

# Category Theory 2022-23

## Lecture 8

25<sup>th</sup> November 2022

# Monoidal closed structure

## Definition (Monoidal closure)

A monoidal category  $\mathcal{C}$  is (left) closed if, for every  $X, Y \in |\mathcal{C}|$ , there is an object  $[X, Y]$  and map

$$[X, Y] \otimes X \xrightarrow{\text{ev}_{X, Y}} Y$$

such that, for every map  $Z \otimes X \xrightarrow{f} Y$ , there exists a unique map  $Z \xrightarrow{\Lambda f} [X, Y]$  such that

$$\begin{array}{ccc} [X, Y] \otimes X & \xrightarrow{\text{ev}} & Y \\ \Lambda f \otimes 1_X \uparrow & \nearrow f & \\ Z \otimes X & & \end{array}$$

## Proposition (alternative formulation of monoidal closure)

A monoidal category is (left) closed if and only if, for every object  $X$ , there is a functor  $[X, -] : \mathcal{C} \rightarrow \mathcal{C}$  together with a natural (in  $Z$  and  $Y$ ) bijection

$$\Lambda : \mathcal{C}(Z \otimes X, Y) \xrightarrow{\cong} \mathcal{C}(Z, [X, Y])$$

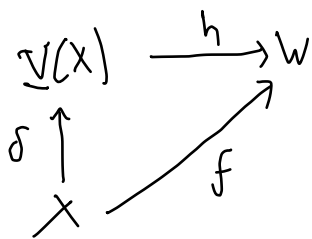
The free vector space  $\underline{V}(X)$  over a set  $X$  is defined by (Fill in the obvious definitions of addition and scalar multiplication)

$\underline{V}(X) :=$  functions  $X \rightarrow K$  with finite support

( $f: X \rightarrow K$  has finite support if  $\{x \in X \mid f(x) \neq 0\}$  is finite)

This is characterised up to linear isomorphism by the following universal property

For any set  $X$ , vector space  $W$  and function  $f: X \rightarrow W$ , there exists a unique linear  $h: \underline{V}(X) \rightarrow W$  such that



(A diagram in Set)

Where  $\delta_x(y) := \begin{cases} 1 & \text{if } y=x \\ 0 & \text{otherwise} \end{cases}$

A more pedantic formulation of the same universal property making the role of the forgetful functor  $\underline{U}: \underline{\text{Vect}} \rightarrow \underline{\text{Set}}$  explicit.

For any set  $X$ , vector space  $W$  and map  $X \xrightarrow{f} \underline{U}W$  in  $\underline{\text{Set}}$ , there exists a unique map  $\underline{V}(X) \xrightarrow{h} W$  in  $\underline{\text{Vect}}$  such that

$$\begin{array}{ccc} \underline{U}(\underline{V}(X)) & \xrightarrow{\underline{U}h} & \underline{U}W \\ \delta \uparrow & \nearrow f & \\ X & & \end{array}$$

(A)

An alternative equivalent statement

There is a functor  $\underline{V}: \underline{\text{Set}} \rightarrow \underline{\text{Vect}}$

together with natural (in  $X$  and  $W$ ) bijections

$$\psi_{X,W}: \underline{\text{Vect}}(\underline{V}X, W) \xrightarrow{\cong} \underline{\text{Set}}(X, \underline{U}W)$$

## Definition of adjunction

An adjunction  $(F, G, \psi)$  between  $\mathcal{C}$  and  $\mathcal{D}$  is given by

- Functors  $F: \mathcal{C} \rightarrow \mathcal{D}$  and  $G: \mathcal{D} \rightarrow \mathcal{C}$
- Natural (in  $X \in |\mathcal{C}|$  and  $Y \in |\mathcal{D}|$ ) bijections

$$\psi_{X,Y}: \mathcal{D}(FX, Y) \xrightarrow{\cong} \mathcal{C}(X, GY)$$

We say that  $F$  is left adjoint to  $G$

$G$  right adjoint  $F$

and we write  $F \dashv G: \mathcal{D} \rightarrow \mathcal{C}$  or  $F \dashv G$ .

## Theorem (Equivalent formulations)

Each of the 3 sets of information below is equivalent to specifying an adjunction  $(F, G, \eta)$  between  $\mathcal{C}$  and  $\mathcal{D}$ .

① A functor  $G: \mathcal{D} \rightarrow \mathcal{C}$ , function  $F: |\mathcal{C}| \rightarrow |\mathcal{D}|$  and family  $(X \xrightarrow{\eta_X} GFx)_{x \in |\mathcal{C}|}$  such that, for any  $X \in |\mathcal{C}|$ ,  $Y \in |\mathcal{D}|$  and  $X \xrightarrow{f} GY$  in  $\mathcal{C}$ , there exists a unique  $Fx \xrightarrow{g} Y$  in  $\mathcal{D}$  s.t.  $GFx \xrightarrow{Gg} GY$  commutes in  $\mathcal{C}$ .

$$\begin{array}{ccc} GFx & \xrightarrow{Gg} & GY \\ \eta_X \uparrow & \nearrow f & \\ X & & \end{array}$$

② A functor  $F: \mathcal{C} \rightarrow \mathcal{D}$ , function  $G: |\mathcal{D}| \rightarrow |\mathcal{C}|$  and family  $(FGY \xrightarrow{\epsilon_Y} Y)_{Y \in |\mathcal{D}|}$  such that, for any  $X \in |\mathcal{C}|$ ,  $Y \in |\mathcal{D}|$  and  $FX \xrightarrow{g} Y$  in  $\mathcal{D}$ , there exists a unique  $X \xrightarrow{f} GY$  in  $\mathcal{C}$  s.t.  $FX \xrightarrow{Ff} FGY$  commutes in  $\mathcal{D}$ .

$$\begin{array}{ccc} FX & \xrightarrow{Ff} & FGY \\ & \searrow g & \downarrow \epsilon_Y \\ & & Y \end{array}$$

③ Functors  $F: \mathcal{C} \rightarrow \mathcal{D}$  and  $G: \mathcal{D} \rightarrow \mathcal{C}$ , and natural transformations  $\eta: 1_{\mathcal{C}} \Rightarrow GF$  and  $\epsilon: FG \Rightarrow 1_{\mathcal{D}}$  s.t.

$$\begin{array}{ccc} FGf & \xrightarrow{\epsilon f} & f \\ F \uparrow \eta & \nearrow 1_F & \\ F & & \end{array}$$

$$\begin{array}{ccc} gFG & \xrightarrow{g\epsilon} & g \\ g \uparrow \eta & \nearrow 1_G & \\ G & & \end{array}$$

Proof outline that ① is equivalent to an adjunction

From adjunction to ①

Let  $(F, G, \psi)$  be an adjunction.

Thus  $G: \mathcal{D} \rightarrow \mathcal{C}$  and  $F: \mathcal{C} \rightarrow \mathcal{D}$  are given

Consider  $\psi_{X, FX}: \mathcal{D}(FX, FX) \xrightarrow{\cong} \mathcal{C}(X, GFX)$

and define  $X \xrightarrow{\ell_X} GFX := \psi(1_{FX})$ .

The naturality of  $\psi_{x,y}$  in  $\mathcal{Y}$  gives us for any map  $FX \xrightarrow{g} Y$  in  $\mathcal{D}$

$$\begin{array}{ccc}
 \mathcal{D}(FX, FX) & \xrightarrow{\psi_{X, FX}} & \mathcal{C}(X, GFX) \\
 \downarrow \scriptstyle i \atop \downarrow \scriptstyle g \circ i & & \downarrow \scriptstyle C(X, Gg) \\
 \mathcal{C}(FX, g) & & \\
 \downarrow & & \\
 \mathcal{D}(FX, Y) & \xrightarrow{\psi_{X, Y}} & \mathcal{C}(X, GY)
 \end{array}
 \quad
 \begin{array}{c}
 \downarrow \scriptstyle h \\
 Gg \circ h
 \end{array}$$

$$\begin{aligned}
 \text{In particular } \psi_{X, Y}(g) &= \psi_{X, Y}(\mathcal{C}(FX, g)(1_{FX})) \\
 &= \mathcal{C}(X, Gg)(\psi_{X, FX}(1_{FX})) = Gg \circ \ell_X.
 \end{aligned}$$

Since  $\psi_{X, Y}$  is a bijection, for any  $X \xrightarrow{f} GY$  in  $\mathcal{C}$ ,  $\psi^{-1}(f)$  is the unique  $g$  s.t.  $f = Gg \circ \ell_X$ .

From ① to adjunction.

Suppose we have  $G: \mathcal{D} \rightarrow \mathcal{C}$ ,  $F: \mathcal{C} \rightarrow \mathcal{D}$  and  $(x \xrightarrow{\zeta_x} GFx)_{x \in \mathcal{C}}$  as in ①.

The functorial action of  $F$  is

$$\begin{array}{ccc}
 X & & \text{the unique } FX \xrightarrow{g} FX' \\
 f \downarrow & \mapsto & \text{such that } GFx \xrightarrow{Gg} GFx' \\
 X' & & \begin{array}{ccc} \zeta_x \uparrow & & \uparrow \zeta_{x'} \\ X & \xrightarrow{f} & X' \end{array}
 \end{array}$$

The bijection  $\psi_{x,y}: \mathcal{D}(FX, y) \xrightarrow{\cong} \mathcal{C}(x, Gy)$  is defined by

$$\psi(Fx \xrightarrow{g} y) := x \xrightarrow{\zeta_x} GFx \xrightarrow{Gg} Gy$$

The diagrammatic property in ① says that this is indeed a bijection.

One then verifies that the functorial action of  $f$  preserves identities and composition and the naturality of  $\psi_{x,y}$ . □



## Examples of adjunctions

The free-vector-space functor  $\underline{V}: \underline{\text{Set}} \rightarrow \underline{\text{Vect}}$  is left adjoint to the forgetful  $\underline{U}: \underline{\text{Vect}} \rightarrow \underline{\text{Set}}$ .

(Property (A) is formulation ① of an adjunction.)

The free-group functor  $\underline{F}: \underline{\text{Set}} \rightarrow \underline{\text{Grp}}$  is left adjoint to the forgetful  $\underline{U}: \underline{\text{Grp}} \rightarrow \underline{\text{Set}}$ .

(Similar statements hold for other free-algebra functors.)

A monoidal category  $\mathcal{C}$  is left closed iff for every object  $X$ , the functor  $(-) \otimes X: \mathcal{C} \rightarrow \mathcal{C}$  has a right adjoint. (The right adjoint is  $[X, -]$ .)

(The definition of left closure is formulation ② of adjunction.)

Any equivalence of categories  $(F, G, \alpha, \beta)$  is an adjunction with  $F \dashv G: \mathcal{D} \rightarrow \mathcal{C}$  and  $G \dashv F: \mathcal{C} \rightarrow \mathcal{D}$ .

(Use formulation ③, exploiting  $\alpha^{-1}$  and  $\beta^{-1}$  as appropriate.)

Consider  $(-)^* : \underline{\text{Vect}}^{\text{op}} \rightarrow \underline{\text{Vect}}$ , the contravariant functor mapping every vector space to its dual.

Then  $(-)^*$  is self adjoint :  $(-)^* \dashv (-)^* : \underline{\text{Vect}}^{\text{op}} \rightarrow \underline{\text{Vect}}$

(There are natural bijections

$$\begin{aligned} \underline{\text{Vect}}(V, W^*) &\cong \underline{\text{Vect}}(V \otimes W, K) \\ &\cong \underline{\text{Vect}}(W, V^*) \cong \underline{\text{Vect}}^{\text{op}}(V^*, W) . \end{aligned}$$

More generally, if  $\mathcal{C}$  is symm. mon. closed then, for every  $Y \in \text{obj } \mathcal{C}$ ,

$[-, Y]$  extends to a functor  $[-, Y] : \mathcal{C}^{\text{op}} \rightarrow \mathcal{C}$

that is self adjoint  $[-, Y] \dashv [-, Y] : \mathcal{C}^{\text{op}} \rightarrow \mathcal{C}$

Exercise • Verify this last point.

- Generalise the statement to categories with (non-symmetric) monoidal structure that are both left and right closed.

### Proposition (Adjoints determined up to isomorphism)

Suppose  $F \dashv G : D \rightarrow C$ .

① If  $F' \dashv G : D \rightarrow C$  then  $F' \cong F : C \rightarrow D$ .

② If  $F \dashv G' : D \rightarrow C$  then  $G' \cong G : D \rightarrow C$ .

### Proof

① we use reformulation ① of adjunction.

The adjunctions  $F \dashv G$  and  $F' \dashv G$  give us  $X \xrightarrow{\zeta_x} GFX$  and  $X \xrightarrow{\zeta'_x} GF'X$ .

Applying reformulation ① we get  $FX \xrightarrow{h_x} F'X$  and  $F'X \xrightarrow{h'_x} FX$  unique s.t.

$$\begin{array}{ccc} GFX & \xrightarrow{Gh_x} & GF'X \\ \zeta_x \uparrow & \nearrow \zeta'_x & \\ X & & \end{array} \quad \text{and} \quad \begin{array}{ccc} GF'X & \xrightarrow{Gh'_x} & GFX \\ \zeta'_x \uparrow & \nearrow \zeta_x & \\ X & & \end{array}$$

It is easy to show that their composites are identities and that  $(h_x)_x$  and  $(h'_x)_x$  are natural.

② Follows from ① because  $F \dashv G : D \rightarrow C$  iff  $G \dashv F : C^{op} \rightarrow D^{op}$ .  
(This 'iff' is immediate from the definition of adjunction.)

### Proposition (Composition of adjunctions)

If  $F \dashv G : D \rightarrow C$  and  $F' \dashv G' : E \rightarrow D$  then  $F'F \dashv GG' : E \rightarrow C$ .

Proof  $C(X, GG'Z) \cong D(FX, G'Z) \cong E(F'FX, Z)$ .

Naturality holds because the composition of natural bijections preserves it.  $\square$

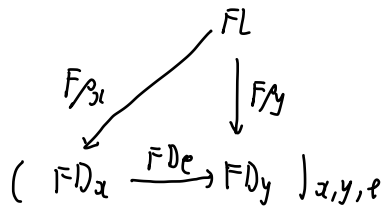
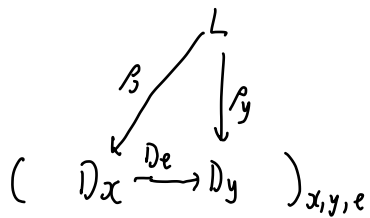
## Definition (Preservation of limits)

A functor  $F: C \rightarrow C'$  preserves the limit of a diagram

$D: G \rightarrow C$  ( $G$  a graph) if, for any limit cone

$(L, (L \xrightarrow{\beta_x} D_x)_{x \in G})$ , it holds that  $(FL, (FL \xrightarrow{F\beta_x} FD_x)_{x \in G})$

is a limit cone in  $C'$  for the diagram  $FD: G \rightarrow C'$ .



$F$  is said to preserve (existing) limits of shape  $G$  if  $F$  preserves the limit of  $D$ , for every diagram  $F: G \rightarrow C$  (that has a limit).

Instances of the above:  $F$  preserves (finite) products,  $F$  preserves pullbacks,  $F$  preserves equalisers,  $F$  preserves finite limits (a.k.a. is left exact) etc.

$F$  is said to preserve (existing) limits if  $F$  preserves (existing) limits of shape  $G$ , for every graph  $G$ .

There are dual definitions of what it means for  $F$  to preserve colimits

Proposition Suppose  $F' \dashv F : \mathcal{C} \rightarrow \mathcal{C}'$ . Then  $F$  preserves limits and  $F'$  preserves colimits.

(Right adjoints preserve limits, left adjoints preserve colimits.)

Proof Let  $(L, (L \xrightarrow{r_x} D_x)_{x \in I})$  be a limit cone for  $D: \mathcal{G} \rightarrow \mathcal{C}$ . We need to show  $(FL, (FL \xrightarrow{Fr_x} FD_x)_x)$  is a limit for  $FD: \mathcal{G} \rightarrow \mathcal{C}'$ .

Let  $(Z, (Z \xrightarrow{r_x} FD_x)_x)$  be a cone for  $FD$ .

Then  $(F'Z, (F'Z \xrightarrow{\psi^{-1}r_x} D_x)_x)$  is a cone for  $D$ .

$$\begin{array}{ccc}
 Z & \xrightarrow{\psi_1} & FL \\
 \downarrow r_x & \swarrow Fr_x & \searrow r_y \\
 & & FD_x \\
 & \searrow Fr_y & \downarrow Fr_y \\
 & & FD_y
 \end{array}
 \quad (FD_x \xrightarrow{FDc} FD_y)_{x,y \in I}$$

$$\begin{array}{ccc}
 F'Z & \xrightarrow{\lambda} & L \\
 \downarrow \psi^{-1}r_x & \swarrow Fr_x & \searrow \psi^{-1}r_y \\
 & & D_x \\
 & \searrow Fr_y & \downarrow Fr_y \\
 & & D_y
 \end{array}
 \quad (D_x \xrightarrow{Dc} D_y)_{x,y \in I}$$

Let  $F'Z \xrightarrow{\lambda} L$  be the unique cone morphism in  $\mathcal{C}$ .

Then  $Z \xrightarrow{\psi_1} FL$  is the unique cone morphism in  $\mathcal{C}'$ .

The statement for colimits follows by duality

because  $F' \dashv F : \mathcal{C} \rightarrow \mathcal{C}' \Leftrightarrow F' \dashv F : \mathcal{C}'^{op} \rightarrow \mathcal{C}^{op}$ .

□

## Application The 'arithmetic' of exponentiation

Suppose throughout that  $\mathcal{C}$  is left monoidal closed.

$$[1, x] \cong x$$

$$\text{cf. } x^1 = x$$

$$[y \otimes z, x] \cong [z, [y, x]]$$

$$x^{yz} = (x^y)^z$$

If  $\mathcal{C}$  has finite products then

$$[x, 1] \cong 1$$

$$1^x = 1$$

$$[x, y \times z] \cong [x, y] \times [x, z]$$

$$(yz)^x = y^x z^x$$

If  $\mathcal{C}$  has finite coproducts then

$$0 \otimes x \cong 0$$

$$0 \multimap x = 0$$

$$(y + z) \otimes x \cong (y \otimes x) + (z \otimes x)$$

$$(y + z)^x = y^x + z^x$$

If  $\mathcal{C}$  has finite products and coproducts then

$$[0, x] \cong 1$$

$$x^0 = 1$$

$$[y + z, x] \cong [y, x] \times [z, x]$$

$$x^{y+z} = x^y x^z$$

N.B., In the case of a cartesian closed category,  $\otimes$  and  $\times$  coincide.

Exercise Prove the above!