

# Communicate Without Errors

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**Definition 0.1** (graph). *A graph  $G$  is a set of vertices with a set of edges connecting pairs of vertices.*

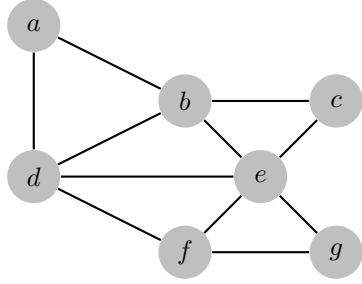


Figure 1: An example of a graph.

**Definition 0.2** ( $\alpha(G)$ ). *Given a graph  $G$ . Given a subgraph of  $G$ ,  $H$ , such that every vertex of  $H$  is not connected in  $G$ . Then  $\alpha(G)$  is the maximum number of vertices of  $H$ .*

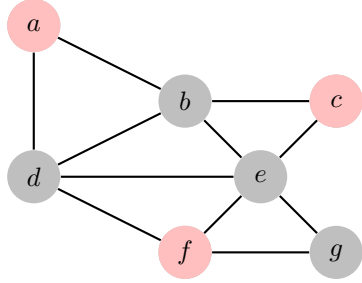


Figure 2: Example of  $\alpha(G)$ . Here  $\alpha(G) = 3$ .

**Definition 0.3** (graph product). *Given two graph  $G$  and  $H$ . The graph product  $G \times H$  is the graph with vertices  $V(G \times H) = V(G) \times V(H)$  in which  $(x, y)$  is adjacent to  $(x', y')$  in  $G \times H$  if and only if  $x$  is adjacent to  $x'$  in  $G$  and  $y$  is adjacent to  $y'$  in  $H$ .*

*A graph  $G$  product itself for  $n$  times will always be denoted by  $G^n$ .*

**Lemma 0.1.** *Given graph  $G$  and  $H$ .*

$$\alpha(G \times H) \geq \alpha(G)\alpha(H) \tag{1}$$

*Proof.* Given graph  $G$  and  $H$ . Let  $G'$  and  $H'$  be subgraph of  $G$  and  $H$  such that no vertex of  $G'$  or  $H'$  is adjacent in  $G$  or  $H$ , respectively. Then  $G' \times H'$  is a subgraph of  $G \times H$  such that no vertex of  $G' \times H'$  is adjacent in  $G \times H$ .  $\square$

**Definition 0.4** (Shannon capacity). *Given a graph  $G$ , the Shannon capacity  $\Theta(G)$  is defined by*

$$\Theta(G) = \lim_{n \rightarrow \infty} \sqrt[n]{\alpha(G^n)} \quad (2)$$

*We will see that this sequence is monotonically increasing and has an upper bound later. So, this sequence will always converge.*

**Definition 0.5** (adjacency matrix). *Given a graph  $G$ , the adjacency matrix  $A$  is defined by*

$$A_{ij} = \begin{cases} 1 & \text{if } i \text{ and } j \text{ are adjacent in } G \\ 0 & \text{otherwise} \end{cases} \quad (3)$$

**Definition 0.6** (orthonormal representation). *Given a graph  $G$  with vertices  $1, 2, \dots, n$ , the orthonormal representation of  $G$  is a set of unit vectors  $\{v_1, v_2, \dots, v_n\}$  such that if  $i$  and  $j$  are not adjacent in  $G$ , then  $v_i$  and  $v_j$  are orthogonal.*

**Lemma 0.2** (existence of orthonormal representation). *Given a graph  $G$  with vertices  $1, 2, \dots, n$ , there exists an orthonormal representation of  $G$ .*

*Proof.* Any orthonormal basis of  $\mathbb{R}^n$  can be used as an orthonormal representation of  $G$ .  $\square$

**Definition 0.7** (tensor product). *Given two vectors  $v = (v_1, \dots, v_n)$  and  $w = (w_1, \dots, w_n)$ , the tensor product  $v \circ w$  is defined by*

$$v \circ w = (v_1 w_1, \dots, v_1 w_n, v_2 w_1, \dots, v_2 w_n, \dots, v_n w_1, \dots, v_n w_n) \quad (4)$$

**Lemma 0.3** (inner product of tensor product). *Given two vectors  $v = (v_1, \dots, v_n)$  and  $w = (w_1, \dots, w_n)$ , the inner product of  $v \circ w$  is defined by*

$$\langle v \circ w, v' \circ w' \rangle = \langle v, v' \rangle \langle w, w' \rangle \quad (5)$$

**Lemma 0.4** (product of orthonormal representation). *Given a graph  $G$  with vertices  $1, 2, \dots, n$ , and a graph  $H$  with vertices  $1, 2, \dots, m$ . Then vectors  $\{v_i \circ w_j\}$  is an orthonormal representation of  $G \times H$ .*

**Definition 0.8** (value of a orthonormal representation). *Given an orthonormal representation  $\{v_1, v_2, \dots, v_n\}$  of a graph  $G$  the value of the representation is defined by*

$$\text{value}(\{v_1, v_2, \dots, v_n\}) = \inf_c \max_i \frac{1}{\langle c, v_i \rangle^2} \quad (6)$$

*which  $i = 1, 2, \dots, n$  and  $c$  is any vector does not orthogonal to  $v_i$ .*

**Definition 0.9** ( $\theta(G)$ ). *Given a graph  $G$ , the  $\theta(G)$  is defined by*

$$\theta(G) = \inf_{\{v_1, v_2, \dots, v_n\}} \text{value}(\{v_1, v_2, \dots, v_n\}) \quad (7)$$

*where  $\{v_1, v_2, \dots, v_n\}$  is an orthonormal representation of  $G$ .*

**Lemma 0.5.** *Given graph  $G$  and  $H$ , then*

$$\theta(G \times H) \leq \theta(G)\theta(H) \quad (8)$$

*Proof.* Let  $\{v_1, v_2, \dots, v_n\}$  and  $\{w_1, w_2, \dots, w_m\}$  be orthonormal representation of  $G$  and  $H$  and  $c_v$  and  $c_w$  such that

$$\max \left\{ \frac{1}{\langle c_v, v_i \rangle^2} : i = 1, 2, \dots, n \right\} = \theta(G) \quad (9)$$

and

$$\max \left\{ \frac{1}{\langle c_w, w_i \rangle^2} : i = 1, 2, \dots, m \right\} = \theta(H) \quad (10)$$

Then

$$\theta(G \times H) \leq \max \left\{ \frac{1}{\langle c_v \circ c_w, v_i \circ w_j \rangle^2} \right\} \quad (11)$$

$$= \max \left\{ \frac{1}{\langle c_v, v_i \rangle^2 \langle c_w, w_j \rangle^2} \right\} \quad (12)$$

$$= \max \left\{ \frac{1}{\langle c_v, v_i \rangle^2} \right\} \max \left\{ \frac{1}{\langle c_w, w_j \rangle^2} \right\} \quad (13)$$

$$= \theta(G)\theta(H) \quad (14)$$

$$(15)$$

□

**Lemma 0.6.**

$$\theta(G) \geq \alpha(G) \quad (16)$$

*Proof.* Let  $\{1, 2, \dots, k\}$  be the set of vertices of  $G$  such that every point is not adjacency in  $G$ . And  $k = \alpha(G)$

Let  $\{v_1, v_2, \dots, v_n\}$  be an orthonormal representation of  $G$  and  $c$  such that

$$\max \left\{ \frac{1}{\langle c, v_i \rangle^2} \right\} = \theta(G) \quad (17)$$

Then

$$1 = c^2 \quad (18)$$

$$\geq \sum_{i=1}^k \langle c, v_i \rangle^2 \quad (19)$$

$$\geq \frac{k}{\theta(G)} \quad (20)$$

□

**Theorem 0.1.** *Given a graph  $G$ , then*

$$\theta(G) \geq \Theta(G) \quad (21)$$

*Proof.*

$$\Theta(G) = \lim_n \sqrt[n]{\alpha(G^n)} \quad (22)$$

$$\leq \lim_n \sqrt[n]{\theta(G^n)} \quad (23)$$

$$\leq \lim_n \sqrt[n]{\theta(G)^n} \quad (24)$$

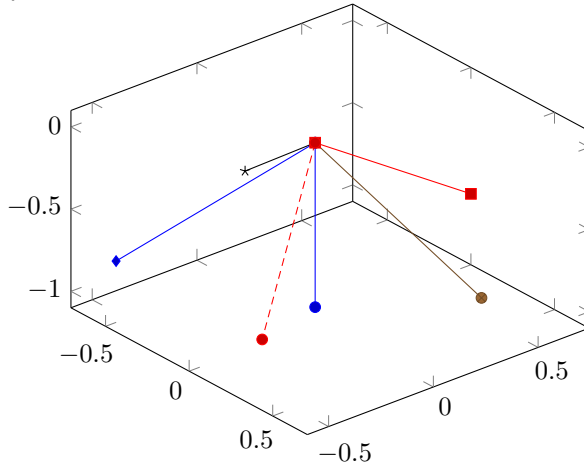
$$= \theta(G) \quad (25)$$

□

**Theorem 0.2.** For odd  $n$ ,

$$\theta(C_n) \geq \frac{n \cos(\pi/n)}{\cos(\pi/n) + 1} \quad (26)$$

*Proof.* Consider an umbrella that has a handle and  $n$  ribs that all have unit length.



Hear, angles between two consecutive ribs are same. Let  $v$  be a rib. And let  $w$  be one of the rib that have the largest angle with  $v$ . Then, we let the angle between  $v$  and  $w$  be  $\pi/2$ .

Then, the  $n$  ribs of such an umbrella form an orthonormal representation of  $\tilde{C}_n$ .

Let  $d$  be the vector represent the handle, and  $v_1, v_2, \dots, v_n$  be the  $n$  ribs. And let  $\gamma$  be the angle between the handle and any rib.

So, by some calculation, we get

$$\theta(C_n) = \max \sum_{i=1}^n (d^T v_i)^2 \quad (27)$$

$$\geq n (\cos(\gamma))^2 \quad (28)$$

$$= n \left( \sqrt{1 - \frac{1}{1 + \cos(\pi/n)}} \right)^2 \quad (29)$$

$$= \frac{n \cos(\pi/n)}{\cos(\pi/n) + 1} \quad (30)$$

□