

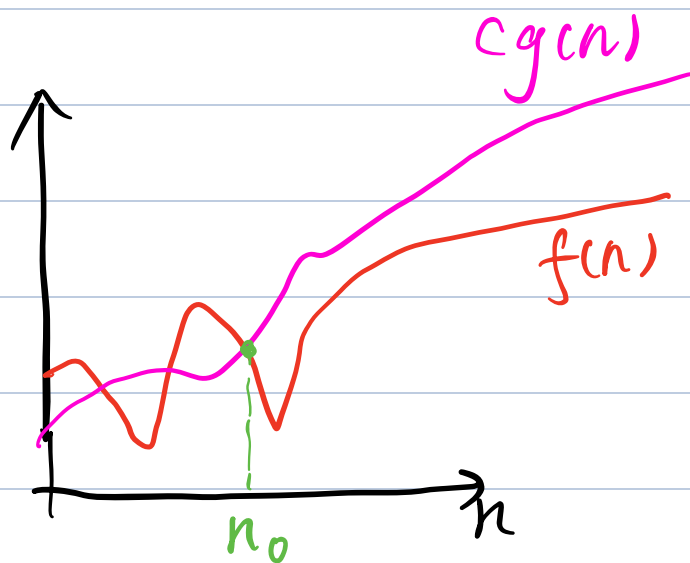
# Growth of function

$O$ ,  $\omega$ ,  $\Theta$ ,  $o$ ,  $w$

## $O$ -notation

$O(g(n)) = \{f(n) : \text{there exist positive constants } c \text{ and } n_0 \text{ such that}$   
 $0 \leq f(n) \leq c g(n)$   
 $\text{for all } n \geq n_0\}$

$g(n)$  is an 渐进性 asymptotic upper bound for  $f(n)$



Ex.  $2n^2 \in O(n^3)$ ,  $c=1$ ,  $n_0=2$

$$2n^2 \leq 1 \cdot n^3, \quad n \geq 2$$

$$2n^2 = O(n^3)$$

$$2n^2 = O(n^2), \quad c, n_0$$

$$0 \leq 2n^2 \leq cn^2, \quad c=3, n_0=1$$

$$\leq 3n^2$$

$$n^2 = O(n^2), \quad n^2 + n = O(n^2)$$

$$0 \leq n^2 + n \leq cn^2, \quad c, n_0$$

$$c=2, n_0=2$$

$$0 \leq n^2 + n \leq 2n^2, \quad \text{if } n \geq 2$$

$$n^2 + 1000n = O(n^2), \quad 1000n^2 + 1000n = O(n^2)$$

$$n = O(n^2), \quad n/1000 = O(n^2)$$

$$\frac{n^2}{\lg \lg n} = O(n^2), \quad n^{1.999999} = O(n^2)$$

$$c, n_0$$

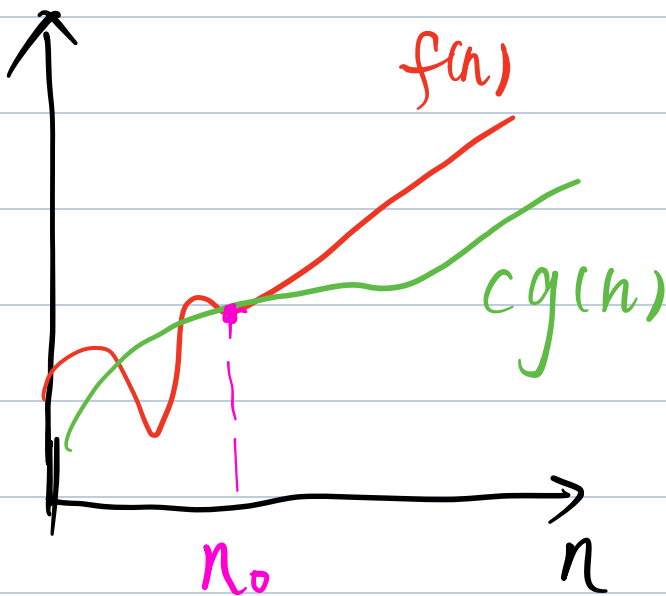
$\Omega$  - notation

$\Omega(g(n)) = \{f(n) : \text{there exist positive constants } c \text{ and } n_0, \text{ such that}$

$$0 \leq c g(n) \leq f(n)$$

for all  $n \geq n_0\}$

$g(n)$  is an asymptotic lower bound for  $f(n)$ .



Ex.  $\sqrt{n} = \Omega(\lg n)$ ,  $c=1$ ,  $n_0=16$

$$n^2 = \Omega(n^2). \quad C = 0.5, \quad n_0 = 1$$

$$0 \leq 0.5n^2 \leq n^2 \quad \text{for } n \geq 1$$

$$n^2 + n = \Omega(n^2), \quad n^2 - 10n = \Omega(n^2)$$

$$n^2 + 10000n = \Omega(n^2), \quad n^2 - 10000n = \Omega(n^2)$$

$$n^3 = \Omega(n^2), \quad n^{2.00001} = \Omega(n^2)$$

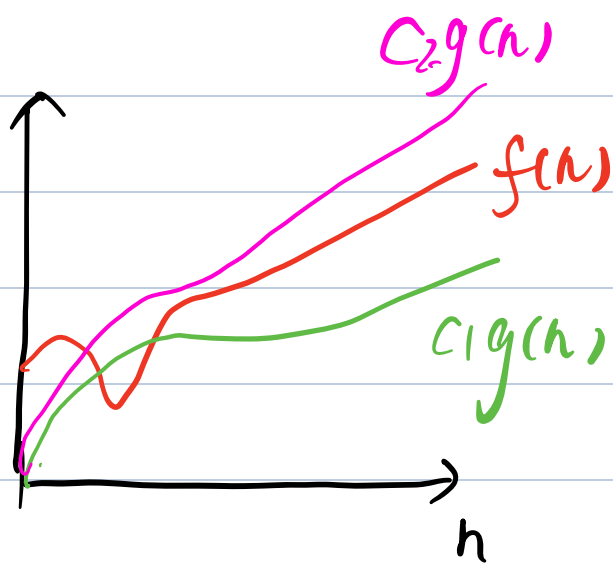
$$n^2 \lg \lg n = \Omega(n^2), \quad 2^{2n} = \Omega(n^2)$$

⊕ - notation

⊕(g(n)) = { f(n) : there exist positive constants  $C_1$ ,  $C_2$  and  $n_0$  such that

$$0 \leq C_1 g(n) \leq f(n) \leq C_2 g(n)$$

for all  $n \geq n_0$  }



$g(n)$  is an asymptotic **tight bound** for  $f(n)$ .

Theorem.

$f(n) = \Theta(g(n))$  if and only if

$f(n) = O(n)$ , and  $f(n) = \Omega(n)$

$$O(1) = O(n)$$

step 1:  $n=1$ ,  $O(1) = O(1)$

step 2: Suppose  $k < n$   
 $O(1) = O(k)$

$$O(n) = O(n-1) + O(1)$$

$$= O(1) + O(1)$$

$$= \underline{O(1)}$$

$$O(n) = O(1)$$

O-notation

$O(g(n)) = \{f(n) : \text{for all constants } c > 0,$   
 there exist a constant  $n_0 > 0$   
 such that

$$0 < f(n) < c g(n)$$

for all  $n \geq n_0\}$

$$\lim_{n \rightarrow \infty} \frac{f(n)}{g(n)} = 0$$

Ex.  $\underline{n^{1.99999}} = O(\overset{g}{n^2})$        $\lim_{n \rightarrow \infty} \frac{n^{1.99999}}{n^2} = 0$

$\uparrow$

$$\frac{n^2}{\ln n} = O(n^2) \quad \cdot \quad \frac{n^2}{10000} \neq O(n^2)$$

$\omega$ -notation

$\omega(g(n)) = \{f(n) : \text{for all constants } c > 0, \text{ there exist } n_0 > 0 \text{ such that } 0 < c g(n) < f(n) \text{ for all } n \geq n_0.\}$

$$\lim_{n \rightarrow \infty} \frac{f(n)}{g(n)} = \infty$$

Ex.  $n^{2.00001} = \omega(n^2)$ ,  $\lim_{n \rightarrow \infty} \frac{n^{2.00001}}{n^2} = \infty$

$$n^2 \lg(n) = \omega(n^2)$$

$$n^2 \neq \omega(n^2)$$

Relation properties:

Transitivity:

$$f(n) = \Theta(g(n)) \text{ and } g(n) = \Theta(h(n)) \Rightarrow f(n) = \Theta(h(n))$$

Same for  $O$ ,  $\Omega$ ,  $\Theta$ ,  $\omega$

Reflexivity:

$$f(n) = \Theta(f(n))$$

Same for  $O$ ,  $\Omega$

Symmetry:

$$f(n) = \Theta(g(n)) \text{ if and only if } g(n) = \Theta(f(n))$$

Transpose Symmetry:  $\text{iff} = \text{iff}$  and only if

$$f(n) = O(g(n)) \text{ iff } g(n) = \Omega(f(n))$$

$$f(n) = o(g(n)) \text{ iff } g(n) = \omega(f(n))$$