

## Quiz #4 CS361 Spring 2017

Name: Alexander Molodyh

### I. True or False (1 Point each for the best 8, answer all)

- ☐ 1. The set of states  $Q$  in a DFA can be empty. When  $Q$  is not empty, the DFA can have more than one start state.
- ☐ 2. Let's  $A$ ,  $p$ , and  $s$  be defined as in the Pumping Lemma, if  $s = xyz$ , then  $s' = xz$  must also be in  $A$ . In addition, the Pumping Lemma for NFA is mostly used to prove that a language is actually regular
- ☐ 3. Any language accepted by a NFA can also be accepted by a DFA.
- ☐ 4. The class of regular language is closed under the concatenation operation.
- ☐ 5. A NFA is considered as a special case of a DFA.
- ☐ 6. If machine  $M$  has two input symbols that transit to  $q_2$  from  $q_1$ , then  $M$  cannot be a DFA.
- ☐ 7. Every CFG has an equivalent NFA, and every DFA has an equivalent CFG.
- ☐ 8. If a DFA has  $n$  states, it cannot accept strings with a length greater than  $n^2$ .
- ☐ 9.  $A = \{0^n 1^n\}$  can be accepted by a PDA.

### II. Short answers (2 points each)

1. Show your state diagram of a NFA or DFA that accepts strings containing 1100, such as 00**11**001.

2. Let  $A = \{a, b\}$  and  $B = \{0, 1\}$  show  $A \cup B$  and  $A \circ B$ . Then explain what is  $A^*$  and show all strings in  $A^*$  that have a length smaller than 3.

### II. Short answers (4 points)

1. Design a DFA or NFA that accepts regular expression  $(01U10)^*$ .
2. Design a PDA that accepts  $E = \{0^i 1^j \mid i > j\}$ . You may have to use the back page.