## From grammar to shared packed parse forest

Semyon Grigorev
Saint Petersburg State University
7/9 Universitetskaya nab.
St. Petersburg, 199034 Russia
semen.grigorev@jetbrains.com

It is possible to build graph structured representation of result which can be interesting for several reasons.

- More user friendly query result represenation.
- Useful for query result exploration and investigation.
- Useful for query debugging.

Example from your article [2] will be used for explanation. Let we intrioduce new notation for nonterminals and terminals for simplification.

$$q_1 = q[A, B]$$
  $q_5 = q[C, E]$   
 $q_2 = q[A, C]$   $q_6 = q[A, D]$   
 $q_3 = q[D, E]$   $q_7 = q[A, E]$   
 $q_4 = q[B, D]$   $q_8 = q[B, E]$ 

$$T_1 = friendOf_{A,B}$$
  $T_4 = friendOf_{B,D}$   
 $T_2 = friendOf_{A,C}$   $T_5 = friendOf_{C,E}$   
 $T_3 = friendOf_{D,E}$ 

Now grammar from [2](pages 6-7) can be represented as below. Let it be named  $G_1$ 

$$\begin{array}{lll} q_1 \to T_1 & q_6 \to q_1 \ q_4 \\ q_2 \to T_2 & q_7 \to q_2 \ q_5 \\ q_3 \to T_3 & q_7 \to q_1 \ q_8 \\ q_4 \to T_4 & q_7 \to q_6 \ q_3 \\ & q_8 \to q_4 \ q_3 \end{array}$$

For each context-free grammar we can create regular tree grammar [1](section 2.4). Moreover if G is a context-free word grammar, then the set of derivation trees of L(G) is a regular tree language. By using algorithm presented in [1](section 2.4) we can build regular tree grammar  $G_2$  for grammar  $G_1$ :

$$\begin{array}{ll} q_1 \to q_1(T_1) & q_6 \to q_6(q_1, q_4) \\ q_2 \to q_2(T_2) & q_7 \to q_7(q_2, q_5) \\ q_3 \to q_3(T_3) & q_7 \to q_7(q_1, q_8) \\ q_4 \to q_4(T_4) & q_7 \to q_7(q_6, q_3) \\ & q_8 \to q_8(q_4, q_3) \end{array}$$

Reduced graph. It is SPPF.

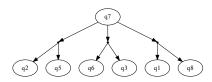


Figure 1: Input graph M

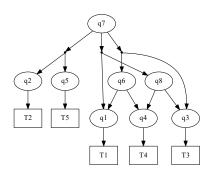


Figure 2: Input graph M

Let return information about edges back.

It is possible to build SPPF with our tool without to CNF transformation.

Query grammar.

 ${\bf Input\ graph.}$ 

Query result

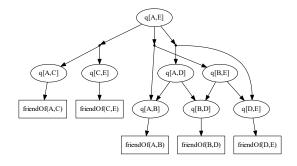


Figure 3: SPPF with information about paths

0: s = L s R
1: s = middle
2: middle = L R

Figure 4: Grammar  $G_1$  for language  $L = \{L^n R^n; n \ge 1\}$ 

## 1. REFERENCES

- [1] Comon, Hubert, et al. "Tree automata techniques and applications." (2007).
- [2] Hellings, Jelle. "Querying for Paths in Graphs using Context-Free Path Queries." arXiv preprint arXiv:1502.02242 (2015).

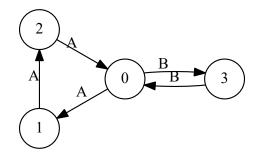


Figure 5: SPPF with information about paths

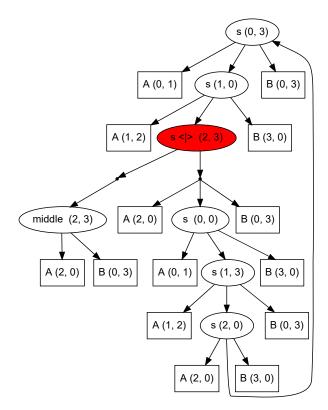


Figure 6: SPPF with information about paths