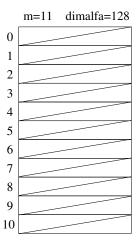
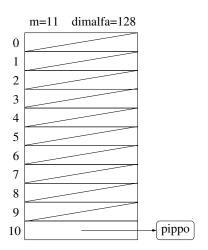
Algoritmi e Strutture Dati Hash (esempi)

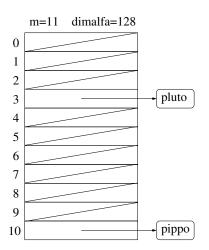
P. Massazza¹

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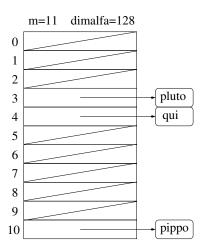




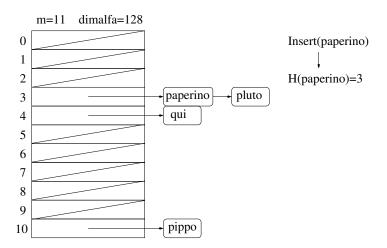


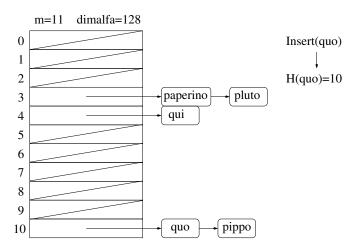






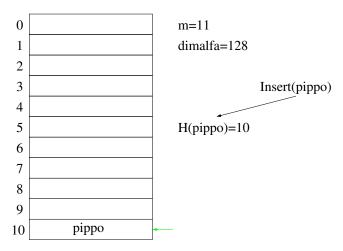


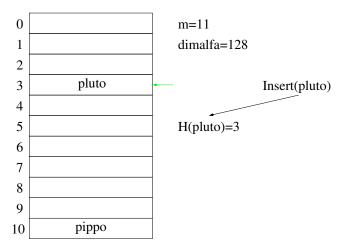


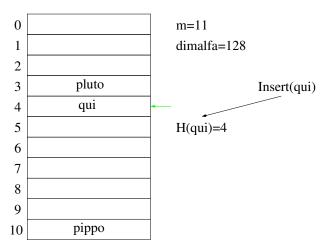


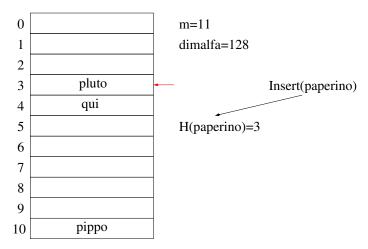
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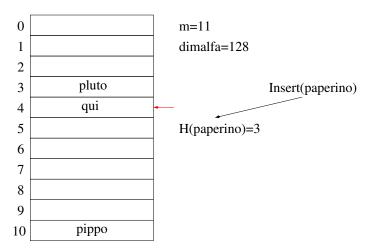
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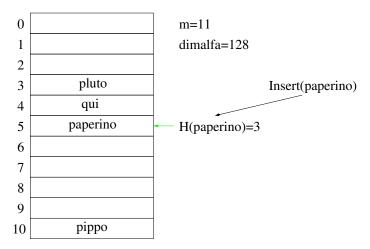


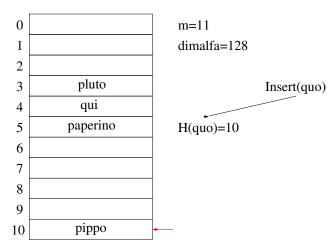


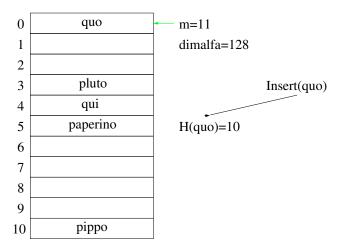


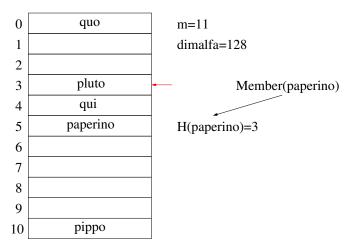


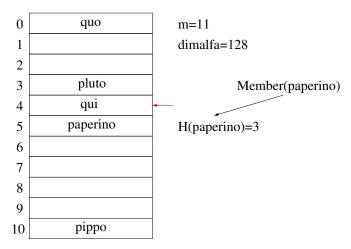


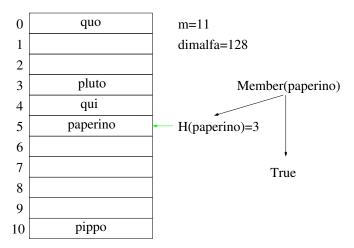


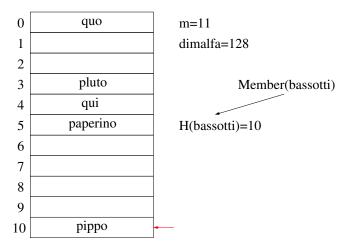


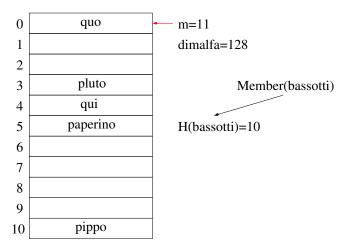


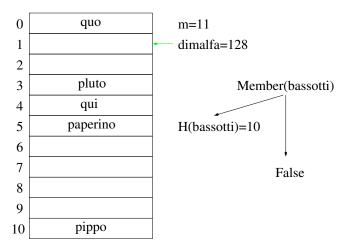


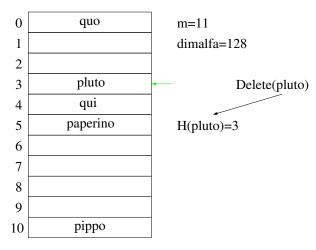






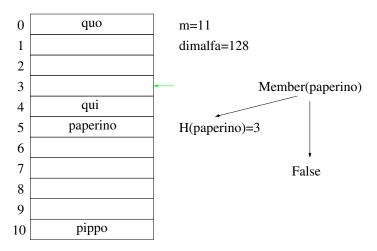


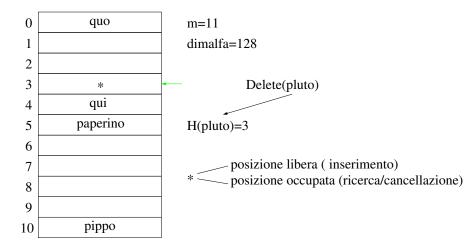


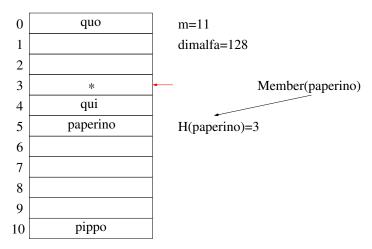


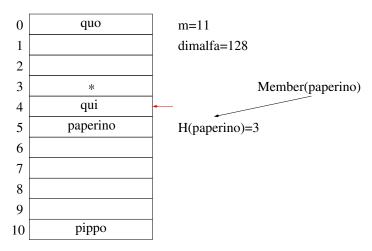
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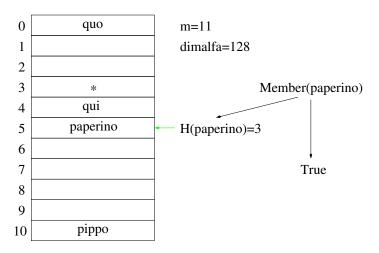
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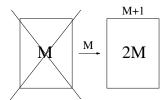


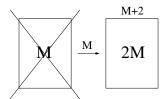
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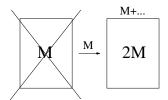
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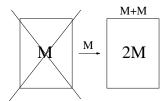
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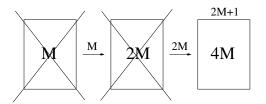
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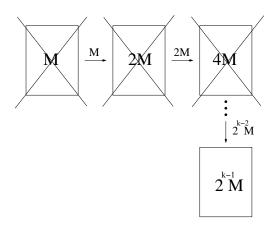


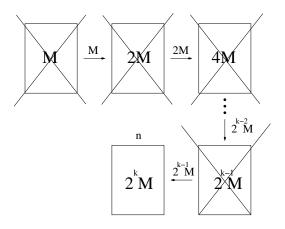


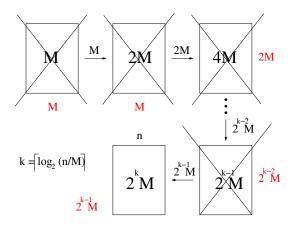












Il costo complessivo di *n* inserimenti è quindi

$$\begin{array}{lll} \textit{M} + \textit{M} \sum_{i=0}^{k-1} 2^{i} + \textit{M} \sum_{j=0}^{k-1} 2^{j} & = & \textit{M} + 2\textit{M}(2^{k} - 1) \\ & = & \textit{M}(2^{k+1} - 1) \quad (\mathsf{nota:} \ k = \lceil \log_{2}(n/\textit{M}) \rceil) \\ & \leq & \textit{M}(2 \cdot 2^{\log_{2} n} - 1) = \textit{M}(2n - 1) = \Theta(n) \end{array}$$

Il costo ammortizzato di un inserimento (costo totale diviso n) è O(1)