

4.1: What are the four steps CPUs use to execute instructions

1. Fetch instruction address from PC
2. Parse instruction for opcode and src/dest for example
3. Do the opcode with ALU (use something other than ALU if modern, FPU maybe?)
4. Store/write result according to instructions

The CPU executes each instruction in a series of small steps. Roughly speaking, the steps are as follows:

1. Fetch the next instruction from memory into the instruction register.
2. Change the program counter to point to the following instruction.
3. Determine the type of instruction just fetched.
4. If the instruction uses a word in memory, determine where it is.
5. Fetch the word, if needed, into a CPU register.
6. Execute the instruction.
7. Go to step 1 to begin executing the following instruction.

4.2:

Dette er gjort fordi at ud af de 9 registre, så er der kun en der må komme ud til B bus, så at vælge registre ud fra 0000-1111. C bus skal skrive til flere registre på en gang, så det er bedre at opdele sådan at hver bit har en betydning. for eksempel hvis vi har 1010 , så skriver vi til MAR og PC

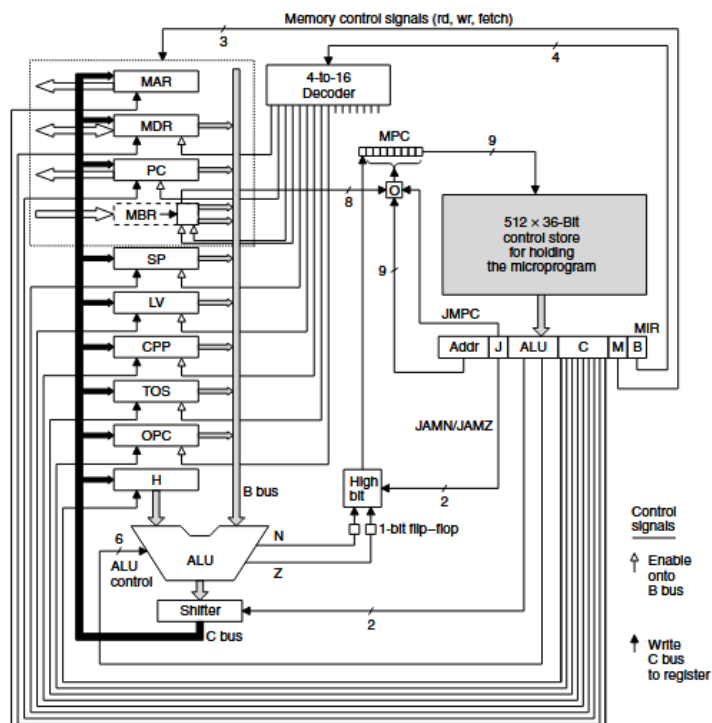


Figure 4-6. The complete block diagram of our example microarchitecture, the Mic-1.

4.4

ILOAD J

ILOAD M

IADD
BIPUSH (1000)
IADD
ISTORE I

Og herefter kan vi bruge omvendt polsk notation

ILOAD J
ILOAD M
BIPUSH (1000)
IADD
IADD
ISTORE I

4.8

Label	Operations	Comments
Main1	PC = PC + 1; fetch; goto (MBR)	MBR holds opcode; get next byte; dispatch
nop1	goto Main1	Do nothing
iadd1	MAR = SP = SP - 1; rd	Read in next-to-top word on stack
iadd2	H = TOS	H = top of stack
iadd3	MDR = TOS = MDR + H; wr; goto Main1	Add top two words; write to top of stack
isub1	MAR = SP = SP - 1; rd	Read in next-to-top word on stack
isub2	H = TOS	H = top of stack
isub3	MDR = TOS = MDR - H; wr; goto Main1	Do subtraction; write to top of stack
iand1	MAR = SP = SP - 1; rd	Read in next-to-top word on stack
iand2	H = TOS	H = top of stack
iand3	MDR = TOS = MDR AND H; wr; goto Main1	Do AND; write to new top of stack
ior1	MAR = SP = SP - 1; rd	Read in next-to-top word on stack
ior2	H = TOS	H = top of stack
ior3	MDR = TOS = MDR OR H; wr; goto Main1	Do OR; write to new top of stack
dup1	MAR = SP = SP + 1	Increment SP and copy to MAR
dup2	MDR = TOS; wr; goto Main1	Write new stack word
pop1	MAR = SP = SP - 1; rd	Read in next-to-top word on stack
pop2		Wait for new TOS to be read from memory
pop3	TOS = MDR; goto Main1	Copy new word to TOS
swap1	MAR = SP - 1; rd	Set MAR to SP - 1; read 2nd word from stack
swap2	MAR = SP	Set MAR to top word
swap3	H = MDR; wr	Save TOS in H; write 2nd word to top of stack
swap4	MDR = TOS	Copy old TOS to MDR
swap5	MAR = SP - 1; wr	Set MAR to SP - 1; write as 2nd word on stack
swap6	TOS = H; goto Main1	Update TOS
bipush1	SP = MAR = SP + 1	MBR = the byte to push onto stack
bipush2	PC = PC + 1; fetch	Increment PC, fetch next opcode
bipush3	MDR = TOS = MBR; wr; goto Main1	Sign-extend constant and push on stack
iload1	H = LV	MBR contains index; copy LV to H
iload2	MAR = MBRU + H; rd	MAR = address of local variable to push
iload3	MAR = SP = SP + 1	SP points to new top of stack; prepare write
iload4	PC = PC + 1; fetch; wr	Inc PC; get next opcode; write top of stack
iload5	TOS = MDR; goto Main1	Update TOS
istore1	H = LV	MBR contains index; copy LV to H
istore2	MAR = MBRU + H	MAR = address of local variable to store into
istore3	MDR = TOS; wr	Copy TOS to MDR; write word
istore4	SP = MAR = SP - 1; rd	Read in next-to-top word on stack
istore5	PC = PC + 1; fetch	Increment PC; fetch next opcode
istore6	TOS = MDR; goto Main1	Update TOS
wide1	PC = PC + 1; fetch;	Fetch operand byte or next opcode
wide2	goto (MBR OR 0x100)	Multiway branch with high bit set
wide_ild1	PC = PC + 1; fetch	MBR contains 1st index byte; fetch 2nd
wide_ild2	H = MBRU << 8	H = 1st index byte shifted left 8 bits
wide_ild3	H = MBRU OR H	H = 16-bit index of local variable
wide_ild4	MAR = LV + H; rd; goto iload3	MAR = address of local variable to push
wide_istore1	PC = PC + 1; fetch	MBR contains 1st index byte; fetch 2nd
wide_istore2	H = MBRU << 8	H = 1st index byte shifted left 8 bits
wide_istore3	H = MBRU OR H	H = 16-bit index of local variable
wide_istore4	MAR = LV + H; goto istore3	MAR = address of local variable to store into
ldc_w1	PC = PC + 1; fetch	MBR contains 1st index byte; fetch 2nd
ldc_w2	H = MBRU << 8	H = 1st index byte << 8
ldc_w3	H = MBRU OR H	H = 16-bit index into constant pool
ldc_w4	MAR = H + CPP; rd; goto iload3	MAR = address of constant in pool

Vi går også igennem MAIN

Så ADD(4) 2*LOAD(6) STORE(7)= 23 clock cycles

$$\frac{1}{4 \cdot 10^{12}}$$

$$2.500000000 \cdot 10^{-13} \cdot 23$$

$$250.00 \times 10^{-15}$$

$$5.750000000 \times 10^{-12}$$