ECE220: Computer Systems and Programming

Past Exam

- 1. This is a closed book exam except for 1 sheet of hand-written notes
- 2. You may not use any personal electronic devices, such as cellphone and calculator
- 3. Absolutely no interaction between students is allowed
- 4. Illegible handwriting will be graded as incorrect

Name:	
NetID:	
Room:	
Question 1 (36 points):	
Question 2 (22 points): A; B; C	
Question 3 (42 points): 3.1; 3.2; 3.3_	; 3.4
Total Score:	

Problem 1 (30 points): Check for Pythagorean Theorem

In this question, you will complete a program which checks the following condition: $C^2 = A^2 + B^2$, where A, B and C are three input numbers

More details:

- 1. The inputs (A, B, C) are already stored in a **stack** as shown below and they are all positive numbers and small enough that there will be no arithmetic overflow even when squared.
- 2. R6 is the stack pointer, pointing at the next available location on the stack. You should implement stack operations directly using R6 rather than calling the PUSH and POP subroutines as in MP2.
- 3. The subroutine SQUARE takes one input from R4 and sets the output in R5 (R5 = R4*R4).
- 4. If $C^2 = A^2 + B^2$, R0 is set to 1. Otherwise, R0 is set to 0 (this part has been done for you).
- 5. Push the value in R0 onto the stack.
- 6. Assume no stack overflow nor underflow.

x3FFA		
x3FFB		
x3FFC		
x3FFD		← R6
x3FFE	A	
x3FFF	В	
x4000	С	

Each line of code is worth 2 points.

.ORIG x3000		NOT R3, R3 ; R3 = $-$ C ²	
; Pop A into R1 in next two lines		ADD R3, R3, #1	
	; line 1	ADD R1, R1, R2 ; R1 = $A^2 + B^2$	
		ADD R1, R1, R3 ; R1 = $A^2 + B^2 - C^2$	
D D: (D2: () 1:	; line 2	BRnp CHECK_DONE	
; Pop B into R2 in next two lines		ADD R0, R0, #1	
	; line 3	CHECK_DONE	11 .
	· line 4	; Push the value in R0 onto the stack in fo two lines	llowing
; Pop C into R3 in next two lines	, inic i	ewo mies	
•	. lim a F		; line 16
	; line 5		; line 17
	; line 6		,
AND R0, R0, #0 ; Set R0 = 0		; Stop the program	
; Set $R1 = R1 * R1$ in following three lines			; line 18
; Must call SQUARE subroutine			
	; line 7	; SQUARE subroutine	
		; Input : R4	
	; line 8	; Output : R5 = R4*R4	
		SQUARE	
	; line 9	ST R3, Save_R3	
; Set R2 = R2 * R2 in following three lines		AND R5, R5, #0 ADD R3, R4, #0	
; Must call SQUARE subroutine		LOOP	
,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,		ADD R5, R5, R3	
	; line 10	ADD R4, R4, #-1	
		BRp LOOP	
	; line 11	ID D2 C D2	
	; line 12	LD R3, Save_R3 RET	
	; IIIIe 12	KE I	
; Set R3 = R3 * R3 in following 3 lines ; Must call SQUARE subroutine		Save_R3 .BLKW #1	
	; line 13		
	1: 1.4		
	, ime 14		
	; line 15		

Problem 2: Debug a Program

Danny implemented a program that reverses the contents of a sequence of memory locations. That is, contents of x4000 is swapped with that of x4009, x4001 with x4008, x4002 with x4007, and so on. The Main program calls the REVERSE subroutine which in turn uses a SWAP subroutine to swap the contents of two memory locations. REVERSE is supposed to return to the Main program once the value of the starting address (in R0) becomes larger than that of the ending address (in R1).

The program compiled fine but it falls into an infinite loop. Help Danny find and fix the bug.

Line	
0	.ORIG x3000
	; Main program
	, Main program
1	LD R0, START
2	LD R1, END
3	JSR REVERSE
4	HALT
	; Reverse Subroutine
	; Input R0, R1; Output - NONE
5	REVERSE
6	ST R0, SAVER0_REVERSE
7	ST R1, SAVER1_REVERSE
8	ST R2, SAVER2_REVERSE
9	ST R3, SAVER3_REVERSE
10	RLOOP
11	JSR SWAP
12	ADD R0, R0, #1
13	ADD R1, R1, #-1
14	NOT R2, R0
15	ADD R2, R2, #1
16	ADD R3, R2, R1
17	BRp RLOOP
18	LD R0, SAVER0_REVERSE
19	LD R1, SAVER1_REVERSE
20	LD R2, SAVER2_REVERSE
21	LD R3, SAVER3_REVERSE
22	RET
	; Swap Subroutine
	; Input R0, R1; Output - NONE
23	SWAP
24	ST R2, SAVER2_SWAP
25	ST R3, SAVER3_SWAP

26	LDR R2, R0, #0
27	LDR R3, R1, #0
28	STR R2, R1, #0
29	STR R3, R0, #0
30	LD R2, SAVER2_SWAP
31	LD R3, SAVER3_SWAP
32	RET
33	START .FILL x4000
34	END .FILL x4009
35	SAVERO_REVERSE .BLKW #1
36	SAVER1_REVERSE .BLKW #1
37	SAVER2_REVERSE .BLKW #1
38	SAVER3_REVERSE .BLKW #1
39	SAVER2_SWAP.BLKW #1
40	SAVER3_SWAP.BLKW #1
41	.END

Part A (7 points). Danny suspects that the RLOOP on lines 10–17 never finishes; i.e., line 18 is never reached. Is he correct in his suspicion or not? Justify your answer (use no more than 30 words).

our Answer:		

line(s)? (use no more than 30 words)			
Your Answer:			
Part C (8 points). Please provide a solution for the problem by modifying of inserting instructions. Maximum 4 lines are allowed. Please be very specific about the lines you are modifying/inserting. e.g., between lines 30 and 31, insert ADD R5, R5, R5			
or change line 30 to ADD R5, R5, R5			
Your Answer:			

Part B (7 points). Why does the program get into an infinite loop and between which

Problem 3: Concepts

3.1 (2 points) LC-3 is a 16-bit system. Now suppose you have a 32-bit system, what is its address space (number of unique memory locations)?

Y	ดน	r A	۱n	SW	ær	•:
	v		111	.7 VI		

- 3.2 (5 points) What are some benefits of using a subroutine in LC-3? (Choose all that apply)
 - a) Hide program details from others
 - b) Separate program from underlying hardware
 - c) Make code more organized
 - d) Create libraries for others to use
 - e) Make debugging easier

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3.3 The following piece of code is given to you to handle I/O. Does it describe interrupt-driven I/O or polling I/O? Justify your choice using no more than 30 words.

	.ORIG x3000		
	КСНЕСК	LDI R0, KBSR BRzp KCHECK	
	DCHECK	LDI RO, KBDR LDI R1, DSR BRzp DCHECK STI RO, DDR	
	HALT	JII Ko, DDK	
	KBSR .FILL KBDR .FILL DSR .FILL DDR .FILL .END	xFE02 xFE04	
Cho	ose one (2 poi	nts):	
Inte	rrupt-driven I/	O Po	olling I/O
Just	ify your choic	e using no more t	than 30 words (3 points):

3.4 Consider the evaluation of postfix expression using a stack in MP2. Assume that all operands in the expression are single-digit. You need to apply the same method used in MP2 to evaluate the postfix expression given below. When an operand is encountered, it should be pushed onto the stack. When an operator is encountered, two numbers should be popped off the stack to perform the arithmetic operation, and the result should be pushed onto the stack.

Postfix expression: 487-3+/2*	
Part A: find the result of this postfix expression	
Your Answer: 487-3+/2* equals	(5 points)

Part B: Given that the stack starts at memory location x4000. Fill out the stack at the following instances. Assume that the stack is empty at the beginning, be sure to include all the values at each instance.

(1) Right before the first arithmetic operation (i.e., before its operands are popped off the stack). (5 points)

x3FFA	ji (e pemie)
x3FFB	
x3FFC	
x3FFD	
x3FFE	
x3FFF	
x4000	

(2) Right after the first arithmetic operation (i.e., after its result has been pushed onto the stack). (5 points)

x3FFA	
x3FFB	
x3FFC	
x3FFD	
x3FFE	
x3FFF	
x4000	

Postfix expression: 487-3+/2*

(3) Right before the second arithmetic operation (i.e., before its operands are popped off the stack). (5 points)

x3FFA	
x3FFB	
x3FFC	
x3FFD	
x3FFE	
x3FFF	
x4000	

(4) Right after the second arithmetic operation (i.e., after its result has been pushed onto the stack). (5 points)

x3FFA	
x3FFB	
x3FFC	
x3FFD	
x3FFE	
x3FFF	
x4000	

(5) Right after the third arithmetic operation (i.e., after its result has been pushed onto the stack). (5 points)

x3FFA	
x3FFB	
x3FFC	
x3FFD	
x3FFE	
x3FFF	
x4000	

Table E.2 The Standard ASCII Table

ASCII			AS	CII		AS	CII		ASCII			
Character	Dec	Hex	Character	Dec	Hex	Character	Dec	Hex	Character	Dec	Hex	
nul	0	00	sp	32	20	@	64	40	*	96	60	
soh	1	01	1	33	21	A	65	41	a	97	61	
stx	2	02		34	22	В	66	42	b	98	62	
etx	3	03	#	35	23	C	67	43	C	99	63	
eot	4	04	\$	36	24	D	68	44	đ	100	64	
enq	5	05	&	37	25	E	69	45	e	101	65	
ack	6	06	<u>&</u>	38	26	F	70	46	f	102	66	
bel	7	07		39	27	G	71	47	g	103	67	
bs	8	80	(40	28	H	72	48	h	104	68	
ht	9	09)	41	29	I	73	49	i	105	69	
1f	10	0A	*	42	2A	J	74	4A	j	106	6A	
vt	11	0B	+	43	2B	K	75	4B	k	107	6B	
ff	12	OC.	,	44	2C	L	76	4C	1	108	6C	
cr	13	0D	-	45	2D	M	77	4D	m	109	6D	
so	14	0E		46	2E	N	78	4E	n	110	6E	
si	15	0F	/	47	2F	0	79	4F	О	111	6F	
dle	16	10	0	48	30	P	80	50	р	112	70	
dc1	17	11	1	49	31	Q	81	51	g	113	71	
dc2	18	12	2	50	32	R	82	52	r	114	72	
dc3	19	13	3	51	33	S	83	53	s	115	73	
dc4	20	14	4	52	34	T	84	54	t	116	74	
nak	21	15	5	53	35	υ	85	55	u	117	75	
syn	22	16	6	54	36	v	86	56	v	118	76	
etb	23	17	7	55	37	W	87	57	w	119	77	
can	24	18	8	56	38	X	88	58	x	120	78	
em	25	19	9	57	39	Y	89	59	у	121	79	
sub	26	1A	:	58	3A	Z	90	5A	z	122	7A	
esc	27	1B	;	59	3B	[91	5B	{	123	7B	
fs	28	1C	<	60	3C	\	92	5C	lì	124	7C	
gs	29	1D	=	61	3D	l ì	93	5D	j	125	7D	
rs	30	1E	>	62	3E	*	94	5E	_	126	7E	
us	31	1F	?	63	3F	_	95	5F	del	127	7F	

NOTES: RTL corresponds to execution (after fetch!); JSRR not shown

$R7 \leftarrow PC, PC \leftarrow M[ZEXT(trapvect8)]$	TRAP 1111 0000 trapvect8 TRAP trapvect8	R7 ← PC, PC ← PC + SEXT(PCoffset11)	JSR 0100 1 PCoffset11 JSR PCoffset11	PC ← BaseR	JMP 1100 000 BaseR 000000 JMP BaseR	((n AND N) OR (z AND Z) OR (p AND P)): $PC \leftarrow PC + SEXT(PCoffset9)$	BR 0000 n z p PCoffset9 BR{nzp} PCoffset9	DR ← SR1 AND SEXT(imm5), Setcc	AND 0101 DR SR1 1 imm5 AND DR, SR1, imm5	DR ← SR1 AND SR2, Setcc	AND 0101 DR SR1 0 00 SR2 AND DR, SR1, SR2	DR ← SR1 + SEXT(imm5), Setcc	ADD 0001 DR SR1 1 imm5 ADD DR, SR1, imm5	DR ← SR1 + SR2, Setcc	ADD 0001 DR SR1 0 00 SR2 ADD DR, SR1, SR2
M[BaseR]+ SEXT(offset6)] ← SR	STR 0111 SR BaseR offset6 STR SR, BaseR, offset6	M[M[PC + SEXT(PCoffset9)]] ← SR	STI 1011 SR PCoffset9 STI SR, PCoffset9	M[PC + SEXT(PCoffset9)] ← SR	ST 0011 SR PCoffset9 ST SR, PCoffset9	DR ← NOT SR, Setcc	NOT 1001 DR SR 111111 NOT DR, SR	DR ← PC + SEXT(PCoffset9), Setcc	LEA 1110 DR PCoffset9 LEA DR, PCoffset9	DR ← M[BaseR + SEXT(offset6)], Setcc	LDR 0110 DR BaseR offset6 LDR DR, BaseR, offset6	DR ← M[M[PC + SEXT(PCoffset9)]], Setcc	LDI 1010 DR PCoffset9 LDI DR, PCoffset9	DR ← M[PC + SEXT(PCoffset9)], Setcc	LD 0010 DR PCoffset9 LD DR, PCoffset9