Connected components

Смирнов Кирилл 371

Определение:

Две вершины u и v называются **связаными** (англ. adjacent), если в графе G существует путь из u в v (обозначение: $u \leadsto v$).

Определение:

Компонентой связности (англ. connected component) называется класс эквивалентности относительно связности.

Пример использования

1. Кластеризация

a. группировка пикселей изображения в группы (aka Pixel connectivity)

2. Social networks

а. Изучения свойств общества посредством представления его в виде графа (aka Social network analysis)

3. Pipeline

а. Определение компонент связности позволяет выстроить конвейеры, в который каждая компонента обрабатывается независимо

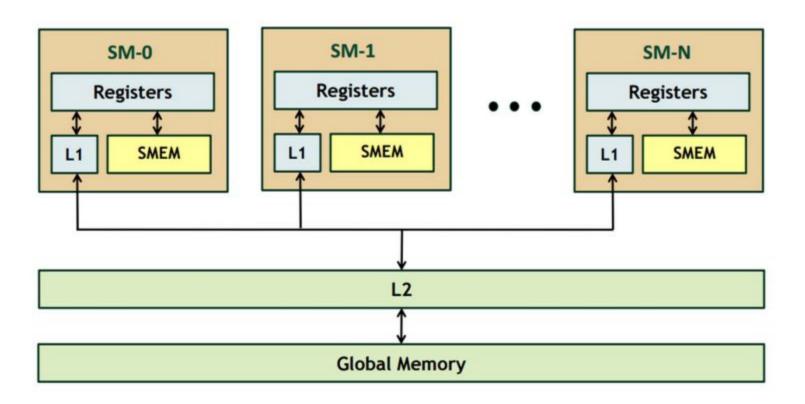
Serial algorithm

```
function doDfs(G[n]: Graph):
    visited = array[n, false]

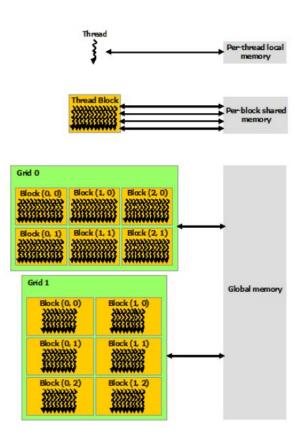
    function dfs(u: int):
        visited[u] = true
        for v: (u, v) in G
            if not visited[v]
            dfs(v)

    for i = 1 to n
        if not visited[i]
            dfs(i)
```

Память GPGPU



Потоки GPGPU



A Fast GPU Algorithm for Graph Connectivity

Jyothish Soman, Kothapalli Kishore, and P J Narayanan IIIT-Hyderabad

Main idea

DSU + Kruskal

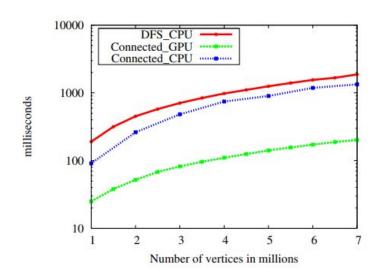
Details

- Hooking
- Pointer jumping

Results:

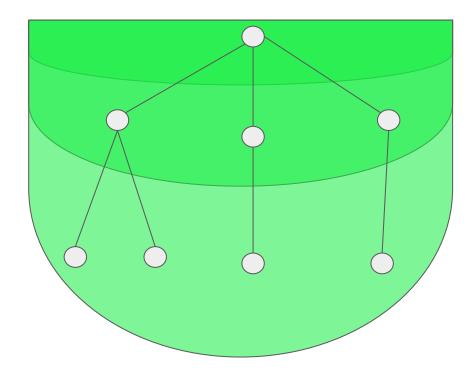
Data set	# vertices, #edges (in M)	Run Time (in ms)
Live journal	4.8, 69	207
Wiki Talk	2.4, 5	12
Citation n/w	3.7, 16.5	127
Road Networks		
California	2, 5.5	27
Pennsylvania	1.0, 3.0	15
Texas	1.4, 3.8	17

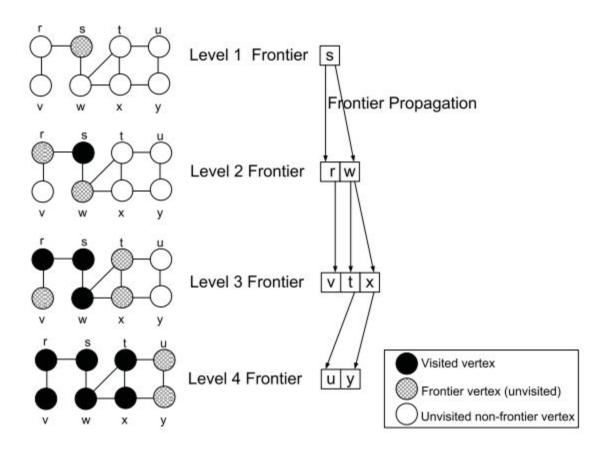
Table III
RUN TIME OF OUR ALGORITHM ON VARIOUS REAL-WORLD INSTANCES.



An Effective GPU Implementation of Breadth-First Search

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Synchronization

- Atomic operations
- Host Device communication

Results:

Table 1: BFS results on regular graphs

#Verte	IIIT-BFS	CPU-BFS	UIUC-BFS	Sp.
1M	462.8ms	$146.7 \mathrm{ms}$	67.8ms	2.2
2M	1129.2ms	311.8ms	$121.0 \mathrm{ms}$	2.6
5M	4092.2ms	1402.2ms	$266.0 \mathrm{ms}$	5.3
7M	6597.5ms	2831.4ms	$509.5 \mathrm{ms}$	5.6
9M	9170.1ms	4388.3ms	$449.3 \mathrm{ms}$	9.8
10M	11019.8ms	5023.0ms	$488.0 \mathrm{ms}$	10.3

Table 2: BFS results on real world graphs

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	#Vertex	IIIT-BFS	CPU-BFS	UIUC-BFS	S_{P}
New York	264,346	79.9ms	41.6ms	19.4ms	2.
Florida	1,070,376	372.0ms	$120.7 \mathrm{ms}$	$61.7 \mathrm{ms}$	2.0
USA-East	3,598,623	1471.1ms	581.4ms	158.5ms	3.
USA-West	6,262,104	$2579.4 \mathrm{ms}$	$1323.0 \mathrm{ms}$	$236.6 \mathrm{ms}$	5.6

Table 3: BFS results on scale-free graphs

#Vertex	IIIT-BFS	CPU-BFS	UIUC-BFS
1M	161.5ms	$52.8 \mathrm{ms}$	100.7ms
5M	1015.4ms	284.0ms	302.0ms
10M	2252.8ms	506.9ms	483.6ms

My results

Languages & tools:

- C++11
- CUDA 7.5

Hardware:

- i5-5200U CPU @ 2.20GHz × 4
- GeForce 920M

Benchmarking

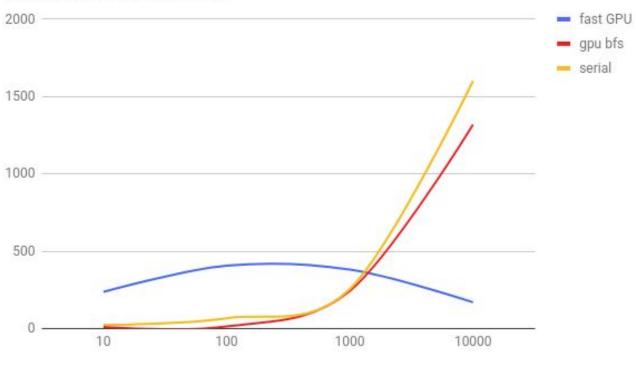
Benchmarking

- In additional to these two algorithms, serial BFS algorithms was implemented to compare GPU and CPU algorithms efficiencies.
- All measurements were done on syntactic date. The generator may be found in gen.cpp
- There are 3 different graph topologies:
 - each node has at least 1 edge (fig. 1)
 - each node has at least 4 edges (fig. 2)
 - o each node has at least 16 edges (fig. 3)
- The graph size range is from 10 nodes to 10000 nodes

Implemented algorithms

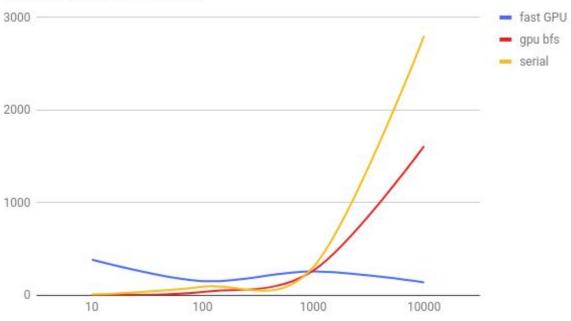
- Fast (see A Fast GPU Algorithm For Graph Connectivity)
- GPU BFS (see An Effective GPU Implementation Of Breadth-First Search)
- Serial (see https://en.wikipedia.org/wiki/Breadth-first_search)

indian, gpu bfs и serial



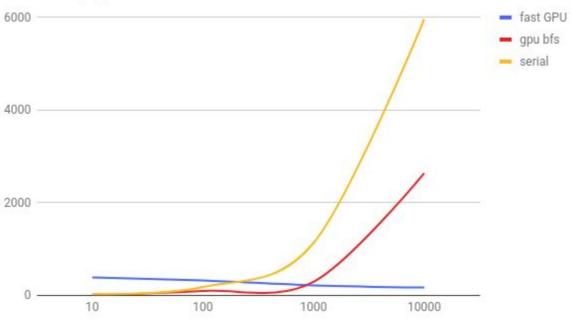
each node has only one edge

indian, gpu bfs и serial



each node has 4 edges

fast GPU, gpu bfs и serial



each node has 16 edges