Midterm CS361 Spring 2017

Name: <u>Alexander Moladyh</u>

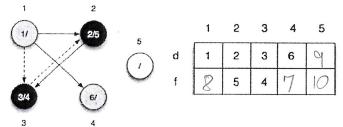
- 1. True or False (2 points each, best 15, answer all)
 - F 1. Quicksort is an excellent example for greedy algorithms and has a best time of O(n log n).
 - f 2. By using Merge Sort, we can sort n^2 elements in $O(\sqrt{nlog}n)$ time.
- \mathcal{T}_3 . If $f(n) = O(n^6)$, then $f(n) = O(n^6 \log n)$.
- 4. The Master Theorem has three cases and it covers all possible recurrence situations.
- \mathcal{T} 5. If $f(n) = \mathbf{O}(n^6)$ and $f(n) = \mathbf{\Omega}(n^6)$ then $f(n) = \Theta(n^6)$.
- **T** <u>F</u> 6. An undirected, connected, and weighted graph may have more than one Minimum Spanning Tree (MST).
 - F 7. Adjacency Matrix cannot be used to represent a DAG.
 - 8. Depth first search and breadth first search generally result in the same "Search tree."
- 9. Like Merge sort, topological sort can sort melements in O(n log n) time.
- 10. Problems that can be solved by Dynamic Programming can, generally speaking, also be solved with a Greedy algorithm.
- 11. Kruskal and Prim algorithm in finding a MST are examples of Greedy algorithms.
- 12. If a connected graph G=(V,E) that satisfying the condition |V| = |E| 1, then it is always a tree.
- T13. For a connected undirected graph G=(V,E), its Breadth-First Tree has |V| -1 edges.
- 14. According to Parenthesis Theorem of DFS, two vertices do not have to be in the same tree.
- 15. Topologic sort can only be performed on DAGs, and the results on the same graph can be different.
- 16. The 0-1 Knapsack Problem has a dynamic programming solution as well as a greedy algorithm solution.

II. Filling in the blank + Multiple choice, circle the best answer on the cases and filling in the complexity blank for each problem (3 points each, best five, answer all)

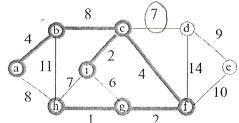
1.	T(n) = 27T(n/3) + n =	2.	$T(n) = 3T(n/3) + n^2 =$	3. $T(n) = 3T(n/3) + 10n =$
	$\Theta(\cancel{n}^{\leq})$. In addition, we		$\Theta(\mathcal{M})$. In addition, we	$\mathbb{Q} \wedge \Theta(\mathcal{D})$. In addition, we
	achieve this by using Master		achieve this by using Master	achieve this by using Master
	Theorem's case?		Theorem's case?	Theorem's case?
	(a. 1) b. 2 c. 3		a. 1 b. 2 (c. 3)	a. 1 (b. 2 (c. 3)
4.	$T(n) = 4T(n/2) + n^2 =$	5.	T(n) = 2T(n/4) + Ign =	6. $T(n) = 8T(n/3) + lgn = \sqrt{93}$
	Θ(1) In addition, we achieve this by using Master		$\Theta(\mathcal{N} \setminus \mathbf{n})$ In addition, we	$\Theta(\mathcal{N}^{\mathcal{S}})$. In addition, we
			achieve this by using Master	achieve this by using Master
	Theorem's case?		Theorem's case?	Theorem's case?
	a. 1 (b. 2) c. 3		a. 1 b. 2 (c. 3)	(a. 1) b. 2 c. 3

III. Finishing the examples (3 points each).

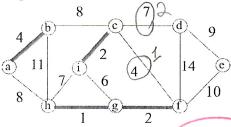
a. For the DFS example below, finish every vertices' starting and ending timestamps if they are still not done as rows d and f in the array.



b. With Prim's algorithm for MST, which edge will be included in the MST? Clearly mark the edge by circling it.



c. Which 2 edges will be included with Kruskal's algorithm for MST? Clearly mark the edges by circling them and make them with I and II.



d. For the BFS example below, which vertex will be discovered next? Clearly mark the vertex by circling it.

