

↓ §o Fixes from last time.

e.g. ↓

R ↓ S ↓ T x

S = R , a R b  $\iff |a - b| \leq 1$

|R<sub>2</sub> , 2R<sub>3</sub> , 1R<sub>3</sub>|

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$R \cup S \times T \setminus$

$S = \mathbb{Z}$        $aRb \Leftrightarrow a|b$

$4|8, 8|4$

$R \times S \cup T \setminus$   
 $S = \mathbb{Z}$        $aRb \Leftrightarrow ab \text{ is odd}$

$S: aRb \Leftrightarrow ab \text{ odd} \Leftrightarrow b \text{ odd} \Leftrightarrow bRa$

$\exists$   $a R b, b R c \Rightarrow a, b, c$  <sup>both</sup> are all odd  
 $\Rightarrow ac$  is odd      Even  
 $\Rightarrow a R c$

$2 R 2$

$2 \cdot 2 = 4$  even not odd

Prop "pf":  $x R y \stackrel{S}{\Rightarrow} y R x \stackrel{T}{\Rightarrow} x R x$   $\blacksquare$   
 $S, T \Rightarrow R$

$$S = \{1, 2, 3\}$$

$$R = \{(1, 1), (2, 1), (1, 3)\}$$

Prop: If a relation is Sym + trans  
and every element of S is relate to something  
then it is reflex.

$$a \mathrel{b} \Leftrightarrow b = a \cdot x$$

$$b \mathrel{c} \Leftrightarrow c = b \cdot y$$

$$C = b \cdot y = a \cdot x \cdot y$$

$$a \mathrel{c}$$



§

## §1 Closures

Def'n: A relation  $R$  dominates  $R'$  if  
 $R' \subseteq R$  as a subset of  $S \times S$ .

Visually:



e.g.]  $S = \mathbb{Z}$ ,  $aRb \Leftrightarrow a < b$   
 $aRb \Leftrightarrow a \leq b$

Def'n: Let  $R$  be a relation on  $S$ , and  
 $P \in \{\text{Symm., reflex., trans.}\}$ . The  $P$ -closure of  
 $R$  is the smallest relation dominating  $R$  which is  $P$ .  
e.g.]  $S = \mathbb{Z}$ ,  $aRb \Leftrightarrow a < b$ , what is the reflexive  
closure? what is the Symm closure? what is the trans closure?

7] Step 1: Find guess  $R'$

Step 2:  $R'$  is P and dominates R

Step 3: If  $R''$  is P and dominates R,  
then  $R'' \geq R'$ .

8] Symm:  $a R b \Leftrightarrow$  (either  $a < b$  or  $b < a$ )  
or  
 $a \neq b$



$$a R b \Leftrightarrow \begin{cases} a R b \\ b R a \end{cases}$$

$$b R a \Leftrightarrow \begin{cases} b R a \\ a R b \end{cases}$$

$$R = \{(a, b) \in S \times S$$

$$R' = \{(a, b) \in S \times S; (a, b) \in R \text{ or } (b, a) \in R\}$$

q]

Prop:  $\neq$  is Symm. closure of  $\subset$ .

Pf: First, observe that  $\neq$  is Symm:

if  $a \neq b$  then  $b \neq a$ , and dominates  $<$

if  $a < b$ ,  $a \neq b$ . Assume, that

$\neq$  is a symm relation dominating  $<$   $\Rightarrow \neq$  dominating

We want to show that if  $a \neq b$

then  $a \neq b$ . Either  $a < b$  or  $b < a$

and as  $\neq$  dominates  $< \Rightarrow a \neq b$  or  $b \neq a$

As  $\neq$  is Symm in either case  $a \neq b$ .  $\blacksquare$

Reflex: Reflexive closure of  $a < b$  is  $a \leq b$

Trans:  $<$        $\times >$

Hm: If  $R$  is a relation, then a P-closure exists

Lemma: If  $\{R_i\}$  is any non-empty collection  
of relations which are P, then  $R = \bigcap R_i$  is P.

$$aR'b \iff aR_i b \text{ for all } i.$$

Pf (transitive): Suppose  $aR^ib$  and  $bR^ic$ .

But  $aR^ib$  means  $(a, b) \in R = \bigcap R_i$ , so

$(a, b) \in R_i$  for all  $i$ , i.e.,  $aR_i b$  for all  $i$ .

Similarly,  $bR_i c$  for all  $i$ . Since  $R_i$  is trans,  
 $aR_i c$  for all  $i$ ; i.e.,  $(a, c) \in \bigcap R_i = R'$ , so  $aR'c$ .  $\square$

12] Pf of Thm: Let  $\mathcal{E}_{R;\mathbb{Z}}$  be the set of all relations dominating  $R$  and which are P. Note,  $S \times S \in \mathcal{E}_{R;\mathbb{Z}}$ .

Set  $R' = \bigcap_i R_i$ . By the lemma,  $R'$  is P and dominates  $R$ .  
Assume  $R''$  is P and dominates  $R$ . Then  $R'' \in \mathcal{E}_{R;\mathbb{Z}}$ .  
So,  $R'' \supseteq \bigcap_i R_i = R'$ , so  $R''$  dominates  $R'$ .  
Note: If  $R'$  and  $P''$  are P-closures of  $R$ , then  $R'' \supseteq R'$ .  
 $\Rightarrow R' \subseteq R''$ ,  $R'' \subseteq R'$   $\Rightarrow R' = R''$ .

IB

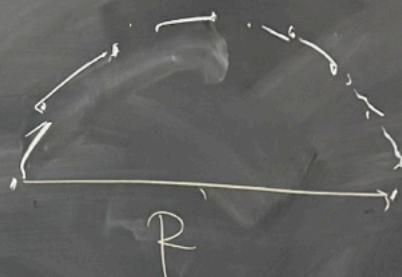
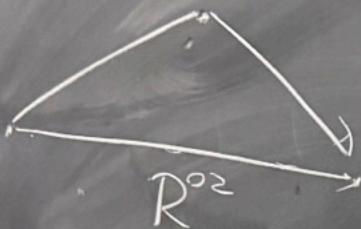
Def'n: For two relations  $R$  and  $R'$  are relations their composition is

$$R \circ R' = \{(a, c) \in S \times S \mid \exists b \text{ s.t. } aRb \text{ and } bR'c\}$$

Visually:



Prop:  $R = R^{\circ 2} \iff R = R^{\circ n} \forall n \iff R^{\text{trans}}$



Prop: The trans closure of  $R$  is

$$\bigcup_{n=1}^{\infty} R^{\circ n} = R'$$

15]

## § 2 FEquivalence relations



$\Delta$

$\Delta \cong \Delta$

Defn: An equivalence relation is a relation which is sym, reflex., trans.

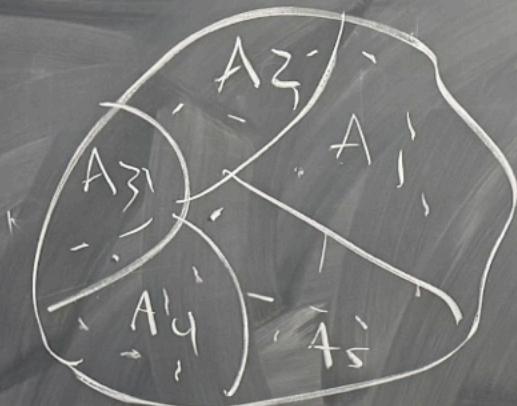
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e.g.]  $S = \{ \text{animals} \}$

$a R b \Leftrightarrow a \text{ and } b \text{ are same species.}$

Defn: A partition of a set  $S$  is a collection  $\{A_i\}$  of subsets of  $S$  s.t.  $S = \bigcup A_i$ ,  $A_i \cap A_j = \emptyset$  for  $i \neq j$ .

Visually:



Thm: Partitions  $\iff$  equivalence relations