

Category Theory 2022-23

Lecture 1

7th October 2022

A category C is given by:

- A collection $|C|$ of objects.
- For every pair X, Y of objects, a collection $C(X, Y)$ of morphisms (maps)
- For every object X an identity morphism $1_X \in C(X, X)$
- For every triple X, Y, Z of objects a composition function $(-) \circ (-) : C(Y, Z) \times C(X, Y) \rightarrow C(X, Z)$

Such that:

- For every X, Y and $X \xrightarrow{f} Y$, $f \circ 1_X = f = 1_Y \circ f$ (identity laws)
- For every $X \xrightarrow{f} Y \xrightarrow{g} Z \xrightarrow{h} W$, $h \circ (g \circ f) = (h \circ g) \circ f$ (associativity)

Exercise The identities are determined uniquely by the composition function.

By the associativity law, chains of morphisms such as

$$X \xrightarrow{f} Y \xrightarrow{g} Z \xrightarrow{h} W$$

yield a unique composite map $X \xrightarrow{h \circ g \circ f} W$.

Equalities between such composites can be expressed as commutative diagrams, e.g.,

$$\begin{array}{ccc} X & \xrightarrow{f} & Y \\ 1_X \downarrow & \searrow f & \downarrow 1_Y \\ X & \xrightarrow{f} & Y \end{array}$$

expresses the identity laws.

Examples

Set The category of sets (and functions).

Objects: sets

Morphisms: $\text{Set}(X, Y)$ = the set of all functions from X to Y (which makes sense as X, Y are sets).

Identities: 1_X = the identity function $x \mapsto x : X \rightarrow X$

Composition: $g \circ f$ = the composite function $x \mapsto g(f(x))$

In the category Set the collection of objects $|\text{Set}|$ is the collection of all sets, which is not itself a set, rather a proper class.

However, for every X, Y , $\text{Set}(X, Y)$ is a set.

A category \mathcal{C} is locally small if, for every $X, Y \in |\mathcal{C}|$, the "hom set" $\mathcal{C}(X, Y)$ is a set.

\mathcal{C} is small if it is locally small and $|\mathcal{C}|$ is also a set.

The category Set is thus locally small but not small.

Grp The category of groups (and homomorphisms)

Objects : groups

Morphisms : $\text{Grp}(G, H) =$ set of homomorphisms from G to H

Identities and Composition : identity functions and function composition (as in Set)

Top The category of topological spaces (and continuous functions)

Objects : topological spaces

Morphisms : $\text{Top}(S, T) =$ set of continuous functions from S to T

Identities and Composition : function identities & composition

Vect_K The category of vector spaces (and linear transformations) over K ^{a field}

Objects : vector spaces over K

Morphisms : $\text{Vect}_K(U, V) =$ set of linear transformations from U to V

Identities and Composition : function identities & composition.

Commonalities

All examples so far are locally small but not small.

In all, morphisms are classes of functions, with identities and composition given by function identities and composition

The following examples are of a different character

Rel The category of (sets and) relations.

Objects: sets

Morphisms: $\text{Rel}(X, Y) = \text{set of relations between } X \text{ and } Y$

(Recall a relation between X and Y is a function $R: X \times Y \rightarrow \{\text{true}, \text{false}\}$.)

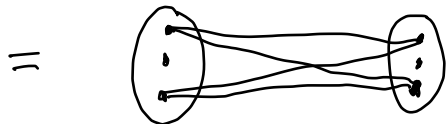
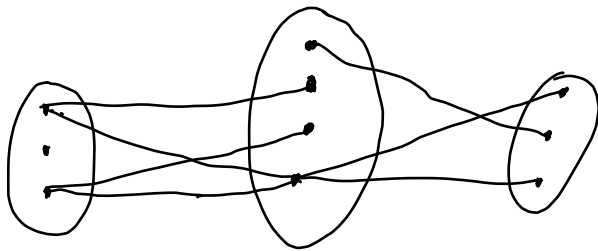
Identities: $1_X = \text{the identity relation} \quad x 1_X x' \Leftrightarrow x = x'$

Composition: $\text{SoR} = \text{relation composition } R; S$

$$x(R; S)z \Leftrightarrow \exists y \quad xRy \text{ and } ySz$$

This is again a locally small but not small category. However composition is not function composition.

E.g. A composition $3 \rightarrow 4 \rightarrow 3$



Let (G, \cdot, e) be any group.

G The group G as a category

Objects: Just one object, $*$

Morphisms: $\underline{G}(*, *) := G$

Identities: $1_* := e$

Composition: $y \circ x := y \cdot x$

More generally, the same construction gives a category M for any monoid (M, \cdot, e) .

G (and M) is a small category.

It has only one object.

Categoryfication:

monoid $:=$ one object category

Let (P, \leq) be any poset (partially ordered set)

P The poset P as a category

Objects: $|P| := P$

Morphisms: $P(x, y) := \begin{cases} \{\ast_{x,y}\} & \text{if } x \leq y \\ \emptyset & \text{otherwise} \end{cases}$

a singleton set

Identities: $1_x := \ast_{x,x}$

Composition: $\ast_{y,z} \circ \ast_{x,y} := \ast_{x,z}$

More generally, the same construction gives a category for any preorder (P, \leq) .

P is a small category.

Moreover, there is at most one morphism between any two objects.

Categorification:

preorder := category in which hom sets have at most one element

Special Morphisms in a category \mathcal{C} the inverse of f

$X \xrightarrow{f} Y$ is an isomorphism (iso) if there exists $Y \xrightarrow{f^{-1}} X$ such that $f^{-1} \circ f = 1_X$ and $f \circ f^{-1} = 1_Y$.

Trivially every identity is an isomorphism. It is its own inverse.

Exercise • If f is an isomorphism then f^{-1} is determined uniquely by f .

• If $X \xrightarrow{f} Y \xrightarrow{g} Z$ are two isomorphisms then $g \circ f$ is also an iso.

In set the isomorphisms are the bijections

Grp

bijective homomorphisms

Top

homeomorphisms

Vect_K

linear isomorphisms

Rel

graphs of bijections

G every morphism is an isomorphism

P (a poset) the only isomorphisms are the identities

Categorification:

group := 1-object category in which every morphism is an isomorphism

poset := category in which hom-sets have at most one morphism and the only isomorphisms are identities.

$X \xrightarrow{f} Y$ is a Monomorphism (mono) if, for every parallel pair $z \xrightarrow{x} X$, $fox = foy \Rightarrow x = y$

Exercises • Every isomorphism is a monomorphism

- If $X \xrightarrow{f} Y \xrightarrow{g} Z$ are two monos then gof is also mono.
- If gof is mono then so is f .
- If gof is iso and g is mono then f and g are both iso.

In Set the monomorphisms are the injections

Top

continuous injections

Grp

injective homomorphisms

Vect

injective linear transformations

P (a poset) every morphism is a monomorphism

Weekly puzzle In Rel :

- Is there a mono from 2 (a 2-element set) to 1 (a singleton)?
- Is there a mono from 3 to 2?
- Characterise the monomorphisms in Rel.