

Generalized LL parsing for context-free constrained path search problem

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ABSTRACT

Graph data model and graph databases are very popular in many different areas such as bioinformatics, semantic web, social networks, etc. One of specific problems of graph databases is a path querying with constraints formulated in terms of formal grammars. There are several solutions, but providing structural representation of query result which is practical for answer understanding and exploration is still a problem. In this paper we propose graph parsing technique which allows us to build such representation with respect to given grammar query for arbitrary context-free grammar and graph. Proposed algorithm is based on generalized top-down parsing algorithm which has cubic worst-case time complexity and linear for LL grammars.

1. INTRODUCTION

Graph data model and graph data bases are very popular in many different areas such as bioinformatics, semantic web, social networks, etc. Extraction of paths which satisfy specific constraints may be useful for investigation of graph structured data and for detection of relations between data items. One specific problem — path querying with constraints — is usually formulated in terms of formal grammars and is called formal language constrained path problem [3].

Classical parsing techniques can be used to solve formal language constrained path problem. It means that such technique can be used for more common problem — “graph parsing”. Graph parsing may be required in graph data base querying, formal verification, string-embedded language processing and another areas where graph structured data is used.

Existing solutions in databases field usually use such parsing algorithms as CYK or Earley (for example [5], [16]). These algorithms have nonlinear time complexity for unambiguous grammars ($O(n^3)$ and $O(n^2)$ respectively). There are such parsing algorithms as GLR and GLL which have cubic worst-case complexity and linear for unambiguous grammars. Also there is no need to transform the input grammar to Chomsky normal form which required for CYK. This fact allows us to avoid grammar size increasing, and hence to improve performance of parsing because it is sensitive to grammar size.

Despite the fact that there is a set of path querying solutions [16, 5, ?], query result exploration is still a challenge [6]. Simplification of complex query debugging is also a problem. Structural representation of query result can be used to solve these problems, and classical parsing techniques pro-

vide such representation — derivation tree — which contains full information about parsed sentence structure in terms of specified grammar.

Graph parsing can also be used to analyze dynamically generated strings or string-embedded languages. Regular approximation for value set of string variable can be represented as a finite automata. Moreover, if we want to check a syntactic correctness of dynamically generated strings, we should check that all generated strings (all paths from start states to final states in the given automata) are correct with respect to given context-free grammar. For example, the grammar of SQL dialect can be used for processing string-embedded SQL. There are some solutions that are addressed to this problem: GLR-based checker of string-embedded SQL queries [2, 4]; relaxed parser of string-embedded languages [20] based on RNGLR parsing algorithm. RNGLR-based algorithm allows to construct derivation forest for all correct paths in the input automata.

In this paper we propose graph parsing technique which allows to construct structural representation of query result with respect to given grammar query. This structure can be useful for query debugging and exploration. Proposed algorithm is based on generalized top-down parsing algorithm — GLL [12] — which has cubic worst-case time complexity and linear for LL grammars on linear input.

2. PRELIMINARIES

In this work we are focused on the parsing algorithm, and not on the data representation, and we assume that the whole input graph can be located in RAM memory in the optimal for our algorithm way.

Also we need to introduce some definitions.

- Context-free grammar $G = (N, \Sigma, P, S)$ where N is a set of nonterminal symbols, Σ is a set of terminal symbols, $S \in N$ is a start nonterminal, and P is a set of productions.
- $\mathcal{L}(G)$ is a language specified by grammar G .
- Directed graph $M = (V, E, L)$ where V is a set of vertices, $L \subseteq \Sigma$ is a set of edge's labels, and $E \subseteq V \times L \times V$. We assume that there are no parallel edges with equal labels: for every $e_1 = (v_1, l_1, v_2) \in E$, $e_2 = (u_1, l_2, u_2) \in E$ if $v_1 = u_1$ and $v_2 = u_2$ then $l_1 \neq l_2$.
- $tag : E \rightarrow L$ is a helper function which allows to get edge's tag.

$$tag(e = (v_1, l, v_2), e \in E) = l$$

- $\oplus : L^+ \times L^+ \rightarrow L^+$ is a concatenation operation.

- Path p in graph M is a list of edges:

$$p = (v_0, l_0, v_1), (v_1, l_1, v_2), \dots, (v_{n-1}, l_{n-1}, v_n) \\ = e_0, e_1, \dots, e_{n-1}$$

where $v_i \in V, e_i \in E, e_i = (v_i, l_i, v_{i+1}), l_i \in L, |p| = n, n \geq 1$.

- Set of paths $P = \{p : p \text{ path in } M\}$ where M is a directed graph.
- $\Omega : p \rightarrow L^+$ is a helper function which allows to get a string produced by the given path.

$$\Omega(p = e_0, e_1, \dots, e_{n-1}, p \in P) = \\ tag(e_0) \oplus \dots \oplus tag(e_{n-1}).$$

As a result, we can state that context-free language constrained path querying means that we get query as grammar G and the result of this query is a set of paths

$$P = \{p | \Omega(p) \in \mathcal{L}(G)\}.$$

Note that P can be an infinite set in some cases, and hence it cannot be explicitly represented. In order to solve this problem, in this paper, we will construct compact data structure representation which stores all elements of P in finite space and allows to extract any of them.

3. MOTIVATING EXAMPLE

In this article we discuss context-free constrained path querying, and one of well-known not regular but context-free language is an language

$$\mathcal{L} = \{A^n B^n; n \geq 1\} = \{AB; AAB; AAAB; \dots\}$$

. This language is a subset of balanced brackets language and in practice may be used to describe many different relations: a n -th generation in parent-child, a correct order of an open-close operations for resources, etc.

Let's a graph $M = (\{0; 1; 2; 3\}, E, \{A; B\})$ is presented in figure 1. Labels on edges represent one of mentioned relations. We want to find all paths p , such that $\Omega(p) \in \{AB; AAB; AAAB; \dots\}$ or $\Omega(p) \in A^n B^n$ where $n \geq 1$. Required language can be specified by grammar G_1 presented in picture 2 where $N = \{s; middle\}$, $\Sigma = \{A; B\}$, and $S = s$.

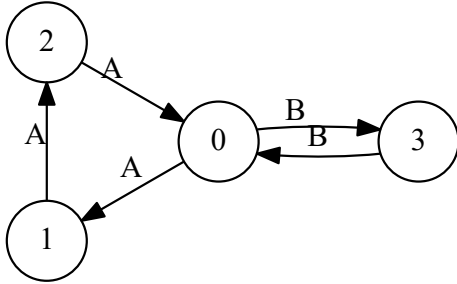


Figure 1: Input graph M

The result of presented query for given graph is an infinite set of paths, hence it cannot be constructed explicitly. Moreover, for some tasks it can be necessary to get the structure

```
0: s = A s B
1: s = middle
2: middle = A B
```

Figure 2: Grammar G_1 for language $L = \{A^n B^n; n \geq 1\}$

of result with respect to given grammar. Later we show how to get finite representation of query result structure in terms of derivation in grammar.

4. GRAPH PARSING ALGORITHM

We propose graph parsing algorithm which allows to create finite representation of parse forest which contains trees for all satisfied paths in graph. Finite representation of result set with structure related to specified grammar may be useful not only for results understanding and processing but also for query debugging especially for complex queries.

Therefore our solution is based on generalized LL (GLL) [12, 1] parsing algorithm which allows to process arbitrary (including left-recursive and ambiguous) context-free grammars with worst-case cubic time complexity and linear for LL grammars.

4.1 Generalized LL Parsing Algorithm

In classical LL algorithm we have a pointer to input (position i) and a pointer to grammar in form $n \rightarrow \alpha \cdot x \beta$ — grammar slot. The parsing may be described as a movement of this pointers from initial position: $i = 0, s \rightarrow \cdot \beta$, where s is start nonterminal. Thus in each step we have two pointers and some possible cases to process them.

1. $n \rightarrow \alpha \cdot x \beta$ when x is a terminal and $x = input[i]$. In this case we should move both pointers right: $i = i + 1$, new slot = $n \rightarrow \alpha x \cdot \beta$.
2. $n \rightarrow \alpha \cdot x \beta$ when x is nonterminal. In this case we should save return address $n \rightarrow \alpha x \cdot \beta$ in stack and move pointer in grammar to position $x \rightarrow \cdot \gamma$.
3. $n \rightarrow \alpha \cdot$. This case means that we finish processing of nonterminal n . We should pop return address from stack and use it as new slot.
4. $s \rightarrow \alpha \cdot$ where s is a start nonterminal of grammar. In this case we should check emptiness of input and report success or failure.

In case 2 we can use *FIRST* set to choose single variant. But sometimes it is not possible to select only one path to continue parsing and it does not allow to use LL parsing algorithm. Generalized LL algorithm handles all possible paths in this case. Instead of immediate processing of all variants, GLL uses descriptors mechanism to store all possible branches and process them sequentially. Descriptor is a quadruple (L, s, j, a) where L is a grammar slot, s is a stack node, j is a position in the input, and a is a node of derivation tree.

The stack in parsing process is used to store return information for the parser — a name of function which would be called when current function will finish computation. As mentioned before, generalized parsers process all possible derivation branches and for every branch parser must store its own stack. It leads to infinite stack growth. Tomita-style

graph structured stack (GSS) [18] allows to combine stacks to solve this problem. In GLL each GSS node contains a pair of position in input and grammar slot.

In order to provide termination and correctness we should avoid duplication of descriptors, and allow to process GSS nodes in arbitrary order. It is necessary to use some additional sets for this.

- R — working set which contains descriptors to process. Algorithm finish when R is empty.
- U — all descriptors was created. Avoid duplication of these.
- P — popped nodes. Allows to process descriptors (and GSS nodes) in arbitrary order without losing new descriptors in step 3.

Instead of explicit code generation used in classical algorithm, we use table version of GLL [?] in order to simplify adaptation to graph processing. As a result, main control function is different from the original one because it should process LL-like table instead of switching between generated parsing functions. Control functions of the table based GLL are presented in listing 1. All other functions are the same as in original algorithm and their description can be found in original article [12] or in Appendix A.

There are more than one tree for ambiguous grammar and generalized algorithms builds all derivation trees. Special data structure — Shared Packed Parse Forest [10] — is used to reduce space required for tree storage.

4.2 Shared packed parse forest

Binarized Shared Packed Parse Forest (SPPF) [15] allows to compress derivation trees with optimal reusing of common nodes and subtrees. Version of GLL which uses this structure for parsing forest representation achieves worst-case cubic space complexity [13].

Let we present an example of SPPF for the input sentence "ABABAB" and ambiguous grammar G_0 (pic 3).

```
0: s = eps
1: s = A s B
2: s = s s
```

Figure 3: Grammar G_0

There are two different leftmost derivations of given sentence in grammar G_0 , hence SPPF should contains two different trees. Result SPPF (fig. 4a) and two trees (fig. 4b and fig. 4c) are presented in figure 4.

Binarised SPPF can be represented as a graph where each node has one of four types which described below with correspondent graphical notation. Let i and j are start and end positions of substring, and tuple of positions (i, j) is an *extension* of node.

- Node with rectangle shape labeled with (i, T, j) is terminal node.
- Node with oval shape labeled with (i, N, j) is nonterminal node. This node denotes that there is at least one derivation for substring α from position i to position j in input string ω such that $N \Rightarrow_G^* \alpha, \alpha = \omega[i..j-1]$. All derivation trees for given substring and nonterminal

Algorithm 1 Control functions of table version of GLL

```
1: function DISPATCHER
2: if  $R.Count \neq 0$  then
3:    $(L, v, i, cN) \leftarrow R.Get()$ 
4:    $cR \leftarrow dummy$ 
5:    $dispatch \leftarrow false$ 
6: else
7:    $stop \leftarrow true$ 
8: function PROCESSING
9:    $dispatch \leftarrow true$ 
10:  switch  $L$  do
11:    case  $(X \rightarrow \alpha \cdot x\beta)$  where  $x = input[i+1]$ 
12:      if  $cN = dummyAST$  then
13:         $cN \leftarrow GETNODET(i)$ 
14:      else
15:         $cR \leftarrow GETNODET(i)$ 
16:         $i \leftarrow i + 1$ 
17:         $L \leftarrow (X \rightarrow \alpha x \cdot \beta)$ 
18:        if  $cR \neq dummy$  then
19:           $cN \leftarrow GETNODEP(L, cN, cR)$ 
20:         $dispatch \leftarrow false$ 
21:    case  $(X \rightarrow \alpha \cdot x\beta)$  where  $x$  is nonterminal
22:       $v \leftarrow CREATE((X \rightarrow \alpha x \cdot \beta), v, i, cN)$ 
23:       $slots \leftarrow pTable[x][input[i]]$ 
24:      for all  $L \in slots$  do
25:         $ADD(L, v, i, dummy)$ 
26:    case  $(X \rightarrow \alpha \cdot)$ 
27:       $POP(v, i, cN)$ 
28:    case  $(S \rightarrow \alpha \cdot)$  when  $S$  is start nonterminal
29:      final result processing and error notification
30: function CONTROL
31: while not  $stop$  do
32:   if  $dispatch$  then
33:     DISPATCHER
34:   else
35:     PROCESSING
```

can be extracted from SPPF by left-to-right top-down graph traversal started from respective node. We use filled nonterminal node labeled with $(\blacktriangleleft (i, N, j))$ for denote that there are more then one derivations from nonterminal N for substring from i to j .

- Packed node with label (i, t, j) where t is a grammar slot. We use dot shape for these nodes and omit label because it is important only for SPPF constriction. Subgraph with root in such node is one variant of derivation in case when parent is nonterminal node with label $(\blacktriangleleft (i, N, j))$.
- Node with rectangle shape and label $(N : \gamma, k)$ is an intermediate node.

Later in our examples we will remove redundant intermediate and packed nodes from the SPPF to simplify it and to decrease the size of structure.

4.3 GLL-based graph parsing

In this section we present such modification of GLL algorithm, that for input graph M , set of start vertices $V_s \subseteq V$, set of final vertices $V_f \subseteq V$, grammar G_1 , it returns SPPF

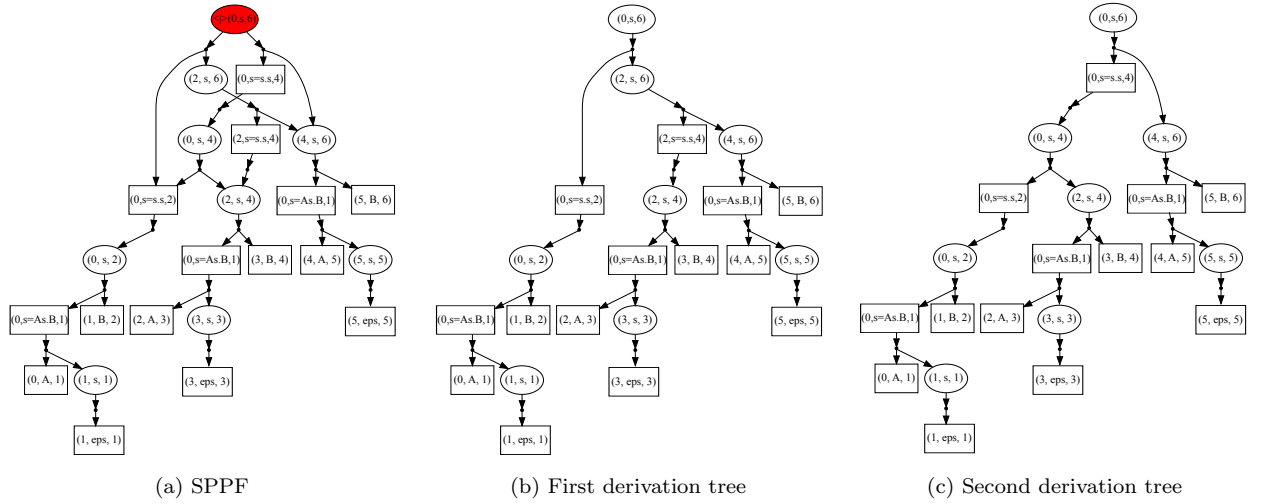


Figure 4: SPPF for sentence "ABABAB" and grammar G_0

which contains all derivation trees for all paths p in M , such that $\Omega(p) \in L(G_1)$, and $p.start \in V_s, p.end \in V_f$.

First of all notice that input string for classical parser is a linear graph, and positions in input is vertices of this graph. This observation can be generalized to arbitrary graph with remark that in the position we have not only one next symbol, but the set of labels of all outgoing edges for given vertex. Thus in order to use GLL for graph parsing we need to use graph vertices as position in input and modify **Processing** function to allow to process more then one "next symbol". Required modifications presented in listing 2. Small modification also are required for initialization of R set: it is necessary to add not only one initial descriptor but the set of descriptors for all vertices in V_s . All other functions can be reused from original algorithm without any changes.

As far as we can specify sets of start and final vertices, our solution can find all paths in graph, all paths from specified vertex, all paths between specified vertices. Also SPPF represents a structure of paths in terms of derivation which allow to get more useful information about result.

Note that termination of proposed algorithm is inherited from basic GLL algorithm. We are working with finite graphs, hence set of positions is finite, and tree construction is not changed. As a result, total number of descriptors is finite, and any of them cannot be added in R twice, hence main loop is finite.

4.4 Complexity

Time complexity estimation in terms of input graph and grammar size is pretty similar to estimation of GLL complexity provided in [13].

LEMMA 1. For any descriptor (L, u, i, w) either $w = \$$ or w has extension (j, i) where u has index j .

PROOF. Proof of this lemma is the same as provided for original GLL in [13] because main function used for descriptors creation was not changed. \square

THEOREM 1. The GSS generated by GLL-based graph parsing algorithm for grammar G on input graph $M = (V, E, L)$ has at most $O(|V|)$ vertices and $O(|V|^2)$ edges.

Algorithm 2 Processing function modified in order to process arbitrary directed graph

```

1: function PROCESSING( )
2:   dispatch  $\leftarrow$  true
3:   switch L do
4:     case  $(X \rightarrow \alpha \cdot x\beta)$  where  $x$  is terminal
5:       for all  $\{e|e \in input.outEdges(i), tag(e) = x\}$  do
6:         new_cN  $\leftarrow$  cN
7:         if new_cN = dummyAST then
8:           new_cN  $\leftarrow$  GETNODET(e)
9:         else
10:          new_cR  $\leftarrow$  GETNODET(e)
11:          L  $\leftarrow$   $(X \rightarrow \alpha x \cdot \beta)$ 
12:          if new_cR  $\neq$  dummy then
13:            new_cN  $\leftarrow$  GETNODEP(L, new_cN, new_cR)
14:            ADD(L, v, target(e), new_cN)
15:     case  $(X \rightarrow \alpha \cdot x\beta)$  where  $x$  is nonterminal
16:       v  $\leftarrow$  CREATE( $(X \rightarrow \alpha x \cdot \beta), v, i, cN$ )
17:       slots  $\leftarrow \bigcup_{e \in input.OutEdges(i)} pTable[x][e.Token]$ 
18:       for all  $L \in slots$  do
19:         ADD(L, v, i, dummy)
20:     case  $(X \rightarrow \alpha \cdot)$ 
21:       POP(v, i, cN)
22:     case _
23:       final result processing and error notification

```

PROOF. Proof is the same as the proof of **Theorem 2** from [13] because structure of GSS was not changed.

\square

THEOREM 2. The SPPF generated by GLL-based graph parsing algorithm on input graph $M = (V, E, L)$ has at most $O(|V|^3 + |E|)$ vertices and edges.

PROOF. Let us estimate the number of nodes of each type.

- **Terminal nodes.** Each of them is labeled with (T, v_0, v_1) , and such label can be created only if there

is such $e \in E$ that $e = (v_0, T, v_1)$. Note, that there are no duplicate edges. Hence there are at most $|E|$ terminal nodes.

- ε **nodes** are labeled with (ε, v, v) , hence there are at most $|V|$ of them.
- **Nonterminal nodes** have labels of form (N, v_0, v_1) , so there are at most $O(|V|^2)$ of them.
- **Indeterminate nodes** have labels of form (t, v_0, v_1) , where t is grammar slot, so there are at most $O(|V|^2)$ of them.
- **Packed nodes** are children of intermediate or nonterminal nodes and have label of form (t, v) where t is a grammar slot $N \rightarrow \alpha \cdot \beta$. There are at most $O(|V|^2)$ parents for packed nodes and each of them can have at most $O(|V|)$ children.

As a result there are at most $O(|V|^3 + |E|)$ nodes in SPPF.

The packed nodes have at most two children so there are at most $O(|V|^3 + |E|)$ edges with source in packed node. Nonterminal and intermediate nodes have at most $O(|V|)$ children and all of them are packed nodes. Thus there are at most $O(|V|^3)$ edges with source in nonterminal or intermediate nodes. As a result there are at most $O(|V|^3 + |E|)$ edges in SPPF.

□

THEOREM 3. *The space complexity of GLL-based graph parsing algorithm for graph $M = (V, E, L)$ is at most $O(|V|^3 + |E|)$.*

PROOF. From theorems 1 and 2 we have that space required for main data structures is at most $O(|V|^3 + |E|)$.

□

THEOREM 4. *The runtime complexity of GLL-based graph parsing algorithm for graph $M = (V, E, L)$ is at most*

$$O\left(|V|^3 * \max_{v \in V} (deg^+(v))\right).$$

PROOF. From Lemma 1 we get that there are at most $O(|V|^2)$ descriptors. Complexity of all functions are the same as in proof of **Theorem 4** from [13] except **Processing** function where we should process not single next input token, but the whole set of outgoing edges. Thus for each descriptor we should examine at most

$$\max_{v \in V} (deg^+(v))$$

edges where $deg^+(v)$ is outdegree of vertex v .

Thus worst-case complexity of proposed algorithm is

$$O\left(V^3 * \max_{v \in V} (deg^+(v))\right).$$

□

From theorem 4 we can get estimations for linear input and for LL grammars: for any $v \in V$ it is true, that $deg^+(v) \leq 1$, so $\max_{v \in V} (deg^+(v)) = 1$ and we get $O(|V|^3)$, as expected. For LL grammars and linear input complexity should be $O(|V|)$ for the same reason as for original GLL.

As discussed in [7] achieving of theoretical complexity required special data structures which can be irrational for practice implementation and it is necessary to find balance between performance, software complexity, and hardware resources. As a result in practice we can get slightly worse performance than theoretical estimation.

Note that result SPPF contains only paths matched specified query, so result SPPF size is $O(|V'|^3 + |E'|)$ where $M' = (V', E', L')$ is a subgraph of input graph M which contains only matched paths. Also note that each specific path can be explored with linear SPPF traversal.

4.5 Example

Let us present a solution of motivating example: grammar G_1 is a query and we want to find all paths in graph M (presented in picture 1) matching this query. Result SPPF for this input is presented in figure 5. Note that presented version does not contains redundant nodes. Each terminal node corresponds to the edge in the input graph: for each node with label (v_0, T, v_1) there is $e \in E : e = (v_0, T, v_1)$. We duplicate terminal nodes only for figure simplification.

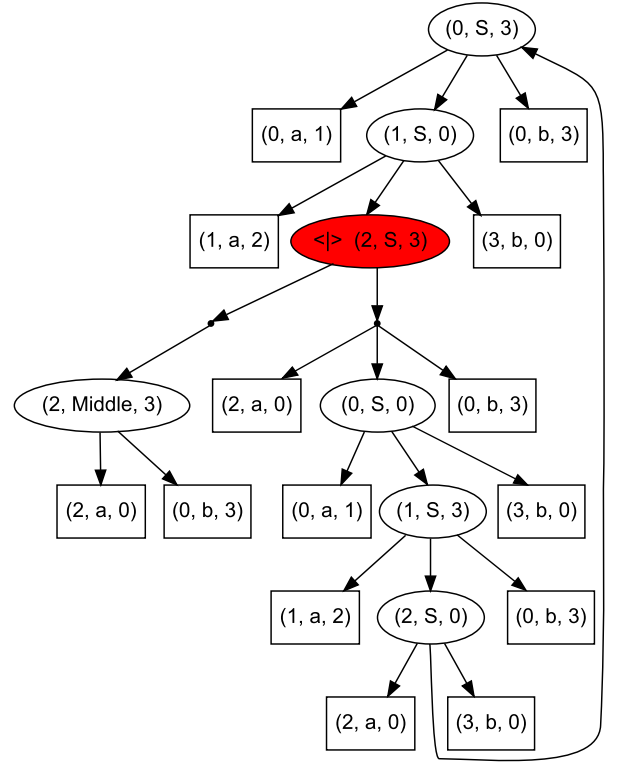


Figure 5: Result SPPF for input graph M (fig. 1) and query G_1 (fig. 2)

As an example of derivation structure usage we can find 'middle' of any path in example above simply by finding correspondent nonterminal *Middle* in SPPF. So we can find out that there is only one common ancestor for all results, and it is a vertex with $id = 0$.

Extensions stored in nodes allow us to check whether path from u to v exists, and to extract it. To extract specified path we need only to traverse SPPF, and it can be done in linear time (in terms of SPPF size).

Lets find paths satisfying specified in G_1 constraints from vertex 0. To do this, we should find vertices with label $(0, S, _)$ in SPPF. We can see that there are two vertices with label matched this pattern: $(0, S, 0)$ and $(0, S, 3)$. At the next step let us to extract corresponded paths from SPPF. In our example there is a cycle in SPPF so there are **at least** two different paths:

$$p_0 = \{(0, a, 1); (1, a, 2); (2, a, 0); (0, b, 3); (3, b, 0); (0, b, 3)\}$$

and

$$p_1 = \{(0, a, 1); (1, a, 2); (2, a, 0); (0, a, 1); (1, a, 2); (2, a, 0); (0, b, 3); (3, b, 0); (0, b, 3); (3, b, 0); (0, b, 3); (3, b, 0)\}.$$

Thus SPPF which was constructed by described algorithm can be useful for query result investigation. But in some cases explicit representation of matched subgraph is preferable, and required subgraph that may be extracted from SPPF trivially by its traversal.

5. EVALUATION

In this section we show that performance of implemented algorithm is in good agreement with theoretical estimations, and that the worst case of time and space complexity can be achieved.

We use two grammars for balanced brackets — ambiguous grammar G_0 (fig. 3) and unambiguous grammar G_2 (fig. 6) — in order to investigate performance and grammar ambiguity correlation.

$$\begin{aligned} 0 : S &\rightarrow a S b S \\ 1 : S &\rightarrow \varepsilon \end{aligned}$$

Figure 6: Unambiguous grammar G_2 for balanced brackets

For input we use complete graphs where for each terminal symbol there is an edge between every two vertices labeled with it. Note that we use only terminal symbols for edges labels. The task we solve in our experiments is to find all paths from all vertices to all vertexes satisfied specified query. Such designed input looks hard for querying in terms of required resources because there are correct path between any two vertices and result set is infinite.

For complete graph $M = (V, E, L)$ we get

$$\max_{v \in V} (deg^+(v)) = (|V| - 1) * |\Sigma|$$

where Σ is terminals of input grammar, hence we should get time complexity at most $O(|V|^4)$ and space complexity at most $O(|V|^3)$.

All tests were performed on a PC with following characteristics:

- OS Name: Microsoft Windows 10 Pro
- System Type: x64-based PC
- CPU: Intel(R) Core(TM) i7-4790 CPU @ 3.60GHz, 3601 Mhz, 4 Core(s), 4 Logical Processor(s)
- RAM: 32 GB

Performance measurement results are presented in figure 7. For time measurement results we have that all two curves can be fit with polynomial function of degree 4 to a high level of confidence with R^2 .

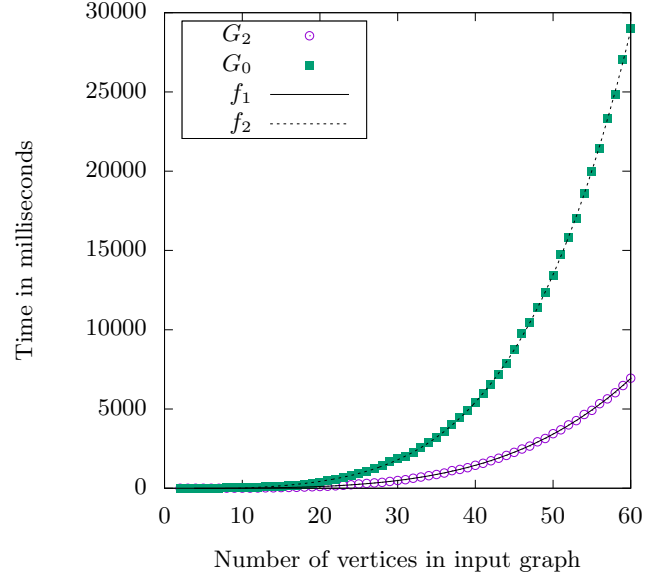


Figure 7: Performance on complete graphs for grammar G_0 and G_2

$$f_1(x) = 0.000495989 * x^4 + 0.001252184 * x^3 + 0.068491746 * x^2 - 0.306749160 * x; R^2 = 0.99996$$

$$f_2(x) = 0.003368883 * x^4 - 0.114919298 * x^3 + 3.161793404 * x^2 - 22.549491142 * x; R^2 = 0.99995$$

Also we present SPPF size in terms of nodes for both G_0 and G_2 grammars 8. As we expected, all two curves are cubic to a high level of confidence with $R^2 = 1$.

6. CONCLUSION AND FUTURE WORK

We propose GLL-based algorithm for context-free path querying which construct finite structural representation of all paths satisfying given constraint. Provided data structure can be useful for result investigation and processing, and for query debugging. Presented algorithm has been implemented in F# [17] and is available on GitHub:<https://github.com/YaccConstructor/YaccConstructor>.

In order to estimate practical value of proposed algorithm we should perform evaluation on a real dataset and real queries. One possible application of our algorithm is metagenomical assembly querying, and we are currently working on this topic.

We are also working on performance improvement by implementation of recently proposed modifications in original GLL algorithm [14, 1]. One direction of our research is generalization of grammar factorization proposed in [14] which may be useful for the processing of regular queries which are common in real world application.

7. REFERENCES

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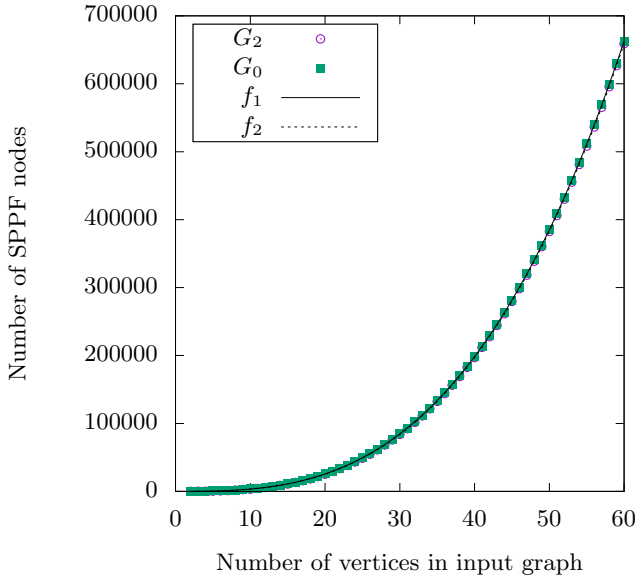


Figure 8: SPPF size on complete graph for grammars G_0 and G_2 on complete graphs

$$f_1(x) = 3.000047 * x^3 + 3.994579 * x^2 + 4.191568 * x; R^2 = 1$$

$$f_2(x) = 3.000050 * x^3 + 2.994338 * x^2 + 4.196472 * x; R^2 = 1$$

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APPENDIX

A. GLL PSEUDOCODE

Main functions of GLL parsing algorithms:

- stack and descriptors manipulation functions 3;
- SPPF construction functions 4.
- R — working set which contains descriptors to process.
- U — all descriptors was created.
- P — popped nodes.

Algorithm 3 Stack and descriptors manipulation functions

```

1: function ADD( $L, v, i, a$ )
2:   if  $(L, v, i, a) \notin U$  then
3:      $U.add(L, v, i, a)$ 
4:      $R.add(L, v, i, a)$ 
5: function POP( $v, i, z$ )
6:   if  $v \neq v_0$  then
7:      $P.add(v, z)$ 
8:     for all  $(a, u) \in v.outEdges$  do
9:        $y \leftarrow GETNODEP(v.L, a, z)$ 
10:       $ADD(v.L, u, i, y)$ 
11: function CREATE( $L, v, i, a$ )
12:   if  $(L, i) \notin GSS.nodes$  then
13:      $GSS.nodes.add(L, i)$ 
14:    $u \leftarrow GSS.NODES.GET(L, i)$ 
15:   if  $(u, a, v) \notin GSS.edges$  then
16:      $GSS.edges.add(u, a, v)$ 
17:   for all  $(u, z) \in P$  do
18:      $y \leftarrow GETNODEP(L, a, z)$ 
19:      $(\_, \_, k) \leftarrow z.lbl$ 
20:      $ADD(L, v, k, y)$ 
return  $u$ 

```

Algorithm 4 SPPF construction functions

```

1: function GETNODET( $x, i$ )
2:   if  $x = \varepsilon$  then
3:      $h \leftarrow i$ 
4:   else
5:      $h \leftarrow i + 1$ 
6:   if  $(x, i, h) \notin SPPF.nodes$  then
7:      $SPPF.nodes.add(x, i, h)$ 
8:     return  $SPPF.nodes.get(x, i, h)$ 
9: function GETNODEP( $(X \rightarrow \omega_1 \cdot \omega_2), a, z$ )
10:   if  $\omega_1$  is terminal or non-nullable nonterminal and  $\omega_2 \neq \varepsilon$  then return  $z$ 
11:   else
12:     if  $\omega_2 = \varepsilon$  then
13:        $t \leftarrow X$ 
14:     else
15:        $h \leftarrow (\rightarrow \omega_1 \cdot \omega_2)$ 
16:      $(q, k, i) \leftarrow z.lbl$ 
17:     if  $a \neq dummy$  then
18:        $(s, j, k) \leftarrow a.lbl$ 
19:        $y \leftarrow findOrCreate SPPF.nodes (n.lbl = (t, i, j))$ 
20:       if  $y$  does not have a child with label  $(X \rightarrow \omega_1 \cdot \omega_2)$  then
21:          $y' \leftarrow newPackedNode(a, z)$ 
22:          $y.chld.add y'$ 
23:         return  $y$ 
24:       else
25:          $y \leftarrow findOrCreate SPPF.nodes (n.lbl = (t, k, i))$ 
26:         if  $y$  does not have a child with label  $(X \rightarrow \omega_1 \cdot \omega_2)$  then
27:            $y' \leftarrow newPackedNode(z)$ 
28:            $y.chld.add y'$ 
29:           return  $y$ 
30:       return  $SPPF.nodes.get(x, i, h)$ 

```
