

Parsing Techniques for Context-Free Path Querying

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Formal language constrained path querying

- Finite directed edge-labelled graph $\mathcal{G} = (V, E, L)$
- The path is a world over L :
$$\omega(p) = \omega(v_0 \xrightarrow{l_0} v_1 \xrightarrow{l_1} \dots \xrightarrow{l_{n-1}} v_n) = l_0 \cdot l_1 \cdot \dots \cdot l_{n-1}$$
- The language \mathcal{L} (over L)

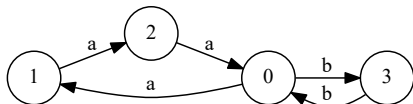
Formal language constrained path querying

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- The language \mathcal{L} (over L)
- Reachability problem: $Q = \{(v_i, v_j) \mid \exists p = v_i \dots v_j, \omega(p) \in \mathcal{L}\}$
- Path querying problem: $Q = \{p \mid \omega(p) \in \mathcal{L}\}$
 - ▶ Single path, all paths, shortest path ...

Context-Free path querying

- \mathcal{L} is a context-free language
- $G_{\mathcal{L}} = (N, \Sigma, R, S)$
- Reachability problem: $Q = \{(v_i, v_j) \mid \exists p = v_i \dots v_j, S \xrightarrow[G_L]{*} \omega(p)\}$
- Path querying problem: $Q = \{p \mid \omega(p) \in \mathcal{L}\}$

Example of CFPQ



Input graph

$S \rightarrow a S b$

$S \rightarrow \text{Middle}$

$\text{Middle} \rightarrow a b$

Query: language $\{a^n b^n \mid n > 0\}$

Paths:

$2 \xrightarrow{a} 0 \xrightarrow{b} 3$

$1 \xrightarrow{a} 2 \xrightarrow{a} 0 \xrightarrow{b} 3 \xrightarrow{b} 0$

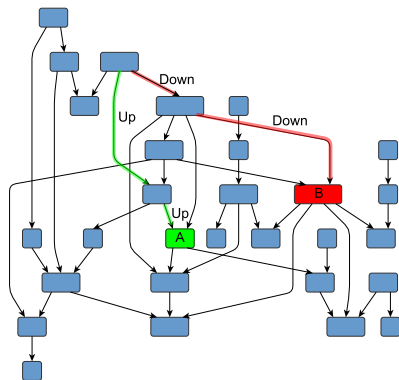
$p_1 = 0 \xrightarrow{a} 1 \xrightarrow{a} 2 \xrightarrow{a} 0 \xrightarrow{b} 3 \xrightarrow{b} 0 \xrightarrow{b} 3$

$p_2 = 0 \xrightarrow{a} 1 \xrightarrow{a} 2 \xrightarrow{a} 0 \xrightarrow{a} 1 \xrightarrow{a} 2 \xrightarrow{a} 0 \xrightarrow{b} 3 \xrightarrow{b} 0 \xrightarrow{b} 3 \xrightarrow{b} 0 \xrightarrow{b} 3 \xrightarrow{b} 0$

...

- Graph data bases querying
Mihalis Yannakakis. “Graph-theoretic methods in database theory.” 1990.
- Static code analysis
Thomas Reps et al. “Precise interprocedural dataflow analysis via graph reachability.” 1995
- ...

Graph data bases querying



Navigation through a graph

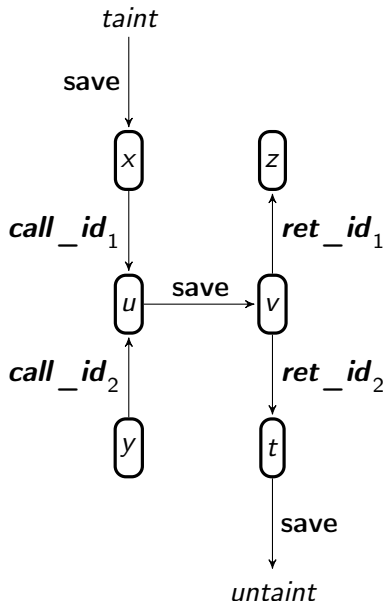
- Are nodes A and B on the same level of hierarchy?
- Is there a path of form $Up^n Down^n$?
- Find all paths of form $Up^n Down^n$ which start from the node A

Context-Free Path Querying

- *Sevon P., Eronen L.* “Subgraph queries by context-free grammars.” 2008
- *Hellings J.* “Conjunctive context-free path queries.” 2014
- *Zhang X. et al.* “Context-free path queries on RDF graphs.” 2016

Static code analysis

```
int id(int u)
{
    v = u;
    return v;
}
int main()
{
    //taint
    int x;
    int z, y;
    //untaint
    int t;
    z = id(x);
    t = id(y);
}
```



Static code analysis (Language Reachability Framework)

- *Thomas Reps et al.* “Precise interprocedural dataflow analysis via graph reachability.” 1995
- *Dacong Yan et al.* “Demand-driven context-sensitive alias analysis for Java.” 2011
- *Jakob Rehof and Manuel Fahndrich.* “Type-base flow analysis: from polymorphic subtyping to CFL-reachability.” 2001

Static code analysis (Language Reachability Framework)

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- *Qirun Zhang and Zhendong Su.* “Context-sensitive data-dependence analysis via linear conjunctive language reachability.” 2017

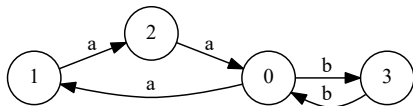
Parsing algorithms for CFPQ

- Structural representation of results
- Number of algorithms with different properties
- Number of theoretical results

Parsing algorithms for CFPQ

- Structural representation of results
- Number of algorithms with different properties
- Number of theoretical results
- Interconnection between different areas

Structural representation of result



Input graph

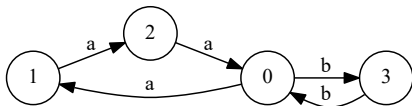
$S \rightarrow a S b$

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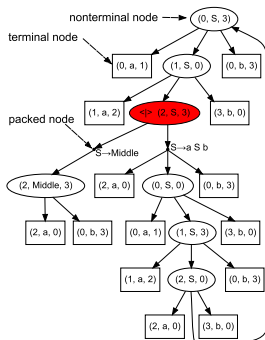
Grammar

Structural representation of result

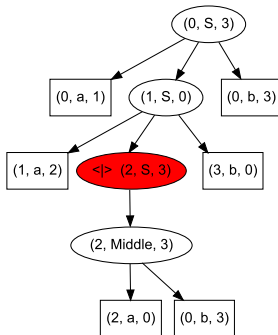


Input graph

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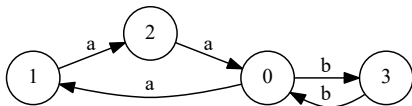


Query result (SPPF)



Tree for p_1

Structural representation of result



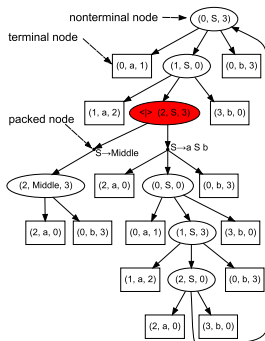
Input graph

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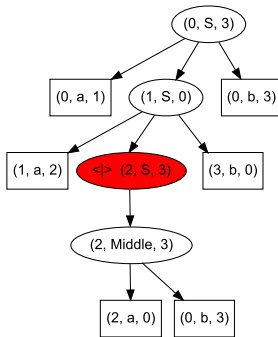
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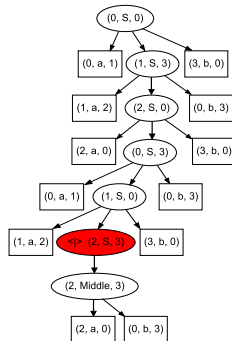
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Query result (SPPF)

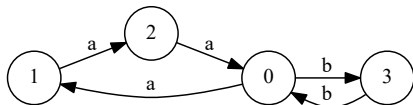


Tree for p_1



Tree for p_2

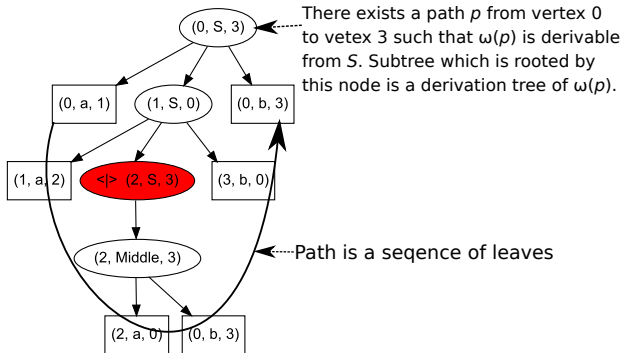
Paths extraction



$S \rightarrow a S b$

$S \rightarrow \text{Middle}$

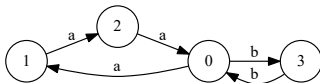
$\text{Middle} \rightarrow a b$



Path: $0 \xrightarrow{a} 1 \xrightarrow{a} 2 \xrightarrow{a} 0 \xrightarrow{b} 3 \xrightarrow{b} 0 \xrightarrow{b} 3$

Bar-Hillel theorem

Context-free languages are closed under intersection with regular languages



Regular language

$S \rightarrow a S b$

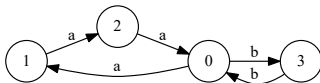
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Context-free language

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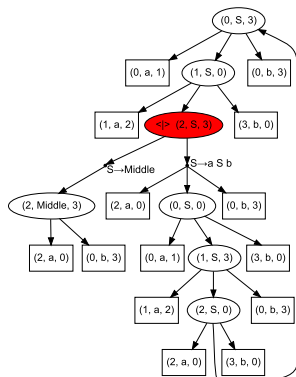
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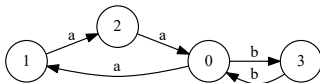
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Context-free language



Bar-Hillel theorem

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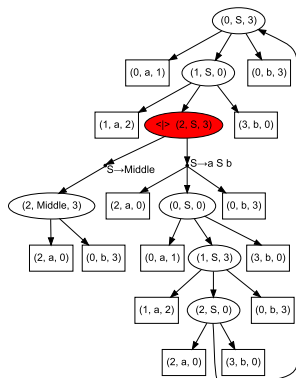
Regular language

$S \rightarrow a S b$

$S \rightarrow \text{Middle}$

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Context-free language



$(0, S, 3) \rightarrow (0, a, 1) (1, S, 0) (0, b, 3)$

$(1, S, 0) \rightarrow (1, a, 2) (2, S, 3) (3, b, 0)$

$(2, S, 3) \rightarrow (2, a, 0) (0, S, 0) (0, b, 3)$

$(2, S, 3) \rightarrow (2, \text{Middle}, 3)$

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$(0, \text{Middle}, 3) \rightarrow (2, a, 0) (0, b, 3)$

- Generalized LR for CFPQ
 - ▶ Based on Right Nulled Generalized LR: *Scott E., Johnstone A.* “Right Nulled GLR Parsers”
 - ▶ *Ekaterina Verbitskaia, Semyon Grigorev, and Dmitry Avdyukhin.* “Relaxed Parsing of Regular Approximations of String-Embedded Languages” 2015

Our experiments

- Generalized LR for CFPQ
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- Generalized LL for CFPQ (**GLL**)
 - ▶ Based on Generalized LL: *Scott E., Johnstone A.* “GLL parsing”
 - ▶ *Semyon Grigorev and Anastasiya Ragozina.* “Context-free path querying with structural representation of result.” 2017

How to integrate query language into general-purpose programming language?

- Transparency
- Compositionality
- Static error checking

Query language integration

How to integrate query language into general-purpose programming language?

- Transparency
- Compositionality
- Static error checking
- String-embedded languages
- ORMs
- Combinators

Combinators for CFPQ

- Implemented in Scala
- Based on Meerkat parser combinator library: *Anastasia Izmaylova, Ali Afroozeh, and Tijs van der Storm*. “Practical, general parser combinators” 2016
- *Ekaterina Verbitskaia, Ilya Kirillov, Ilya Nozkin, Semyon Grigorev*. “Parser Combinators for Context-Free Path Querying” 2019

Supported combinators

Combinator	Description
$a \sim b$	sequential parsing: a then b
$a \mid b$	choice: a or b

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$a \sim b$	sequential parsing: a then b
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$a ?$	optional parsing: a or nothing
$a *$	repetition of zero or more a
$a +$	repetition of at least one a
$a \wedge f$	apply f function to a if a is a token
$a \wedge \wedge$	capture output of a if a is a token
$a \& f$	apply f function to a if a is a parser
$a \& \&$	capture output of a if a is a parser

A set of functions for edges and vertices values handling.

```
def LV(labels: String*) =  
  V(e => labels.forall(e.hasLabel))  
def outLE(label: String) = outE(_.label() == label)  
def inLE (label: String) = inE ( _.label() == label)
```

Basic example

Is there a path from vertex 0 to vertex 3 which has form $a^n b^n$?

```
val Query : Nonterminal
    = syn (LV("0") ~ S ~ LV("3"))
```

```
val S: Nonterminal
    = syn (
        | "a" ~ S ~ "b"
        | "a" ~ "b"
    )
```

Example of generalization

```
def sameGen( brs ) =  
  reduceChoice(  
    brs.map { case ( lbr , rbr ) =>  
      lbr ~ syn( sameGen( brs ).? ) ~ rbr }
```

Example of generalization

```
def sameGen( brs ) =  
  reduceChoice(  
    brs.map { case ( lbr , rbr ) =>  
      lbr ~ syn( sameGen( brs ).? ) ~ rbr } )  
  
val query1 = syn( sameGen( List( ( "a" , "b" ) ) ) ) )  
  
val query2 = syn(  
  sameGen( List( ( p1 , p2 ) , ( "(" , ")" ) ) ) ~ p3 )
```


Example of values handling

Actors who played in some film

In Cypher

```
MATCH (m:Movie {title: 'Forrest_Gump'})  
      <-[:ACTS_IN]-(a:Actor)  
RETURN a.name, a.birthplace;
```

In Meerkat

```
val query =  
  syn((  
    (LV("Movie")::V(_.title == "Forrest_Gump")) ~  
    inLE("ACTS_IN") ~  
    syn(LV("Actor") ^  
      (e => (e.name, e.birthplace)))) &&  
  executeQuery(query, input)
```

Limitations

- Overhead for the regular constraints
- Not exactly clear how to compute arbitrary semantics for the paths
 - ▶ Paths can be lazily extracted, but in what order?
 - ▶ Is it possible to compute some semantics in case of cycles?

Boolean Matrix Multiplication for CFPQ

Transitive Closure

- Subset multiplication, $N_1, N_2 \subseteq N$
 - ▶ $N_1 \cdot N_2 = \{A \mid \exists B \in N_1, \exists C \in N_2 \text{ such that } (A \rightarrow BC) \in P\}$
- Subset addition: set-theoretic union.
- Matrix multiplication
 - ▶ Matrix of size $|V| \times |V|$
 - ▶ Subsets of N are elements
 - ▶ $c_{i,j} = \bigcup_{k=1}^n a_{i,k} \cdot b_{k,j}$
- Transitive closure
 - ▶ $a^{cf} = a^{(1)} \cup a^{(2)} \cup \dots$
 - ▶ $a^{(1)} = a$
 - ▶ $a^{(i)} = a^{(i-1)} \cup (a^{(i-1)} \times a^{(i-1)}), \quad i \geq 2$

The algorithm

Algorithm Context-free recognizer for graphs

```
1: function CONTEXTFREEPATHQUERYING( $D, G$ )
2:    $n \leftarrow$  the number of nodes in  $D$ 
3:    $E \leftarrow$  the directed edge-relation from  $D$ 
4:    $P \leftarrow$  the set of production rules in  $G$ 
5:    $T \leftarrow$  the matrix  $n \times n$  in which each element is  $\emptyset$ 
6:   for all  $(i, x, j) \in E$  do ▷ Matrix initialization
7:      $T_{i,j} \leftarrow T_{i,j} \cup \{A \mid (A \rightarrow x) \in P\}$ 
8:   while matrix  $T$  is changing do
9:      $T \leftarrow T \cup (T \times T)$  ▷ Transitive closure  $T^{cf}$  calculation
10:  return  $T$ 
```

Boolean Matrix Multiplication for CFPQ

- The matrix for nonterminal is a set of boolean matrices
- Matrices multiplication can be implemented efficiently by using modern hardware and high-performance libraries

Performance comparison setup

We use graphs from the classical set of ontologies: *skos*, *foaf*, *univ-bench*, *wine*, *pizza*, etc.

Queries are classical variants of the same-generation query

$$\begin{array}{ll} S \rightarrow \text{subClassOf}^{-1} S \text{ subClassOf} & S \rightarrow B \text{ subClassOf} \\ S \rightarrow \text{type}^{-1} S \text{ type} & S \rightarrow \text{subClassOf} \\ S \rightarrow \text{subClassOf}^{-1} \text{subClassOf} & B \rightarrow \text{subClassOf}^{-1} B \text{ subClassOf} \\ S \rightarrow \text{type}^{-1} \text{type} & B \rightarrow \text{subClassOf}^{-1} \text{subClassOf} \end{array}$$

Query 1

Query 2

Performance comparison results

№	#V	#E	Query 1 (ms)			Query 2 (ms)	
			CYK ¹	GLL	GPGPU	GLL	GPGPU
1	144	323	1044	10	12	1	1
2	129	351	6091	19	13	1	0
3	131	397	13971	24	30	1	10
4	179	413	20981	25	15	11	9
5	337	834	82081	89	32	3	6
6	291	685	515285	255	22	66	2
7	341	711	420604	261	20	45	24
8	671	2604	3233587	697	24	29	23
9	733	2450	4075319	819	54	8	6
10	6224	11840	–	1926	82	167	38
11	5864	19600	–	6246	185	46	21
12	5368	20832	–	7014	127	393	40

¹Zhang, et al. "Context-free path queries on RDF graphs."

Performance comparison results

Graph	Scipy	M4RI	GPU4R	GPU_N	GPU_Py	CuSprs
G5k-0.001	10.352	0.647	0.113	0.041	0.216	5.729
G10k-0.001	37.286	2.395	0.435	0.215	1.331	35.937
G10k-0.01	97.607	1.455	0.273	0.138	0.763	47.525
G10k-0.1	601.182	1.050	0.223	0.114	0.859	395.393
G20k-0.001	150.774	11.025	1.842	1.274	6.180	-
G40k-0.001	-	97.841	11.663	8.393	37.821	-
G80k-0.001	-	1142.959	88.366	65.886	-	-

Directions for research: engineering part

- Develop parallel and distributed algorithms
 - ▶ Distributed GLL(GLR)-based algorithms
 - ▶ Distributed matrix multiplication algorithms
 - ▶ Efficient implementation of sparse boolean matrices multiplication algorithms
- Adopt other parsing algorithms
 - ▶ Brzozowski's derivatives
 - ▶ Derivatives for graph querying: *Maurizio Nole and Carlo Sartiani*. "Regular path queries on massive graphs." 2016
 - ▶ Derivatives for context-free parsing: *Matthew Might, David Darais, and Daniel Spiewak*. "Parsing with derivatives: a functional pearl." 2011.
- Utilize other classes of languages for constraints specification
- Investigate incremental queries evaluation

Directions for research: theoretical part

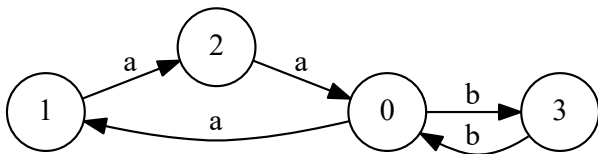
- Time complexity of GLL(GLR)-based algorithms for different classes of grammars
 - ▶ Current result for GLL-based: $O\left(|V|^3 * \max_{v \in V}(\deg^+(v))\right)$
- Theoretical lower bound for CFPQ
 - ▶ Is it possible to reduce CFPQ to PolyLog BMM?
 - ★ Our result is $O(|V|^2|N|^3(BMM(|V|) + BMU(|V|)))$
 - ▶ $\tilde{O}(n^\omega)$ solution for Dyck with one type of brackets: *Phillip G. Bradford*. “Efficient Exact Paths For Dyck and semi-Dyck Labeled Path Reachability.” 2018.
 - ▶ Derivatives for context-free parsing: *Matthew Might, David Darais, and Daniel Spiewak*. “Parsing with derivatives: a functional pearl.” 2011.

BMM-based algorithm: the Worst Case

Input graph: two cycles connected via a shared node

- first cycle has $2^k + 1$ edges labeled a
- second cycle has 2^k edges labeled b

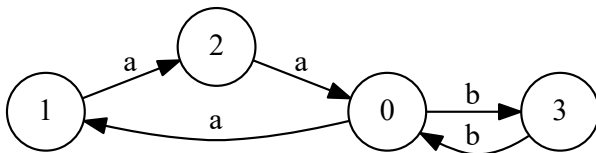
$$\begin{array}{ccc} S & \rightarrow & a S b \\ | & & a b \end{array}$$



The Worst Case: Step by Step

$$\begin{array}{lcl} S & \rightarrow & A B \\ & | & A S_1 \\ S_1 & \rightarrow & S B \\ A & \rightarrow & a \\ B & \rightarrow & b \end{array}$$

$$T_0 = \begin{pmatrix} \emptyset & \{A\} & \emptyset & \{B\} \\ \emptyset & \emptyset & \{A\} & \emptyset \\ \{A\} & \emptyset & \emptyset & \emptyset \\ \{B\} & \emptyset & \emptyset & \emptyset \end{pmatrix}$$



The Worst Case: Step by Step

$$T_0 \times T_0 = \begin{pmatrix} \emptyset & \emptyset & \emptyset & \emptyset \\ \emptyset & \emptyset & \emptyset & \emptyset \\ \emptyset & \emptyset & \emptyset & \{S\} \\ \emptyset & \emptyset & \emptyset & \emptyset \end{pmatrix}$$

$$T_1 = T_0 \cup (T_0 \times T_0) = \begin{pmatrix} \emptyset & \{A\} & \emptyset & \{B\} \\ \emptyset & \emptyset & \{A\} & \emptyset \\ \{A\} & \emptyset & \emptyset & \{S\} \\ \{B\} & \emptyset & \emptyset & \emptyset \end{pmatrix}$$

The Worst Case: Step by Step

$$T_2 = \begin{pmatrix} \emptyset & \{A\} & \emptyset & \{B\} \\ \emptyset & \emptyset & \{A\} & \emptyset \\ \{A, \mathbf{S_1}\} & \emptyset & \emptyset & \{S\} \\ \{B\} & \emptyset & \emptyset & \emptyset \end{pmatrix}$$

The Worst Case: Step by Step

$$T_3 = \begin{pmatrix} \emptyset & \{A\} & \emptyset & \{B\} \\ \{S\} & \emptyset & \{A\} & \emptyset \\ \{A, S_1\} & \emptyset & \emptyset & \{S\} \\ \{B\} & \emptyset & \emptyset & \emptyset \end{pmatrix}$$

The Worst Case: Step by Step

$$T_4 = \begin{pmatrix} \emptyset & \{A\} & \emptyset & \{B\} \\ \{S\} & \emptyset & \{A\} & \{\mathbf{S_1}\} \\ \{A, S_1\} & \emptyset & \emptyset & \{S\} \\ \{B\} & \emptyset & \emptyset & \emptyset \end{pmatrix}$$

The Worst Case: Step by Step

$$T_5 = \begin{pmatrix} \emptyset & \{A\} & \emptyset & \{B, S\} \\ \{S\} & \emptyset & \{A\} & \{S_1\} \\ \{A, S_1\} & \emptyset & \emptyset & \{S\} \\ \{B\} & \emptyset & \emptyset & \emptyset \end{pmatrix}$$

The Worst Case: Step by Step

$$T_6 = \begin{pmatrix} \{\mathbf{S_1}\} & \{A\} & \emptyset & \{B, S\} \\ \{S\} & \emptyset & \{A\} & \{S_1\} \\ \{A, S_1\} & \emptyset & \emptyset & \{S\} \\ \{B\} & \emptyset & \emptyset & \emptyset \end{pmatrix}$$

The Worst Case: Step by Step

$$T_7 = \begin{pmatrix} \{S_1\} & \{A\} & \emptyset & \{B, S\} \\ \{S\} & \emptyset & \{A\} & \{S_1\} \\ \{A, S_1, \mathbf{S}\} & \emptyset & \emptyset & \{S\} \\ \{B\} & \emptyset & \emptyset & \emptyset \end{pmatrix}$$

The Worst Case: Step by Step

$$T_8 = \begin{pmatrix} \{S_1\} & \{A\} & \emptyset & \{B, S\} \\ \{S\} & \emptyset & \{A\} & \{S_1\} \\ \{A, S_1, S\} & \emptyset & \emptyset & \{S, \mathbf{S_1}\} \\ \{B\} & \emptyset & \emptyset & \emptyset \end{pmatrix}$$

The Worst Case: Step by Step

$$T_9 = \begin{pmatrix} \{S_1\} & \{A\} & \emptyset & \{B, S\} \\ \{S\} & \emptyset & \{A\} & \{S_1, \mathbf{S}\} \\ \{A, S_1, S\} & \emptyset & \emptyset & \{S, S_1\} \\ \{B\} & \emptyset & \emptyset & \emptyset \end{pmatrix}$$

The Worst Case: Step by Step

$$T_{10} = \begin{pmatrix} \{S_1\} & \{A\} & \emptyset & \{B, S\} \\ \{S, \mathbf{S_1}\} & \emptyset & \{A\} & \{S_1, S\} \\ \{A, S_1, S\} & \emptyset & \emptyset & \{S, S_1\} \\ \{B\} & \emptyset & \emptyset & \emptyset \end{pmatrix}$$

The Worst Case: Step by Step

$$T_{11} = \begin{pmatrix} \{S_1, S\} & \{A\} & \emptyset & \{B, S\} \\ \{S, S_1\} & \emptyset & \{A\} & \{S_1, S\} \\ \{A, S_1, S\} & \emptyset & \emptyset & \{S, S_1\} \\ \{B\} & \emptyset & \emptyset & \emptyset \end{pmatrix}$$

The Worst Case: Step by Step

$$T_{12} = \begin{pmatrix} \{S_1, S\} & \{A\} & \emptyset & \{B, S, \mathbf{S_1}\} \\ \{S, S_1\} & \emptyset & \{A\} & \{S_1, S\} \\ \{A, S_1, S\} & \emptyset & \emptyset & \{S, S_1\} \\ \{B\} & \emptyset & \emptyset & \emptyset \end{pmatrix}$$

The Worst Case: Step by Step

$$T_{13} = \begin{pmatrix} \{S_1, S\} & \{A\} & \emptyset & \{B, S, S_1\} \\ \{S, S_1\} & \emptyset & \{A\} & \{S_1, S\} \\ \{A, S_1, S\} & \emptyset & \emptyset & \{S, S_1\} \\ \{B\} & \emptyset & \emptyset & \emptyset \end{pmatrix}$$