Category Theory 2022-23 Lecture 11 16th December 2022

In Set objects are determined up to isomorphism by their global points $X \cong Y \iff Set(1,X) \cong Set(1,Y)$ The corresponding property does not hold in an arbitrary category with terminal object. (Exercise find counterexamples) In an arbitrary category C we need to consider generalised points of X Maps Z -> X where Z ranges over all of 10/

Generalised points form a contravariant functor in ? We assume e is locally Small We write $\underline{y} \times := c(-,x) : c^{or} \rightarrow (\underline{sct})$ for this functor, the representable functor from X (sometimes representable presheaf) Proposition 1 ((crollery of the Yoneda Lemma to Follow) Objects are determined by generalisal points; 1.e., $yX \cong yY \quad \text{in} \left[c^{op}, \underline{set} \right] \Leftrightarrow X \cong Y \text{ in } C$ [Cor, Set] is well-defined as a locally small category with

[Cor, Set] is well-defined as a locally small category with a class of objects if and only if C is small. In the case of a locally small C, the collection of objects of [cor, set] lies outside the real of Sets and classes. Nevertheless one can make some of yX = yy, which refers to the hon class between 2 objects.

The Mapping from object X to representable yX 13 itself a covariant functor passimer. assume C is small the Yoneda functor y: c > [cor, set]) One can equivalently obtain the above functor from our original hom Functor C(-,-): $C^{op} \times C \longrightarrow Set$ using the exponential property of functor categoria $C^{op} \times C \longrightarrow \underline{\text{Set}}$ c -> [eor, Set] Proposition 2 (cooling of the Yaneda lemma to Follow)

The Yoneda functor is full and faithful

Application Let G be a group. Y: G -> [Gop, set] is (Full and) Faithful. Since G has only one object, U:= X M dx: [60, set] -> set is faithful. So Uy: G - set is faithful. Since function preserve isos this gives a faithful Functor Uy: 6 - Set iso i.e., we have embedded a in a symmetric group, proving Cayley's theoren. The argument works in ignorance! One does not need to know that [600, set] is isomorphic to the category of right Gactions (cf. Week] putzle). One does not need to know that yx is the

transitive right action of G on itself.

The Yoneda lemma

For any $F: C^{op} \to Set$ and $X \in Icl$ $[C^{op}, Set](yX, F) \cong FX$ naturally in X and F.

Proof idea

The required bijections are $w \in FX \mapsto (2f)(w)_{x}$ $x : yx \Rightarrow F \mapsto \alpha_{x}(1x)$

One must then verify that these are Mutual inverses and the naturality property. This is routine.

Proof of Prop 2: y is full & faithful The Yoneda lemma gives $[C^{op}, Set](yx, yy) \cong yy(x) = C(x,y)$ Where the right-to-left bijection is $X \xrightarrow{2} X \mapsto (\xi \xrightarrow{1} X \mapsto \partial \circ t)$ which is the marphism action of y Proof of prop1: AX = AA () X = A. I is full and faithful It therefore creates isos. F: C-) D creates isos if, for any XIYE |C| and iso g: FX -> FY in D, there exists a unique fix -> y s.t ff = g; moreover, this unique f i) an iso. For a <u>snall</u> category C, the Yoneda Functor gives a full & faithful embedding of C into [Cor, set] its category of <u>presheaves</u>.

(other standard notation: Ĉ, Psh(C))

Psh(C)

preserves
existing
limits

(but not colimits)

Psh(C)

complete,
complete,
cocomplete,
cartesian closed,
a (Goothendiecu) topos

Limits and colimits in PS4(a) are computed pointwise as in set

E.g., given presheaves F and G an G.

Derine the product presheaf $F \times G$ by: $(F \times G)(X) := F \times X G \times Y$ product in let

 $Fx : f \longrightarrow Fx : f^{X}$ $Fx : f \longrightarrow Fx : f^{X}$

Psh(C) is cartesian clased. Proof outline Given presheaves F, G we need to find an exponential presheaf [F, G]. Suppose such a preshock exists, then it must enjoy the following properties $[F,G](X) \cong PSL(G)(YX, [F,G])$ (YorldG)

 \cong PSL(a) ($\forall X \times F$, G) (defining Property of [F,G])

Theorem

 $[F,G](X) := PSH(C)(YX \times F,G)$ The corresponding morphism action is determined by the naturality of the above bijection.

Accordingly, we define

One now needs to check that the preshess so defined indeed satisfies the properties required of [F,G].

This is left as an exercise (for the enthusiastic only)

Example Psh(G) (i.e. $[G^{op}, Set]$ of Week 7-puttle) $Psh(G) \cong Category of right G-actions (Week 3 puttle)$ Let A, B be presheared corresponding to right G-actions (A, a), (B, a), since X is

right G-actions $(A, *_A)$, $(B, *_B)$ since x is symmetric it we calculate [A, B] in Psh (G).

We calculate [A, B] in Psh (G).

Let $(A \times Supped)$ the order (A, B) $(A \times Y \mid X)$, (B) $(A \times Y \mid X)$, (B) (B)

 $\cong \mathcal{B}^A$ (set of all functions $A \to B$). For the last bisection, any function $f:A \to B$

determines $\hat{f} \cdot A \times G \rightarrow B$ by $\hat{f}(\alpha, g) := f(\alpha \times g') \cdot_B g$

We show that F is equivariant $\widehat{f}((a,g) \cdot h) = \widehat{f}(a \cdot h, g \cdot h)$ = f(a.h.h-1.g-1).g.h $= f(a \cdot g^{-1}) \cdot g \cdot h = \hat{f}(a,g) \cdot h$ The mapping f In f gives the required bisection from BA to G-Actir (AxG, 13), with inverse (and (a,e)) for equivariant & (e the group identity) The preshed structure on [A,B] (see A) Corresponde to the following right action on G-Atr (AxG,B) For equivariant $\phi: A \times G \longrightarrow B$ and $g \in G$ $\phi \cdot g : (a,h) \mapsto \phi(a,h,g)$ (B) Via tle bijection $\Psi: f \mapsto \widetilde{f}$, re obtain the following isomorphic action on \mathcal{B}^A $(f \cdot g)(a) = (\psi^{-1}(\psi(f) \cdot g))(a) = (\hat{f} \cdot g)(a, e)$ $\stackrel{\text{(g)}}{=} \hat{f}(a, g) = f(a \cdot g^{-1}) \cdot g.$

This explicitly defines the exponential [A,B] in GACTR.