

$$1) \quad i) \quad 2RTT + RTT + (3 \times 9L/R + 2L/R) \div \frac{1}{2} + 2 \times 9Q$$

$$= 3RTT + 58 L/R + 18Q$$

$$ii) \quad 11RTT + 58 L/R + 18Q$$

$$iii) \quad 20RTT + 58 L/R + 18Q$$

$$iv) \quad 12RTT + 34 L/R + 9Q$$

$$v) \quad 8RTT + 22 L/R + 5Q$$

$$2) \quad i) \quad \text{linkrate: } \frac{375 \text{ k}}{6 \text{ mil}} = 0.0595 \text{ seconds}$$

$$T_B: \quad \frac{120}{60} \times 0.0595 = 0.119$$

$$\text{Avg access delay: } \frac{0.0595}{(1 - 0.119)} = 0.06754$$

$$\text{Total: } 0.06754 \times 2 \times 4 = 6.06754 \text{ seconds}$$

$$2) \text{ ii) new link delay : } \frac{0.075}{1-(0.4)(0.15)} = 0.07979 \text{ seconds}$$

$$\text{If miss} \Rightarrow 0.07979 + 2 + 4 = 6.07979 \text{ seconds,,}$$

$$\text{Total average Time} \Rightarrow (0.6)(0) + (0.4)(6.07979) = 2.431916 \text{ seconds,,}$$

Reduced from 6.088235 seconds to 2.431916 seconds,,

(ii)

$$T = \frac{375}{5M} = 0.075$$

$$T_b = \frac{120 \times 0.3}{60} \times 0.075 = 0.045$$

$$\text{delay} = \frac{T}{1-T_b} = 0.07853$$

$$T = \frac{375}{10M} = 0.0375$$

$$T_b = \frac{120 \times 0.7}{60} \times 0.0375 = 0.0525$$

$$\text{delay} = \frac{T}{1-T_b} = 0.03958$$

$$\begin{aligned} \text{Total time} &= 0.07853 \times 0.3 + 0.03958 \times 0.7 + 6 \\ &= 6.051265,, \end{aligned}$$

$$3) \quad 12 S/R > RTT > 4 S/R ;$$

$$2RTT + 4S/R$$

$$RTT + 4S/R$$

$$RTT + 16S/R + 28S/R$$

$$= 4RTT + 52S/R$$

$$RTT \geq 12 S/R :$$

$$2RTT + 4S/R$$

$$RTT + 4S/R$$

$$RTT + 4S/R$$

$$RTT + 28S/R$$

$$= 5RTT + 40S/R$$

4) i) Sender Events:

Send data octets 1-1000

Send data octets 1001-2300

Receive ack for 1000

Receive ack for 2300

Receiver Events:

advertise window = 2300

ack up to 1000, window = 1000

ack up to 2300, window = 0

i) Sender Events:

Send data octets 2501, 3500

Send data octets 3501, 4500

Receive ack for 3500

Receiver ack for 4500

Receiver ack for 4500

Receiver Events.

Application reads 1000 octets

ack up to 2500, window = 2500

ack up to 3500, window = 1500

ack up to 4500, window = 500

Application reads 1000 octets

ack up to 4500, window = 1500

5) i) TCP slow start is operating at $[4, 7]$
because the window sizes are doubling up.

ii) TCP congestion avoidance is operating at $[1, 3]$ and $[8, 16]$
because of the linear increments to the window size

iii) After the 3rd transmission, packet loss is recognized by timeout
and the window size is reset to 1.

iv) After the 16th transmission, segment loss is recognized by a triple
duplicate Ack where window size is halved

v) Maximum possible initial size on first transmission round is
30, when TCP congestion avoidance started working.

vi) Since a timeout was detected, the congestion size of
32 was reduced to 1 at the 4th transmission round.
So the threshold becomes 16 //

vii) Since triple duplicate Ack was detected, the congestion size of
18 got reduced to 9 during the 11th transmission round.
So the threshold is 9 //

viii) The cws will be 3 and threshold of 3 at 16th