Where did we leave off?

```
lemma radius_pos (a b : Pt) (α : Circ)
  (a_cent_of_α : Center_of_Circ a α)
  (b_on_α : Pt_on_Circ b α) :
  0 < dist a b := by</pre>
```

Main idea: axiom

```
Pts_of_Circ : \forall (\alpha : Circ), \exists (b c : Pt), Pt_on_Circ b \alpha \land Pt_on_Circ c \alpha \land b \neq c
```

tells us that there are *two distinct* points on any circle! Which will somehow imply that a circle can't consist of a single point (i.e. have radius zero)

Question: Why do we even want to have two distinct points on a circle? Why not assume we only have one?

Challenge question: Given a circle with no center, how to find the center???

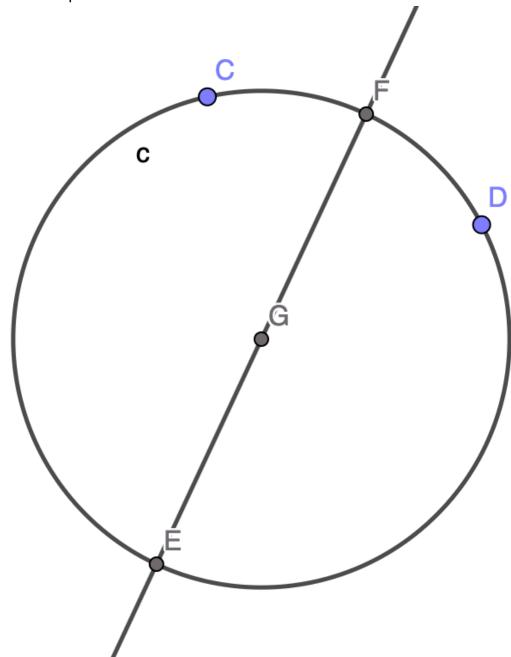
[Hint: start with two distinct points on the circle!]

Solution: Step 1: get two distinct points on the circle

Step 2: by straightedge and compass, draw the perpendicular bisector of those two points

Step 3: by some process (???) this line intersects the circle in two points, E and F

Step 4: the midpoint of E and F is the center



Back to original problem of positivity of radius:

By dist_nonneg, either dist a b = 0 or 0 < dist a b -- and in the latter case, we're done. Formally, first let's record:

```
have : 0 ≤ dist a b := dist_nonneg a b
```

And then we want to prove

```
have : 0 = dist a b v 0 < dist a b := by
```

and this is a basic statement about the real numbers, not something about geometry, so the name for the theorem that closes this goal is not known to us. But it's such a basic fact that it must be somewhere in the Lean Mathlib library. The way to search for a theorem that you suspect is exact ly something already in the library is to write exact? When you do that, if that theorem already exists, you'll see in the goal state something like:

```
Try this: exact Decidable.eq_or_lt_of_le this
```

And when you click on the exact, it replaces your text with the text that solves the goal. Next step: We now have a hypothesis:

```
this : 0 = dist a b v 0 < dist a b
```

We can't use constructor because (1) not in the goal, and (2) not Λ . We want to break the logic by cases according to which case it is. That's called: by_cases. More generally, if we have a hypothesis H: P v Q. Then writing

```
cases this with
| inl dist_eq =>
    sorry
| inr dist_notEq => exact dist_notEq
```

splits the goal into two goals, one assuming that P is true, and another assuming that Q is true. In our case, Q is exactly the original goal, so exact dist_notEq solves everything. We need to see what happens in case P, where the game board says:

```
dist_eq : 0 = dist a b

- 0 < dist a b</pre>
```

We're obviously in a state of contradiction. If we can prove false, then we can prove whatever we want (we must've stumbled into a contradiction). Note: we're trying to prove that the radius is always positive, so this is the part of the argument that says the radius can't be zero. To replace any goal by false, the tactic is called exfalso.

...

```
If you have H1: X = Y and H2: Y = Z, then have H3: X = Z := H1. trans H2
```

. . .

We've gotten to the point where we proved e = f and $e \ne f$. Recall that the latter is by definition: $e = f \rightarrow False$.

```
import Mathlib
class EuclideanPlane where
  Pt: Type
  Line : Type
  Circ: Type
  Pt_on_Line : Pt → Line → Prop
  Line_of_Pts : ∀ a b : Pt, ∃ L : Line, (Pt_on_Line a L) ∧ (Pt_on_Line b L)
  Pt_on_Circ : Pt → Circ → Prop
  Center_of_Circ : Pt → Circ → Prop
  Circ_of_Pts : \forall a b : Pt, \exists \alpha : Circ, (Center_of_Circ a \alpha) \land (Pt_on_Circ b
α)
  dist : Pt \rightarrow Pt \rightarrow \mathbb{R}
  dist_symm : ∀ (a b : Pt), dist a b = dist b a
  dist_nonneg : \forall (a b : Pt), 0 \le dist a b
  dist pos def : \forall (a b : Pt), dist a b = 0 \leftrightarrow a = b
  Circs inter : Circ → Circ → Prop
  Pt_in_Circ : Pt → Circ → Prop
  Circs_inter_iff : \forall (\alpha \beta : Circ), Circs_inter \alpha \beta \leftrightarrow
     (∃ (a b : Pt), Pt_on_Circ a α ∧ Pt_in_Circ a β ∧
       Pt_on_Circ b \beta \Lambda Pt_in_Circ b \alpha \Lambda a \neq b)
  Pt_on_Circ_iff : \forall (a b c : Pt) (\alpha : Circ), Center_of_Circ a \alpha \rightarrow
     Pt on Circ b \alpha \rightarrow (Pt on Circ c \alpha \leftrightarrow dist a c = dist a b)
  Pt_in_Circ_iff : \forall (a b c : Pt) (\alpha : Circ), Center_of_Circ a \alpha \rightarrow
     Pt_on_Circ b \alpha \rightarrow (Pt_in_Circ c \alpha \leftrightarrow dist a c < dist a b)
  Pts_of_Circ : \forall (\alpha : Circ), \exists (b c : Pt), Pt_on_Circ b \alpha \land Pt_on_Circ c \alpha
\Lambda b \neq c
variable [EuclideanPlane]
namespace EuclideanPlane
def IsEquilateralTriangle (a b c : Pt) : Prop := (dist a b = dist b c) Λ
(dist a b = dist a c)
-- An easier helper lemma:
lemma dist_self (a : Pt) : dist a a = 0 := by
```

```
have := (dist_pos_def a a).2
  apply this
  -- now the goal is to prove `⊢ a = a` and that's solved by
  -- the reflexive property of the equal sign, called `rfl`
  rfl
-- Another helper lemma: the "radius" is always strictly positive
lemma radius pos (a b : Pt) (\alpha : Circ)
    (a_cent_of_\alpha : Center_of_Circ a \alpha)
   (b on \alpha: Pt on Circ b \alpha):
    0 < dist a b := by</pre>
  -- idea:
  -- by `dist_nonneg`, either `0 = dist a b` or `0 < dist a b`
  have : 0 \le \text{dist a b} := \text{dist nonneg a b}
  have : 0 = dist a b v 0 < dist a b := by
    exact Decidable.eq_or_lt_of_le this
  cases this with
  | inl dist_eq =>
   exfalso
  -- Now we're arguing by contradiction (we want to prove `false`!)
    have := Pts_of_Circ \alpha
    obtain (e, f, e_on_\alpha, f_on_\alpha, e_ne_f) := this
  -- Using `Pts_of_Circ`, we can `obtain` two new points, say, `e` and
  -- `f` that are distinct and both on `\alpha`.
    have a_eq_b : a = b := (dist_pos_def a b).1 dist_eq.symm
  -- If `dist a b = 0` then `a = b` (by `dist_pos_def`).
  -- But `Pt on Circ iff` says that radii are well-defined
  -- that is, `dist a e = dist a b = 0 = dist b e` So
    have := (Pt_on_Circ_iff \ a \ b \ e \ \alpha \ a_cent_of_\alpha \ b_on_\alpha).1 \ e_on_\alpha
    have := this.trans dist_eq.symm
    have a_eq_e : a = e := (dist_pos_def a e).1 this
    have := (Pt_on_Circ_iff \ a \ b \ f \ \alpha \ a_cent_of_\alpha \ b_on_\alpha).1 \ f_on_\alpha
    have : dist a f = 0 := this.trans dist eq.symm
    have a_{eq}f : a = f := (dist_pos_def a f).1 this
    have e_eq_f : e = f := a_eq_e.symm.trans a_eq_f
  -- `b = e` (again by `dist_pos_def`). The same argument proves
```

```
-- `b = f`, which by transitivity proves `e = f`.
  -- Which is a contradiction!
  -- Hint: e \neq f by defintion means: e = f \rightarrow false.
    exact e_ne_f e_eq_f
  | inr dist_notEq => exact dist_notEq
  -- if the second case, then we're done. So we just want to rule
  -- out the first case
-- Let's make a helper lemma: given a circle \alpha and a point \alpha that's the
center of the circle, `a` is in the interior of the circle
lemma in_Circ_of_center (a : Pt) (\alpha : Circ) (a_cent_\alpha : Center_of_Circ a \alpha)
    Pt in Circ a \alpha := by
  have := Pts_of_Circ \alpha
  obtain (b, c, b_on_\alpha, c_on_\alpha, b_ne_c) := this
  have := (Pt_in_Circ_iff \ a \ b \ a \ \alpha \ a_cent_\alpha \ b_on_\alpha).2
  apply this
  have dist_aa : dist a a = 0 := dist_self a
  sorry -- HOMEWORK
theorem PropIp1 (a b : Pt) : ∃ (c : Pt), IsEquilateralTriangle a b c := by
  have Step1 : \exists (\alpha : Circ), Center_of_Circ a \alpha \wedge Pt_on_Circ b \alpha :=
    Circ_of_Pts a b
  obtain (\alpha, a\_center\_of\_\alpha, b\_on\_\alpha) := Step1
  have Step2 : \exists (\beta : Circ), Center_of_Circ b \beta \Lambda Pt_on_Circ a \beta :=
    Circ_of_Pts b a
  obtain (β, b_center_of_β, a_on_β) := Step2
  have Step3 : Circs inter \alpha \beta := by
    have := (Circs_inter_iff \alpha \beta).2
    apply this
    use b, a
    constructor
    exact b_on_α

    constructor

         have := (Pt_in_Circ_iff b a b β b_center_of_β a_on_β).2
```

```
sorry -- Pt_in_Circ b β

    constructor

       exact a_on_β
       · sorry -- Pt_in_Circ a α
have Step4 : \exists (c d : Pt), Pt_on_Circ c \alpha \land Pt_on_Circ d \alpha \land
  Pt_on_Circ c \beta \Lambda Pt_on_Circ d \beta \Lambda c \neq d :=
    Pts_of_Circs_inter \alpha \beta Step3
obtain (c, d, c_on_\alpha, d_on_\alpha, c_on_\beta, d_on_\beta, c_ne_d) := Step4
have Step5 : dist a c = dist a b :=
  (Pt_on_Circ_iff a b c \alpha a_center_of_\alpha b_on_\alpha).1 c_on_\alpha
have Step6 : dist b a = dist b c :=
  (Pt_on_Circ_iff b c a β b_center_of_β c_on_β).1 a_on_β
use c
unfold IsEquilateralTriangle
constructor
· convert Step6 using 1
  have := dist_symm a b
 exact this
have Step7 := Step5.symm
 exact Step7
```

Homework: Prove in_Circ_of_center