

CLOUD COMPUTING CONCEPTS

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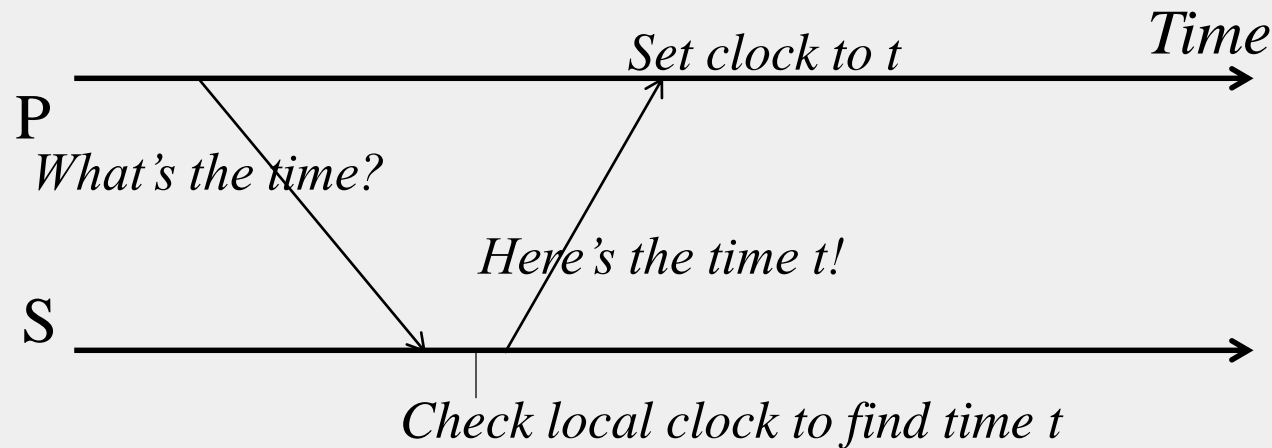
TIME AND ORDERING

Lecture B

CRISTIAN'S ALGORITHM

BASICS

- External time synchronization
- All processes P synchronize with a time server S

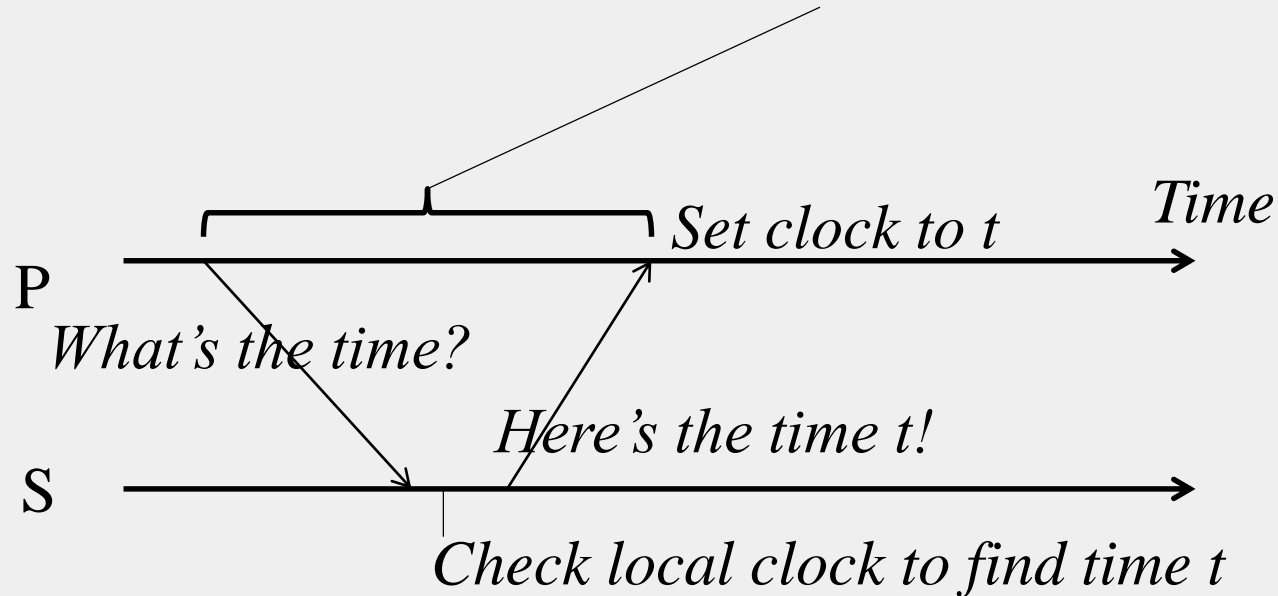


WHAT'S WRONG

- By the time response message is received at P, time has moved on
- P's time set to t is inaccurate!
- Inaccuracy a function of message latencies
- Since latencies unbounded in an asynchronous system, the inaccuracy cannot be bounded

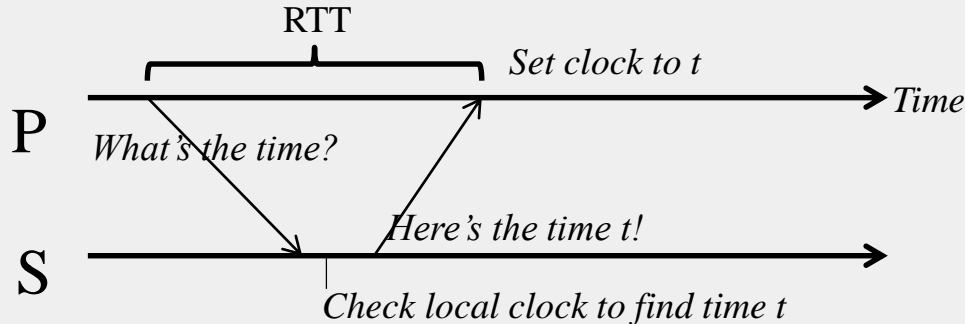
CRISTIAN'S ALGORITHM

- P measures the round-trip-time RTT of message exchange



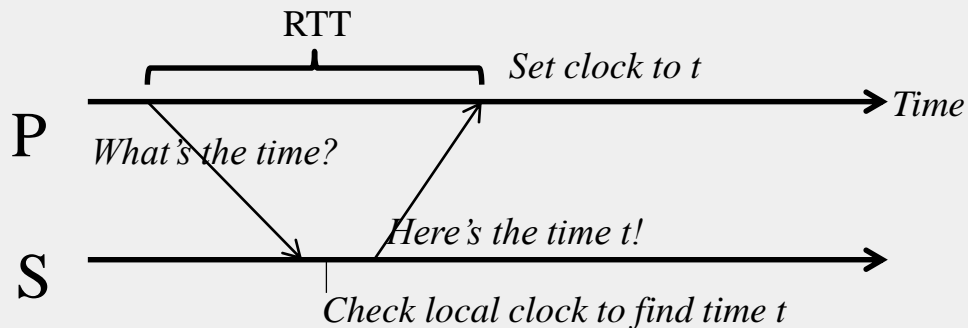
CRISTIAN'S ALGORITHM (2)

- **P** measures the round-trip-time RTT of message exchange
- Suppose we know the minimum $P \rightarrow S$ latency min1
- And the minimum $S \rightarrow P$ latency min2
 - min1 and min2 depend on operating system overhead to buffer messages, TCP time to queue messages, etc.



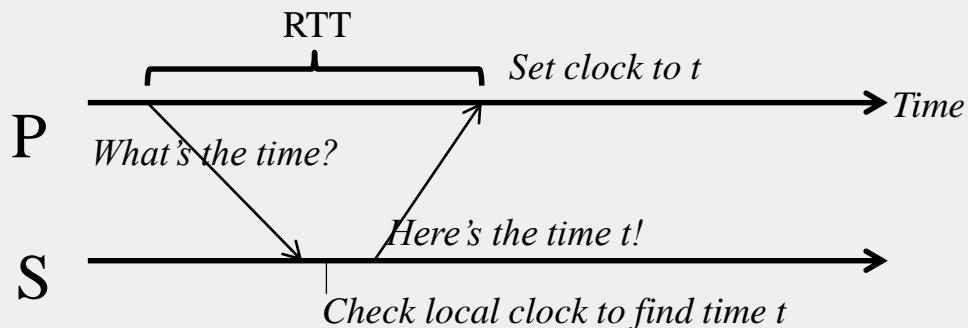
CRISTIAN'S ALGORITHM (3)

- P measures the round-trip-time RTT of message exchange
- Suppose we know the minimum $P \rightarrow S$ latency $\min1$
- And the minimum $S \rightarrow P$ latency $\min2$
 - $\min1$ and $\min2$ depend on Operating system overhead to buffer messages, TCP time to queue messages, etc.
- The actual time at P when it receives response is between $[t+\min2, t+\text{RTT}-\min1]$



CRISTIAN'S ALGORITHM (4)

- The actual time at P when it receives response is between $[t + \min_2, t + \text{RTT} - \min_1]$
- P sets its time to halfway through this interval
 - To: $t + (\text{RTT} + \min_2 - \min_1) / 2$
- Error is at most $(\text{RTT} - \min_2 - \min_1) / 2$
 - Bounded!



GOTCHAS

- **Allowed to increase clock value but should never decrease clock value**
 - May violate ordering of events within the same process
- **Allowed to increase or decrease speed of clock**
- **If error is too high, take multiple readings and average them**