

# Category Theory 2022-23

## Lecture 6

11<sup>th</sup> November 2022

## Reformulation of limit

The category of D-cones in  $\mathcal{C}$  has

- Objects : D-cones
- Morphisms from  $(Z', (Z' \xrightarrow{r'_u} Du)_{u \in I})$   
to  $(Z, (Z \xrightarrow{r_u} Du)_{u \in I})$

are maps  $Z' \xrightarrow{h} Z$  in  $\mathcal{C}$  s.t.  $r_u \circ h = r'_u \quad \forall u \in I$

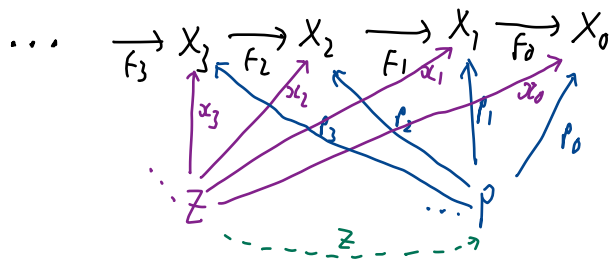
A limit D-cone is then just a terminal object in the category of D-cones.

Limit D-cones are determined up to isomorphism, because terminal objects are determined up to isomorphism!

# Projective limits

The projective limit of an infinite sequence of maps  $(X_{n+1} \xrightarrow{f_n} X_n)_{n \geq 0}$  in  $\mathcal{C}$  is given by  $P \in |\mathcal{C}|$  and  $(p_n \xrightarrow{f_n} X_n)_{n \geq 0}$  such that:

- for all  $n$ ,  $p_n = f_n \circ p_{n+1}$ ; and
- for any family  $(z \xrightarrow{x_n} X_n)_{n \geq 0}$  such that  $x_n = f_n \circ x_{n+1} \forall n$ , there exists a unique  $z \xrightarrow{z} P$  such that  $p_n \circ z = x_n \forall n$ .



Set has projective limits

$$P := \{ (x_n)_{n \geq 0} \mid \forall n \ f_n(x_{n+1}) = x_n \}$$

$$p_n := (x_n)_{n \geq 0} \mapsto x_n : P \rightarrow X_n$$

A category  $\mathcal{C}$  is said to be (small-) complete if every small diagram in  $\mathcal{C}$  has a limit cone.

Theorem The following are equivalent.

(1)  $\mathcal{C}$  is complete

(2)  $\mathcal{C}$  has products and equalisers

Proof idea

(1)  $\Rightarrow$  (2) is trivial.

For (2)  $\Rightarrow$  (1) let  $D: G \rightarrow \mathcal{C}$  be a diagram

The limit cone is the components of the equaliser of

$$\prod_{u \in |G|} D_u \begin{array}{c} \xrightarrow{\lambda_1} \\ \xrightarrow{\lambda_2} \end{array} \prod_{(u,v,e) \in \{(u,v,e) \mid u,v \in |G|, e \in G(u,v)\}} D_v$$

where  $\lambda_1 = (D \circ \pi_u)_{(u,v,e)}$  and  $\lambda_2 = (\pi_v)_{(u,v,e)}$   $\square$

A category is said to be finitely complete if every finite diagram has a limit cone.

Theorem The following are equivalent

- (1)  $\mathcal{C}$  is finitely complete
- (2)  $\mathcal{C}$  has finite products and equalisers
- (3)  $\mathcal{C}$  has terminal object and pullbacks

Proof idea (1)  $\Rightarrow$  (2) and (1)  $\Rightarrow$  (3) are immediate.  
(2)  $\Rightarrow$  (1) is proved in same way as last theorem.

For (3)  $\Rightarrow$  (2), first construct the binary product  $X \times Y$  as the pullback of  $X \xrightarrow{!} 1 \xleftarrow{!} Y$ .

Then construct the equaliser of  $X \xrightleftharpoons[g]{f} Y$  as the pullback of  $X \xrightarrow{(f,g)} Y \times Y \xleftarrow{(1_Y, 1_Y)} Y$ .



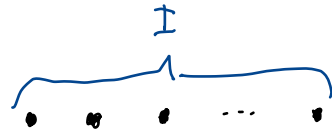
A colimit in  $\mathcal{C}$  of a diagram  $D: \mathcal{G} \rightarrow \mathcal{C}$   
 is a limit in  $\mathcal{C}^{op}$  of  $D: \mathcal{G}^{op} \rightarrow \mathcal{C}^{op}$

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Type of colimit

Diagram shape

$I$ -indexed coproducts



Initial object



Coequaliser



Pushout



Direct limit

e.g.



$\mathcal{C}$  is finitely cocomplete if every finite diagram has a colimit.

$\mathcal{C}$  is small cocomplete (or just cocomplete) if every small diagram has a colimit.

Theorem The following are equivalent

- $\mathcal{C}$  is finitely cocomplete
- $\mathcal{C}$  has finite coproducts and coequalisers
- $\mathcal{C}$  has initial object and pushouts

Theorem The following are equivalent

- $\mathcal{C}$  is cocomplete
- $\mathcal{C}$  has coproducts and coequalisers

Both theorems are immediate from the corresponding theorems about limits by duality

## Coproducts in Set

The coproduct of  $(X_i)_{i \in I}$  is  $\sum_{i \in I} X_i := \{(i, x) \mid i \in I, x \in X_i\}$

together with the cocore

$$(X_i \xrightarrow{in_i} \sum_{i \in I} X_i)_{i \in I}$$

$$in_i := x \mapsto (i, x)$$

## Universal property

For any core  $(X_i \xrightarrow{f_i} Z)_{i \in I}$

there exists a unique  $\sum_{i \in I} X_i \xrightarrow{h} Z$  such that

$$\begin{array}{ccc} \sum_{i \in I} X_i & \xrightarrow{h} & Z \\ \uparrow in_i & \nearrow f_i & \\ X_i & & \end{array} \quad \text{commutes } \forall i \in I$$

namely  $h := (i, x) \mapsto f_i(x)$ .



In Vect<sub>K</sub> finite coproducts coincide with finite products. (Vect<sub>K</sub> has biproducts.)

Given vector spaces  $V, W$ , the coproduct is

$$V \xrightarrow{v \mapsto (v, 0)} V \times W \xleftarrow{(0, w) \mapsto w} W$$

Indeed given linear  $V \xrightarrow{f} U$  and  $W \xrightarrow{g} U$

the map  $h := (v, w) \mapsto f(v) + g(w)$  is the unique

linear map such that the diagram below commutes

$$\begin{array}{ccccc}
 V & \xrightarrow{(-, 0)} & V \times W & \xleftarrow{(0, -)} & W \\
 & \searrow f & \downarrow h & \swarrow g & \\
 & & U & & 
 \end{array}$$

A coequaliser is a colimit for

$$X \begin{matrix} \xrightarrow{f} \\ \xrightarrow{g} \end{matrix} Y$$

i.e., an object  $Q$  together with  $Y \xrightarrow{q} Q$  s.t.  $q \circ f = q \circ g$

Satisfying:

for any  $Y \xrightarrow{z} Z$  s.t.  $z \circ f = z \circ g$

there exists a unique  $Q \xrightarrow{w} Z$  s.t.  $z = w \circ q$

$$\begin{array}{ccccc} X & \begin{matrix} \xrightarrow{f} \\ \xrightarrow{g} \end{matrix} & Y & \xrightarrow{q} & Q \\ & & & \searrow z & \downarrow w \\ & & & & Z \end{array}$$

Every  $q$  arising as a coequaliser is epi.

Morphisms that arise as coequalisers are called regular epis.

## Coequalisers in Set

The coequaliser in Set of  $X \begin{matrix} \xrightarrow{f} \\ \xrightarrow{g} \end{matrix} Y$

is  $Y \xrightarrow{q} Q$  defined as follows

$$Q := Y / \sim$$

where  $\sim$  is the smallest equivalence relation such that  $f(x) \sim g(x)$ , for every  $x \in X$ .

$$q := y \mapsto [y]_{\sim}$$

(every element is mapped to its equivalence class)

## Relations in a category

A representative for a relation between  $X$  and  $Y$  in  $\mathcal{C}$  is a jointly monic pair

$$R \begin{array}{c} \xrightarrow{\Gamma_1} X \\ \searrow \Gamma_2 \\ Y \end{array}$$

Recall joint monicity means, for any  $Z \xrightarrow{x} R$   
if  $\Gamma_1 \circ x = \Gamma_2 \circ x$  and  $\Gamma_2 \circ x = \Gamma_2 \circ y$  then  $x = y$ .

(If  $\mathcal{C}$  has products this is  $\equiv$  to  $R \xrightarrow{(\Gamma_1, \Gamma_2)} X \times Y$  is mono.)

Given relations  $R \begin{array}{c} \xrightarrow{\Gamma_1} X \\ \searrow \Gamma_2 \\ Y \end{array}$  and  $R' \begin{array}{c} \xrightarrow{\Gamma'_1} X \\ \searrow \Gamma'_{2'} \\ Y \end{array}$

$R \subseteq R'$  means there exists a (necessarily unique) diagonal making the two triangles commute

$$\begin{array}{ccc} R & \xrightarrow{\Gamma_1} & X \\ \downarrow \Gamma & \searrow & \uparrow \Gamma' \\ Y & \xleftarrow{\Gamma_{2'}} & R' \end{array}$$

$R \equiv R'$  means  $R \subseteq R'$  and  $R' \subseteq R$

Given  $R \begin{matrix} \xrightarrow{r_1} \\ \xrightarrow{r_2} \end{matrix} \begin{matrix} X \\ Y \end{matrix}$

We say  $z \begin{matrix} \xrightarrow{x} \\ \xrightarrow{y} \end{matrix} \begin{matrix} X \\ Y \end{matrix}$  are related by  $R$  ( $x R y$ )

If there exists a (necessarily unique) commuting diagonal

$$\begin{array}{ccc} z & \xrightarrow{x} & X \\ & \searrow & \uparrow r_1 \\ y & & R \\ & \swarrow & \downarrow r_2 \\ & & Y \end{array}$$

$R \begin{matrix} \xrightarrow{r_1} \\ \xrightarrow{r_2} \end{matrix} \begin{matrix} X \\ X \end{matrix}$  is an equivalence relation if

$$\forall z \quad \forall z \begin{matrix} \xrightarrow{x} \\ \xrightarrow{y} \end{matrix} \begin{matrix} X \\ Y \end{matrix} \quad x R y$$

$$\forall z \quad \forall x, y : z \begin{matrix} \xrightarrow{x} \\ \xrightarrow{y} \end{matrix} \begin{matrix} X \\ Y \end{matrix} \quad x R y \Rightarrow y R x$$

$$\forall z \quad \forall x, y, z : z \begin{matrix} \xrightarrow{x} \\ \xrightarrow{y} \end{matrix} \begin{matrix} X \\ Y \end{matrix} \quad x R y \wedge y R z \Rightarrow x R z$$

The kernel pair of  $X \xrightarrow{f} Y$  is the  
pullback  $R \rightrightarrows X$  of  $f$  along itself

$$\begin{array}{ccc} R & \xrightarrow{\pi_1} & X \\ \pi_2 \downarrow & \lrcorner & \downarrow f \\ X & \xrightarrow{f} & X \end{array}$$

### Exercise

- 1)  $f$  is mono  $\Leftrightarrow$  it has  $X \rightrightarrows_{\pi_x} X$  as its kernel pair.
- 2) Any kernel pair is an equivalence relation

## The symbiotic relationship between Coequalisers and kernel pairs

1) If  $Y \xrightarrow{q} Q$  is regular epi and  $q$  has a kernel pair  $R \rightrightarrows Y$  then  $q$  is the coequaliser of  $\pi_1, \pi_2$

2) If  $R \rightrightarrows Y$  is a kernel pair and it has a coequaliser  $Y \xrightarrow{q} Q$  then  $R \rightrightarrows Y$  is the kernel pair of  $q$

A diagram

$$R \rightrightarrows Y \xrightarrow{q} Q$$

in which  $\pi_1, \pi_2$  is the kernel pair of  $q$  and  $q$  is the coequaliser of  $\pi_1, \pi_2$  is called an exact fork.