

<b>HOMEWORK 6</b>	<b>Transaction Processing</b>
<b>Due Wed, Nov 11 at 11:30 pm</b>	<b>Objectives: Introduction to Transaction Processing Concepts</b>

## CHAPTER 20: Introduction to Transaction Processing Concepts and Theory

**20.16.** Add the operation commit at the end of each of the transactions  $T_1$  and  $T_2$  in Figure 20.2, and then list all possible schedules for the modified transactions. Determine which of the schedules are recoverable, which are cascadeless, and which are strict.

Note: In general, given  $m$  transactions with number of operations  $n_1, n_2, \dots, n_m$ , the number of possible schedules is:  $(n_1 + n_2 + \dots + n_m)! / (n_1! * n_2! * \dots * n_m!)$ , where  $!$  is the factorial function. In our case,  $m = 2$  and  $n_1 = 5$  and  $n_2 = 3$ , so the number of possible schedules is:

$$(5+3)! / (5! * 3!) = 8! / (5! * 3!) = 8 * 7 * 6 * 5 * 4 * 3 * 2 * 1 / (5 * 4 * 3 * 2 * 1 * 3 * 2 * 1) = 56.$$

You don't need to list all 56 possible schedules; only list 2 strict, 2 recoverable, 2 non-recoverable, and 2 cascadeless schedules.

**20.17.** List all possible schedules for transactions  $T_1$  and  $T_2$  in Figure 20.2, and determine which are conflict serializable (correct) and which are not.

**20.23.** Consider the three transactions  $T_1$ ,  $T_2$ , and  $T_3$ , and the schedules  $S_1$  and  $S_2$  given below. Draw the serializability (precedence) graphs for  $S_1$  and  $S_2$ , and state whether each schedule is serializable or not. If a schedule is serializable, write down the equivalent serial schedule(s).

$T_1$ :  $r_1(X)$ ;  $r_1(Z)$ ;  $w_1(X)$ ;

$T_2$ :  $r_2(Z)$ ;  $r_2(Y)$ ;  $w_2(Z)$ ;  $w_2(Y)$ ;

$T_3$ :  $r_3(X)$ ;  $r_3(Y)$ ;  $w_3(Y)$ ;

$S_1$ :  $r_1(X)$ ;  $r_2(Z)$ ;  $r_1(Z)$ ;  $r_3(X)$ ;  $r_3(Y)$ ;  $w_1(X)$ ;  $w_3(Y)$ ;  $r_2(Y)$ ;  $w_2(Z)$ ;  $w_2(Y)$ ;

$S_2$ :  $r_1(X)$ ;  $r_2(Z)$ ;  $r_3(X)$ ;  $r_1(Z)$ ;  $r_2(Y)$ ;  $r_3(Y)$ ;  $w_1(X)$ ;  $w_2(Z)$ ;  $w_3(Y)$ ;  $w_2(Y)$ ;

**20.24.** Consider schedules  $S_3$ ,  $S_4$ , and  $S_5$  below. Determine whether each schedule is strict, cascadeless, recoverable, or nonrecoverable. (Determine the strictest recoverability condition that each schedule satisfies.)

$S_3$ :  $r_1(X)$ ;  $r_2(Z)$ ;  $r_1(Z)$ ;  $r_3(X)$ ;  $r_3(Y)$ ;  $w_1(X)$ ;  $c_1$ ;  $w_3(Y)$ ;  $c_3$ ;  $r_2(Y)$ ;  $w_2(Z)$ ;  $w_2(Y)$ ;  $c_2$ ;

$S_4$ :  $r_1(X)$ ;  $r_2(Z)$ ;  $r_1(Z)$ ;  $r_3(X)$ ;  $r_3(Y)$ ;  $w_1(X)$ ;  $w_3(Y)$ ;  $r_2(Y)$ ;  $w_2(Z)$ ;  $w_2(Y)$ ;  $c_1$ ;  $c_2$ ;  $c_3$ ;

$S_5$ :  $r_1(X)$ ;  $r_2(Z)$ ;  $r_3(X)$ ;  $r_1(Z)$ ;  $r_2(Y)$ ;  $r_3(Y)$ ;  $w_1(X)$ ;  $c_1$ ;  $w_2(Z)$ ;  $w_3(Y)$ ;  $w_2(Y)$ ;  $c_3$ ;  $c_2$ ;