Discrete Mathematics for Computing



Ch 1.5 Nested Quantifiers

- Two Quantifiers P and Q are nested
 - One is within the scope of the other
 - Propositional Function
 - Math and Computer Science



$$\forall x Q(x)$$

where Q(x) is $\exists y P(x, y)$ where P(x,y) is x + y = 0



 Commutative law for Addition of real numbers

$$\forall x \forall y (x + y = y + x)$$

For all real numbers x and y,

$$x + y = y + x$$

Associative law for Addition of real numbers

$$\forall x \forall y \forall z (x + (y + z) = (x + y) + z)$$

For all real numbers x, y and z,

$$x + (y + z) = (x + y) + z$$

Additive Inverse

$$\forall x \exists y (x + y = 0)$$

For every real number x, there is a real number y

$$x + y = 0$$

Translate into English the statement

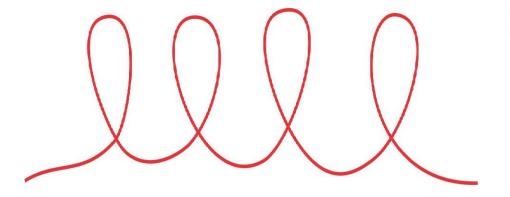
$$\forall x \forall y ((x > 0) \land (y < 0) \rightarrow (xy < 0))$$

- For real numbers x and y, if x is positive and y is negative, then xy is negative.
- The product of a positive real number and a negative real number is always a negative real number.



- Mathematical Statements
- Multiple Quantifications of Propositional Functions
 - Involve more than one variable
 - Order of Quantifiers is relevant





```
for ( i = 0; i < rows; i++) {
    for ( j = 0; j < columns; j++) {
        System.out.print(aryNumbers[i][j] + " ");
    }
    System.out.println( "" );
}</pre>
```



- Quantification as Loops
 - Nested Loops

$$\forall x \forall y P(x, y)$$

- Loop through the values of x
- For each value of x, loop through the values of y.



 $\exists x \forall y P(x,y)$

- We loop through values of x
- Until we find an x for which P(x,y) is always T
- when we loop through all values of y
- If we never hit such an x, F



```
\forall x \forall y P(x, y)
```

When true?

P(x,y) is true for every pair x,y.

When false?

There is a pair x, y for which P(x,y) is false.

 $\forall x \exists y P(x, y)$

When true?

For every x there is a y for which P(x,y) is true.

When false?

There is an x such that P(x,y) is false for every y.



$$\exists x \forall y P(x, y)$$

When true?

There is an x for which P(x,y) is true for every y.

When false?

For every x there is a y for which P(x,y) is false.

$$\exists x \exists y P(x, y)$$

When true?

There is a pair x, y for which P(x,y) is true.

When false?

P(x,y) is false for every pair x, y.



■ Example: Let P(x,y) be the statement x+y=y+x. What are the truth values of the quantification $\forall x \forall y P(x,y)$ where the domain for all variables consists of all real numbers?

The quantification $\forall x \forall y P(x, y)$ denotes the proposition "For all real numbers x, for all real numbers y, x + y = y + x." P(x,y) is T for all real numbers x and y,

$$\Rightarrow \forall x \forall y P(x, y)$$
is **T**



■ Example: Let Q(x,y) be the statement x + y = 0. What are the truth values of the quantification $\exists y \forall x Q(x, y)$ where the domain for all variables consists of all real numbers?

The quantification $\exists y \forall x Q(x, y)$ denotes the proposition "There is a real number y such that for every real number x, Q(x,y)."

$$\Rightarrow \exists y \forall x Q(x, y)$$
 is **F**



■ Example: Let Q(x,y) be the statement x + y = 0. What are the truth values of the quantification $\forall x \exists y Q(x, y)$ where the domain for all variables consists of all real numbers?

The quantification $\forall x \exists y Q(x, y)$ denotes the proposition "For every real number x, there is a real number y such that Q(x,y)."

 $\Rightarrow \forall x \exists y Q(x, y) \text{ is } \mathsf{T}$

ORDER OF QUANTIFIERS IMPORTANT!



Example: Translate the statement into logical expressions.

"The sum of two positive integers is always positive."

Rewrite: "For every two integers, if these integers are positive, then the sum of these integers is positive."

Introduce variables x & y: "For all positive integers x and y, x + y is positive."

 $\forall x \forall y ((x>0) \land (y>0) \rightarrow (x+y>0))$, domain all integers. $\forall x \forall y (x+y>0)$, domain all positive integers.

Example: Translate the statement into logical expressions.

"Every real number except zero has a multiplicative inverse."

Translate to a logical expression: "For every real number x, if $x \ne 0$, then there is a real number y such that xy = 1."

 $\forall x ((x \neq 0) \rightarrow \exists y (xy = 1)), domain real numbers$

 Example: Translate the statement into logical expressions using nested quantifiers.

"There is a woman who has taken a flight on every airline in the world."

Let P(w,f) be "w has taken f"

Let Q(f,a) be "f is a flight on a"

$$\exists w \forall a \exists f (P(w, f) \land Q(f, a))$$

Domain for 'w' consists of all the women in the world, domain for 'f' all the airplane flights, domain for 'a' all airlines.



■ Example: What is the negation of the following statement $\forall x \exists y (x = -y)$

```
\neg \forall x \ P(x) \quad \text{where } P(x) = \exists y \ (x = -y)
\equiv \exists x \ \neg P(x)
\equiv \exists x \ \neg \exists y (x = -y)
\equiv \exists x \ \forall y \ \neg (x = -y)
\equiv \exists x \ \forall y \ (x \neq -y)
```

Negating Nested Quantifiers

Example: Use quantifiers to express the statement that "There does not exist a woman who has taken a flight on every airline in the world."

Use De Morgan's laws for Quantifiers, move the negation

$$\neg \exists w \forall a \exists f (P(w, f) \land Q(f, a))$$

$$\equiv \forall w \neg \forall a \exists f (P(w, f) \land Q(f, a))$$

$$\equiv \forall w \exists a \neg \exists f (P(w, f) \land Q(f, a))$$

$$\equiv \forall w \exists a \forall f \neg (P(w, f) \land Q(f, a))$$

$$\equiv \forall w \exists a \forall f \neg (P(w, f) \lor \neg Q(f, a))$$

For every woman there is an airline such that for all flights, this woman has not taken that flight or that flight is not on this airline.



Practical Example

- Prolog Based on first-order predicate logic
- Developed in 1970 by Colmerauer & Roussel (Marseilles) and Kowalski (Edinburgh) + others.
- Used in Artificial Intelligence, databases, expert systems.
- Program = a bunch of axioms
- Run your program by:
 - Enter a series of facts and declarations
 - Pose a query
 - System tries to prove your query by finding a series of inference steps

