VII. Container with Most Water Medium
int max Area Cint[] heights):
left = 0, right = len(heights)-1, max=0//set initial pointers & initial max
WHILE left < right: // While the pointers do not meet
diff = right - left, cap = min (heights [left], heights [right])
amount = diff * cap//calculate amount
IF amount > max: max = amount // update max if necessary
IF heights [left] < heights [right]: left+t // increment left if right is greater
ELSE: right // decrement right otherswise
Return max Complexity: O(n)
713. Subarray Product Less Than K Medium
int count Subarrays (int[] nums, int k):
prev = 1, beg=0, total=0 // default previous product = 1, beg of window =0
FOR 1=0; 1 < len (nums); i++:
IF nums[i] >= k://num[i] > k => product * nums[i] > k
IF $nums[i] >= k : //num[i] \ge k => product * nums[i] \ge k$ Prev = 1, beey = i+1 //reset prev product \land beg to next continue // skip rest of this ilention
Prev=1, beey = i+1 // recet prev product 1 beg to next Continue // skip rest of this iteration
<u>'</u>

138. Find All Anagrams in a String Easy
boolean is Anagram (Map <integer) anagram="" check="" function="" helper="" if="" is="" one,="" one.equals(two)<="" return="" td="" to="" two):=""></integer)>
List < Integer> find Anagrams (String s, p): ret = [] // initialize return variable
IF len(s) < len(p): return ret // Not possible to contain anagrams
Map < Integer > letters = {3}, temp = {3} // define the maps
FOR char c IN s:
etters[c] = etters[c]+ IF c IN etters ELSE 0 // map each etter
FOR 1=0; i < len(p); itt:
temp[sti]] = temp[sti]] + IF sti] IN temp ELSE O//check first
FOR i=0; i < len(s) - len(p) +1; i++: // check all windows by sliding
IF is Anagram (letters, temp):
ret.add (i) / add index if is anagram
temp [S[i]] // remove first element in the window
IF i < len(s) - len(ρ): // check if we have not reached last window
temp [Stitlen(p)]] It / add next element to end of window
Return ret/return list containing all the indices
Complexity: O(n)

V 209. Minimum Size Subarray Sum Medium
int count Subarrays (int[] nums, int target)
Min = MAX, sum = 0, beg = 0 // initial min is constant MAX
FOR 1=0; 1 < len (nums); 1++:
IF nums[i] > target: return // Return 1 is a single greater or equal
Sum t= nums[i]//update sum
IF sum < target: continue // skip rest of this iteration if sum < target
cur = 1-beg + 2// length of window + 1 (to account for apcoming decrement)
WHILE sum = target:
sum -= nums [beg+], cur//chop window & decrement cur
min = cur < min ? cur: min // update min it necessary
Return min == MAX ? 0: min // return min it min changed, otherwise 0
Complexity () (11)