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Quantum Abstract Interpretation

Seminar for the Introduction to Quantum Computing course

Università di Pisa Dipartimento di Informatica

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Roadmap



Introduction

As quantum computing advances, we would like to have some means to prove correctness properties on quantum programs, *especially* since quantum programming is counterintuitive.



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No: **exponential** space and time cost.



$$n_{qubits}=1$$

$$|0\rangle\,\langle 0|$$

$$\begin{pmatrix} 1 & 0 \\ 0 & 0 \end{pmatrix}$$

 $2^2 = 4$ complex numbers



$$n_{qubits} = 2$$

$$|00\rangle\langle00|$$

 $2^4 = 16$ complex numbers



$$n_{qubits} = 4$$

 $|0000\rangle \langle 0000|$

 $2^8 = 64$ complex numbers



$$n_{qubits} = 300$$

$$\left|0\right>^{\otimes_{300}}\left<0\right|^{\otimes_{300}}$$

?????



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$$\left|0\right>^{\otimes_{300}}\left<0\right|^{\otimes_{300}}$$

77777

 $2^{600} = 414951556888099295851240786369116115101244623224243689999 \\ 56573296906528114129081463997070489471037942881978866113007 \\ 89182395151075411775307886874834113963687061181803401509523685376$

Bigger than the number of atoms in the universe.



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Solution: abstract interpretation



Ingredients

- Abstract domain
 - Abstraction function
 - Concretization function
 - Abstract operations
- Assertions



Density Matrix

Instead of dealing with a state $|\phi\rangle$ in vector form, we use their *density matrix*:

 $\left|\phi\right\rangle \left\langle \phi\right|$ (For a pure state)



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Example:

$$\ket{eta_{00}} = rac{\ket{00} + \ket{11}}{\sqrt{2}} \ \ket{eta_{00}} raket{eta_{00}} = rac{1}{2} egin{pmatrix} 1 & 0 & 0 & 1 \ 0 & 0 & 0 & 0 \ 0 & 0 & 0 & 0 \ 1 & 0 & 0 & 1 \end{pmatrix}$$



Reduced Density Matrix

Suppose we have a composite quantum system $AB = A \otimes B$, and we want to consider a state $|\phi\rangle \in AB$ with respect to the subsystem A.

$$\begin{aligned} A &= \mathbb{C}^{2^n} \times \mathbb{C}^{2^n} \\ B &= \mathbb{C}^{2^m} \times \mathbb{C}^{2^m} \\ AB &= (\mathbb{C}^{2^n} \times \mathbb{C}^{2^n}) \otimes (\mathbb{C}^{2^m} \times \mathbb{C}^{2^m}) \end{aligned}$$

$$Tr_B[\rho]:AB o A$$
 $Tr_A[\rho]:AB o B$ $Tr_A[\rho]=\sum_{v=0}^{2^m}(I_A\otimes\langle v|)
ho(I_A\otimes|v
angle)$ $Tr_A[\rho]=\sum_{v=0}^{2^n}(\langle v|\otimes I_B)
ho(|v
angle\otimes I_B)$

