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"Foundation of HPC" course



Outline



Memory Padding

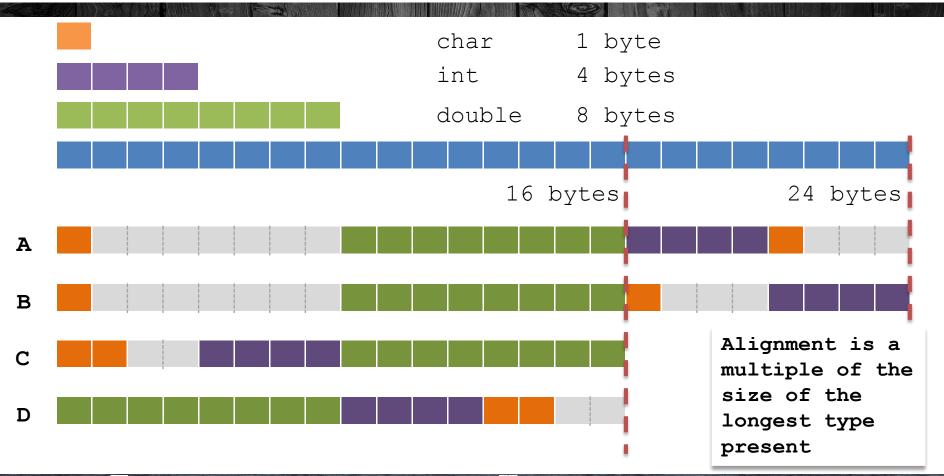


Overhead of Memory Allocation

```
typedef struct
                                             typedef struct
 char
          char field;
                                                       char_field;
                                              char
 double double_field;
                                              char
                                                       char_field2;
          int_field;
                                                        int_field;
 int
                                              int
                                              double double_field;
 char
          char field2:
                                             } STRUCT C:
} STRUCT_A;
typedef struct
                                             typedef struct
 char
          char field;
                                              double double field;
 double double_field;
                                              int
                                                        int_field;
          char field2:
                                                       char field;
 char
                                              char
 int
          int_field;
                                              char
                                                       char field2:
} STRUCT_B;
                                             } STRUCT_D;
```

Is there any difference among the above C structures?

...just have a look at the live results...



```
gcc ... -fpack struct[=n]
```

An instruction to pack them all an in the bitfield chain them down.

To be used carefully: it generates code binary-incompatible with code generated without the option (offset are different) and sub-optimal code.

Normally used for non-default binary interface (reduces the data stream). If given, *n* must be a (small) power of two.

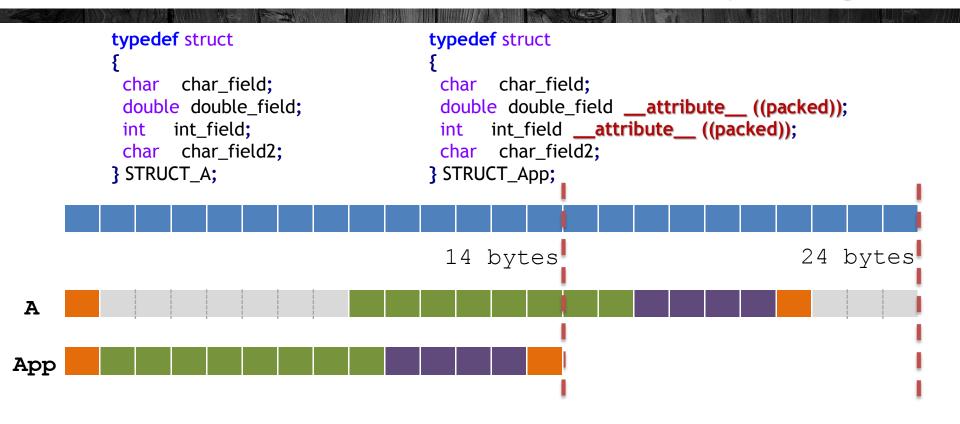
```
attribute ((packed));
```

Inline in the structure definition, at each field you want not to waste any byte. The same words of caution than above hold.

```
typedef struct
 char char field;
 double double field;
      int field:
 char char_field2;
} STRUCT_A;
```

```
typedef struct
 char char field;
 double double field attribute ((packed));
      int field:
 char char_field2;
} STRUCT_Ap;
```

```
typedef struct
                                           typedef struct
          char char_field;
                                                   char_field;
                                            char
                                             double double_field __attribute__ ((packed));
          double double field;
                int field:
                                                  int field;
                                             int
          int
          char char_field2;
                                             char char_field2;
         } STRUCT_A;
                                           } STRUCT_Ap;
                                                                    20 bytes 24 bytes
A
Ap
```



...just have a look at the live results...

General remark

in order to have best memory access performances, it is usually better to have data aligned to the natural alignment of your machine, which typically is 32 or 64 bits.

```
attribute ((aligned (n)));
```

You may achieve this in your data structures by using

System's malloc() allocator returns memory addresses multiples of 8 or 16 for 32- and 64-bits systems.

If you different needs (for instance: using AVXN instructions), you can use special calls:

```
void * posix memalign (void **memptr, size t alignment, size t size)
void * aligned alloc (size t alignment, size t size)
 introduced in ISO C11 and hence may have better portability to non-POSIX systems
```

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Overhead of Memory Allocation

Everytime you allocate some memory chunk on the heap, the system has to keep track of it, in order to

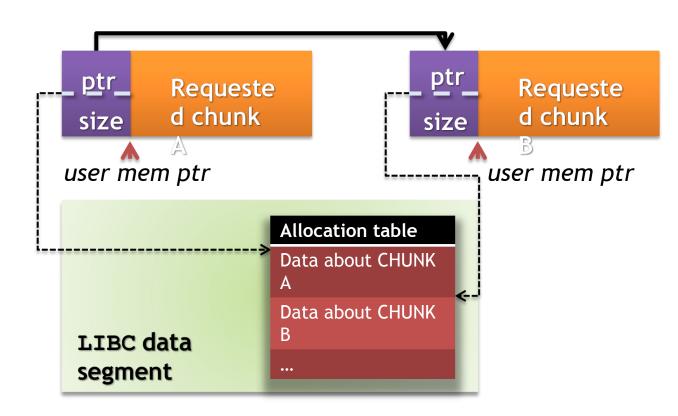
- maintain a precise account of used and available memory
- be able to "register back" that chunk into the free memory when you don't need it anymore (you free() it).

So, in addition to allocate *padded* memory layouts, which amounts to consume more memory than you formally request, there is also an additional "administrative" overhead whose amount is strongly system- and implementation-dependent (normally few bytes).

LIBC prepends an header to the chunk of data it allocates upon your request.

In that header, whose dimension is implementationdependent, there are several data it uses to check the consistency of memory chain and to recover the data about that same allocation.

Every detail is strongly implementation-dependent!



Allocating a huge amount of small data (like in text-book linked-list techniques) might be not the optimal strategy:

- you incur in some waste of memory due to padding, depending on the layout of data:
- for sure in a larger amount of data needed to trace every chunk you require.

How much large is the latter overhead? Again, that is implementation-dependent.

Let's measure it on your computer and for different implementation (i.e. different compilers and libc)

# of allocated blocks	Requested Mem	AdditionallyAl located	Ptr Overhead	System overhead
10M	400 MB	24 %	80 MB	18 %
1M	400 MB	1.3 %	8 MB	3.2 %
100k	400 MB	0.09 %	0.8 MB	0.27 %
10k	400 MB	0.0%	80 K	0.05%
1k	400 MB	0.12 %	8 K	0.125 %
100	400 MB	0.033 %	0.8 K	0.034%
10	400 MB	0.006 %	0.1 K	0.006 %