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Thesis

Title

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To someone very special...

Acknowledgments

Abstract

This thesis is about the synchronized flight of a formation of multirotors which executes a mission. The mission is characterized by a trajectory for every drone which composes the formation. The formation must be able to deal with unforeseen events, which can compromise the outcome of the mission. In order to do that, the drones must exchange information between each others, therefore a support network is needed.

The purpose of this thesis is to present the algorithm for synchronized formation flight called “Consensus algorithm” and implement it in a simulated environment and in a real system with heterogeneous drones. We want to verify that the theoretical results can be applied in a real distributed system, with not ideal network performances.

In the first part, we explain the algorithm used during the experimental work, while in the following chapter the software and hardware used are presented. In the last chapters we show in detail the structure of the software and the experiment conducted. In particular, many simulations will be presented to confirm the quality of the “Consensus algorithm”. Finally, a comparison between the simulated results and the ones obtained in the real environment is proposed.

Sommario

La tesi tratta di volo sincronizzato di una formazione di multirotori che eseguono una missione, caratterizzata da una traiettoria per ogni drone che compone la formazione. La formazione deve essere in grado di reagire a eventi inaspettati, che possono compromettere l'esito della missione stessa. Per raggiungere tale scopo, i droni devono scambiare informazioni, attraverso una rete di supporto.

L'obbiettivo della tesi è di presentare l'algoritmo utilizzato per il volo sincronizzato, chiamato "Algoritmo di Consenso" e implementarlo sia in un ambiente simulato, che in un sistema reale, composto da macchine eterogenee. Si vuole verificare che i risultati teorici possano essere applicati in un sistema reale distribuito, dotato di una rete di comunicazione dalle prestazioni non ideali.

Nella prima parte della tesi, verrà esposto l'algoritmo usato durante il lavoro sperimentale, mentre nel capitolo seguente saranno presentati l'hardware e il software. Negli ultimi capitoli saranno evidenziati nel dettaglio la struttura del software e gli esperimenti condotti. In particolare, saranno mostrate alcune simulationi, in modo da confermare la bontà dell'"Algoritmo di Consenso". Infine, nell'ultimo capitolo, sarà proposta una comparazione tra i risultati delle simulazioni e quelli ottenuti mediante l'implementazione in un sistema reale.

Contents

Acknowledgments	5
Abstract	7
Sommario	9
List of figures	16
Introduction	17
1 Consensus theory	23
1.1 Parametrized trajectory	25
1.2 Virtual time	28
1.3 Consensus law	30
1.4 Network topology	33

1.5 Convergence property	34
2 System architecture	37
2.1 Hardware	37
2.1.1 Pixfalcon	38
2.1.2 Intel Edison	40
2.1.3 RaspberryPi Zero	40
2.1.4 Motive Optitrack	41
2.2 Software	43
2.2.1 Ground station	44
2.2.2 Raspberry Pi Zero	45
2.2.3 Intel Edison	45
2.2.4 Both companions	45
2.2.5 Pixfalcon	46
2.2.6 Simulation	46
2.2.7 Additional software	47
3 Consensus node	49
3.1 Start and stop services	51

3.2 Consensus variable callback	53
3.3 Local position callback	54
4 Simulation results	61
5 Experimental results	63
Conclusions	65

List of Figures

2.1	Pixfalcon board	39
2.2	Edison board	40
2.3	RaspberryPi Zero board	41
2.4	Motive Optitrack screenshot	42
2.5	Optitrack marker	43
2.6	Ant drone	43
2.7	Hexa drone	44
3.1	Input and output of the node	51
3.2	Custom service structure	56
3.3	Custom message structure	57
3.4	Consensus variable class	58
3.5	Class used to manage the position and velocity of the drones	59

3.6 Class used to manage a generic trajectory	60
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Introduction

These years are characterized by the development of the robotics. The growing applications of mobile robots and drones in many fields has brought important increasing in the number of robotic vehicles. Many commercial solutions are available, but, in many cases, they offer a product which needs an expert pilot. For this reason and due to the high prices, the diffusion of mobile robots is limited. In the latest years, these problems were mitigated and we now find lower prices and easier-to-use products. This will bring a growing non professional diffusion of robotic vehicles which, hopefully, will populate our houses and help us in our life. The most important application of UAVs are the ones related to exploration and data collection. Indeed, drones are fundamental in inspection of unknown areas, such as forests or unreachable terrains. Also the monitoring of industrial artifacts constitutes a valuable application of UVAs. For instance, solar panels, wind turbines, high-voltage cables or tubes can be supervised using aerial vehicles. The multimedia production is the field in which the drones are most present. Indeed, most of the machines are equipped with high-definition cameras in order to provide professional videos.

The power consumptions of the machines are significantly reduced in the last

years and we are now able to build smaller and lighter vehicles. In particular, in the aerial field, there are commercial products which weight about 200g with a flight time of 20 minutes. These improvements allow us to develop more advanced features, such as trajectory planning, obstacles avoidance, formation flight. The development of autonomous UAVs is fundamental in application such as humanitarian response around the world. When a natural calamity happens, the UAVs can play a crucial role, delivering essential goods or finding missing people. Moreover, the operations can be done without risking the safety or the life of the rescuers, because they can act remotely or plan an autonomous mission.

The trajectory planning is essential for the development of autonomous vehicles, because the generation of a feasible trajectory for a mission is a requirement, otherwise no mission can be done. The trajectory generation is an hard problem, because of its computational complexity. Indeed, many suboptimal algorithm have been developed in order to reduce the complexity. There are different classes of planning algorithm and they can be summarized as follows:

- Artificial Potential Field: these algorithms assign a value of the potential for every point in the map. The goal has the lowest (highest) value and the obstacles have high (low) values. Then, the robot tries to descend (climb) the potential, in order to reach the goal.
- Sampling Based Planning: random samples are used to find a path from the starting point to the goal. Advanced random samples techniques can be used in order to reduce the complexity of the algorithm.
- Grid Based Planning: these kind of algorithm overlay a grid on the map, so every configuration corresponds with a grid pixel. The robot can move from

one grid pixel to any adjacent grid pixels as long as that grid pixel does not contain an obstacle.

- Reward-Based Planning: the robot can apply different actions in every state of the world. The outcome of an action can be not deterministic and after every action, the robot gains a reward. The objective of the robot is to maximize the sum of the rewards.

Another important aspect is the obstacle avoidance. The trajectory usually is planned offline, when the mission is not started yet. So, the trajectory does not take into account the fact that an unplanned obstacle might interfere with the mission. In this case another algorithm is implemented in order to run online. This kind of algorithms are called obstacle avoidance algorithms and they replan locally the trajectory, in order to get rid of unforeseen events. These algorithms need the drone to be equipped with proximity sensors, which provide information about the surrounding environment. Different methodologies have been developed for obstacles avoidance. The main ones are summarized below without details.

- Artificial Potential Field: as before, the idea of this approach is that obstacles exert a virtual repulsive force, given by the potential field, to push away the robot from them and the goal position generates a virtual attractive force to guide the robot to it.
- Virtual Force Field method (VFF): is the combination of the Artificial Potential Field with the concept of Probabilistic Occupancy Grid maps.
- Fuzzy Controller for Obstacle Avoidance: as the name says, a fuzzy controller is used to derive the variables for the vehicle's orientation and accel-

eration control, depending on the current perception of the robot's sensors

- Vector Field Histogram method (VFH): improved version of the VFF method.
- VFH+, VFH* method: improved versions of the VFH method, but more computational costly. The most recent one is the VFH method, which uses an A* search.
- Traversability Field Histogram (TFH): The local path planner bases on the VFH concept, but is extended by the information provided within the Traversability Map.
- Dynamic Window Approach: the algorithm takes into account the dynamic and kinematic constraints of the robot and it is similar to the VFH+ algorithm. The method is called dynamic window approach and considers only admissible velocities which can be reached within the next time interval and which allow the robot to stop safely.
- Curvature Velocity Space method: This method chooses a point in the linear-angular velocity space which satisfies some constraints and maximizes an objective function. This objective function tries to move the robot close to the commanded direction at the highest feasible speed, while travelling the trajectory with the largest clearance from obstacles.
- Beam Curvature method: improvement for the curvature velocity space method.
- Nearness diagram Navigation: it is similar to the VFH method, but uses a polar histogram to derive actions to be taken for the robot.

When multiple machines have to fly in a formation, a synchronization mechanism is needed. Indeed, if one of the drones has to deviate from the planned trajectory (for instance, because of an unplanned obstacle), the formation must be preserved. To reach the synchronization, the vehicles must communicate with each others and exchange information. Notions from graph theory are needed to analyze the behaviour of the system and demonstrate the convergence properties.

This work is focused on the presentation of a distributed synchronization algorithm, which we call “Consensus Algorithm”. Every drone must exchange its mission progression with its neighbors and must adjust its mission progression based on the information received from the other vehicles. The algorithm is a distributed algorithm and it is executed on every vehicle.

In the first chapter, we will show the theoretical aspects of the “Consensus Algorithm” and then, in the second chapter, we will provide a general overview of the hardware and software used to deploy the algorithm. In the third chapter we will show the software implementation of the main part of the system, providing classes and snippets of code. At the end, we will explain in the fourth and fifth chapters the simulated and experimental results of our work. We will apply the algorithm to a real system and we will present our results and its performances.

Chapter 1

Consensus theory

The literature about the “Consensus” has grown significantly in the latest years because of the increasing presence of autonomous vehicles. This, in accordance with the new improvements in robotics, has brought to a growing interest in consensus between multiple agents which have to accomplish a mission in cooperative or adversarial scenarios.

The “Consensus” theory has its roots in Graph Theory and Automatics and it can be used to coordinate a mission in order to achieve the synchronization between the vehicles even during unforeseen events which force one or more components of the mission to change the planned trajectory or task. In this case, the other components identify this variation and they act to preserve the synchronization.

The “Consensus” algorithm is a distributed algorithm which can sometimes be simulated in a centralized fashion because of the reduced computational power

of the machines involved in the mission, which are equipped with low power hardware to account for the crucial issue of power consumption. In this scenario, a centralized server simulates the algorithm and communicates with all the machines.

The most common “Consensus” application is a spatial and timing consensus: in this implementation, each vehicle of the formation has to travel along a specified trajectory and the completion of the mission occurs when all the agents reach the final positions of their spatial paths. The algorithm has to guarantee that the difference between the ending time at which all the vehicles finish their tasks is minimized and asymptotically goes to zero, when the execution time goes to infinity and no other unforeseen event happens. We also consider formations of Unmanned Aerial Vehicles, but the key concepts can be freely applied to other categories of robots when they are able to follow a trajectory. This study is focused on this kind of application, and further simulation results and experimental achievements are presented in chapters 4 and 5.

In the next sections we refer to the main work of Venanzio Cichella [5] in the field of UAV consensus, in order to provide an homogeneous state of the art about the “Consensus”. The paper considers all the details needed to build a “Consensus” system. The main components of this kind of system are listed and explained below:

- Parametrized trajectory
- Virtual time
- Consensus law

- Network topology

1.1 Parametrized trajectory

The trajectory is a spatial path with an associated timing law and it is used to identify the position of the center of mass of our agent at a given time. First of all, the trajectory used during our experiments are polynomials curves. In particular, we use Bézier curves, because of the existence of many computationally efficient algorithms designed for this kind curves, such as algorithms to efficiently compute the minimum distance between two Bezier curves and the existence of a closed-form solution for the arc lengths of the paths. We can start with the following definition of a generic trajectory $\mathbf{p}_{d,i}(t_d)$ for N vehicles:

$$\mathbf{p}_{d,i} : [0, t_{d,i}^f] \rightarrow \mathbb{R}^4, \quad i = 1, 2, \dots, N \quad (1.1)$$

where $t_d \in [0, T_d]$, with $T_d := \max\{t_{d,1}^f, \dots, t_{d,N}^f\}$, is the time variable of the trajectory, while $t_{d,i}^f \in \mathbb{R}^+$ are the individual final mission times of the vehicles obtained during the planning phase. Usually all these final times are equal and therefore $t_{d,1}^f = \dots = t_{d,N}^f = T_d$, but we introduced the notation for the sake of completeness. Obviously, the trajectories need to be collision free and must comply with spatial and temporal constraints due to the dimension of the vehicle and its maximum velocities and accelerations. The range of the function is \mathbb{R}^4 because, in the case of UAVs, we want to control x, y, z, yaw . The trajectory can

account also for *roll* and *pitch*, and therefore, the function might take images in \mathbb{R}^m , where m is the number of dimensions considered. In the following sections, we will refer only to the position and *yaw*, but a more general theory can also be developed.

We now parametrize the trajectory using a dimensionless variable $\zeta_i \in [0, 1]$, related to the time t_d . In this way we can specify a function $\theta(\cdot)$, which represents the timing law associated with the spatial path $\mathbf{p}_{d,i}(\zeta_i)$. We can specify this timing law using the dynamic relation of the form:

$$\theta(t_d) = \frac{d\zeta_i}{dt_d} \quad (1.2)$$

where $\theta(t_d)$ is a smooth and positive (the parameter increases when the time increases) function.

As it allows a one-to-one correspondence between the time variable t_d and the parameter ζ_i , an analytical expression for the function $\zeta_i(t_d)$ is desirable. Using the timing law defined in (1.2), the map $\zeta_i(t_d)$ is given by the integral:

$$\zeta(t_d) = \int_0^{t_d} \theta_i(\tau) d\tau \quad (1.3)$$

Usually, all these functions are defined as polynomials in order to make quicker

and easier the evaluation process, since multiplication and addition are the basic operations in a digital processor. However, it is not mandatory to use them and a generic shape for the functions can be designed.

For the spatial path we use quintic Bezier curves for the x , y and z coordinates, while we use a third degree polynomial for the *yaw* of the drone since we need less elaborated shapes. Since we are using drones with a decoupled control of x , y , z and *yaw*, the polynomials are completely independent.

The timing law associated to these curves is a Bezier third order polynomial. So, in our case, we have five polynomials which have to be evaluated in real time in order to provide the set points. The rate at which we evaluate the trajectory is 5Hz .

The trajectory which describes the mission of a single drone is shown in the following system.

$$\begin{cases} x_i(\zeta_i) = \sum_{j=0}^5 \bar{x}_{i,j} b_j^5(\zeta_i) \\ y_i(\zeta_i) = \sum_{j=0}^5 \bar{y}_{i,j} b_j^5(\zeta_i) \\ z_i(\zeta_i) = \sum_{j=0}^5 \bar{z}_{i,j} b_j^5(\zeta_i) \\ yaw_i(\zeta_i) = \sum_{j=0}^3 y\bar{w}_{i,j} b_j^3(\zeta_i) \\ \zeta_i(t_d) = \sum_{j=0}^3 \bar{\zeta}_{i,j} b_j^3(t_d) \end{cases}$$

$\bar{x}_{i,j}$, $\bar{y}_{i,j}$, $\bar{z}_{i,j}$, $y\bar{w}_{i,j}$ and $\bar{\zeta} \in \mathbb{R}$ are the desired control points of the spatial

path and b_n^m are the (up to degree m) Bernstein basis polynomials. We do not enter in details about the Bezier curves and Bernstein basis polynomials, because it is not the objective of this work.

We have defined all the elements of a trajectory and we do not go into detail about the trajectory generation phase. Further information about boundary conditions and flyable trajectory, which satisfy the dynamic constraints of the vehicles, can be found in [5] and is extensively analyzed in [4], [24], [26].

1.2 Virtual time

Given N collision free trajectories, we want each vehicle to follow a *virtual target*, moving along the path computed offline by the trajectory generation algorithm. The objective can be achieved introducing a *virtual time*, γ_i , which is used to evaluate the trajectory and can be adjusted online to reach the synchronization even when external disturbances occur. Thus, the position of the i^{th} virtual target is denoted by $\mathbf{p}_{d,i}(\gamma_i(t))$ and the i^{th} vehicle tries to follow it, by reducing to zero a suitably defined error vector using control inputs.

Considering the trajectory $p_{d,i}(t_d)$ produced by the trajectory generation algorithm, we consider the virtual time γ_i as a function of time t , which relates the actual time t to mission planning time t_d .

$$\gamma_i : \mathbb{R}^+ \rightarrow [0, T_d], \quad \text{for all } i = 1, 2, \dots, N \quad (1.4)$$

We can now define the virtual target's position, velocity and acceleration, which have to be followed by the i^{th} vehicle at time t

$$\begin{aligned}\mathbf{p}_{c,i}(t) &= \mathbf{p}_{d,i}(\gamma_i(t)) \\ \mathbf{v}_{c,i}(t) &= \dot{\mathbf{p}}_{d,i}(\gamma_i(t), \dot{\gamma}_i(t)) \\ \mathbf{a}_{c,i}(t) &= \ddot{\mathbf{p}}_{d,i}(\gamma_i(t), \dot{\gamma}_i(t), \ddot{\gamma}_i(t))\end{aligned}\tag{1.5}$$

With the above formulation, if $\dot{\gamma} = 1$, then the speed profile of the virtual target is equal to the desired speed profile computed at trajectory generation level. Indeed, if $\dot{\gamma} = 1$, for all $t \in [0, T_d]$, with $\gamma_i(0) = 0$, it implies that $\gamma_i(t) = t_d$ for all t and thus:

$$\mathbf{p}_{c,i}(t) = \mathbf{p}_{d,i}(\gamma_i(t)) = \mathbf{p}_{d,i}(t) = \mathbf{p}_{d,i}(t_d)$$

In this particular case, the desired and commanded trajectories coincide in every time instant and also the velocity profiles coincide with the ones chosen at the trajectory generation time. If instead $\dot{\gamma}_i > 1$, it implies a faster execution of the mission; on the other hand, $\dot{\gamma}_i < 1$ implies a slower one.

We can now normalize γ_i in order to have a range which is $[0, 1]$. We simply need to divide all by T_d . In this way, we could use Bezier curves in order to

represent the spatial path. This kind of curves offer interesting properties for computing minimum distances between two of them and allow the computation of smooth trajectories. In any case, it is not mandatory to use them and a general function could be used instead.

The second order derivative of γ_i , $\ddot{\gamma}_i$, is a free parameter used to achieve the consensus. In the next section we will introduce the control law which commands its evolution during time and we will explain how it is possible to implement a distributed algorithm.

1.3 Consensus law

Now, we formally state the path following problem. We define as $\mathbf{p}_i(t) \in \mathbb{R}^4$ the position of the center of mass of the i^{th} agent and its yaw and since $\mathbf{p}_{c,i}(t)$ describes the commanded pose to be followed by the agent at time t , the errors are defined as:

$$\begin{aligned}\mathbf{e}_{p,i}(t) &= \mathbf{p}_{c,i}(t) - \mathbf{p}_i(t) \in \mathbb{R}^4 \\ \mathbf{e}_{v,i}(t) &= \mathbf{v}_{c,i}(t) - \dot{\mathbf{p}}_i(t) \in \mathbb{R}^4\end{aligned}\tag{1.6}$$

Then, the objective reduces to that of regulating the error defined in (1.6) to a neighbourhood of zero. This task is solved with an autopilot capable of following the set points computed from the desired trajectory at specified instances of time.

The virtual time is the parameter used to reach consensus between multiple vehicles. In fact, since the the trajectories are parametrized by γ_i , the agents are synchronized at time t when:

$$\gamma_i(t) - \gamma_j(t) = 0 \quad \text{for all } i, j \in 1, \dots, N, \quad i \neq j \quad (1.7)$$

We can also control the rate of progression of the mission using a parameter $\dot{\gamma}_d \in \mathbb{R}$, which represents the velocity of the virtual time with respect to the real time. All the agents share this variable and they proceed at the same rate of progression if:

$$\dot{\gamma}_i(t) - \dot{\gamma}_d(t) = 0 \quad \text{for all } i \in 1, \dots, N \quad (1.8)$$

Adjusting $\dot{\gamma}_d$, we can decide the speed of the mission: for instance, if we set $\dot{\gamma}_d = 1$ and (1.7) and (1.8) are satisfied for all the vehicles, then the mission is executed at the speed originally planned in the trajectory generation phase. If instead we use $\dot{\gamma}_d > 1$ or $\dot{\gamma}_d < 1$ we carry out the mission faster or slower. This term can be changed in real time in order to avoid moving objects or unplanned obstacles, which make it necessary to change one of the path of the agents. For the purpose of “Consensus”, the parameter is only a reference command, rather than a control input.

Now, we introduce the coordination control law which regulates the evolution of $\ddot{\gamma}_i(t)$ during the time and determines $\gamma_i(t)$:

$$\begin{aligned}\ddot{\gamma}_i(t) &= \ddot{\gamma}_d(t) - b(\dot{\gamma}_i(t) - \dot{\gamma}_d(t)) - a \sum_{j \in \aleph_i} (\gamma_i(t) - \gamma_j(t)) - \bar{\alpha}_i(\mathbf{e}_{p,i}(t)) \\ \dot{\gamma}_i(0) &= \dot{\gamma}_d(0) = 1 \\ \gamma_i(0) &= \gamma_d(0) = 0\end{aligned}\tag{1.9}$$

where a and b are positive coordination control gains, while $\bar{\alpha}_i(\mathbf{e}_{p,i}(t))$ is defined as:

$$\bar{\alpha}_i(\mathbf{e}_{p,i}(t)) = \frac{\mathbf{v}_{c,i}^T(t) \mathbf{e}_{p,i}(t)}{||\mathbf{v}_{c,i}(t)|| + \epsilon}\tag{1.10}$$

with ϵ being a positive design parameter, $\mathbf{e}_{p,i}$ the position error vector defined in (1.6) and \aleph_i the set of the neighbors which can communicate with the i^{th} vehicle (we will see details later). In the equation (1.9) we have four terms. The feedforward term, $\ddot{\gamma}_d$, allows the virtual target to follow the acceleration profile of γ_d . The second term, $-b(\dot{\gamma}_i(t) - \dot{\gamma}_d(t))$, reduces the error between the speed profile imposed by $\dot{\gamma}_d(t)$ and $\dot{\gamma}_i(t)$, which corresponds to the control objective given in (1.8). In particular, if $\dot{\gamma}_d(t)$ is one, then the virtual target converges to the desired speed profile chosen in the trajectory generation phase. The third term, $-a \sum_{j \in \aleph_i} (\gamma_i(t) - \gamma_j(t))$, ensures that all the vehicles are coordinated with

their neighbors as specified in (1.7). Finally, the fourth term, $-\bar{\alpha}_i(\mathbf{e}_{p,i}(t))$, is a correction term used to take into account for the path following errors of the agent. Indeed, if the vehicle is behind its target, the term is not zero and the target slows down in order to wait for the real vehicle.

With this control law, we want our vehicles to be synchronized and to proceed at a desired rate of progression, in order to accomplish the mission even when some unforeseen disturbances occur during the execution phase.

1.4 Network topology

To achieve time-coordination objective, agents must exchange information over a supporting communication network. To analyze the information flow, we need to consider some tools from algebraic graph theory, whose key concepts can be found in [2].

We assume that a vehicle i exchanges information with only a subset of all vehicles, denoted as $\aleph_i(t)$. We assume that arcs of the network are bidirectional and that there are no network delays. The information exchanged is composed by the virtual time of the agents, $\gamma_i(t)$.

The topology of the graph, $\Gamma(t)$, that represents the communication network, must comply with the following constraint in order to guarantee the convergence

of the “Consensus” algorithm:

$$\frac{1}{NT} \int_t^{t+T} QL(\tau)Q^T d\tau \geq \mu I_{N-1}, \quad \text{for all } t \geq 0 \quad (1.11)$$

where $L(t) \in \mathbb{R}^{N \times N}$ is the Laplacian of the graph $\Gamma(t)$ and $Q \in \mathbb{R}^{(N-1) \times N}$ is a matrix such that $Q1_N = 0$ and $QQ^T = I_{N-1}$, with 1_N being a vector in \mathbb{R}^N whose components are all 1s. In (1.11), the parameters $T > 0$ and $\mu \in (0, 1]$ represent a measure of the level of connectivity of the communication graph. This condition requires the graph $\Gamma(t)$ to be connected only in an integral sense, not pointwise in time. Therefore, even if the graph were disconnected during the mission at some interval of time, the convergence of the ”Consensus” algorithm would still be possible. With this condition, we can capture also packets dropouts, loss of communication and switching topologies, which can all occur during the mission, but these events do not necessary break the convergence property.

1.5 Convergence property

The control law given by (1.9) guarantees that the error of the “Consensus” algorithm converges to zero exponentially. It can be shown that the maximum convergence rate is given by the sum of the convergence rate of the path following

error and the term

$$\frac{a}{b} \frac{N\mu}{T(1 + (a/b)NT)^2} \quad (1.12)$$

which depends on the control gains a and b , the number of vehicles N and the quality of service of the communication network, characterized by the parameters T and μ . If we fix the gains and the number of the vehicles, the convergence rate depends only on the amount of information which the agents exchange each other over time.

Chapter 2

System architecture

In this chapter, we will describe the hardware and software architecture of the system. Most of the software is completely decoupled from the hardware part and can be run on different machines, allowing portability and reusability. On the contrary, the code is developed only for the specific hardware; mostly the code which is more related to the specific functionalities of the hardware itself. We will have an overview of the hardware used and then we will show a general software architecture and its deployment on the machines involved.

2.1 Hardware

The hardware used to run the system is heterogeneous and we will show in detail the machines involved in the project.

A desktop pc, whose characteristics are listed below, will be our ground station:

- Processor: Intel Core(TM) i5-6500 CPU @ 3.20GHz
- Memory: 16 GB
- Network: Intel Gigabit CT Network Adapter
- Storage: 230 GB

On this machine, we will run a virtual machine which has the following specifications:

- Processor: 1 Core
- Memory: 4 GB
- Storage 25 GB
- Network: Virtual adapter

The flying vehicles involved in the mission are of two kinds, but both adopt the same general configuration, even with different hardware. Indeed, they are equipped with a flight control unit connected to a companion microcomputer by the serial port. The microcomputer communicates with the ground station through the Wi-Fi connection.

2.1.1 Pixfalcon

The flight control unit adopted is the Pixfalcon, (figure 2.1) which belongs to the family of the Pixhawk [18].

Its specifications are the following:



Figure 2.1: Pixfalcon board

- Main System-on-Chip: STM32F427
 - CPU: 180 MHz ARM Cortex M4 with single-precision FPU
 - RAM: 256 KB SRAM (L1)
 - Failsafe System-on-Chip: STM32F100
 - CPU: 24 MHz ARM Cortex M3
 - RAM: 8 KB SRAM
 - Wifi: ESP8266 external
 - GPS: U-Blox 7/8 (Hobbyking) / U-Blox 6 (3D Robotics)
 - Connectivity:
 - 1x I2C
 - 1x CAN (2x optional)
 - 1x ADC
 - 4x UART (2x with flow control)
 - 1x Console
 - 8x PWM with manual override

- 6x PWM / GPIO / PWM input
- S.BUS / PPM / Spektrum input
- S.BUS output

2.1.2 Intel Edison



Figure 2.2: Edison board

The companion computers are of two types. The first kind is the Intel Edison [14] (figure 2.2) which is a general purpose computer with the following specifications:

- Atom 2-Core (Silvermont) x86 @ 500 MHz
- Memory: LPDDR3 1 GB
- Storage: 4 GB EMMC

2.1.3 RaspberryPi Zero

The second kind of companion is the RaspberryPi Zero [10] (figure 2.3), with the following specifications.

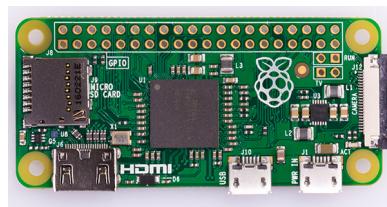


Figure 2.3: RaspberryPi Zero board

- Processor:
 - Broadcom BCM2835
 - contains an ARM1176JZFS (ARM11 using an ARMv6-architecture core)
- Memory: 512MB LPDDR2 SDRAM
- USB On-The-Go port
- Mini HDMI
- 40pin GPIO header
- CSI camera connector

2.1.4 Motive Optitrack

The motion capture system is Motive Optitrack [20], we use eight Optitrack Prime 13 cameras [21], arranged on a square as show in the figure 2.4. The cameras define a volume, which is a cube of 5 meters per side, and the drones can fly without obstacles inside it.

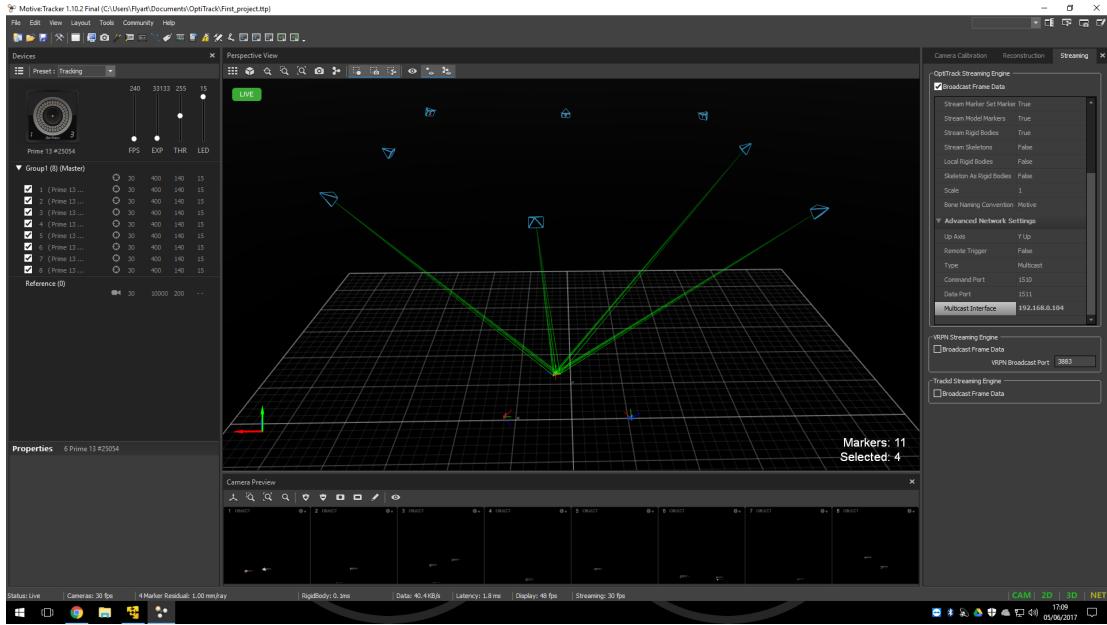


Figure 2.4: Motive Optitrack screenshot

The cameras are connected to a Netgear Prosafe 28PT GE POE [22] switch through Gigabit Ethernet cables.

In order to be seen by the Optitrack system, every drone must be equipped with markers (figure 2.5) with different configurations, which differentiate the drones from each other.

The two models of drones which we will use are presented in the figures 2.6 and 2.7. The first one is the Ant model, a tiny 200g drone, which is equipped with the Raspberry Pi Zero and the Pixfalcon. The second one is the Hexa model. It is provided with the Intel Edison board and the Pixfalcon board. It is larger and its diameter is about 40cm long. As we can see from the figures, the Ant model has four propellers, while the Hexa has six.

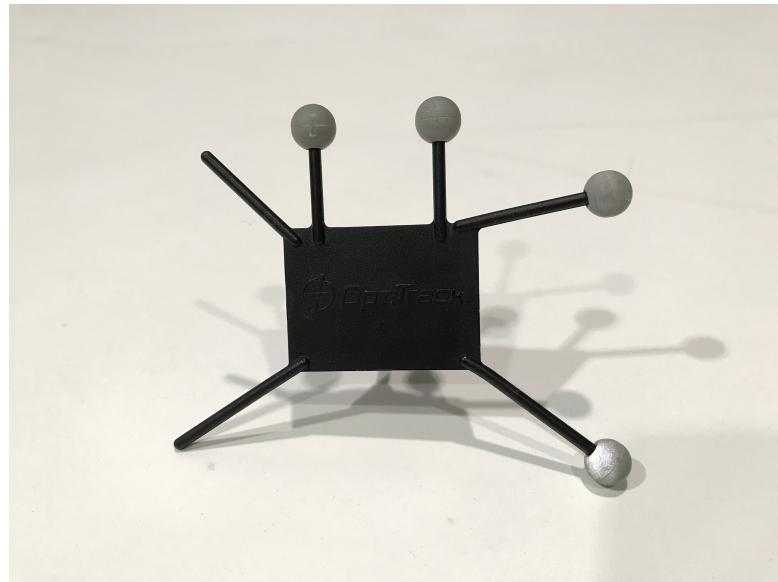


Figure 2.5: Optitrack marker



Figure 2.6: Ant drone

2.2 Software

We now list the software adopted to execute the algorithm in a real environment and in the simulated one.

The principal software used to manage the distributed architecture is ROS Kinetic Kame [7]. ROS is a robotic middleware with a structure which is mainly



Figure 2.7: Hexa drone

publisher-subscriber and it can manage more machines in a distributed environment. The central part of the ROS architecture is a node called ROS core, which manages the topics of the system and the subscriptions. The ROS core offers also other functionalities, such as the Parameter Server or the possibility to advertise services. The Parameter Server is a central infrastructure, which is responsible for storing configuration parameters loaded by the nodes of the system. These parameters can be retrieved by other nodes and used if necessary. Instead, a ROS service is a sort of remote function call. One node can advertise the service, which can be called by any other node. The call is synchronous, so the caller is blocked until the callee has executed its callback function. The ROS architecture is based on queues, threads and callback functions, but most of the provided tools hide part of the implementations of the distributed environment.

2.2.1 Ground station

The ground station runs Windows 10 Pro [19] and the software used to virtualize a Desktop machine is VMware [25]. On the virtual machine is installed Ubuntu

16.04 LTS [3] in order to run software needed and available only for Unix systems.

On Windows operative system we launch the Motive Optitrack software [20], which allows to calibrate and control the cameras for position tracking. It then provides the streaming of the positions of the markers identified by the cameras and sends it to the Ubuntu operative system using a multicast IP address. Here, the information is converted by a ROS node and sent through the ROS topics, which are read by the drones. In this way, each drone knows exactly its position. This conversion node is an open source node called Mocap which can be found on GitHub [13]. On Ubuntu side, we launch the ROS core, which manages all the ROS nodes and topics.

2.2.2 Raspberry Pi Zero

The Raspberry Pi Zero executes a dedicated version of Debian operative system, which is Raspbian. The version used is Raspbian Jessie 4.4 [9].

2.2.3 Intel Edison

The Intel Edison runs a version of Debian called Jubilinux, at version 0.1.1 [6].

2.2.4 Both companions

Both companions, the Raspberry and the Edison, are provided with ROS Kinetic and both have to execute some ROS nodes in order to communicate with the

other drones.

Both of them run Mavros nodes [16], which can be downloaded from GitHub and manage the conversion of the information taken from the ROS topics to the serial port and vice versa. Indeed, the ROS messages are converted into Mavlink messages and sent through the serial to the Pixfalcon autopilot. The same is done for the Mavlink messages from the autopilot, which are published on ROS topics.

The second kind of ROS node run by the companions is a custom consensus node, which loads the desired trajectory and sends the next set point to the Mavros node. This node will be analyzed in detail in the chapter 3.

2.2.5 Pixfalcon

The Pixfalcon autopilot is flashed with PX4 Pro Autopilot [17], an open source firmware downloadable from GitHub. The release used is the v1.5.5.

2.2.6 Simulation

The simulation part is developed with the utilities provided by the PX4 firmware. In particular, the physical engine used is Gazebo [8], which is an Open Source project.

2.2.7 Additional software

We use Matlab R2016_B [15] to process the data, to plot the graphs and to validate some theoretical results.

This document is written in L^AT_EX [23] and the code IDE used is Atom [1], while the versioning control platform used are GitHub [11] and GitLab [12].

Chapter 3

Consensus node

In this chapter, we will examine the structure of the consensus node: we will see its main components, then, we will show some snippets of code in order to make clearer which parts are involved.

As shown in the general architecture, the consensus node is developed as a ROS node which subscribes and publishes messages to different topics. Moreover, the node offers some ROS services used to start and stop the trajectory following algorithm or the consensus one.

The structure of the node takes into account the main architectural patterns used in the software development field and it was designed to allow the maximum degree of usability and customization, even though one of the most important metric taken into account is the efficiency of the code, because it has to be executed on machines with limited amount of resources.

First of all, the node functionalities are enclosed into a C++ class which

initializes all the ROS elements and prepares the node to receive the start and stop commands. The initialization is done by the class constructor when the object is created. In order to apply the consensus dynamic equation, we need the current position of the UAV. The Px4 board already publishes the estimated local position on a topic; so, we need to subscribe to that topic to retrieve the messages with the needed information. Second, we want to publish the consensus variable of the drone in the topic used by all the other UAVs, because, having obtained the others' consensus variables from the same topic, we are able to compute the proportional consensus error. Third, we need the next set point and the next desired velocity profile, because we want to compute the position error and weight it for the target velocity. At the end, we compute the acceleration of the consensus parameter using the consensus equation and we publish the next set point. We will see the details through the code. All these elements can be summarized and shown in figure 3.1.

The subscription of a ROS topic works with a callback function, which accepts as parameter the pointer to the new message. Since in our case we have multiple subscriptions and we must advertise the start and stop services, we need to implement a multithreading architecture which takes care of the concurrent accesses to the state of our object. The number of thread used is three and these are their functions:

- Start and stop services
- Consensus variable callback
- Local position callback

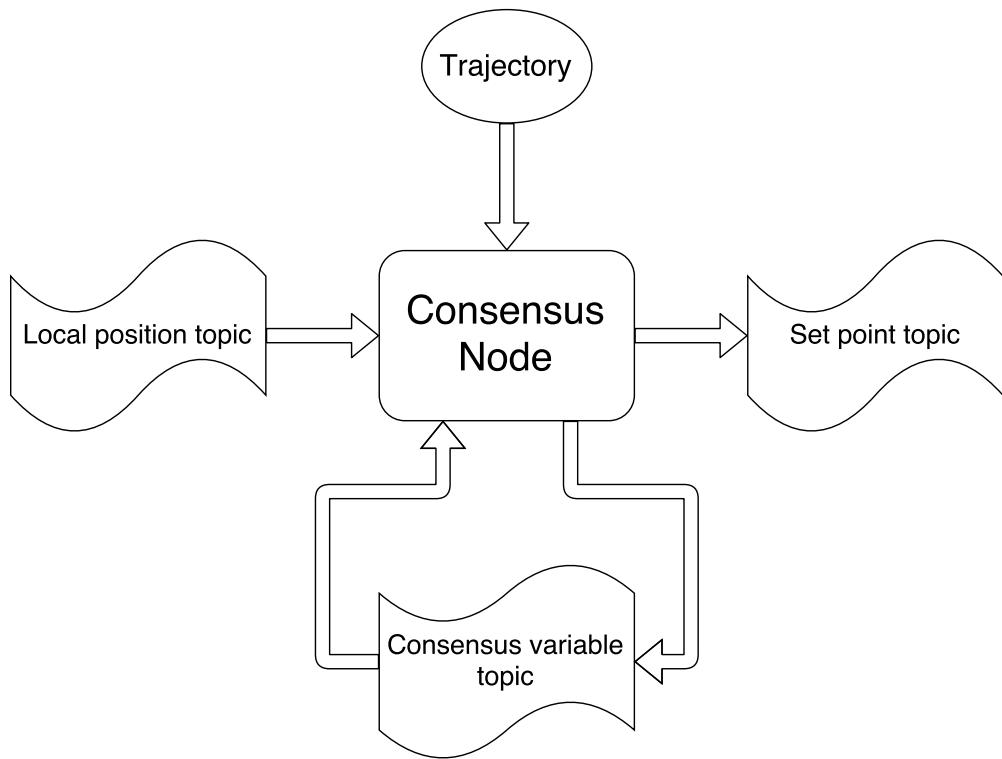


Figure 3.1: Input and output of the node

3.1 Start and stop services

First of all, the node can work in two different modes:

- trajectory-only
- consensus

In the trajectory-only mode, the node computes the next set point and sends it to the UAV, without taking care if there are others UAVs in the mission. Its only objective is to follow the trajectory and reach its final position trying to respect the time constraints imposed by the trajectory. Instead, in the consensus mode, the node does the same computation as before, but it also publishes its consensus

variable and reads all the other ones. It then considers this information and adjusts its variable. The consensus mode includes the trajectory-only, and can be started even if the trajectory-only is already started, while the opposite is not true. When you stop the consensus mode, also the trajectory-only is stopped. When the mission is accomplished, the current active mode is stopped automatically, so that you can freely restart one of the two modes without the need of stopping the previous one.

The services are implemented using the ROS Service class which manages the whole infrastructure needed for calling the service. The call to the service is synchronous and the caller is blocked until the service function is terminated. In this case, we offer two services: one for starting and stopping the trajectory-only mode and the other for the consensus mode. It is possible to customize the service call in order to pass different number and types of arguments to the service function and the response can also be defined. For the two services we have defined the same parameters that are shown in the figure 3.2.

The message is composed by two parts: the request and the response. In the request we need a boolean field in order to know if we want to start or stop the algorithm. The response consists in a boolean variable which represents the success of the operation and an exit code which identifies the eventual problems occurred. The constants for the exit codes are directly specified in the definition of the service.

The trajectory-only service starts or stops the thread which, taken a local position, computes the next set point; while the consensus one starts or stops the same thread as before and the one which retrieves the consensus variables from

the other quadrotors.

3.2 Consensus variable callback

The thread responsible for collecting the consensus variables of all the other UAVs, is managed by ROS and it executes a callback function when a new message is published on a specific topic. This topic is used by all the drones to publish their consensus variable, γ_i , and it accepts a custom message which contains only a string with the name of the owner of the variable and the value itself. The message has also a header, which contains general information such as timestamp or message id. The structure of the message can be seen in the figure 3.3.

The callback function receives the information from the topic and updates a local view of the variables of the neighbors. This information has a timeout validity, because we do not want to consider too old values. Indeed, if we consider too old values and a problem in the network causes a loss of packets, our drone might think which the other drones have a significantly different values of the consensus variables and so, wait for them. On the contrary, it is better to discard the values and remove the neighbors after a timeout.

In order to store the information, we use a support class, thread safe, which provides a procedure to check if the variable is expired or not. The signatures of the methods of the class are presented in the figure 3.4. We use a container to store the values of the neighbors and we always check if the value is expired or not before using it.

These variables are used in the consensus law as γ_j of the neighbors and are used to compute the proportional error. We can see that the the expiration interval of the values can model the fact that the network topology can change. Indeed, if a link between two drones vanished because, for instance, they are too far from each other, after a timeout (which it is equal to the expiration time), the neighbor will be removed from the container and the drone will not take into account the old neighbor. Even a failure of a drone is ignored using the timeout: if a machine has a critical problem and does not send its consensus variable, the other drones remove it from their neighbors and, therefore, they can continue their mission without problems.

3.3 Local position callback

In the position callback function we apply the consensus law. First of all, we store the actual position of the drone in an object of a custom class, which signature is presented in figure 3.5. In this class, we also include, besides x, y and z, the yaw of our vehicle. Indeed, all the operations defined over the class consider also the orientation.

Now, we compute the synchronization term which corresponds to the sum of the difference between the current γ_i and all the γ_j of the neighbors. In order to do this, we iterate over a container, which stores the values, and we incrementally form the synchronization term.

The next step is to form the $\bar{\alpha}_i$ term. First of all, we need the next set point, which must be obtained evaluating the trajectory with the actual value of γ_i .

Our trajectory is represented by a class, and its signature is shown in the figure 3.6. Now, we simply compute the *position_error* as $set_point - position$. We also need the desired velocity, which can be obtained using the suitable function of the trajectory class. At this point, we have all the terms needed to compute $\bar{\alpha}_i$ as defined in -TODO-.

Since we have all the elements, it is time to apply the consensus law and find $\ddot{\gamma}_i$. We simply need to have the coefficients a and b , as shown in -TODO-, and the references $\ddot{\gamma}_d$ and $\dot{\gamma}_d$.

One of the last step needed is the update of $\dot{\gamma}_i$ using $\ddot{\gamma}_i$ and γ_i using $\dot{\gamma}_i$. We compute the interval of time, dt , between the last update and the current update and we do the math as:

```
dgamma += ddgamma * dt ;  
gamma += dgamma * dt ;
```

At the end, we need to publish the value of γ_i to the right topic only if we are operating in consensus mode, otherwise we ignore it. An operation which is always needed is the publication of the set point message to the autopilot of the UAV, in order to allow it to follow the trajectory and reach its final destination.

```
# Command for enabling or disabling
# trajectory and consensus tasks

#####
## Request fields
bool cmd
_____
#####
### exit code constants
# everything ok
uint8 SUCCESS = 0
# trajectory following already
# active and consensus activated
uint8 ESCALATION = 1
# nothing to do
uint8 ALREADY_DONE = 2
# other errors
uint8 FAILURE = 3

#####
# Response fields
bool success
uint8 exit_code
```

Figure 3.2: Custom service structure

```
Header header
string owner
float64 gamma
```

Figure 3.3: Custom message structure

```

class GammaParameter {
    private:
        std::string owner;
        double gamma;
        double acquisition_time; //In seconds
        std::mutex mtx;
        static double expiration_interval; //In seconds

    public:
        GammaParameter();
        GammaParameter(const GammaParameter& gp);
        GammaParameter(std::string owner,
                      double gamma,
                      double acquisition_time);

        ~GammaParameter();

        static void setExpirationInterval(double expiration_interval);
        static double getExpirationInterval();
        GammaParameter& setOwner(std::string owner);
        GammaParameter& setGamma(double gamma, double acquisition_time);
        GammaParameter& setData(std::string owner,
                               double gamma,
                               double acquisition_time);

        std::string getOwner();
        double getGamma();
        bool isExpired();
        GammaParameter& getData(std::string *owner_ptr,
                               double * gamma_ptr,
                               bool * exp_ptr);
        GammaParameter& operator= (const GammaParameter &gp);
};


```

Figure 3.4: Consensus variable class

```

class DronePose {

private:
    double x, y, z, yaw;

public:
    DronePose();
    DronePose(double x, double y, double z, double yaw);
    DronePose(const DronePose &dp);
    ~DronePose();

    double getX() const;
    double getY() const;
    double getZ() const;
    double getYaw() const;
    void setX(double x);
    void setY(double y);
    void setZ(double z);
    void setYaw(double yaw);

    double module() const;
    DronePose& operator= (const DronePose &dp); //Assignment
    DronePose operator+ (const DronePose &dp) const; //Sum
    DronePose& operator+= (const DronePose &dp);
    DronePose operator- (const DronePose &dp) const; //Difference
    DronePose& operator-= (const DronePose &dp);
    DronePose operator- () const; //Unary minus
    double operator* (const DronePose &dp) const; //Scalar product
    DronePose operator* (double scal) const; //Product with constant
    DronePose& operator*=(double scal);
    DronePose operator/ (double scal) const; //Ratio with constant
    DronePose& operator/=(double scal);
};


```

Figure 3.5: Class used to manage the position and velocity of the drones

```
class Trajectory {
private:
    std::vector<TrajectorySegment *> segments;
    std::string drone_id;

    double filterAndConvertTime(double t) const;
    TrajectorySegment* getRightSegment (double t) const;

public:
    double min_time, max_time;

    bool loadXML( std::string dorne_ns , char * document );
    void cleanAll();

    DronePose evaluateNED(double t) const;
    DronePose evaluateENU(double t) const;
    DronePose operator[] (double t) const; //Default return ENU

    DronePose evaluateVelNED(double t) const;
    DronePose evaluateVelENU(double t) const;
};
```

Figure 3.6: Class used to manage a generic trajectory

Chapter 4

Simulation results

Chapter 5

Experimental results

Conclusions

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