

CLOUD COMPUTING CONCEPTS with Indranil Gupta (Indy)

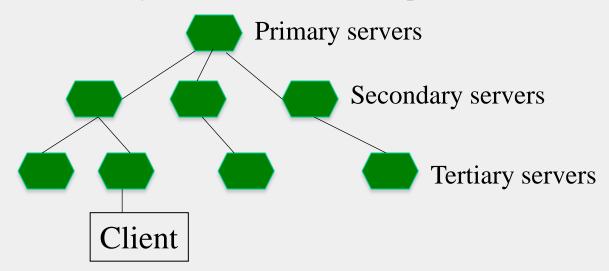
TIME AND ORDERING

Lecture C

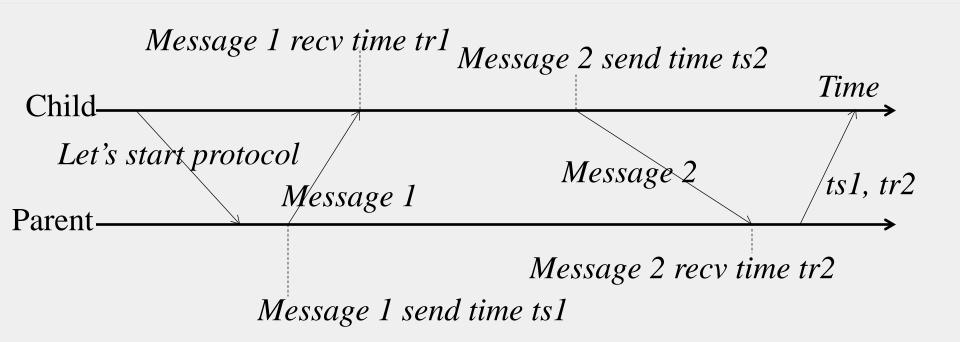
NTP

NTP = NETWORK TIME PROTOCOL

- NTP servers organized in a tree
- Each client = a leaf of tree
- Each node synchronizes with its tree parent



NTP PROTOCOL



WHAT THE CHILD DOES

- Child calculates *offset* between its clock and parent's clock
- Uses ts1, tr1, ts2, tr2
- Offset is calculated as

$$o = (tr1 - tr2 + ts2 - ts1)/2$$

Why o = (trl - trl + tsl - tsl)/2?

- Offset o = (tr1 tr2 + ts2 ts1)/2
- Let's calculate the error
- Suppose real offset is *oreal*
 - Child is ahead of parent by *oreal*
 - Parent is ahead of child by -oreal
- Suppose one-way latency of Message 1 is *L1* (*L2* for Message 2)
- No one knows *L1* or *L2*!
- Then

$$tr1 = ts1 + L1 + oreal$$

 $tr2 = ts2 + L2 - oreal$

WHY o = (trl - tr2 + ts2 - ts1)/2?

• Then

$$tr1 = ts1 + L1 + oreal$$

 $tr2 = ts2 + L2 - oreal$

• Subtracting second equation from the first

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oreal = (tr1 - tr2 + ts2 - ts1)/2 + (L2 - L1)/2
=> oreal = o + (L2 - L1)/2
=> |oreal - o| < |(L2 - L1)/2| < |(L2 + L1)/2|
- Thus, the error is bounded by the round-triptime
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AND YET...

- We still have a non-zero error!
- We just can't seem to get rid of error
 - Can't, as long as message latencies are non-zero
- Can we avoid synchronizing clocks altogether and still be able to order events?