Evaluation of Matrix-Based Context-Free Path Quering Algorithm

Nikita Mishin
Iaroslav Sokolov
Egor Spirin
mishinnikitam@gmail.com
sokolov.yas@gmail.com
egor@spirin.tech
Saint Petersburg State University
St. Petersburg, Russia

Vladimir Kutuev Egor Nemchinov Sergey Gorbatyuk vladimir.kutuev@gmail.com nemchegor@gmail.com sergeygorbatyuk171@gmail.com Saint Petersburg State University St. Petersburg, Russia Semyon Grigorev semen.grigorev@jetbrains.com Saint Petersburg State University St. Petersburg, Russia JetBrains Research St. Petersburg, Russia

ABSTRACT

Resantly proposed matrix-based algorithm for context-free path querying (CFPQ) evaluation looks very promising because it offloads the most performance-critical parts to boolean matrices multiplication (BMM). By this way it is possible to utilize modern hardware and software to achive highe performance CFPQ evaluation. We collect and publish a dataset which includes both real data and syntatica data for worst cases. We provide a number of implementations of matrix-based CFPQ algorithm which utilize different software and hardware features, such as GPGPUs and high-perfprmance libraries for BMM, and compare performance of them. Our evaluation shows that !!!

CCS CONCEPTS

• Information systems \rightarrow Query languages for non-relational engines; • Theory of computation \rightarrow Grammars and context-free languages; Parallel computing models; • Computing methodologies \rightarrow Massively parallel algorithms; • Computer systems organization \rightarrow Single instruction, multiple data.

KEYWORDS

Transitive closure, context-free path querying, graph databases, context-free grammar, GPGPU, matrix multiplication, boolean matrix

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1 INTRODUCTION

Language-constrained path querying [?], and particulary Context-Free Path Querying (CFPQ) [?] widely used for graph-structured data analysis in such areas as biological data analysis, RDF, network analysis. Real-world data is huge, so performance of CFPQ evaluation is important for practical tasks. Resently active research !!! [?]

One of the most promising algorithm is a matrix-based algorithm, proposed by Rustam Azimov. This algorith is pretty simple for implementation and allow to utilize massive-parallel hardware for CFPQs evaluation by offloading the most critical computattions to matrices multiplication. CPU/GPGPU based implementation. Investigate and compare.

There is no publically available standartized dataset for algorithms evaluation. We collect some data and propose possible candidate for it.

Research question: comparison of differend implementations of matrix-based CFPQ. We implement and compare performance.

We make the following contributions in this paper.

- (1) Implementation. Source code is available on GitHub:!!!
- (2) Dataset for evaluation. Real and syntethic data. Available. Data format. Reference values.
- (3) Evaluation. We show that !!!

2 MATRIX-BASED ALGORITHM FOR CFPQ

Matrix-based algoritm for CFPQ was proposed by Rustam Azimov [4]. This algorithm can be expressed in few lines of code in terms of matrices operations, and it is a sufficient advantage for implementation. It was shown that GPGPU utilization for queryes evaluation can significantly improve performance in comparison with other implementations [4] even float matrices used instead of boolean matrices.

Pseudocode of the algorithm is presented in listing 1.

Algorithm 1 Context-free path quering algorithm

```
1: function contextFreePathQuerying(D, G)
        n \leftarrow the number of nodes in D
 3:
        E \leftarrow the directed edge-relation from D
        P \leftarrow the set of production rules in G
 4:
        T \leftarrow the matrix n \times n in which each element is \emptyset
 5:
                                                     ▶ Matrix initialization
        for all (i, x, j) \in E do
 6:
             T_{i,j} \leftarrow T_{i,j} \cup \{A \mid (A \rightarrow x) \in P\}
 7:
 8:
        while matrix T is changing do
 9:
             T \leftarrow T \cup (T \times T)
                                          ▶ Transitive closure calculation
10:
        end while
11:
        return T
12:
13: end function
```

Here D=(V,E) be the input graph and $G=(N,\Sigma,P)$ be the input grammar. Each cell of the matrix T contains the set of nonterminals such that $N_k \in T[i,j] \iff \exists p=v_i\dots v_j$ —path in D, such that $N_k \stackrel{*}{\Longrightarrow} \omega(p)$, where $\omega(p)$ is a word formed by labels along path p. Thus, this algorithm solves reachability problem, or, according Hellings [6], process CFPQs by using relational query semantics.

As you can see, performance-critical part of this algorithm is a matrix multiplication. Note, that the set of nonterminals is finite, we can represent the matrix T as e set of boolean matrices: one for each nonterminal. In this case the matrix updeta operation be $T_{N_i} \leftarrow T_{N_i} + (T_{N_j} \times T_{N_k})$ for each production $N_i \rightarrow N_j \ N_k$ in P. Thus we can reduse CFPQ to boolean matrices multiplication. After such transfromation we can apply the next optimization: we can skip update if there are no changes in the matrices T_{N_j} and T_{N_k} at the previous iteration.

Thus, the most important part is efficient implementation of operations over boolean matrices, and in this work we compare effects of utilization of different approaches to matrices multiplication. All our implementations are based on the optimized version of the algorithm.

3 IMPLEMENTATION

We implement matrix-based algorithm fo CFPQ by using a number of different programming languages and tools. Our goal is to investigate effects of the next features of implementation.

- GPGPU utilization. It is mell-known thet GPGPUs are sutable for matrices operations, but performance of whole solution depends on task details: overhead on data transferring may negate effect of parallel computations. Moreover, it is believed that GPGPUs is not sutable boolean calculations [?]. Can GPGPUs utilization for CFPQ improve performance in comparison with CPU version?
- Existing libraries utilization is a good practice in software engeneering. Is it possible to achaive highe performance by using existing libraries for matrices operations or we need to create own solution to get more control?
- Low-level programming. GPGPU programming is traditionally low-level programming by using C-based languages (CUDA C, OpenCL C). On the other hand, there are number of approaches to create GPGPU-based solution by ysing

- such high-level languages as a Python. Can we get highperformance solution by using such approaches?
- Sparce matrices. Real graphs often are sparse, but not always. Is it sutable to use sparse matrix representation for CFPQ?

We provide next implementations for investigation.

- CPU-based solutions
- [Scipy] Saprse matrices multiplication by using Scipy [8] in Python programming language.
- [M4RI] Dense matrices multiplication by using m4ri¹ [1] library which implements 4 russian method [3] in C language. This library choosen because it is one of performnat implementation of 4 russian method [2].
 - GPGPU-based solutions

 ${\bf [GPU4R]}\,$ Manual implementation of 4 russian metod in CUDA C.

[GPU_N] Manual implementation of naïve boolean matrix multiplication in CUDA C.

[GPU_Py] Manual implementation of naïve boolean matrix multiplication in Pyton by using numba compiler².

Generic notes on optimizations. Notes on data transferring. On matrix changes tracking (we should multiply pair of matrices only if one of them changed in last iteration)

4 DATASET DESCRIPTION

We create and publish a dataset for CFPQ algorithms evaluation. This dataset contains both the real data and syntetic data for different cpecific cases, such as theoretical worst case, or matrices representation specific worst cases.

Our goal is querying algorithms evaluation, not a graph storages or graph databases evaluation, so all data is presented in text-based format to simplify usage in different environments. Grammars are in Chomscy Normal Form and are stored in the files with yrd extension. Each line is a rule in form of triple or pair. The example of grammar representation is presented in figure 1

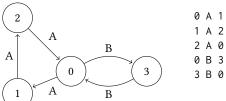
Figure 1: Example of grammar representation in the yrd file

Graphs are represented as a set of triples (edges) and are stored in the files with txt extension. Example of graph is presented in figure 2.

Each case is a pair of set of graphs and set of grammars: each query (grammar) should be applied to each graph. Cases are placed

 $^{^{1}}$ Actually we use pull request which is not merged yet: https://bitbucket.org/malb/m4ri/pull-requests/9/extended-m4ri-to-multiplication-over-the/diff. The original library implements operations over GF(2), and this pull request conteains operations over boolean semiring

²Numba is a JIT compiler which supports GPGPU for subset of Python programming. Offical page: http://numba.pydata.org/. Access date: 03.05.2019



(b) Representation of the input

(a) Example of input graph D_1 graph D_1 in txt file

Figure 2: Example of graph representation in txt file

in folders with case-specific name. Grammars and graph are placed in subfolders with names Grammars and Matrices respectively.

It is known that varians of the *same generation query* ?? are classical example of queryes that are context-free but not regular, so we use this type of queryes in our evaluation. The dataset includes data for next cases.

[RDF] The set of real RDF files (ontologies) from [9] and two variants of the same generation query (figures ??).

[Worst] Theoretical worst case for CFPQ time complexity which is proposed by Hellings [7]: graph is a two cycles of coprime lengths with single common vertex. First cycle labelled by open bracket and the second cysle is labelled by close bracked. Query is a grammar for A^nB^n language (grammar G_1 , figure 1).

[Full] The case when input graph is sparce, but result is a full graph. Such case may be a hard for sparse matrices representation. As an input graph we use a cycle all edges of which is labelled by the same token. As a queryes we use two grammars which describe arbitrary repetition of a token: unambiguous and highly ambiguous grammar (figure ??).

[**Sparse**] Sparse graphs from [5] which emulates realistic sparse data. Query is a same generation query.

5 EVALUATION

We evaluate all described implementations on all data and queries from presented dataset.

For evaluation we use PC with the next characteristics.

- OS
- CPU
- RAM
- GPU
 - Name
 - GHz
 - RAM
 - Threads
 - ...
- Libs versions
- Pyton runtime

Compiler options, Python runtime, etc.

Results of evaluation are presented in tables below. Time is mesured in milliseconds. Time for each algorithm is an averege time of 10 runs.

First is a **[RDF]** dataset. Results are presented in a table 1. We can see, that in this case !!!!

Results of theoretical worst case ([Worst] dataset) is presented ni table ??.

In this case !!!!! In this case !!!!!

Next is a **[Sparse]** datasset. Results are presented in table ??. For such type of graphs !!!!

The last dataset is a [Full], and results a shown in table 4 Finally, we can coculude that

- On GPU utilization
- On Existinng libraries
- On Low-level programming
- On sparse matrices

6 CONCLUSION AND FUTURE WORK

We provide a number of implementations of matrix-based algorithm for context-free path querying, collect a dataset for evaluation and provide results of evaluation of our implementation on collected dataset. Our evaluation shouws that!!!

log cfg. datalog -> cfg or to boolean/conjunctive

First direction for future research is a more detailed CFPQ algorithms investigation. We should do more evaluation on sparse matrices on GPGPUs.

Also it is nesessary to implement and evaluate solutions for graphs which is not fit in RAM. There is a set of technics for huge matrices multiplication. Is it possible to dopt it for CFPQ

Another direcion is a dataset improvement. First of all, it is necessary to collect more data, and more grammars/queries. Especially it would be interesting to add to dataset more real graphs and more real queryes. Secondly, it is nesessary to !!! data format to be able to evaluate different algorithms. Collaboration with community is required.

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Table 1: RDFs querying results

RDF	Query 1						Query 2						
	Scipy	M4RI	GPU4R	GPU_N	GPU_Py	CuSprs	Scipy	M4RI	GPU4R	GPU_N	GPU_Py	CuSprs	
rdf1	1	2	2	2	2	2	2	2	2	2	2	2	
rdf1	1	2	2	2	2	2	2	2	2	2	2	2	
rdf1	1	2	2	2	2	2	2	2	2	2	2	2	

Table 4: Full querying results

Graph size	Query 1							Query 2						
	Scipy	M4RI	GPU4R	GPU_N	GPU_Py	CuSprs	Scipy	M4RI	GPU4R	GPU_N	GPU_Py	CuSprs		
rdf1	1	2	2	2	2	2	2	2	2	2	2	2		
rdf1	1	2	2	2	2	2	2	2	2	2	2	2		
rdf1	1	2	2	2	2	2	2	2	2	2	2	2		

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