

# Extended Context-Free Grammars Parsing with Generalized LL

Artem Gorokhov and Semyon Grigorev

Saint Petersburg State University  
7/9 Universitetskaya nab.  
St. Petersburg, 199034 Russia  
gorohov.art@gmail.com  
semen.grigorev@jetbrains.com

**Abstract.** Parsing is important step of static program analysis because it allows to get structural representation of code which is necessary for analysis. One of important source of languages' syntax specification is official documentation which often use EBNF to make grammars more compact and expressive. Moreover, grammar in documentation may be ambiguous. These facts leads to problems with parsers creation automation by using parser generators, because the most parsing algorithms and tools require transformation to more simple form of grammar (BNF, CNF) and can not handle ambiguous grammars. There are number of works on ECFG processing without transformation, but no any can handle ambiguous CFG grammar. On the other hand, there is an generalized LL parsing algorithm, which can handle arbitrary grammars in  $O(n^3)$ , and LL grammars in  $O(n)$ , but this algorithm can not handle grammars in EBNF. We propose modification of generalized LL algorithm which can handle arbitrary grammar in form which close to EBNF (ECFG) without transformations. Also we implement parser generator based on described algorithm and we demonstrate that proposed modifications allow to get parsing performance improvement.

**Keywords:** Parsing, GLL, SPPF, EBNF, ECFG, RRPG, Recursive Automata

## 1 Introduction

Static program analysis is usually performed over a structural representation of code and parsing is a classical way to get such representation. Parser generators are often used to automate parser creation: these tools allow to derive parser from grammar. It decreases amount of efforts required for syntax analyzer creation and maintenance.

There are a wide range of parsing techniques and algorithms (CYK, LR(k), LALR(k), LL, etc) and parser generation tools, which based on them. The most practical parsing algorithms are LL(k)- and LR(k)-based algorithms. The LL family is more intuitive than LR and can provide better error diagnostic. LL(1) parsers are the most practical, but not powerful enough, moreover LL(k) for any

$k$  is not enough to process some languages: there are LR, but not LL languages. Also left and hidden left recursion in grammars is a problem for LL-based parsers. At the same time there is a common problem for both LL- and LR-based tools: handling of arbitrary ambiguous grammars. All these facts restrict class of grammars which can be handled. It makes parser creation difficult. In order to solve these problems generalized LL (GLL) [14] was proposed. This algorithm handles arbitrary context free grammar, even ambiguous and (hidden)left-recursive. Worst-case time and space complexity of GLL is cubic in terms of input size and for LL(1) grammars it demonstrates linear time and space complexity.

Extended BNF (EBNF) [19] is a useful format of grammar specification because it allows to make description of language syntax more expressive and compact. This formalism often used in documentation, which is one of main source of information for parsers developers. So, it is necessary to have a parser generator which supports grammar in EBNF. But classical parsing algorithms requires BNF, and as a result, parser generators requires BNF too. It is possible to convert EBNF to BNF but with this conversion we loose the structure of main grammar and resulting trees are for the BNF grammars.

In order to provide ability to process grammar in ELL, ELR [3–5, 7–9, 11, 12] and other can process EBNF but they do not deal with ambiguities in grammars.

At the same time, algorithm for left factorized grammars processing was introduced in [16]. Factorization means that there are no two productions for one nonterminal with equal prefixes (look at fig 1 for example). Shown, that factorization can reduce memory usage and increase performance which achieved by reusing common parts of rules for one nonterminal. Purposed idea can be used for processing grammars in EBNF with exception of same effects.

To summarise, it is possible to simplify language description required for parser generation in case a parser generator is based on generalized algorithm which can handle grammars in ECFG. In this work we present modified generalized LL parsing algorithm which handles grammars in ECFG without transformations. We show that changes of basic algorithm are very native for GLL nature. Also we demonstrate that proposed modifications allow to get parsing performance and memory usage improvement.

## 2 Extended Context-Free grammars

Parser generators widely use Extended CFG form: right parts of productions are regular expressions over union alphabet  $\Sigma \cup N$ .

**Definition 1** *An **extended context-free grammar** (ECFG) [9] is a tuple  $(N, \Sigma, P, S)$ , where  $N$  and  $\Sigma$  are finite sets of nonterminals and terminals,  $S \in N$  is the start symbol, and  $P$  (the productions) is a map from  $N$  to regular expressions over alphabet  $N \cup \Sigma$ .*

It is possible to transform ECFG to CFG [8], but this transformation leads to grammar size increase and change in grammar structure: new nonterminals addition is required during transformation. As a result, parsing performs not in

terms of user defined grammar. This fact leads to such problem: parser build structural representation not by the original grammar but by transformed and it may differ from expected. There are algorithms for parsing ECFG without transformations, based on different classical algorithms: ELL(k) and ELR(k) [8] parsers, Early-style parsers [?]. Some of them point out a problem with parsing conflicts [], and none of them work with arbitrary ECFG. Generalized parsing algorithms can handle arbitrary grammars and in this paper we will show how to use them for parsing with arbitrary ECFG.

### 3 Generalized LL Parsing Algorithm

Generalized parsing algorithms (GLL and GLR) was purposed to perform syntax analysis by arbitrary context-free grammar. Unlike the GLR, GLL algorithm [14] is rather intuitive and allows to perform better syntax error diagnostic. As an output of GLL we get Shared Packed Parse Forest(SPPF) [15] that represents all possible derivations of input string.

Work of the GLL algorithm based on descriptors, it allows to handle all possible derivations. Descriptor is a four-element tuple  $(L, i, T, S)$  that can uniquely define state of parsing process.  $L$  is a grammar slot — pointer to position in grammar of the form  $(S \rightarrow \alpha \cdot \beta)$ ,  $i$  — position in input,  $T$  — already built SPPF root,  $S$  — current Graph Structured Stack(GSS) [?] node.

In initial state we have descriptors that describe start positions in grammar and input, dummy tree node and bottom of GSS. On each step algorithm processes first descriptor in queue and makes actions depending on the grammar and input. If there are any ambiguity algorithm will queue descriptor for all cases to handle them all.

There are table based approach [13] which allows to generate only tables for given grammar instead of full parser code. The idea is similar to one in original article and main function uses same tree construction and stack processing functions. Pseudo code can be found in appendix A. Note that we do not include the check for first/follow sets in this paper.

### 4 Extended CFG GLL Parsing

In this section we will show an application of Extended Context-Free Grammars(ECFG) in automaton and corresponding GLL-style parsers.

#### 4.1 Factorization

In order to improve performance Elizabeth Scott and Adrian Johnstone offered support of factorised grammars in GLL [16]. The idea is to automatically factorize grammars and use them for parser generation.

The algorithm creates and queues new descriptors depending on current parse state that we get from unqueued descriptor. In case descriptor has been already

created it does not add it to queue. For this purpose we have a set of **all** created descriptors. Thus reducing a number of possible descriptors decreases the parse time and required memory.

Factorization decreases the number of grammar slots. Consider example from the paper [16] on fig. 1.

$$\begin{array}{ll}
 S ::= a \ a \ B \ c \ d & \\
 \begin{array}{l} | a \ a \ c \ d \\ | a \ a \ c \ e \\ | a \ a \end{array} & S ::= a \ a \ (B \ c \ d \mid c \ (d \mid e) \mid \varepsilon) \\
 & \text{(b) Production } P_0' \\
 \text{(a) Production } P_0
 \end{array}$$

**Fig. 1.** Example of factorization

Production  $P_0$  factorises to  $P_0'$ . Second is much compact and contains much less possible slots, so parser creates less descriptors. It gives significant performance improvement on some grammars.

## 4.2 Automata

The idea of factorisation has evolved to use of automata and their minimization.

First, we should define recursive automata [17].

**Definition 2** *Recursive automaton (RA)  $R$  is a tuple  $(\Sigma, Q, S, F, \delta)$ , where  $\Sigma$  is a set of terminals,  $Q$  — set of states of  $R$ ,  $S \in Q$  — start state,  $F \in Q$  — set of final states,  $\delta : Q \times (\Sigma \cup Q) \rightarrow Q$  — transition function.*

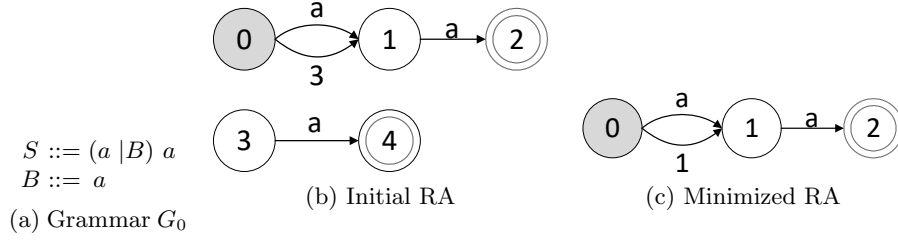
The only difference between RA and FSA is that in RA transition can be labeled either by terminal ( $\in \Sigma$ ) or by state ( $\in Q$ ). Further in this paper we will call transitions by elements from  $Q$  as nonterminal transitions and by terminal as terminal transitions.

ECFG can be converted to recursive automata. In conversion of ECFG to RA we will use following fact: right parts of ECFG are regular expressions over alphabet of terminals and nonterminals. Thus for each right-hand side of grammar productions we can build a finite state automaton using Thompson's method [18]. To transform the set of produced automata to RA we need to eliminate  $\varepsilon$ -transitions and replace transitions by nonterminals with transitions labeled by start states of corresponding to nonterminal FSA. An example of constructed recursive automaton for grammar  $\Gamma_0$  (fig. 2a) is given on fig. 2b, state 0 is start state.

Decrease of quantity of automaton states decreases number of GLL descriptors, as it was with factorization. Thus to increase performance of parsing we can minimize the number of states in produced automata.

First, RA should be converted to deterministic RA using the algorithm for FSA described in [2]. Then John Hopcroft's algorithm [10] can be applied to RA to minimize the number of states. An example for grammar  $G_0$  is shown on fig. 2c.

Note: later we will need a nonterminal names to build a SPPF, for this purpose we define function  $\Delta : Q \rightarrow N$  where  $N$  is nonterminal name.



**Fig. 2.** Example of automaton

### 4.3 Input processing

An GLL idea is to move through grammar and input simultaneously, creating multiple descriptors for the case of ambiguity.

Just as we can move through grammar slots we can move through states of automaton. Grammar slot in descriptor changes to state in RA. The problem is that in automaton we have nondeterministic choice because there can be many transitions to other states. Consider such significant cases:

- transition by current input terminal to final state exists
- transition by current input terminal exists and it leads to state that is not final
- nonterminal transition exists

All of them should be handled and this leads to nondeterminism. For the last case we just can call create function for each state. But for the terminal cases we need to add descriptor that describes next position to queue without checking it's existence in descriptor elimination set. Thus we use descriptors queue to handle nondeterminism in states, while original algorithm uses it to handle ambiguity in grammars.

```

function ADD( $S, u, i, w$ )
  if ( $S, u, i, w$ )  $\notin U$  then
     $U.add(S, u, i, w)$ 
     $R.add(S, u, i, w)$ 
    
```

Function **add** queues descriptor if it was not already created.

```

function CREATE( $S_{call}, S_{next}, u, i, w$ )
    
```

```

 $A \leftarrow \Delta(S_{call})$ 
if ( $\exists$  GSS node labeled  $(A, i)$ ) then
   $v \leftarrow$  GSS node labeled  $(A, i)$ 
  if (there is no GSS edge from  $v$  to  $u$  labeled  $(S_{next}, w)$ ) then
    add a GSS edge from  $v$  to  $u$  labeled  $(S_{next}, w)$ 
    for  $((v, z) \in \mathcal{P})$  do
       $(y, N) \leftarrow \text{getNodes}(S_{next}, u.nonterm, w, z)$ 
      if  $N \neq \$$  then
         $(\_, \_, h) \leftarrow N$ 
        pop $(u, h, N)$ 
         $(\_, \_, h) \leftarrow y$ 
        add $(S_{next}, u, h, y)$ 
  else
     $v \leftarrow$  new GSS node labeled  $(A, i)$ 
    create a GSS edge from  $v$  to  $u$  labeled  $(S_{next}, w)$ 
    add $(S_{call}, v, i, \$)$ 
return  $v$ 

```

Function **create** is called when we meet nonterminal transition. It performs necessary operations with GSS and checks if there are already built SPPF for current input position and nonterminal.

```

function POP( $u, i, z$ )
  if  $((u, z) \notin \mathcal{P})$  then
     $\mathcal{P}.add(u, z)$ 
    for all GSS edges  $(u, S, w, v)$  do
       $(y, N) \leftarrow \text{getNodes}(S, v.nonterm, w, z)$ 
      if  $N \neq \$$  then
        pop $(v, i, N)$ 
      if  $y \neq \$$  then
        add $(S, v, i, y)$ 

```

**Pop** function is called when we reach final state. It queues descriptors for all outgoing edges from current GSS node.

```

function PARSE
   $R.add(StartState, newGSSnode(StartNonterminal, 0), 0, \$)$ 
  while  $R \neq \emptyset$  do
     $(C_S, C_U, C_i, C_N) \leftarrow R.Get()$ 
     $C_R \leftarrow \$$ 
    if  $(C_N = \$) \& (C_S \text{ is pop state})$  then
       $eps \leftarrow \text{getNodeT}(\varepsilon, C_i)$ 
       $(\_, N) \leftarrow \text{getNodes}(C_S, C_U.nonterm, \$, eps)$ 
      pop $(C_U, C_i, N)$ 
    for each  $transition(C_S, label, S_{next})$  do
      switch  $label$  do
        case  $Terminal(x)$  where  $(x = input[i])$ 
           $R \leftarrow \text{getNodeT}(x, C_i)$ 

```

```

(y, N) ← getNodes( $S_{next}$ ,  $C_U.nonterm$ ,  $C_N$ ,  $R$ )
if  $N \neq \$$  then
    pop( $C_U, i + 1, N$ )
     $R.add(S_{next}, C_U, i + 1, y)$ 
case  $Nonterminal(S_{call})$ 
    create( $S_{call}, S_{next}, C_U, C_i, C_N$ )

```

The main function **parse** handles queued descriptor and checks all transitions from current state to be appropriate for current input terminal, or calls create function when meets nonterminal transitions.

#### 4.4 Parse forest construction

Result of the parsing process is structural representation of input — tree, or parse forest for the case of many derivation variants.

First, we should define derivation trees for recursive automaton: it is an ordered tree whose root labeled with start state, leaf nodes are labeled with a terminals or  $\varepsilon$  and interior nodes are labeled with nonterminals  $A$  and have a sequence of children that corresponds to transition labels of path in automaton that starts from the state  $\Delta(A)$ . More formal.

**Definition 3** *Derivation tree of sentence  $\alpha$  for the recursive automaton  $R = (\Sigma, Q, S, F, \delta)$ :*

- *Ordered rooted tree. Root labeled with  $\Delta(S)$*
- *Leaves are terminals  $\in \Sigma$*
- *Nodes are nonterminals  $\in \Delta(Q)$*
- *Node with label  $N_i \in \Delta(q_i)$  has children  $l_0 \dots l_n (l_i \in \Sigma \cup \Delta(Q))$  iff exists path  $q_i \xrightarrow{l_0} \dots \xrightarrow{l_n} q_m, q_m \in F$ .*

RA is ambiguous if there exist string that have more than one derivation trees. We work with arbitrary grammars, thus our RA can be ambiguous and we can define SPPF that can represent all possible derivation trees. It is similar to SPPF for grammars described in [15]. SPPF contains symbol nodes, packed nodes and intermediate nodes.

Packed nodes are of the form  $(S, k)$ , where  $S$  is a state of automaton. Symbol nodes have labels  $(X, i, j)$  where  $X \in \Sigma \cup \Delta(Q) \cup \varepsilon$ . Intermediate nodes have labels  $(S, i, j)$ , where  $S$  is a state of automaton.  $i$  is position in input before leftmost leaf terminal,  $j$  — position after rightmost leaf.

Packed node necessarily has right child — symbol node, and optional left child — symbol or intermediate node. Nonterminal and intermediate nodes may have several packed children. Terminal symbol nodes are leaves.

Use of intermediate and packed nodes leads to binarization of SPPF and thus the space complexity is  $O(n^3)$ .

**function** `getNodeT( $x, i$ )` did not change

We defined function **getNodes** which can construct two nodes: intermediate and nonterminal (at least one of them, at most both). It uses modified function

**getNodeP** that takes additional argument: state or nonterminal name. Symbol in returned SPPF node will be this argument's value.

```

function GETNODES( $S, A, w, z$ )
  if ( $S$  is pop state) then
     $x \leftarrow \text{getNodeP}(S, A, w, z)$ 
  else
     $x \leftarrow \$$ 

  if ( $w = \$$ ) & not ( $z$  is nonterminal node and it's extents are equal) then
     $y \leftarrow z$ 
  else
     $y \leftarrow \text{getNodeP}(S, S, w, z)$ 
  return ( $y, x$ )

function GETNODEP( $S, L, w, z$ )
  ( $\_, k, i$ )  $\leftarrow z$ 
  if ( $w \neq \$$ ) then
    ( $\_, j, k$ )  $\leftarrow w$ 
     $y \leftarrow$  find or create SPPF node labelled ( $L, j, i$ )
    if ( $\nexists$  child of  $y$  labelled ( $S, k$ )) then
       $y' \leftarrow \text{new packedNode}(S, k)$ 
       $y'.addLeftChild(w)$ 
       $y'.addRightChild(z)$ 
       $y.addChild(y')$ 
    else
       $y \leftarrow$  find or create SPPF node labelled ( $L, k, i$ )
      if ( $\nexists$  child of  $y$  labelled ( $S, k$ )) then
         $y' \leftarrow \text{new packedNode}(S, k)$ 
         $y'.addRightChild(z)$ 
         $y.addChild(y')$ 
  return  $y$ 

```

## 5 Evaluation(under construction)

Left factorization vs EBNF

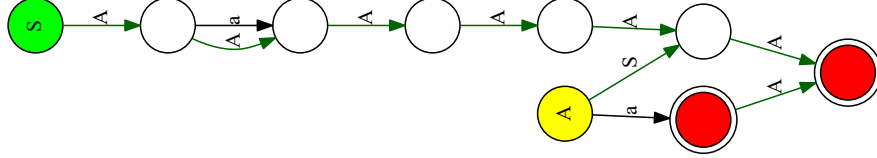
Small demo example (message to Scott)

$$\begin{aligned}
 S &::= A A A A A A \\
 &\quad | A a A A A A \\
 A &::= S A \mid a A \mid a
 \end{aligned}$$

**Fig. 3.** Grammar  $G_0$ .



We have compared our parsers built on factorized grammar and on minimized automaton. Grammar  $G_0$ (fig. 3) was used for the tests, it has long “common tail” which is not unified with factorization. FSA built for this grammar presented on fig. 4.



**Fig. 4.** Minimized automaton for grammar  $G_0$

Explanation of slots difference: for BNF, for factorized, for ECFG

Description of input. Short info about PC.

Note: SPPF construction was disabled while testing.

Length	Time, seconds		Descriptors		GSS Nodes		GSS Edges	
	factorized	minimized	factorized	minimized	factorized	minimized	factorized	minimized
100	0.206	0.127	52790	38530	200	200	42794	28534
200	1.909	1.54	215540	157030	400	400	175544	117034
300	8.844	7.125	488290	355530	600	600	398294	265534
400	25.876	21.707	871040	634030	800	800	711044	474034
500	60.617	51.245	1363790	992530	1000	1000	1113794	742534
1000	842.779	768.853	5477540	3985030	2000	2000	4477544	2985034
	Average gain: 19%		Average gain: 27%		Average gain: 0%		Average gain: 33%	

**Table 1.** Experiments results.

Table 1 shows that in general minimized version works 19% faster, uses 27% less descriptors and 33% less GSS edges. Also we use this automaton approach in metagenomic assemblies parsing and it gives visible performance increase. A bit more discussion on evaluation.

Examples of SPPF. May be some nontrivial cases:  $s -j a^* a^*$  and so on

## 6 Conclusion and Future Work

Described algorithm implemented in F# as part of the YaccConstructor project. Source code available here: [1].

Proposed modification can not only increase performance, but also decrease memory usage. It is critical for big input processing. For example, Anastasia Ragozina in her master’s thesis [13] shows that GLL can be used for graph

parsing. In some areas graphs can be really huge: assemblies in bioinformatics ( $10^8 \dots$ ). Proposed modification can improve performance not only in case of classical parsing, but in graph parsing too. We perform some tests that shows performance increasing in metagenomic analysis, but full integration with graph parsing and formal description is required.

One of way to specify any useful manipulations on derivation tree (or semantic of language) is an attributed grammars [?]. YARD supports it but our algorithm is not. So, attributed grammar and semantic calculation is a future work.

Yet another question is possibility of unification our results with tree languages: our definition of derivation tree for ECFG is quite similar to unranked tree and SPPF is similar to automata for unranked trees [6]. Theory of tree languages seems more mature than theory of general SPPF manipulations.

## References

1. Yaccconstructor project repository. <https://github.com/YaccConstructor/YaccConstructor>.
2. A. V. Aho and J. E. Hopcroft. *The design and analysis of computer algorithms*. Pearson Education India, 1974.
3. H. Alblas and J. Schaap-Kruseman. An attributed ell (1)-parser generator. In *International Workshop on Compiler Construction*, pages 208–209. Springer, 1990.
4. L. Breveglieri, S. C. Reghizzi, and A. Morzenti. Shift-reduce parsers for transition networks. In *International Conference on Language and Automata Theory and Applications*, pages 222–235. Springer, 2014.
5. A. Bruggemann-Klein and D. Wood. The parsing of extended context-free grammars. 2002.
6. H. Comon, M. Dauchet, R. Gilleron, C. Löding, F. Jacquemard, D. Lugiez, S. Tison, and M. Tommasi. Tree automata techniques and applications. 2007.
7. R. Heckmann. An efficient ell (1)-parser generator. *Acta Informatica*, 23(2):127–148, 1986.
8. S. Heilbrunner. On the definition of elr (k) and ell (k) grammars. *Acta Informatica*, 11(2):169–176, 1979.
9. K. Hemerik. Towards a taxonomy for ecfg and rrpq parsing. In *International Conference on Language and Automata Theory and Applications*, pages 410–421. Springer, 2009.
10. J. Hopcroft. An  $n \log n$  algorithm for minimizing states in a finite automaton. Technical report, DTIC Document, 1971.
11. S.-i. Morimoto and M. Sassa. Yet another generation of lalr parsers for regular right part grammars. *Acta informatica*, 37(9):671–697, 2001.
12. P. W. Purdom Jr and C. A. Brown. Parsing extended lr (k) grammars. *Acta Informatica*, 15(2):115–127, 1981.
13. A. Ragozina. Gll-based relaxed parsing of dynamically generated code. Masters thesis, SpBU, 2016.
14. E. Scott and A. Johnstone. Gll parsing. *Electronic Notes in Theoretical Computer Science*, 253(7):177–189, 2010.
15. E. Scott and A. Johnstone. Gll parse-tree generation. *Science of Computer Programming*, 78(10):1828–1844, 2013.

16. E. Scott and A. Johnstone. Structuring the gll parsing algorithm for performance. *Science of Computer Programming*, 125:1–22, 2016.
17. I. Tellier. Learning recursive automata from positive examples. *Revue des Sciences et Technologies de l'Information-Série RIA: Revue d'Intelligence Artificielle*, 20(6):775–804, 2006.
18. K. Thompson. Programming techniques: Regular expression search algorithm. *Commun. ACM*, 11(6):419–422, June 1968.
19. N. Wirth. Extended backus-naur form (ebnf). *ISO/IEC*, 14977:2996, 1996.

## A GLL pseudocode

```

function ADD( $L, u, i, w$ )
  if ( $(L, u, i, w) \notin U$ ) then
     $U.add(L, u, i, w)$ 
     $R.add(L, u, i, w)$ 

function CREATE( $L, u, i, w$ )
  ( $X ::= \alpha A \cdot \beta$ )  $\leftarrow L$ 
  if ( $\exists$  GSS node labeled ( $A, i$ )) then
     $v \leftarrow$  GSS node labeled ( $A, i$ )
    if (there is no GSS edge from  $v$  to  $u$  labeled ( $L, w$ )) then
      add a GSS edge from  $v$  to  $u$  labeled ( $L, w$ )
      for ( $(v, z) \in \mathcal{P}$ ) do
         $y \leftarrow \text{getNodeP}(L, w, z)$ 
         $\text{add}(L, u, h, y)$  where  $h$  is the right extent of  $y$ 
  else
     $v \leftarrow$  new GSS node labeled ( $A, i$ )
    create a GSS edge from  $v$  to  $u$  labeled ( $L, w$ )
    for each alternative  $\alpha_k$  of  $A$  do
       $\text{add}(\alpha_k, v, i, \$)$ 
  return  $v$ 

function POP( $u, i, z$ )
  if ( $(u, z) \notin \mathcal{P}$ ) then
     $\mathcal{P}.add(u, z)$ 
    for all GSS edges ( $u, L, w, v$ ) do
       $y \leftarrow \text{getNodeP}(L, w, z)$ 
       $\text{add}(L, v, i, y)$ 

function GETNODET( $x, i$ )
  if ( $x = \varepsilon$ ) then
     $h \leftarrow i$ 
  else
     $h \leftarrow i + 1$ 
   $y \leftarrow$  find or create SPPF node labelled ( $x, i, h$ )
  return  $y$ 

function GETNODEP( $X ::= \alpha \cdot \beta, w, z$ )
  if ( $\alpha$  is a terminal or a non-nullable nonterminal) & ( $\beta \neq \varepsilon$ ) then

```

```

    return  $z$ 
else
    if  $(\beta = \varepsilon)$  then
         $L \leftarrow X$ 
    else
         $L \leftarrow (X ::= \alpha \cdot \beta)$ 
     $(-, k, i) \leftarrow z$ 
    if  $(w \neq \$)$  then
         $(-, j, k) \leftarrow w$ 
         $y \leftarrow$  find or create SPPF node labelled  $(L, j, i)$ 
        if  $(\nexists \text{ child of } y \text{ labelled } (X ::= \alpha \cdot \beta, k))$  then
             $y' \leftarrow \text{new packedNode}(X ::= \alpha \cdot \beta, k)$ 
             $y'.addLeftChild(w)$ 
             $y'.addRightChild(z)$ 
             $y.addChild(y')$ 
        else
             $y \leftarrow$  find or create SPPF node labelled  $(L, k, i)$ 
            if  $(\nexists \text{ child of } y \text{ labelled } (X ::= \alpha \cdot \beta, k))$  then
                 $y' \leftarrow \text{new packedNode}(X ::= \alpha \cdot \beta, k)$ 
                 $y'.addRightChild(z)$ 
                 $y.addChild(y')$ 
    return  $y$ 

function DISPATCHER
    if  $R \neq \emptyset$  then
         $(C_L, C_u, C_i, C_N) \leftarrow R.Get()$ 
         $C_R \leftarrow \$$ 
         $dispatch \leftarrow false$ 
    else
         $stop \leftarrow true$ 

function PROCESSING
     $dispatch \leftarrow true$ 
    switch  $C_L$  do
        case  $(X \rightarrow \alpha \cdot x\beta)$  where  $(x = input[C_i] \parallel x = \varepsilon)$ 
             $C_R \leftarrow \text{getNodeT}(x, C_i)$ 
            if  $x \neq \varepsilon$  then
                 $C_i \leftarrow C_i + 1$ 
             $C_L \leftarrow (X \rightarrow \alpha x \cdot \beta)$ 
             $C_N \leftarrow \text{getNodeP}(C_L, C_N, C_R)$ 
             $dispatch \leftarrow false$ 
        case  $(X \rightarrow \alpha \cdot A\beta)$  where  $A$  is nonterminal
            create $((X \rightarrow \alpha A \cdot \beta), C_u, C_i, C_N)$ 
        case  $(X \rightarrow \alpha \cdot)$ 
            pop $(C_u, C_i, C_N)$ 

function PARSE

```

```
while not stop do  
  if dispatch then  
    dispatcher()  
  else  
    processing()
```