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Lab 4: Sort Analysis

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This program implements five sorting methods, using four variations of quicksort and natural merge sort. The versions of quicksort pick different locations to create pivot points and may use insertion sort to recombine them. The different methods of picking mid points causes the algorithm to vary depending on the length and ordering of the data. By running different length tests with varied ordering, we are able to analyze the sorting efficiency based on the number of comparisons the algorithm makes as well as exchanges of elements during its runtime. This data shows us that as the number of elements increases, so does the number of comparisons and exchanges but at different rates depending on the sort and order. For example, ascending and descending orders make quicksort based approaches very inefficient and quicksort tends to have extremely high numbers of comparisons for ordered data which can be costly just to determine that the items are already in order.

A good example is quicksort using first pivot running on ascending and descending data. When ascending the sort required only 49 exchanges for 50 numbers but the descending order required 674. This large step up in exchanges illustrates the inefficiencies some sorts can experience just by ordering. Choosing the median pivot was better than using the first number as a pivot. And the difference in threshold for insertion sort, 100 vs 50, made little difference for large files. As the sizes of files grew, the effect that ordering had on quicksort grew as well. Overall the most important factors as the size grew were the pivot choice for quick sort, with median being preferred over insert, over first pivot, and the order of the data. Comparatively, natural merge sort was more resistant to ordering changes and file sizes. Natural merge sort operated more efficiently compared to a straight merge sort which recursively divides data and uses extra space to hold the arrays. By choosing an in place sort that uses a stack, a quick structure that lends itself well to popping on and off ordered sequences of data, the natural merge

sort ran fairly consistently and efficiently but crucially it used far less space than a straight merge sort would have. Natural merge sort is also faster for data that has partial order, compared to straight merge which performs similarly regardless of order. Natural merge sort is more consistent overall compared to quicksort, and would be faster than straight merge sort regardless of order or size.

Justification for Design Decisions

Initially I tried implementing the sorts using recursion but ran into trouble running them with large amounts of data. I was getting recursion errors, that max depth had been reached, leading to most of the sorts failing to complete. I went back and reimplemented the sorts using iterative approaches by using the MyStack class made in previous labs to act as a stack structure. Choosing a stack was a natural fit given our experience in previous labs, as well as for the task, the stack was a good choice for time and space efficiency of operations.

What I Learned/What I would Do Differently

Implementing 5 recursive sorts resulted in a lot of recursion failures, so I would first attempt iterative designs over recursion until I am sure the method would benefit from recursion. It has been much harder to debug recursive problems than iterative ones.

Sort Name	Order	Size	Comparisons	Exchanges
Quicksort first pivot	random	50	277	119
Quicksort first pivot	random	1000	10569	5653
Quicksort first pivot	random	2000	23399	12554
Quicksort first pivot	random	5000	68270	35059
Quicksort first pivot	random	10000	157753	80463
Quicksort first pivot	ascending	50	1225	49
Quicksort first pivot	ascending	1000	499500	999
Quicksort first pivot	ascending	2000	1999000	1999
Quicksort first pivot	ascending	5000	12497500	4999
Quicksort first pivot	ascending	10000	49995000	9999
Quicksort first pivot	descending	50	1225	674
Quicksort first pivot	descending	1000	499500	250999
Quicksort first pivot	descending	2000	1999000	1001999
Quicksort first pivot	descending	5000	12497500	6254999

Sort Name	Order	Size	Comparisons	Exchanges
Quicksort first pivot	descending	10000	49995000	25009999
Quicksort insert 100	random	50	725	727
Quicksort insert 100	random	1000	23416	20537
Quicksort insert 100	random	2000	48713	42335
Quicksort insert 100	random	5000	127638	105268
Quicksort insert 100	random	10000	274484	218722
Quicksort insert 100	ascending	50	49	49
Quicksort insert 100	ascending	1000	494550	999
Quicksort insert 100	ascending	2000	1994050	1999
Quicksort insert 100	ascending	5000	12492550	4999
Quicksort insert 100	ascending	10000	49990050	9999
Quicksort insert 100	descending	50	1225	1274
Quicksort insert 100	descending	1000	499500	253449
Quicksort insert 100	descending	2000	1999000	1004449
Quicksort insert 100	descending	5000	12497500	6257449
Quicksort insert 100	descending	10000	49995000	25012449
Quicksort insert 50	random	50	725	727
Quicksort insert 50	random	1000	16172	12778
Quicksort insert 50	random	2000	31958	24297
Quicksort insert 50	random	5000	92209	66858
Quicksort insert 50	random	10000	207732	146252
Quicksort insert 50	ascending	50	49	49
Quicksort insert 50	ascending	1000	498275	999
Quicksort insert 50	ascending	2000	1997775	1999
Quicksort insert 50	ascending	5000	12496275	4999
Quicksort insert 50	ascending	10000	49993775	9999
Quicksort insert 50	descending	50	1225	1274
Quicksort insert 50	descending	1000	499500	251599
Quicksort insert 50	descending	2000	1999000	1002599
Quicksort insert 50	descending	5000	12497500	6255599
Quicksort insert 50	descending	10000	49995000	25010599
Quicksort median pivot	random	50	218	173
Quicksort median pivot	random	1000	9107	5570
Quicksort median pivot	random	2000	20861	12116
Quicksort median pivot	random	5000	60807	37644
Quicksort median pivot	random	10000	134190	81497
Quicksort median pivot	ascending	50	193	148
Quicksort median pivot	ascending	1000	7987	4960
Quicksort median pivot	ascending	2000	17964	10916

Sort Name	Order	Size	Comparisons	Exchanges
Quicksort median pivot	ascending	5000	51822	30713
Quicksort median pivot	ascending	10000	113631	66421
Quicksort median pivot	descending	50	245	187
Quicksort median pivot	descending	1000	14205	9001
Quicksort median pivot	descending	2000	32880	20499
Quicksort median pivot	descending	5000	98004	60501
Quicksort median pivot	descending	10000	218585	133464
Natural merge sort	random	50	217	269
Natural merge sort	random	1000	8737	9738
Natural merge sort	random	2000	19395	21396
Natural merge sort	random	5000	56820	61825
Natural merge sort	random	10000	123668	133673
Natural merge sort	ascending	50	161	213
Natural merge sort	ascending	1000	5052	6053
Natural merge sort	ascending	2000	11104	13105
Natural merge sort	ascending	5000	33356	38361
Natural merge sort	ascending	10000	71712	81717
Natural merge sort	descending	50	133	185
Natural merge sort	descending	1000	4932	5933
Natural merge sort	descending	2000	10864	12865
Natural merge sort	descending	5000	29804	34809
Natural merge sort	descending	10000	64608	74613

Figure 1. Comparisons and Exchanges for 5 different sorts

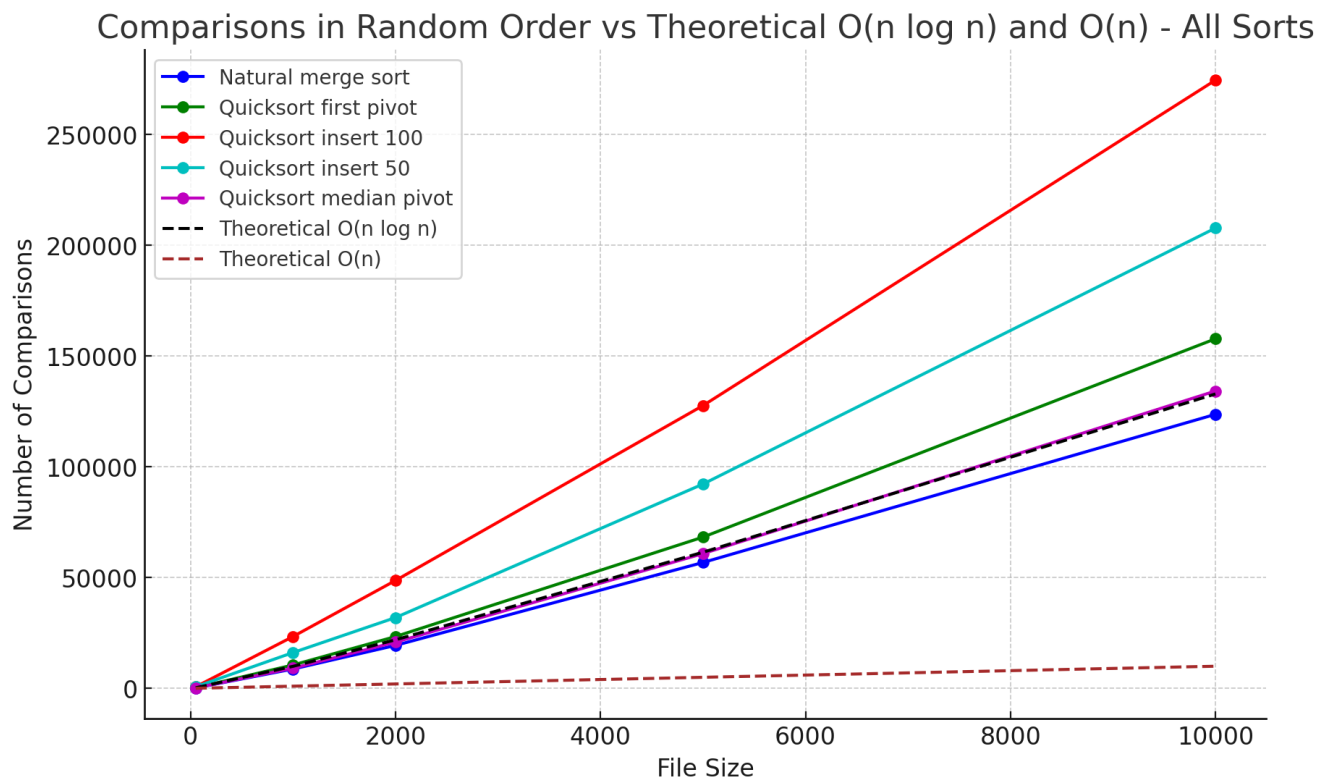


Figure 2. Chart comparing random ordered performance

We can see from the chart above that natural merge sort performed the best, followed closely by quicksort using the median pivot. Randomly ordered data shows an average case for quicksort, making it a better visual comparison but this doesn't reflect the performance in all cases. Looking at $n \log n$ and n lines, we can see how natural merge sort falls under $n \log n$ for random data, making it a good sort.

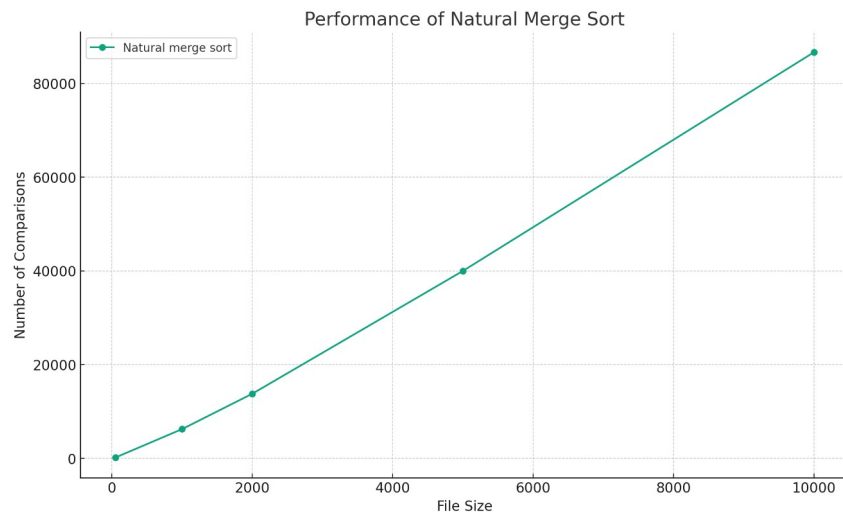


Figure 3. Natural merge sort average comparisons

Natural merge sort shows a nearly linear correlation to file size.

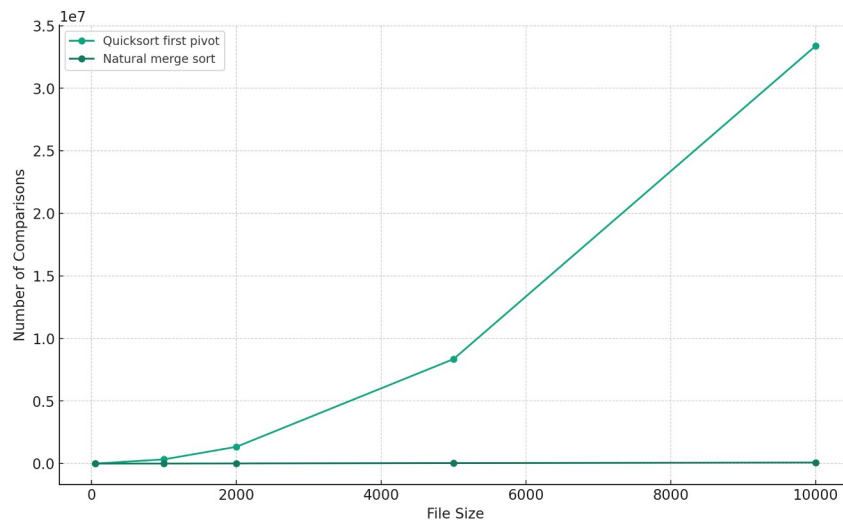


Figure 4. Natural merge and quicksort first pivot comparison

Comparatively, the number of comparisons becomes much greater at the same file size for quicksort, with noticeable growth as the file size grows.