Modular I/O Reasoning in DimSum

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Modular I/O Reasoning

```
int echo () :=
  let c := getc();
  putc(c);
  return 0;
```

reads from increasing sequence

```
{...}
echo
{...}
```

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$$\{\lambda\ \mathsf{es},\ \ulcorner \mathsf{es} = [\ ] \urcorner\}$$

$$\texttt{echo}$$

$$\{\lambda\ \mathsf{v},\ \ulcorner \mathsf{v} = 0 \urcorner\}$$

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reads from increasing sequence

$$\Big\{\lambda\ \mathsf{es},\ \ulcorner \mathsf{es} = [\ ]\urcorner * \{\lambda\ \mathsf{es},\ \ulcorner \mathsf{es} = [\ ]\urcorner\}\texttt{getc}\{\lambda\ \mathsf{v},\ \{\lambda\ \mathsf{es},\ \ulcorner \mathsf{es} = \mathsf{v}\urcorner\}\texttt{putc}\{\_\}\ \}\Big\}$$

$$\texttt{echo}$$

$$\{\lambda\ \mathsf{v},\ \ulcorner \mathsf{v} = 0\urcorner\}$$

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$$\{\lambda\ \mathsf{es},\ \ulcorner \mathsf{es} = [\ ] \urcorner * \exists\ \mathsf{v},\ \mathsf{P\ v} * (\mathtt{getc\_spec}\ \mathsf{P}) * (\mathtt{putc\_spec}\ \mathsf{P})\}$$

$$\mathtt{echo}$$

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reads from increasing sequence

$$\{\lambda\ \mathsf{es},\ \ulcorner \mathsf{es} = [\,]\urcorner * \exists\ \mathsf{v},\ \mathsf{P}\ \mathsf{v}\}\ \texttt{getc}\ \{\lambda\ \mathsf{ret},\ \ulcorner \mathsf{ret} = \mathsf{v}\urcorner * \mathsf{P}\ (\mathsf{v} + 1)\}$$

$$\{\lambda\ \mathsf{es},\ \exists\ \mathsf{v},\ \ulcorner \mathsf{es} = \mathsf{v}\urcorner * \mathsf{P}\ (\mathsf{v} + 1)\}\ \texttt{putc}\ \{\lambda\ \mathsf{ret},\ \mathsf{P}\ (\mathsf{v} + 1)\}$$

$$\{\lambda\ \mathsf{es},\ \ulcorner \mathsf{es} = [\,]\urcorner * \exists\ \mathsf{v},\ \mathsf{P}\ \mathsf{v} * (\texttt{getc\_spec}\ \mathsf{P}) * (\texttt{putc\_spec}\ \mathsf{P})\}$$

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$$\{\lambda\ \mathsf{v},\ \ulcorner \mathsf{v} = 0\urcorner\}$$

DimSum

- Decentralized/language-agnostic multi-language semantics
  - No fixed source
  - No fixed set of languages
  - No fixed memory model
  - No fixed notion of linking
- Notion of semantic linking: $\oplus$
  - Link semantic components (modules) rather than syntactic
  - Link programs with specifications (abstract program)
- Program semantics as LTS, interaction via synchronization
- Reason locally in terms of interaction of two modules

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Proofs in DimSum

$$\boxed{[\![\mathsf{echo}_{\mathsf{rec}} \quad \oplus \quad \mathsf{getc}_{\mathsf{spec}}]\!]} \preceq \quad [\![\mathsf{echo}_{\mathsf{spec}}]\!]$$

Proofs in DimSum

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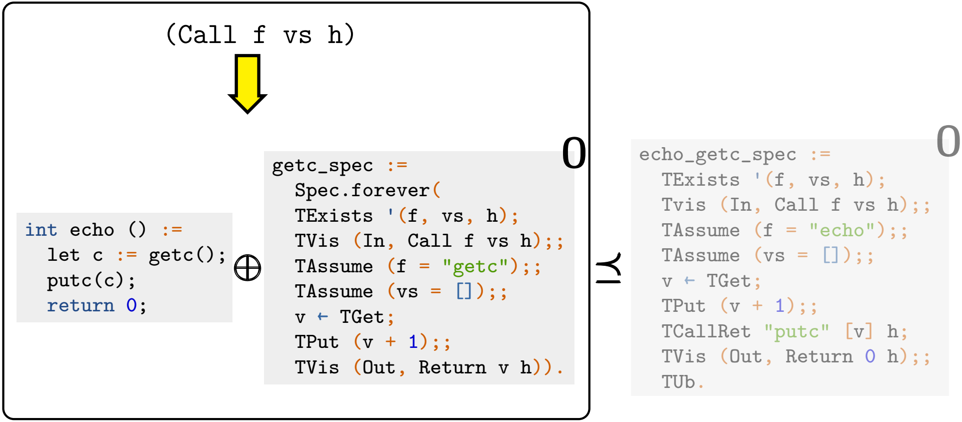
Proofs in DimSum



(Call f vs h)

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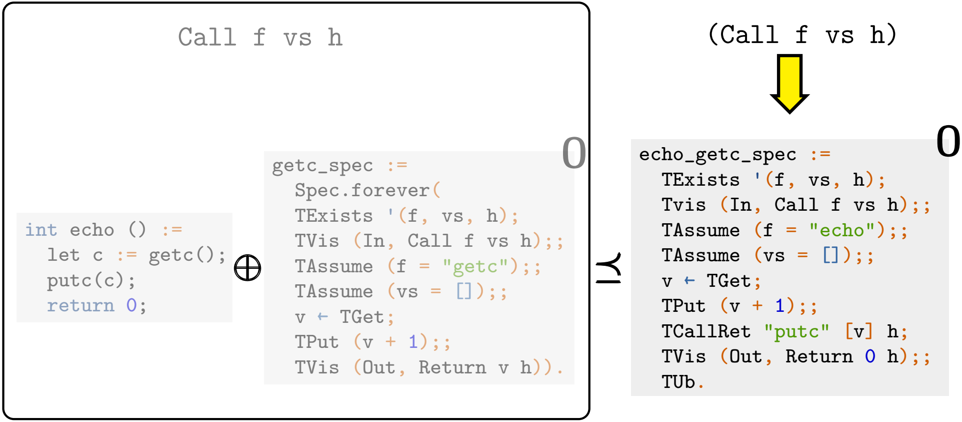

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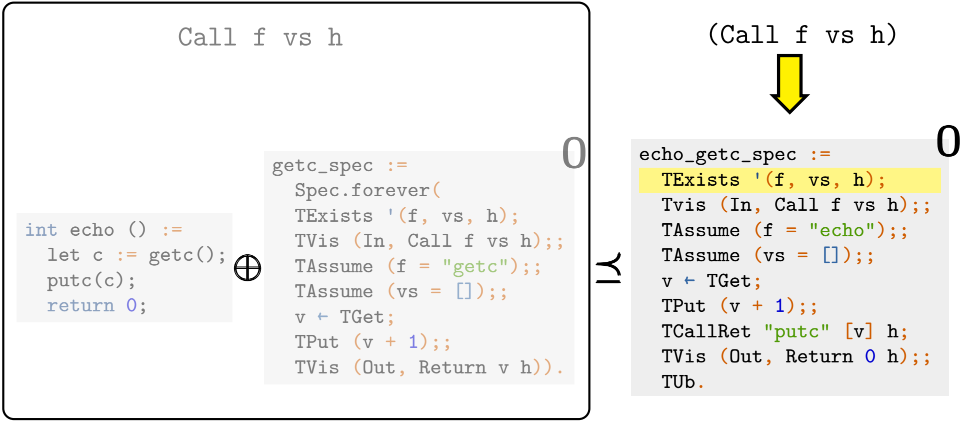

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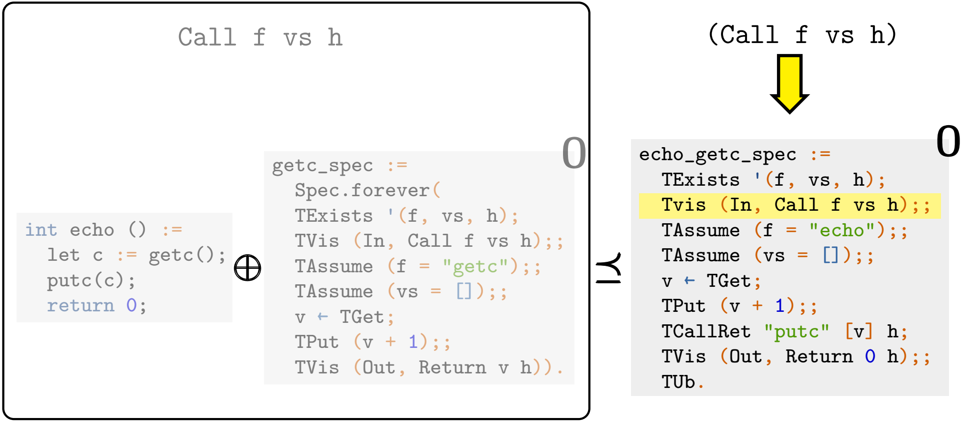

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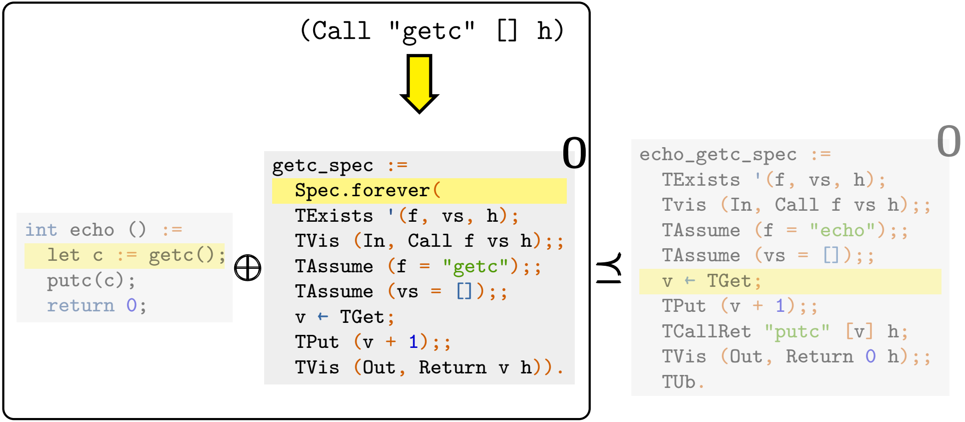

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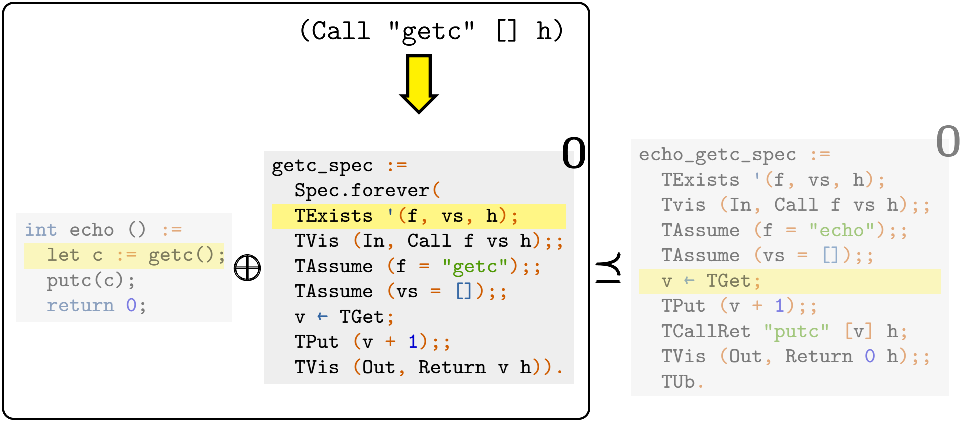

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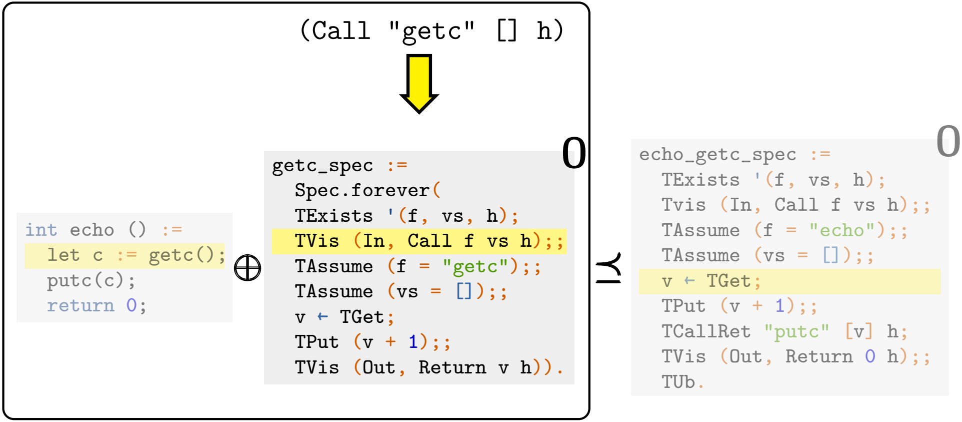

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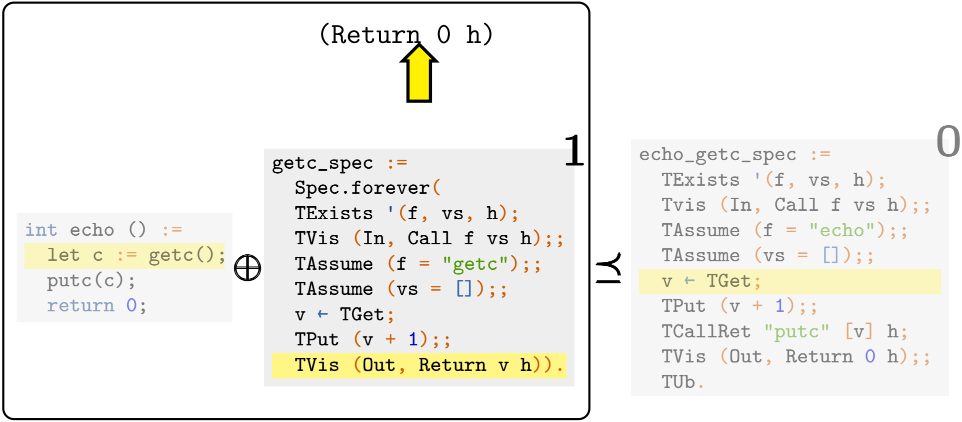

```
(Return 0 h)
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⊕

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⪯

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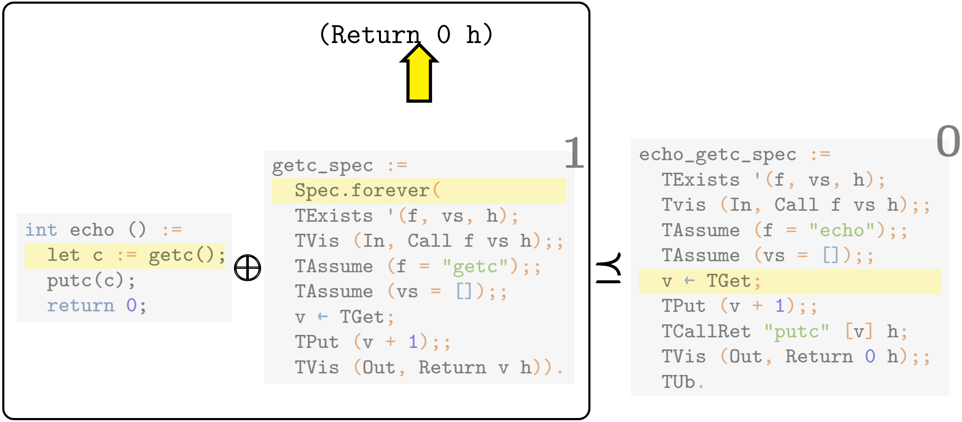
Proofs in DimSum



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$\oplus$

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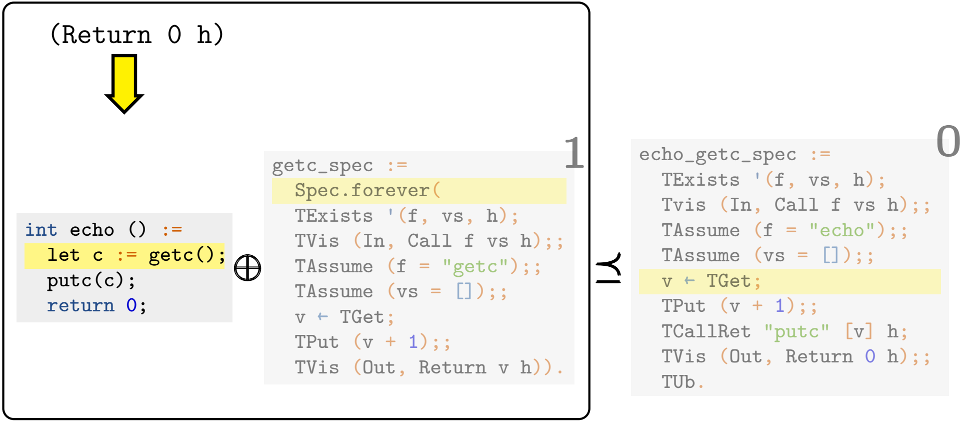
Proofs in DimSum



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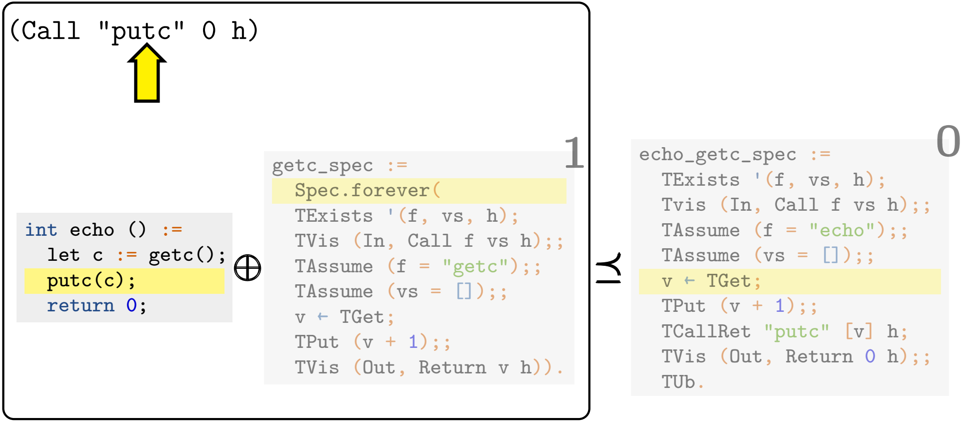
Proofs in DimSum



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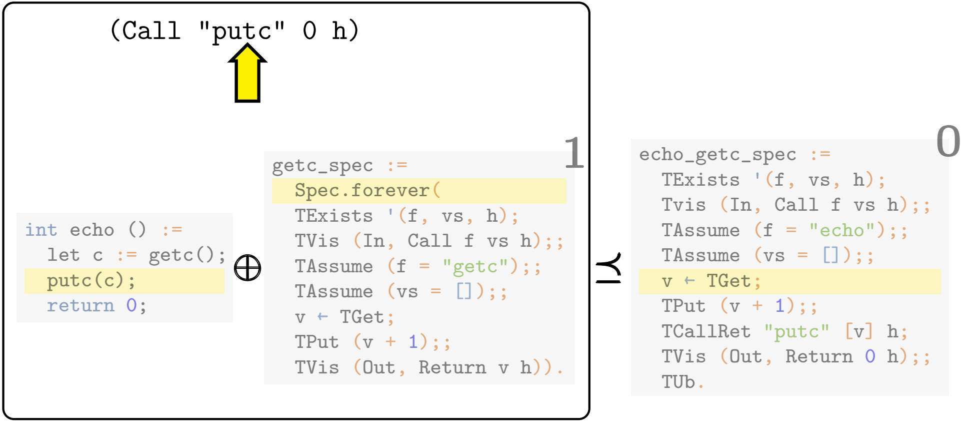
Proofs in DimSum



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Proofs in DimSum

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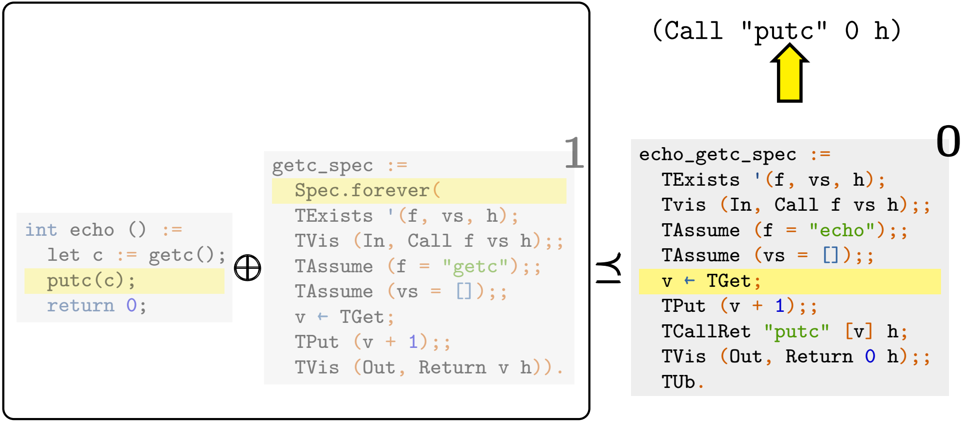
Proofs in DimSum



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Proofs in DimSum

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≼

(Call "putc" 0 h)

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Proofs in DimSum

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Proofs in DimSum

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⪯

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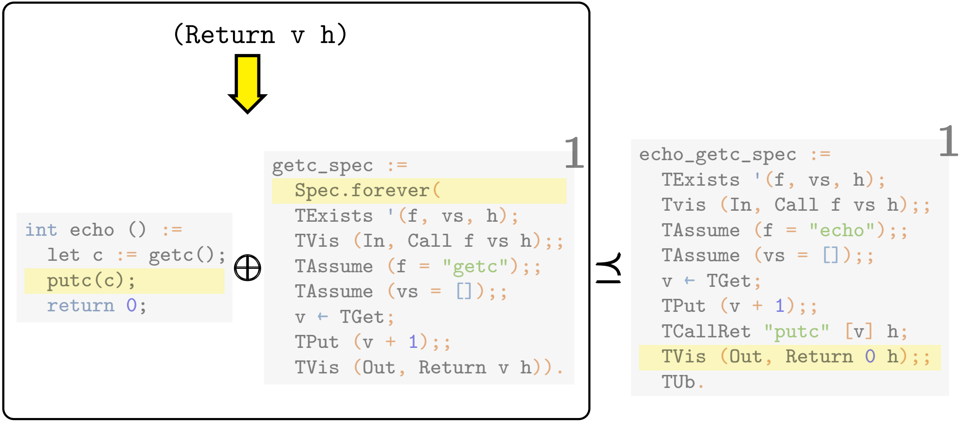
Proofs in DimSum



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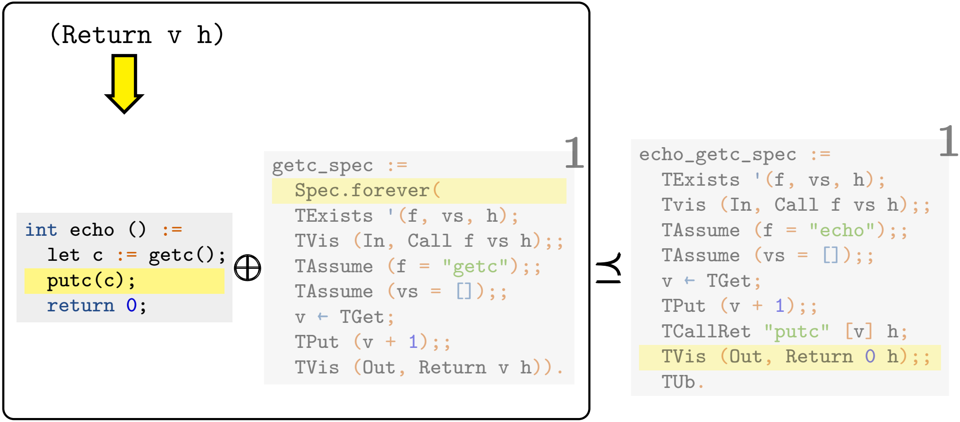

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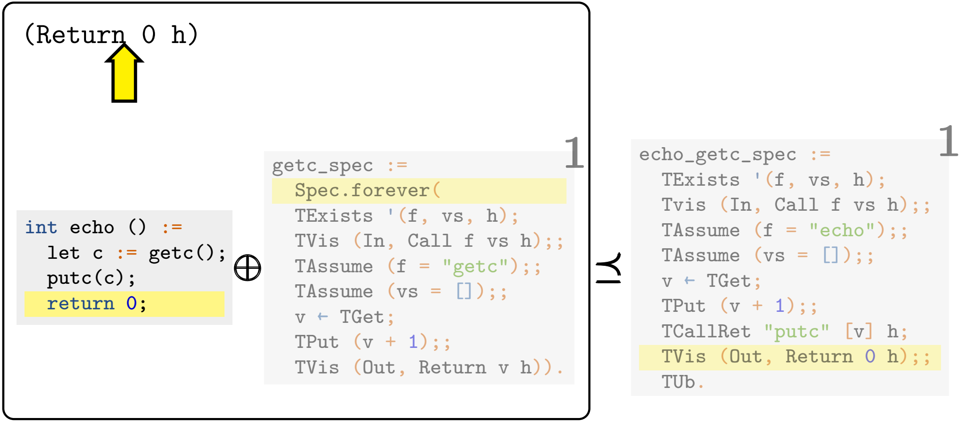

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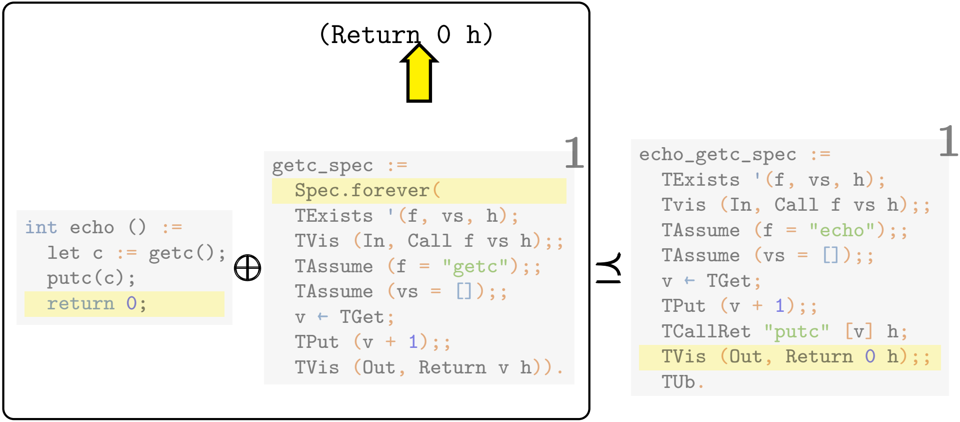

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Towards Hoare-style Reasoning

$$[\![\mathsf{echo}_{\mathsf{rec}} \oplus \mathsf{getc}_{\mathsf{spec}}]\!] \quad \preceq \quad [\![\mathsf{echo}_{\mathsf{spec}}]\!]$$

Towards Hoare-style Reasoning

$[\![\mathsf{echo}_{\mathsf{rec}} \oplus \mathsf{getc}_{\mathsf{spec}}]\!] \approx\!> (\boldsymbol{\lambda} \kappa_t \, \sigma_t,$

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$$\llbracket \mathsf{echo}_{\mathsf{rec}} \oplus \mathsf{getc}_{\mathsf{spec}} \rrbracket \approx\!> (\boldsymbol{\lambda} \kappa_t \sigma_t, \llbracket \mathsf{echo}_{\mathsf{spec}} \rrbracket \approx\!> (\boldsymbol{\lambda} \kappa_s \sigma_s,$$

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Towards Hoare-style Reasoning

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$$[\![\mathsf{echo}_{\mathsf{rec}} \oplus \mathsf{getc}_{\mathsf{spec}}]\!] \approx> \prod_s$$

$$[\![\mathsf{echo}_{\mathsf{rec}}]\!] \approx> \prod_{\oplus}(\Pi_s)$$

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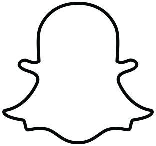

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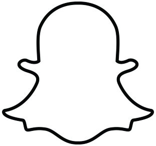

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$$[\![\text{echo}_{\text{rec}}]\!] \approx\!> \prod$$

Towards Hoare-style Reasoning

$$[\![\mathsf{echo}_{\mathsf{rec}} \oplus \mathsf{getc}_{\mathsf{spec}}]\!] \approx> (\boldsymbol{\lambda} \kappa_t \sigma_t, [\![\mathsf{echo}_{\mathsf{spec}}]\!] \approx> (\boldsymbol{\lambda} \kappa_s \sigma_s, \kappa_t = \kappa_s * \sigma_t \preceq \sigma_s))$$

$$[\![\mathsf{echo}_{\mathsf{rec}} \oplus \mathsf{getc}_{\mathsf{spec}}]\!] \approx> \prod_s$$

$$[\![\mathsf{echo}_{\mathsf{rec}}]\!] \approx> (\boldsymbol{\lambda} \kappa_l \sigma_l, \texttt{if}\,\_\,\texttt{then}\ [\![\mathsf{getc}_{\mathsf{spec}}]\!] \approx> \ldots\ \texttt{else}\ \Pi_s \kappa_l \sigma)$$

$$[\![\mathsf{echo}_{\mathsf{rec}}]\!] \approx> \prod_{\bigoplus} \left(\gamma_{\sigma_1}, \gamma_{\sigma_2}, \gamma_{\sigma_\oplus}, \Pi_s(\gamma_{\sigma_{spec}})\right)$$

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$$\mathsf{TGT\ Call}\ \text{“echo”}\ \mathsf{es}\ @\ \Pi\ \{\{\Phi\}\}$$

Towards Hoare-style Reasoning

$$\llbracket \mathsf{echo}_{\mathsf{rec}} \oplus \mathsf{getc}_{\mathsf{spec}} \rrbracket \approx\!> (\boldsymbol{\lambda} \kappa_t \sigma_t, \llbracket \mathsf{echo}_{\mathsf{spec}} \rrbracket \approx\!> (\boldsymbol{\lambda} \kappa_s \sigma_s, \kappa_t = \kappa_s * \sigma_t \preceq \sigma_s))$$

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TGT Call “echo” es @ Π $\{\{\Phi\}\}$

$$\gamma_{\sigma_1} \rightsquigarrow \sigma_1, \gamma_{\sigma_2} \rightsquigarrow \sigma_2, \gamma_{\sigma_\oplus} \rightsquigarrow \sigma_\oplus, \gamma_{\sigma_{spec}} \rightsquigarrow \sigma_{spec}$$

Rotation Project

- Get familiar with DimSum (and Iris)
- Experiment with new reasoning style
- Proof refinements for multiple example programs

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```
int echo () :=
  let c := getc();
  putc(c);
  return 0;
```

```
int getc (l) :=
  let _ := read(l);
  return *l;
```

```
void echo () :=
  let c := getc();
  putc(c);
  echo ();
```

```
global pos = 0;
read (l, c) :=
  if (c <= 0) {
    return 0;
  } else {
    l <- *pos;
    pos <- *pos + 1;
    ret = read(l + 1, c -1);
    return ret + 1;
  }
```

```
int echo () :=
  let c := getc();
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(Modular) Hoare-style Reasoning

```
TGT Call "echo" es @ Π
  PRE |-∗: es POST_e, ⌜es = []⌝ ∗
    TGT Call "getc" es @ Π
      PRE |-∗: es POST, ⌜es = []⌝ ∗
      POST |∗: ret,
        TGT Call "putc" es @ Π
          PRE |-∗: es POST, ⌜es = [v]⌝ ∗
          POST |∗: _,
  POST_e |∗: ret, ⌜ret = 0⌝.
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(Modular) Hoare-style Reasoning

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TGT Call "getc" es @ Π
  PRE |-∗: es POST, ∃ v, P v ∗ ⌜es = []⌝ ∗
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Example 1

$$[\![\mathsf{echo}_{\mathsf{rec}} \oplus \mathsf{getc}_{\mathsf{spec}}]\!] \preceq [\![\mathsf{echo}_{\mathsf{spec}}]\!]$$

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getc_spec :=
  Spec.forever(
  TExists '(f, vs, h);
  TVis (In, Call f vs h);;
  TAssume (f = "getc");;
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  v ← TGet;
  TPut (v + 1);;
  TVis (Out, Return v h)).
```

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Lemma sim_getc_spec `{!specGS} Π Φ :
  switch Π
    PRE |-∗: κ σ1 POST,
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    POST Tgt _ _ |∗: σ' Π',
      ∃ v, ⌜f = "getc"⌝ ∗ ⌜es = []⌝ ∗ spec_state v ∗ ⌜σ' = σ1⌝ (∗)
  switch Π'
    PRE |-∗: κ σ POST,
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$$[\![\mathsf{echo}_{\mathsf{rec}} \oplus \mathsf{getc}_{\mathsf{spec}}]\!] \preceq [\![\mathsf{echo}_{\mathsf{spec}}]\!]$$

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Lemma sim_getc fns Π_l Π_r PL σi :
  "getc" ↪ None -∗
  PL σi -∗
  ⌜σi.1 ≡ getc_spec⌝ -∗
  ⌜σi.2 = 0⌝ -∗
  □ switch_linked_fixed Tgt Π_l Π_r
      PRE |-∗: σ_l POST, ∃ h v σg, PL σg ∗
      POST (ERCall "getc" [] h) σg |∗: σr Π_r',
    switch_link Tgt Π_r'
      Pre |-∗: σ_r' POST, ∃ h'
      POST (ERReturn (ValNum v) h') _ σ_l |∗: _ Π_l',
        ⌜Π_l' = Π_l⌝ ∗ PL σ_r' ==∗
  ∃ P, P 0 ∗ □ rec_fn_spec_hoare Tgt Π_l "getc" (getc_fn_spec P).
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  ⌜σi.1 ≡ getc_spec⌝ -∗
  ⌜σi.2 = 0⌝ -∗
  □ switch Π_l
      PRE |-∗: κ σ0 POST, ∃ h v σg, PL σg ∗
      POST Tgt _ _ |∗:  σi0 Πi, ⌜σi0 = σg⌝ ∗ ⌜Πi = Π_r⌝ ∗
    switch Πi
      PRE |-∗: κ' σ POST0, ∃ e' : rec_ev, ⌜κ' = Some (Incoming, e')⌝ ∗
      POST0 Tgt _ _ |∗: σr Πr, ⌜σr = σ⌝ ∗ ⌜e' = ERCall "getc" [] h⌝ ∗
    switch Πr
      PRE |-∗: κ0 σ1 POST1, ∃ h', ⌜κ0 = Some (Outgoing, ERReturn v h')⌝ ∗
      POST1 Tgt _ _ |∗: σi1 Πi0, ⌜σi1 = σ0⌝ ∗
    switch Πi0
      PRE |-∗: κ'0 σ2 POST2, ∃ e'0, ⌜κ'0 = Some (Incoming, e'0)⌝ ∗
      POST2 Tgt _ _ |∗: σr0 Πr0,
        ⌜σr0 = σ2⌝ ∗ ⌜e'0 = ERReturn v h'⌝ ∗ ⌜Πr0 = Π_l⌝ ∗ PL σ1 ==∗
  ∃ P, P 0 ∗ □ rec_fn_spec_hoare Tgt Π_l "getc" (getc_fn_spec P).
```

Outcome & Takeaways

- Lemma for `TCallRet`
- Keep Πs the same - new lemmas for linking
- Balance between Abstraction and Information
- Balance between Hacking and Thinking

ToC

1. Background

2. DimSum Proof

3. Hoare Style

4. Examples