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Quantum Abstract Interpretation

Seminar for the **Introduction to Quantum Computing** course

Università di Pisa
Dipartimento di Informatica

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Roadmap

Introduction

As quantum computing advances, we would like to have some means to prove correctness properties on quantum programs, *especially* since quantum programming is counterintuitive.

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No: **exponential** space and time cost.

Example

$$n_{qubits} = 1$$

$$|0\rangle \langle 0|$$

$$\begin{pmatrix} 1 & 0 \\ 0 & 0 \end{pmatrix}$$

$$2^2 = 4 \text{ complex numbers}$$

Example

$$n_{qubits} = 2$$

$$|00\rangle \langle 00|$$

$$\begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \end{pmatrix}$$

$2^4 = 16$ complex numbers

Example

$$n_{qubits} = 3$$

$$|000\rangle \langle 000|$$

$$\begin{pmatrix} 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \end{pmatrix}$$

$2^6 = 64$ complex numbers

Example

$$n_{qubits} = 300$$

$$|0\rangle^{\otimes 300} \langle 0|^{\otimes 300}$$

?????

Example

$$n_{qubits} = 300$$

$$|0\rangle^{\bigotimes 300} \langle 0|^{\bigotimes 300}$$

$$2^{600} = 4149515568880992958512407863691161151012446232242436899995657329690652811412908146399707048947103794288197886611300789182395151075411775307886874834113963687061181803401509523685376$$

Bigger than the number of atoms in the universe.

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Solution: abstract interpretation

Ingredients

- Abstract domain
  - Abstraction function
  - Concretization function
  - Abstract operations
- Assertions

Density Matrix

Instead of dealing with a state $|\phi\rangle$ in vector form, we use their *density matrix*:

$\rho_\phi = |\phi\rangle \langle\phi|$ (For a pure state)

- positive semi-definite
- $Tr(\rho) = 1$
- projection ($P = P^\dagger = P^2$)

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Example:

$$|\beta_{00}\rangle = \frac{|00\rangle + |11\rangle}{\sqrt{2}}$$

$$\begin{aligned}
\rho_{\beta_{00}} &= |\beta_{00}\rangle\langle\beta_{00}| = \frac{1}{2}(|00\rangle + |11\rangle)(\langle 00| + \langle 11|) \\
&= \frac{1}{2}(|00\rangle\langle 00| + |00\rangle\langle 11| + |11\rangle\langle 00| + |11\rangle\langle 11|) \\
&= \frac{1}{2}\begin{pmatrix} 1 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ 1 & 0 & 0 & 1 \end{pmatrix}
\end{aligned}$$

Reduced Density Matrix

Suppose we have a composite quantum system $AB = A \otimes B$, and we want to focus our attention on a state $|\phi\rangle \in AB$ with respect to the subsystem A.

$$A = \mathbb{C}^{2^n} \times \mathbb{C}^{2^n}$$
$$B = \mathbb{C}^{2^m} \times \mathbb{C}^{2^m}$$
$$AB = (\mathbb{C}^{2^n} \times \mathbb{C}^{2^n}) \otimes (\mathbb{C}^{2^m} \times \mathbb{C}^{2^m})$$

$$Tr_B[\rho] : AB \to A \qquad Tr_A[\rho] : AB \to B$$

$$Tr_B[\rho] = \sum_{v=0}^{2^m} (I_A \otimes \langle v|)\rho(I_A \otimes |v\rangle) \qquad Tr_A[\rho] = \sum_{v=0}^{2^n} (\langle v| \otimes I_B)\rho(|v\rangle \otimes I_B)$$

Example

$$Tr_B[\rho] = \sum_{v=0}^{2^m} (I_A \otimes \langle v|)\rho(I_A \otimes |v\rangle) \qquad \rho_{\beta_{00}} = \frac{1}{2}\begin{pmatrix}1 & 0 & 0 & 1\\ 0 & 0 & 0 & 0\\ 0 & 0 & 0 & 0\\ 1 & 0 & 0 & 1\end{pmatrix}$$

$$\begin{aligned}Tr_B[\rho_{\beta_{00}}] =&(I_2 \otimes \langle 0|)\rho_{\beta_{00}}(I_2 \otimes |0\rangle) + (I_2 \otimes \langle 1|)\rho_{\beta_{00}}(I_2 \otimes |1\rangle)\\ =&\frac{1}{2}\begin{pmatrix}1 & 0 & 0 & 0\\ 0 & 0 & 1 & 0\end{pmatrix}\begin{pmatrix}1 & 0 & 0 & 1\\ 0 & 0 & 0 & 0\\ 0 & 0 & 0 & 0\\ 1 & 0 & 0 & 1\end{pmatrix}\begin{pmatrix}1 & 0\\ 0 & 0\\ 0 & 1\\ 0 & 0\end{pmatrix} +\\ &\frac{1}{2}\begin{pmatrix}0 & 1 & 0 & 0\\ 0 & 0 & 0 & 1\end{pmatrix}\begin{pmatrix}1 & 0 & 0 & 1\\ 0 & 0 & 0 & 0\\ 0 & 0 & 0 & 0\\ 1 & 0 & 0 & 1\end{pmatrix}\begin{pmatrix}0 & 0\\ 1 & 0\\ 0 & 0\\ 0 & 1\end{pmatrix}\end{aligned}$$

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Loss of precision

Computing a reduced density matrix **discards information**!

$$\rho_{\beta_{00}} = \frac{|00\rangle\langle 00| + |00\rangle\langle 11| + |11\rangle\langle 00| + |11\rangle\langle 11|}{2} \qquad \text{(Pure state)}$$

$$\rho_2 = \frac{|00\rangle\langle 00| + |01\rangle\langle 01| + |10\rangle\langle 10| + |11\rangle\langle 11|}{4} \qquad \text{(Mixed state)}$$

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$$Tr_B[\rho_{\beta_{00}}] = \frac{|0\rangle\langle 0| + |1\rangle\langle 1|}{2} = Tr_B[\rho_2]$$

The partial traces of two different initial states can be equal.

Moreover, for a state $\rho \in A \otimes B$, even if we know $Tr_B[\rho]$ and $Tr_A[\rho]$, we cannot uniquely determine ρ.

Abstract Domain

$$D = \mathbb{C}^{2^n} \otimes \mathbb{C}^{2^n}, \quad S = (s_1, ..., s_m), \quad 1 \leq m \leq 2^n, \quad s_i \subseteq [n]$$

$$AbsDom(S) = \left\{ (P_{s_1}, ..., P_{s_m}) \mid P_{s_i} \text{ is a projection in } \mathbb{C}^{2^{|s_i|}} \otimes \mathbb{C}^{2^{|s_i|}} \right\}$$