

1.

Proof: Suppose EQ_{CFG} is undecidable. Let TM H be its decider:

$H(\langle A, B \rangle) = \{ \text{accept, if } A \text{ accepts } L(A) = L(B). \text{ Decline, if } A \text{ does not accept } L(A) = L(B) \}$

Make a new TM D:

D = on input $\langle N \rangle$, for TM N

1. Run H on $\langle N, \langle N \rangle \rangle$.

2. If H accepts, reject. If H rejects, accept.

$D(\langle D \rangle) = \{ \text{accept, if } D \text{ does not accept } \langle D \rangle. \text{ reject, if } D \text{ accepts } \langle D \rangle \}$

$D \text{ accepts } \langle D \rangle, \text{ so long as } D \text{ does not accept } \langle D \rangle. \rightarrow \text{TM } D \text{ cannot exist} \rightarrow \text{TM } H \text{ cannot exist} \rightarrow EQ_{CFG} \text{ is undecidable}$

2. Proof: Suppose $HALT - \varepsilon_{TM} = \{ \langle M \rangle : M \text{ is a TM and halts on blank input (i.e. } \varepsilon) \}$ is undecidable.

Let TM H be its decider:

$H(\langle M \rangle) = \{ \text{accept, if } M \text{ accepts. Decline, if } M \text{ does not halt on blank input} \}$

Make a new TM D:

D = on input $\langle N \rangle$, for TM N

1. Run H on $\langle N, \langle N \rangle \rangle$.

2. If H accepts, reject. If H rejects, accept.

$D(\langle D \rangle) = \{ \text{accept, if } D \text{ does not accept } \langle D \rangle. \text{ reject, if } D \text{ accepts } \langle D \rangle \}$

$D \text{ accepts } \langle D \rangle, \text{ so long as } D \text{ does not accept } \langle D \rangle. \rightarrow \text{TM } D \text{ cannot exist} \rightarrow \text{TM } H \text{ cannot exist} \rightarrow HALT - \varepsilon_{TM} \text{ is undecidable}$

3. Proof: Suppose $EVEN_{TM} = \{ \langle M \rangle : M \text{ is a TM and } L(M) \text{ contains every string of even length} \}$ is undecidable. Let TM H be its decider:

$H(\langle M \rangle) = \{ \text{accept, if } M \text{ accepts. Decline, if } L(M) \text{ contains every string of even length} \}$

Make a new TM D:

D = on input $\langle N \rangle$, for TM N

1. Run H on $\langle N, \langle N \rangle \rangle$.

2. If H accepts, reject. If H rejects, accept.

$D(\langle D \rangle) = \{ \text{accept, if } D \text{ does not accept } \langle D \rangle. \text{ reject, if } D \text{ accepts } \langle D \rangle \}$

$D \text{ accepts } \langle D \rangle, \text{ so long as } D \text{ does not accept } \langle D \rangle. \rightarrow \text{TM } D \text{ cannot exist} \rightarrow \text{TM } H \text{ cannot exist} \rightarrow EVEN_{TM} \text{ is undecidable}$