

# CS411

## SQL/Mongo

Trigger:

```
create trigger [trigger_name]
[before | after]
{insert | update | delete}
on [table_name]
REFERENCING
NEW ROW AS NewTuple
[for each row]
[trigger_body]

INSERT INTO Drinks(name)
VALUES(NewTuple.drink);

// Use [NEW / OLD] and [ROW / TABLE] AS <name>.
// INSERT statements imply a new tuple (for row-level)
// or new set of tuples (for statement-tuple)
// DELETE implies an old tuple or set of tuples
// UPDATE implies both (the row before the update, and
// the row after the update)
```

Procedure:

```
CREATE PROCEDURE Result()
BEGIN
    DECLARE done int default 0;
    DECLARE cursor_container VARCHAR(10);
    DECLARE temp VARCHAR(255);
    DECLARE cur CURSOR FOR SELECT DISTINCT netid FROM
        Enrollments;
    DECLARE CONTINUE HANDLER FOR NOT found SET done =
        1;

    DROP TABLE IF EXISTS NewTable;

    CREATE TABLE NewTable(
        NetId VARCHAR(10),
        Course_Load_Status VARCHAR(255),
        PRIMARY KEY (NetId)
    );

    OPEN cur;
    REPEAT
        FETCH cur INTO cc;
        IF (query1) > 12 THEN SET temp = 'Maximum';
        ELSEIF (query2) = 12 THEN SET temp='Full';
        ELSE SET temp='Minimum';
        END IF;
        INSERT IGNORE INTO NewTable
        VALUES (cc, temp);
    UNTIL DONE
    END REPEAT;
    CLOSE cur;
    SELECT FROM NewTable;END
```

Mep Reduce:

```
let mapper = function() {
    for(i=0; i < this.actors.length; i++) {
        emit(this.actors[i], this.ratings);
    }
};

let reducer = function(stack_id, values_arr) {
    sum = 0;
    for (i = 0; i < values_arr.length; i++) {
        sum += values_arr[i];
    }
    if (values_arr.length) {
        avg = sum/values_arr.length;
        return avg;
    } else {return 0;}
};

let mapReduceOutput = db.Movies.mapReduce(mapper,
    reducer, {out:"times", query: {release_year:{$lt:
    2000}}});
let updateRes = db.times.update({}, {$rename: {"value":
    "avg_ratings"}}), false, true);

db.times.find();
```

Sample Mongo Query

```
db.Movies.aggregate([
    {$match: {$and: [{ratings: {$gte: 9}}, {
        release_year: {$gt: 2000}}]}},
    {$unwind: "$actors"},
    {$group: {_id:"$actors", numMovies:{$sum: 1}}},
    {$sort : {numMovies : -1}},
    {$project: {_id:0, actorId:"$_id", numMovies:1}}
]);
```

## ER Models

Rectangle means entity set  
Diamond represents a relationship set between two entity sets  
Oval is an attribute name.  
A triangle with "isa" points from a subclass to a parent class  
 $a \rightarrow b$  means there is at most 1 b for an a, but the b can have 0 to infity as. For example, a is the employee and b is the department. (many-to-one)  
 $a \leftrightarrow b$  means at most one entity is connected to at most one entity. For example, a is the employee and b is a laptop id. (one-to-one)  
 $a \dashrightarrow b$  means there are no constraints. Could be anything! (many-to-many)  
 $a \rightarrow b$  means that there is exactly one b for an a.

## Converting Model to DDL

# NOTE: The primary key of table a is a\_pk  
Multiple attributes with the same name: check if all those names are duplicate. If so, you can add those attributes as 1 attribute to that entity. Otherwise create a different entity called attr-assignment that includes the primary key of the first entity and the attribute value. For example, Employee(SSN, E-Name, Office, Office) becomes Employee(SSN, E-Name) and Office-Assignment(SSN, Office), and make SSN the foreign key.  
Many-to-Many relationship. Suppose there is a many-to-many relationship from a to b called c. Then create a new table called c where c\_pk = a\_pk, b\_pk where a\_pk and b\_pk are foreign keys  
Many-to-One relationship. Suppose there is a many-to-one relationship from a to b called c. Here a has only one b associated with it. Then create a new table called c where c\_pk = a\_pk where a\_pk and b\_pk are foreign keys. You can also simply add b\_k to a as an attribute. If c has attributes, do it the former way.  
Weak Entity Set. Suppose you have a weak entity set a and an entity set b that is connected. Now, when you create b, simply add b\_pk as part of the primary key of a. Then follow the many-to-X rules. Then add FOREIGN KEY(b\_pk) REFERENCES b, ON DELETE CASCADE  
Subclasses: Suppose b and c are subclasses of a. Simply create a table for a-c and a-b.

## Normal Forms

# BCNF (lossless)  
1. Find every FD (use the process of finding the full set of FDs)  
2. Check that A of every non-trivial FD in the form  $A \rightarrow B$  is a superkey, meaning it uniquely identifies every attribute in the table. If so, you are done!  
3. If not, compute  $A^+$ .  
4. Decompose it by chopping off everything that the left side does not uniquely determine and put it in a separate table with the left side attributes. This new table has no FDs. So you have table  $R_1 = A^+$ ,  $R_2 = A$  union  $(R - A^+)$ .  
5. For each table, repeat from 1 to 4.  
# 3NF (lossless and dependency preserving)  
1. Find the key  
2. Get a minimal basis G of given FDs  
3. For each FD  $A \rightarrow B$  in the minimal G, use AB as the schema of a new relation  
4. If none of the schemas from step 2 is a superkey, add another relation whose schema is a key for the original relation.