ECE 391 Exam 2, Fall 2014

Monday, November 10, 2014

Name:	
NetID:	
Discussion Section:	
10 a.m. (Eric Badger)	
11 a.m. (Matt Tischer)	
1 p.m. (Matt Tischer)	
2 p.m. (Yixiao Nie)	
○ 3 p.m. (Ying Chen)	

- Be sure that your exam booklet has 10 pages.
- Write your net ID at the top of each page.
- This is a closed book exam.
- $\bullet\,$ You are allowed two $8\,{}^1\!/_2\times11"$ sheet of notes.
- Absolutely no interaction between students is allowed.
- Show all of your work.
- Don't panic, and good luck!

Page	Points	Score
3	9	
4	7	
5	10	
6	4	
7	5	
8	8	
9	5	
10	14	
Total:	62	

here, where he requires you students to displace all his fear. If you finish his quest and finish it well, you will be rewarded with points and all will be swell. So grab your quill, pencil, or pen, and finish every task, requested by Ben.

absolute block numbers (4kB per block) N+D0 2 3 1 N-1 + 0D-1index nodes (inodes) data blocks boot # dir. entries 4Bblock length in B 4B# inodes (N) 4B0th data block # 4Bdata blocks (D) 4B 1st data block # 52B reserved etc. file name 32B file type 4Binode# 4B 64B dir. entries 24B reserved

Notes:

1KB = 1024 B

You may leave expressions in unsimplified form.

Assume only regular files and directories (e.g. no devices).

(a) (2 points) What is the maximum number of files that this file system can supp	ort?
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(b) (2 points) What is the maximum size of any single file in this file system?

(c) (5 points) Assume that your file system contains the maximum number of files and that each of those files is 4B. What fraction of the total filesystem size is used for actual data?

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(d) (2 points) Why do the reserved bytes exist in the boot block as well as in each of the directory entries?

(e) (5 points) Describe the process of reading a file from this file system step-by-step

(f) (5 points (bonus)) Implementing write on this filesystem would be inefficient. Explain why and suggest an additional file system data structure that would help.

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Question 2: Paging Tradeoffs	14	· po	oin	ts
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(a) (10 points) Ben Bitdiddle wants to explore the tradeoffs between different (hypothetical) paging schemes. Help him fill in the tables below. One row is filled in for you as an example. Assume that the index bits are divided equally between all levels and that the size of one entry in a paging structure is always 4 bytes.In the last row, you may assume that s is a power of 2 and that n and s are chosen such that the number of index bits (per level) is an integer.

# Levels	Size of Pages	# Page Offset Bits	# Index Bits (/level)
2	4kB	12	10
20	4kB		
1	1GB		
2	4B		
n	S		

# Levels	Size of Pages	Size of One Entry	# Entries (/level)	Size of Paging
				Data (/level)
2	4kB	4B	1024	4kB
20	4kB	4B		
1	1GB	4B		
2	4B	4B		
n	S	4B		

(b)	(2 points)	Why would v	we choose the	e paging s	scheme	with one	level of 1	GB 1	pages o	ver the	scheme	with 20
	levels of 4	kB pages?										

(c) (2 points) Why would we choose the paging scheme with two levels of 4B pages over the scheme with one level of 1GB pages?

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Virtual address	Data
0xc80000	matrix[0][0]
0xc80004	matrix[0][1]
0xc80008	matrix[0][2]
0xc81000	matrix[1][0]
0xc81004	matrix[1][1]

Table 1: matrix array storage in memory

Consider the following code:

```
int x, y;
result = 0;
for (x = 0; x < 1024; x++) {
  for (y = 0; y < 4; y++) {
    result += matrix[0][x]*matrix[y][x];
  }
}</pre>
```

(a) (5 points) Show the state of the TLB after each iteration of the inner loop, i.e., after the result += ... instruction has been executed. For each TLB translation entry, write down the virtual address of the corresponding page. Assume that the TLB starts out empty, and that the variables x, y, and result are stored in registers and thus do not require memory accesses; likewise, instruction fetches do not use the TLB. The table captures the first six iterations, with the first column already filled in for you.

round 1	round 2	round 3	round 4	round 5	round 6
(x=0, y=0)	(x=0, y=1)	(x=0, y=2)	(x=0, y=3)	(x=1, y=0)	(x=1, y=1)
0xc80000					
empty					
empty					
1 2					
	(x=0, y=0) 0xc80000 empty	(x=0, y=0) (x=0, y=1) 0xc80000 empty	(x=0, y=0) (x=0, y=1) (x=0, y=2) 0xc80000 empty	(x=0, y=0) (x=0, y=1) (x=0, y=2) (x=0, y=3) 0xc80000 empty	(x=0, y=0) (x=0, y=1) (x=0, y=2) (x=0, y=3) (x=1, y=0) 0xc80000 empty

(b) (3 points) What is the number of TLB misses that will occur during the execution of the entire program (i.e., 4×1024 rounds).?

(c) (5 points) Rewrite this code to perform the same calculation with fewer TLB misses

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(d)	(2)	points)	How many	TLB	misses	will	occur	during	the	execution	of v	your revised	code?
()	\ -	0011100)	110 111411		11110000	* * * * * * *	~~~				~	,	· ·

(e) (3 points) Your partner, Ben Bitdiddle, has configured the scheduler so that there is a context switch after every inner loop iteration. He notes, with pride, that your optimized code and the original now both incur the same number of TLB misses. Explain why.

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(a) (4 p	points) Give one example each of a IO-bound and a compute	-bound process.
1-) (4	which of the transfer of the line of the l	to and other?
ы) (4 р	points) Which of the two processes should have higher priori	ty and wny?
c) (2 r	points) How long is a typical time slice (quantum)? Circle the	e correct choice.
1	A. 10–100 ns	
	B. 10–100 ms C. 10–100 s	
d) (4 p	points) Describe one advantage and one disadvantage of mak	ing a quantum larger.