
- π_n is the identity
- equivariance follows from symmetric equivariance

Examples

- Σ : Σ_n , symmetric groups
- T: 1, trivial groups
- B: B_n , braid groups
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Non-Examples

- A: A_n, alternating groups
- H: H_n, hyperoctahedral groups
- C: C_n, cyclic groups
- $R: C_2, \pi_n$ order-reversing /

Fiedorowicz–Loday

Most skew-simplicial groups, despite similarities.

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∧-Operads

- $T: \mathcal{C} \to \mathcal{C}$ a monad
- a **strength** is a natural transformation + axioms

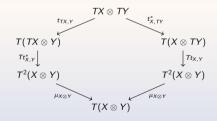
equivalent when
$${\cal C}$$
 sym. mon. cat

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$$t_{X,Y}^* \colon TX \otimes Y \to T(X \otimes Y)$$

• T is **commutative** if this commutes:



Data

- an action operad Λ
- an operad P
- actions $P(n) \times \Lambda(n) \rightarrow P(n)$
- satisfying:

$$\mu^{P}(p \cdot g; q_{1}, \dots, q_{n}) = \mu^{P}(p; q_{g^{-1}(1)}, \dots, q_{g^{-1}(n)}) \cdot \delta^{\Lambda}(g)$$

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Monads

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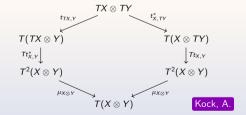
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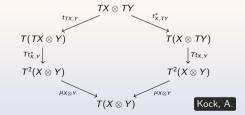


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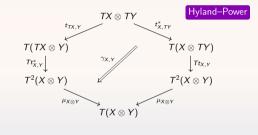
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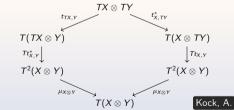
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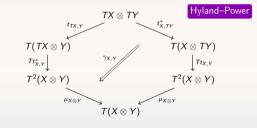
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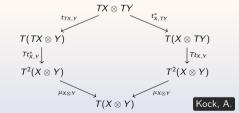
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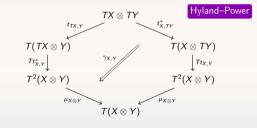
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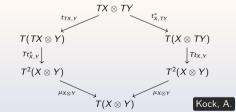
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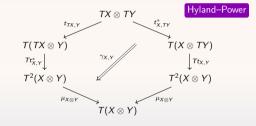
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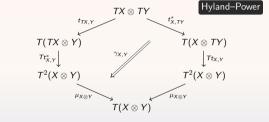
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$\textbf{Pseudo-Commutative} \ \Lambda\textbf{-Operads}$

can characterise Λ -operads with ps-commutativity

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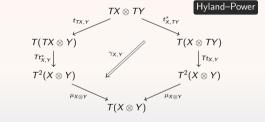
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- $\bullet \ \, \textit{T} \,\, \mathsf{ps\text{-}commutative} \Rightarrow \, \textit{T\text{-}}\mathbf{Alg}_{ps} \,\, \mathsf{is} \,\, \mathsf{pseudo\text{-}closed} \,\,$
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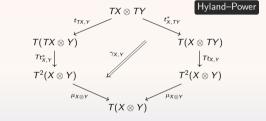
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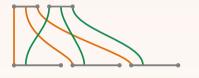
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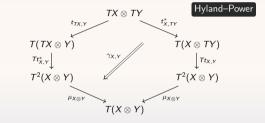
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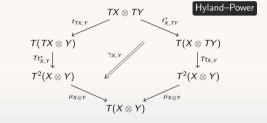
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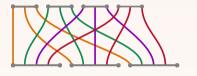
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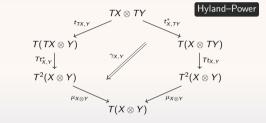
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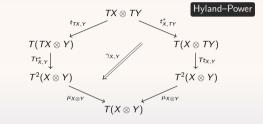
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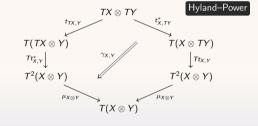


• these permutations correspond to matrix transpose

$$\begin{bmatrix} (x_1,y_1) & (x_1,y_2) & (x_1,y_3) \\ (x_2,y_1) & (x_2,y_2) & (x_2,y_3) \\ (x_3,y_1) & (x_3,y_2) & (x_3,y_3) \\ (x_4,y_1) & (x_4,y_2) & (x_4,y_3) \\ (x_5,y_1) & (x_5,y_2) & (x_5,y_3) \end{bmatrix}^T = \begin{bmatrix} (x_1,y_1) & (x_2,y_1) & (x_3,y_1) & (x_4,y_1) & (x_5,y_1) \\ (x_1,y_2) & (x_2,y_2) & (x_3,y_2) & (x_4,y_2) & (x_5,y_2) \\ (x_1,y_3) & (x_2,y_3) & (x_3,y_3) & (x_4,y_3) & (x_5,y_3) \end{bmatrix}$$

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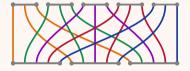
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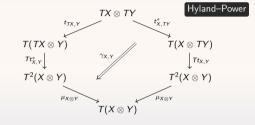
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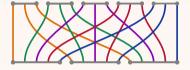
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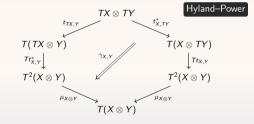
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lots of axioms

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Examples/Non-Examples



Guillou-May-Merling-Osorno

- Σ: symmetric strict monoidal categories
- B: braided strict monoidal categories



Non-Examples

• J: cactus/coboundary strict monoidal categories



- arXiv:1508.04050 Operads, tensor products, and the categorical Borel construction
- arXiv:1312.5910 Operads with general groups of equivariance, and some 2-categorical aspects of operads in Cat