

TLS, CERTS, KERBEROS

CS 361S

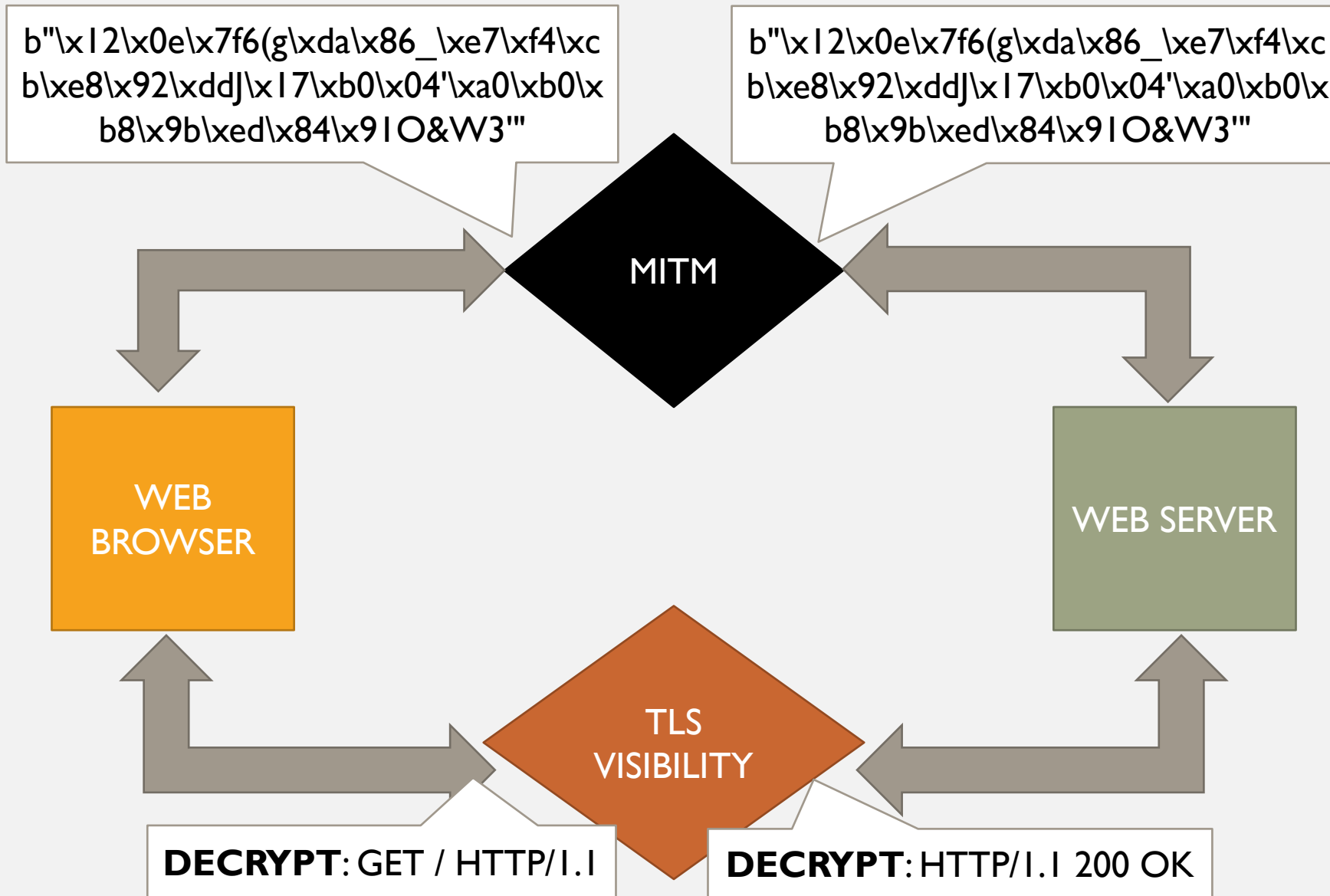
Fall 2020

Seth James Nielson

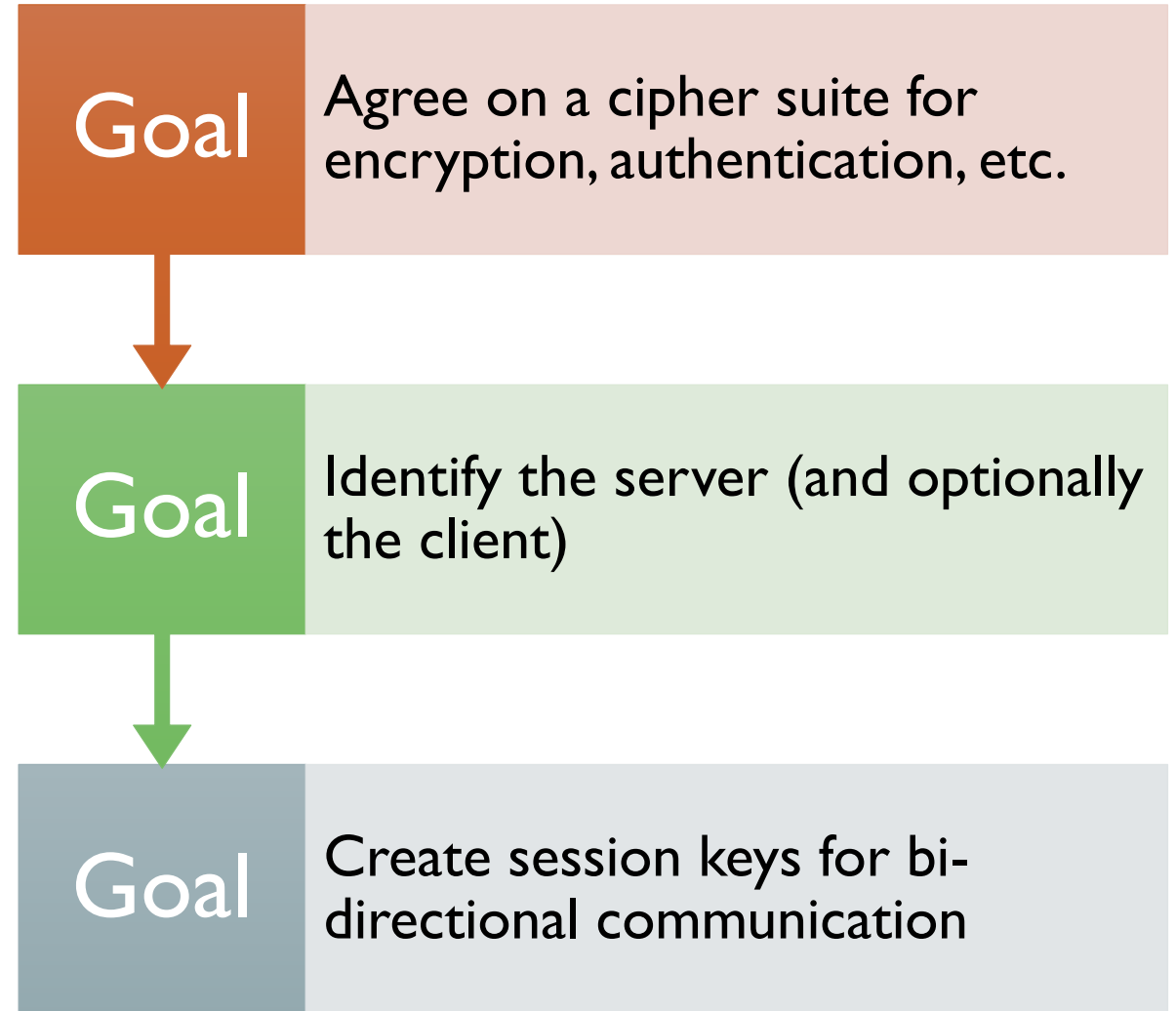
The background of the slide is a dark, high-contrast image. On the left side, there is a close-up of a combination lock with several dials visible, each showing numbers. The lock is metallic and appears to be part of a larger device. To the right and below the lock, there is a blurred view of electronic circuitry, showing various components like resistors and capacitors on a printed circuit board. The overall lighting is dim, with the lock and circuitry providing a technical, security-related aesthetic.

TLS VISIBILITY































- TLS is designed to provide END-TO-END “security”
- MITM should NOT be able to read/modify/forgo data
- TLS Visibility “breaks” this for “authorized” purposes



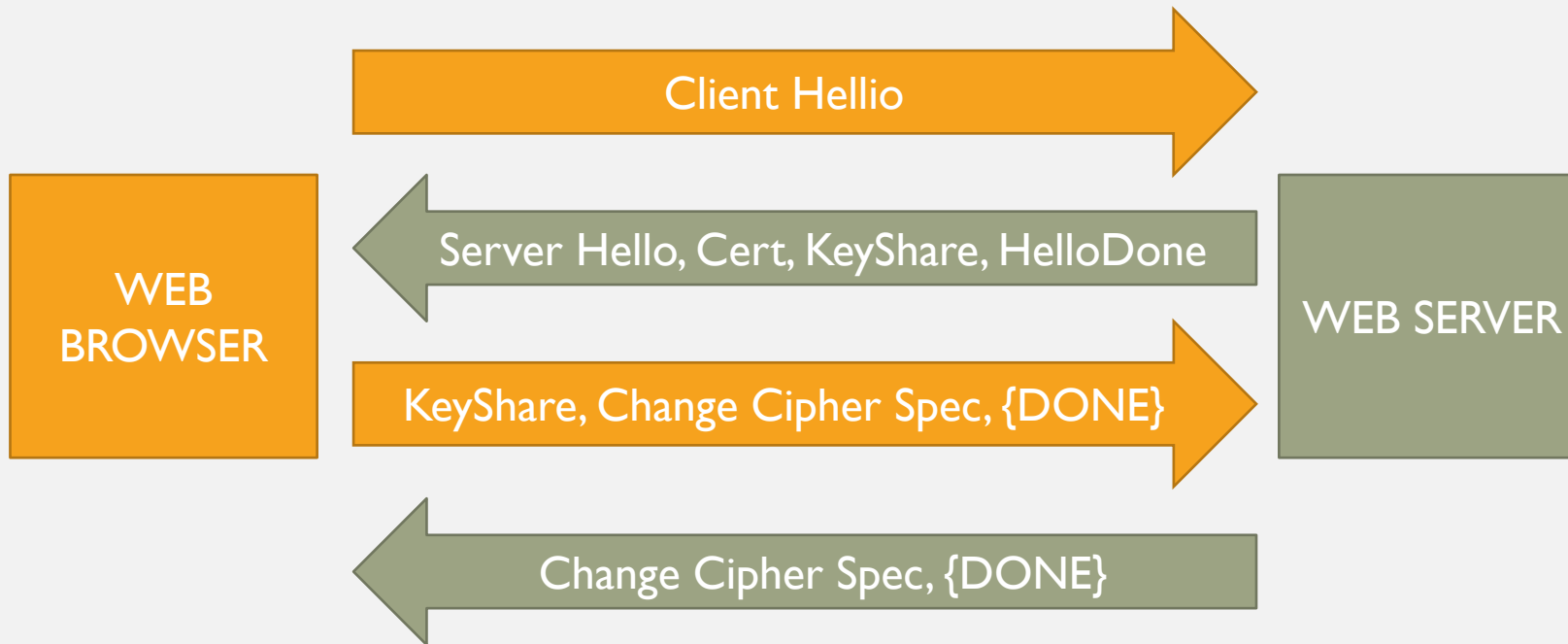
TLS 1.2 HANDSHAKE



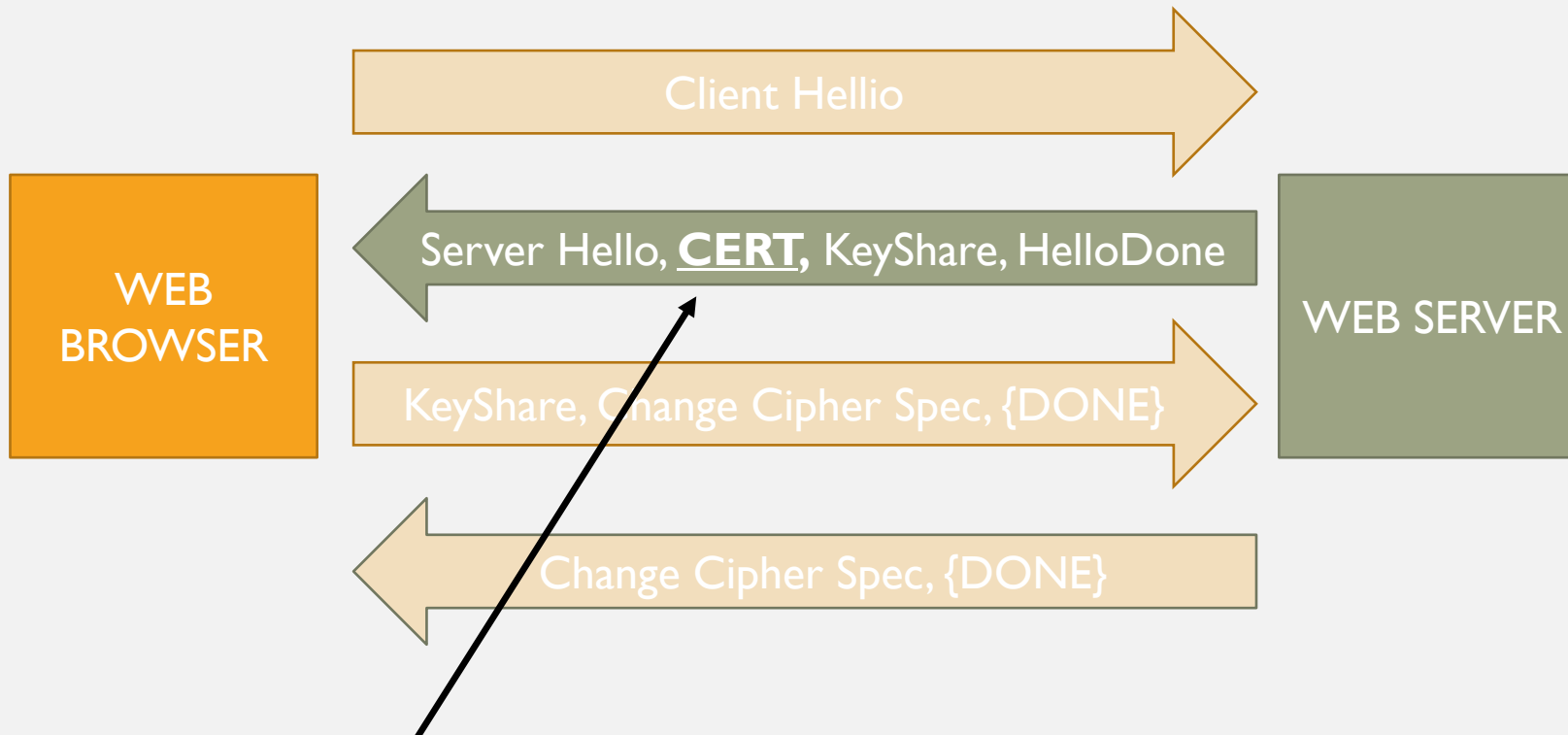
TLS 1.2 HANDSHAKE REVIEW

| Step | Client | Direction | Message | Direction | Server |
|------|--|---|---------------------|---|---|
| 1 |  | | Client Hello |  |  |
| 2 |  |  | Server Hello | |  |
| 3 |  |  | Certificate | |  |
| 4 |  |  | Server Key Exchange | |  |
| 5 |  |  | Server Hello Done | |  |
| 6 |  | | Client Key Exchange |  |  |
| 7 |  | | Change Cipher Spec |  |  |
| 8 |  | | Finished |  |  |
| 9 |  |  | Change Cipher Spec | |  |
| 10 |  |  | Finished | |  |

END-TO-END HANDSHAKE VISUALIZATION #2



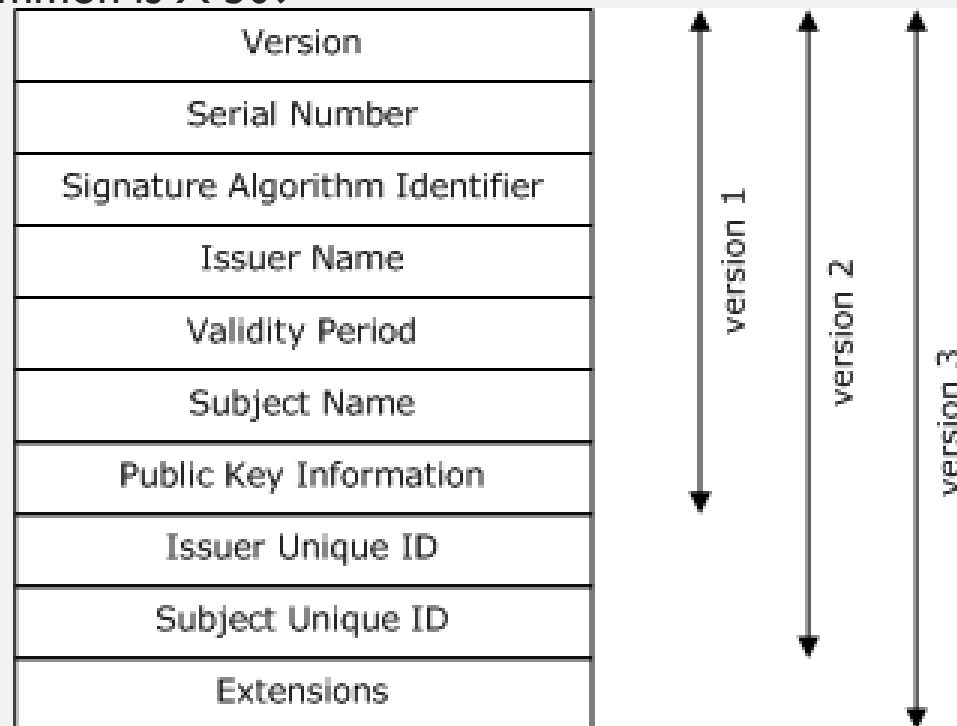
AUTHENTICATION



The “Certificate” message includes ONE OR MORE certificates.

WHAT IS A CERTIFICATE?

- TLS specification (RFC) doesn't specify cert or cert verification
- The most common is X 509



CERTIFICATE VERIFICATION

WEB
BROWSER

Verify “amazon.com” is the URL
Verify the validity period
(Other Verification)
Who issued the cert?

CERTIFICATE

Subject CN: amazon.com
Not Valid Before: 2001
Not Valid After: 2030
Issued By: amazon CA
Signature Blob: <sig>

PUBLIC KEY PRIVATE KEY



PRIVATE KEY

CERTIFICATE CHAINS

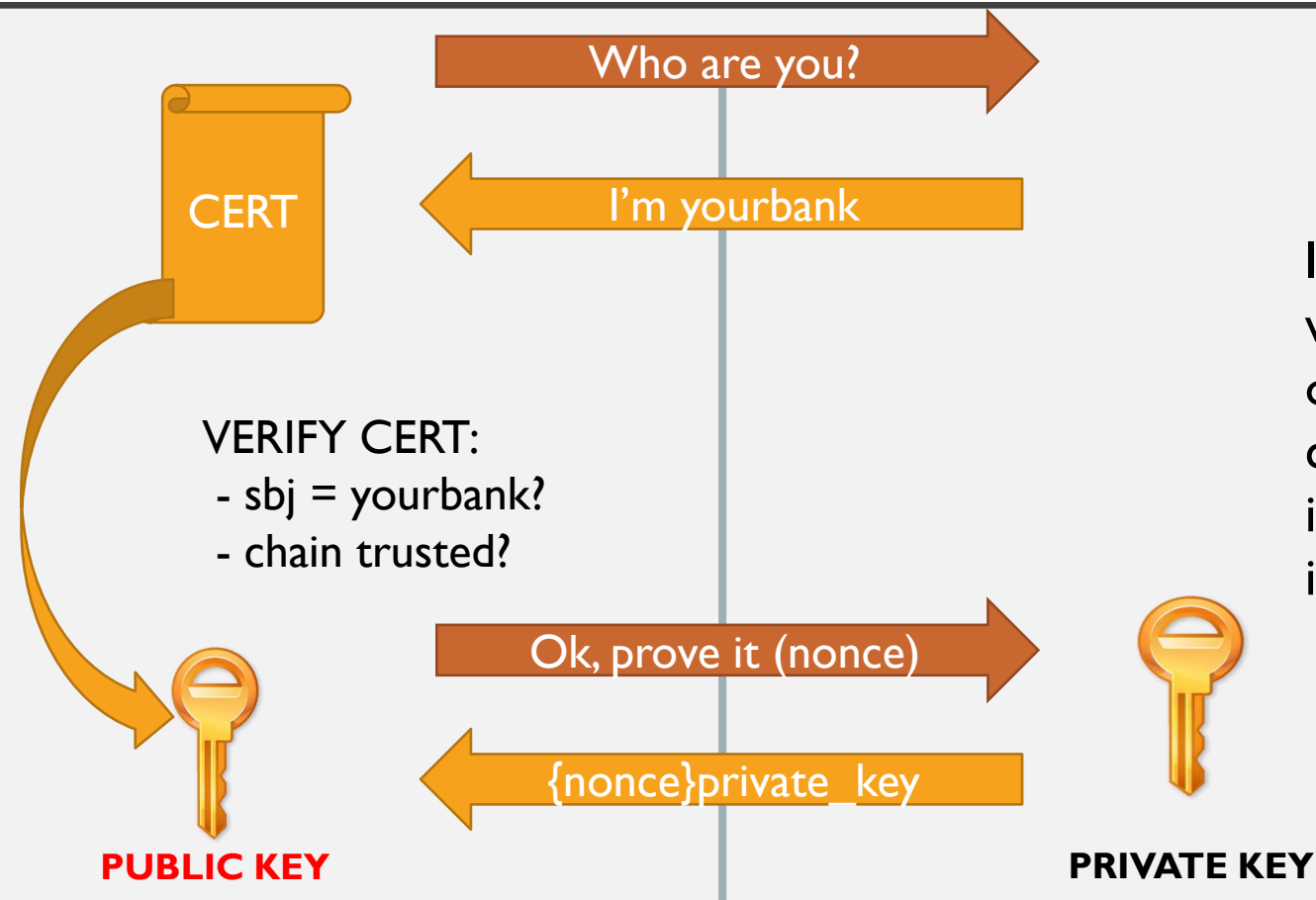
The certificate for the Host may be signed by an INTERMEDIATE Certificate Authority

Because the web browser probably doesn't have this intermediate cert, the TLS handshake includes both certificates.

Subject CN: amazon CA
...
Issued By: GlobalSign
Signature Blob: <sig>

Subject CN: amazon.com
...
Issued By: amazon CA
Signature Blob: <sig>

PROVING IDENTITY



In TLS, the “nonce” is just wrapped up with the other data, such as the client hello, which is all included in the final hash in the finished message.

ROOT CA CERTIFICATES

Certificate chains MUST
have a ROOT



A Root Certificate is SELF
SIGNED



Browsers trust a set of root
certificates
AXIOMATICALLY



Certificate chains must have
a trust chain to one of these
roots.

TRUSTING DIFFIE HELLMAN

Recall that DH keys are EPHEMERAL

The Server's cert includes a long-term public key

The Server's DH key is signed by this key pair

IF the client trusts the cert, THEN it can validate the DH key

TLS BULK TRANSPORT

Both Client
and Server
derive keys

Encryption
keys AND
MAC keys

MAC's ensure
continuous
authentication

WHEN A TLS
MESSAGE IS
RECEIVED:

The sender is “proved” by
the MAC

The MAC is “proved” via
MAC key derived from DH

Server’s DH key “proved”
authentic by cert signature

Certificate “proved”
authentic by chain to trusted
root

IT ALL DEPENDS ON THE CERT

IF a browser trusts
MY certificate to be
Amazon's certificate

- THEN the browser
will trust my DH
public key

IF the browser trusts
my DH public key

- THEN the browser
will derive the same
MAC key I do

IF the browser
derives the same
MAC key I do

- THEN the browser
will believe my
messages are from
Amazon

CERTIFICATE REVOCATION

- How do you revoke a certificate?
- Difficult: so long as the cert is properly signed, it is believed
- You can publish certificate revocation lists:
 - Uses just serial number
 - So make sure your serial numbers are actually unique!
 - But, until the new CRL is received, bad cert still usable

ONLINE CERTIFICATE STATUS PROTOCOL (OCSP)

- Certificates were designed to be used offline
- However, modern security constraints often necessitate OCSP
- Client can ask a server ('OCSP Responder') about a cert
 - Server can respond "Good", "Revoked", "Unknown"
 - Response is signed; however, **vulnerable to replay attacks!**
 - An extension permits nonces, but often not used for efficiency
 - Also, potential privacy loss
 - But, more efficient and timely than CRL

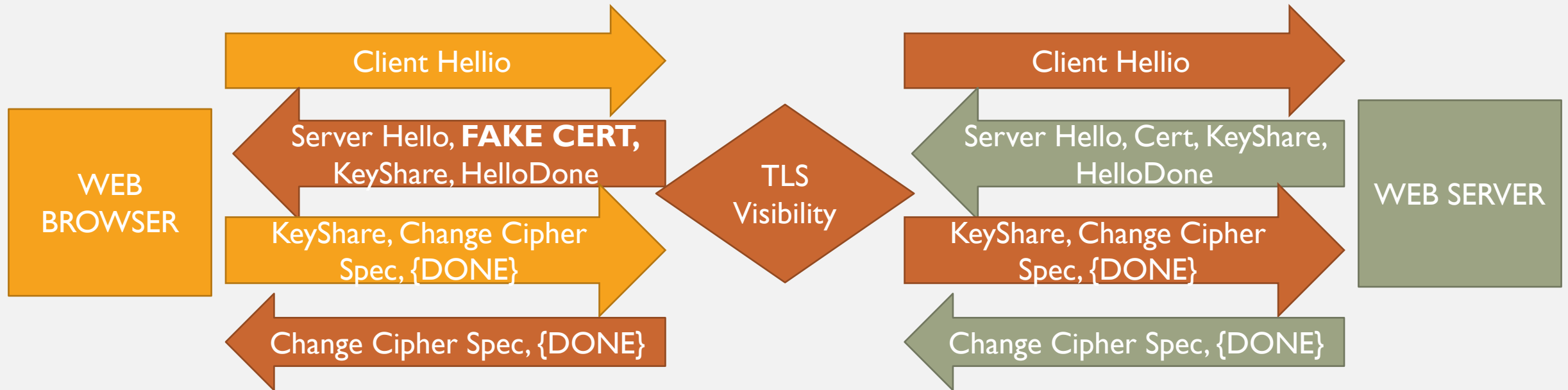
OTHER ALTERNATIVES TO TRUST?

- Sadly, there is no known way to create trust out of thin air
- In almost every case, there must be a trust basis:
 - Out-of-band communication (e.g., in real life)
 - Evolutionary trust over time with long-term identifiers
 - Third parties, including CA's, authentication/reputation servers
 - Crowds, such as distributed ledger

TLS VISIBILITY

- Typically, a browser/client **MUST** have a new root CA installed
- This root CA is a self-signed certificate from the Visibility appliance
- The appliance can now generate **ANY** cert and the browser believes it!
- We will discuss the huge security concerns in a later lecture

TLS VISIBILITY HANDSHAKE VISUALIZATION



KERBEROS

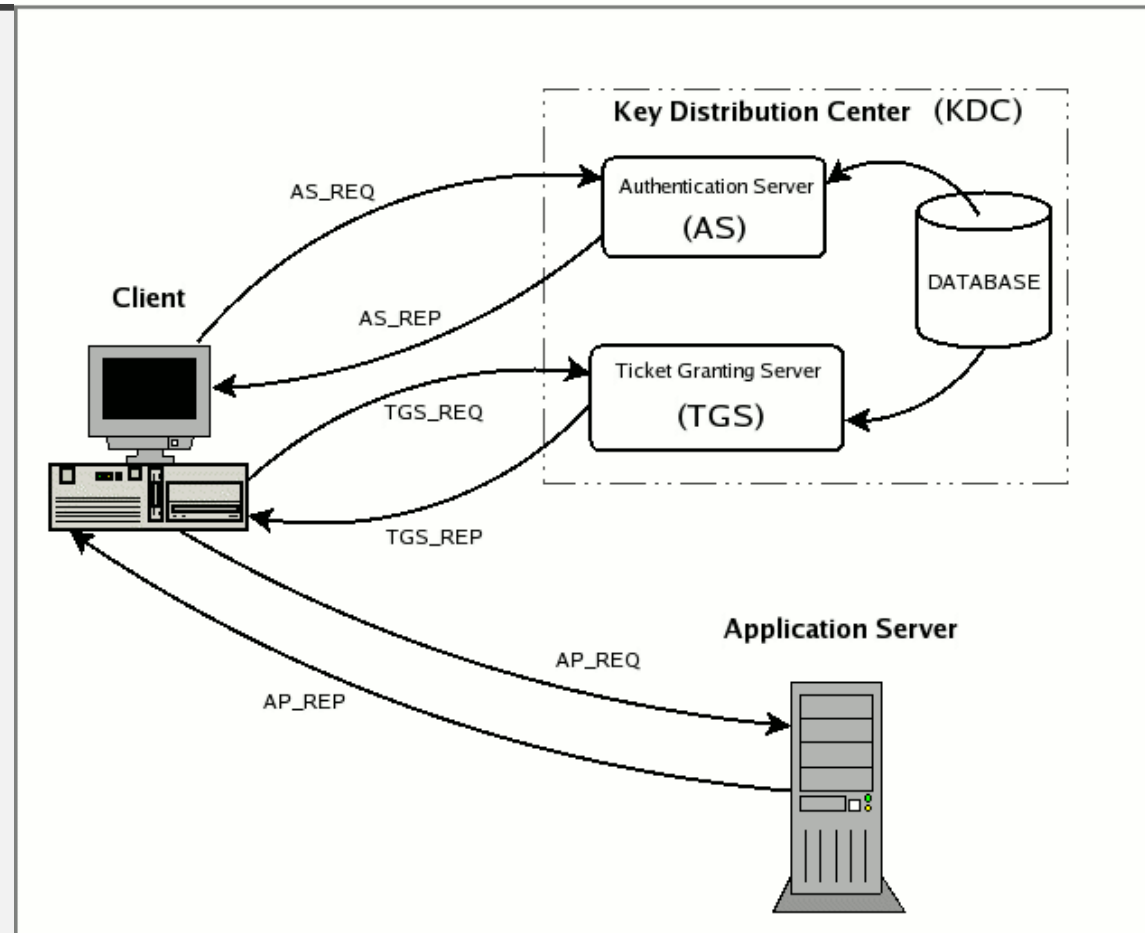
Kerberos vs TLS

- Kerberos uses trusted authentication server
- Must be online.
- If compromised, entire system compromised
- Mutual authentication, confidentiality

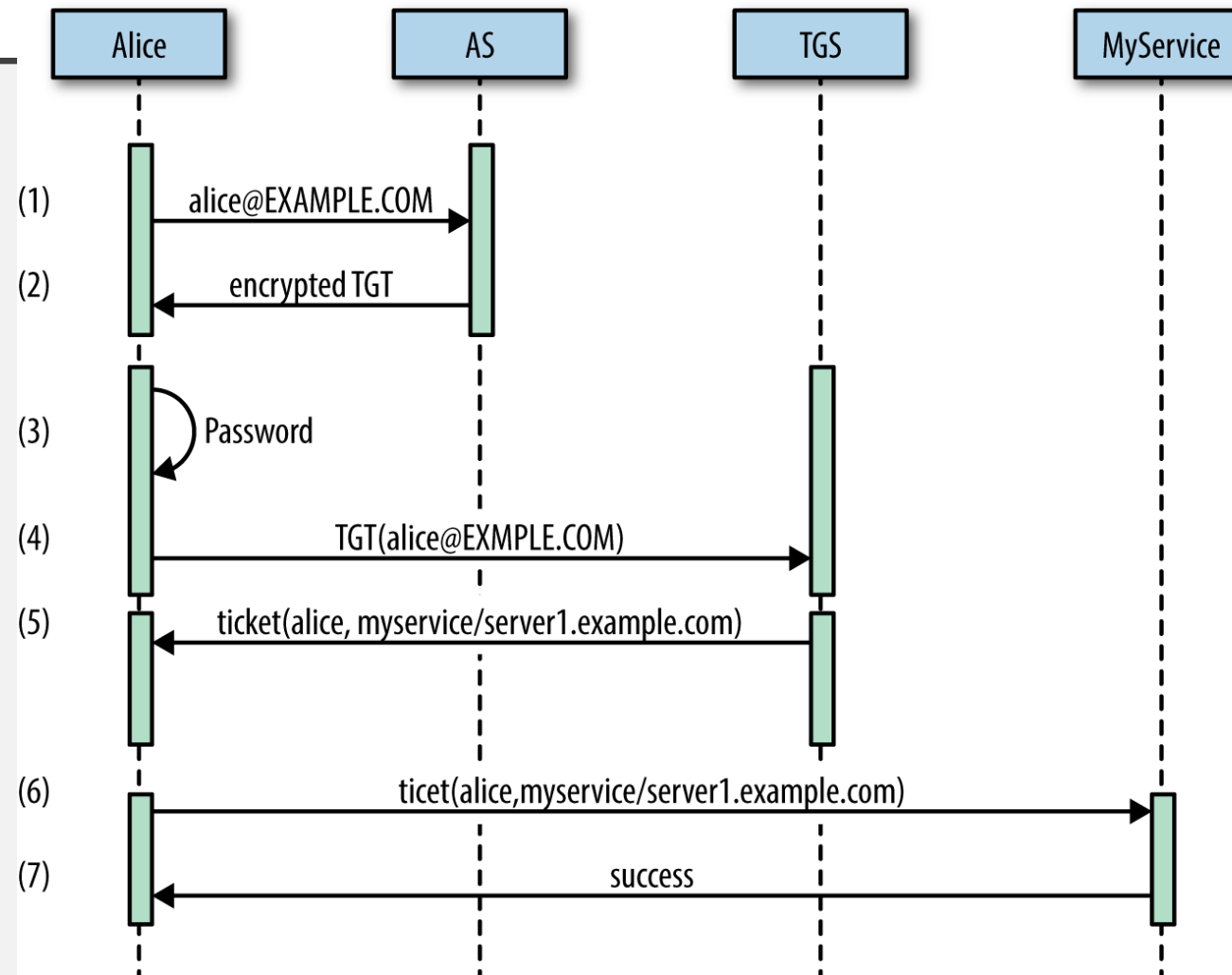
Basic components:

- Authentication Server
- Key Distribution Server (KDS)
- Ticket Granting Service
- Service Server

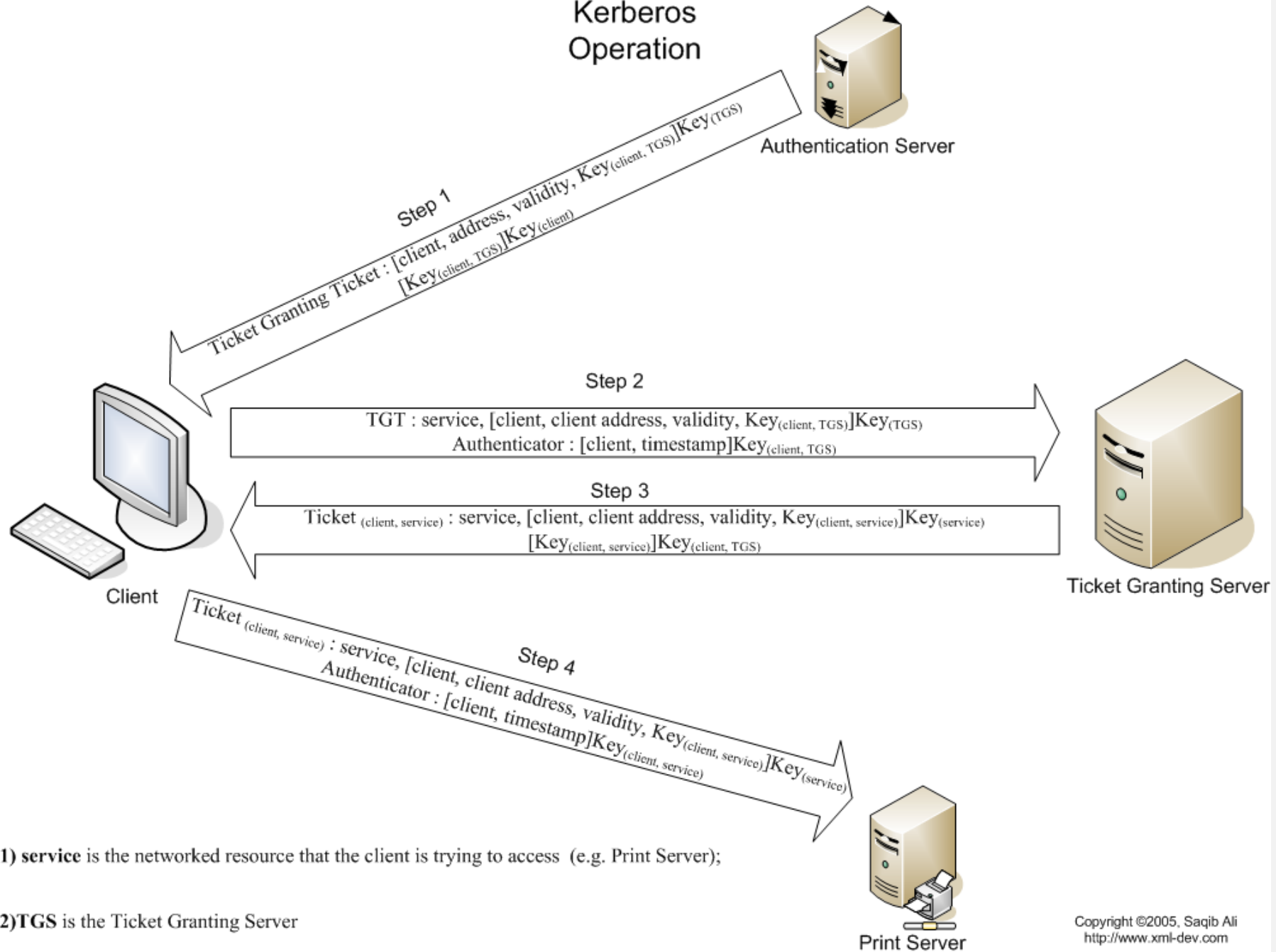
KERBEROS COMMUNICATION



KERBEROS PROTOCOL



Kerberos Operation



PROTOCOL PRINCIPLES

Note that the user's key never goes over the wire

Note that pre-encrypted messages can be sent.

- AS sends a message to A that only TGT can decrypt
- Thus, TGT knows that the message sent by A MUST come from AS

How scalable is this system?