

1a.) One feasible solution is  $(1,1)$  with value of 5.  
It is feasible because it is within the feasible region created by constraints.

$$1b.) m = \frac{y_2 - y_1}{x_2 - x_1} = \frac{3 - 1}{0 - 2} = \frac{2}{-2} = -1$$

$$y - y_1 = m(x - x_1)$$

$$y - 3 = -1(x - 0)$$

$$y - 3 = -x$$

$$x + y = 3 \text{ - test } (0,0) \text{ for inequality}$$

$$0 \leq 3$$

$$x + y \leq 3$$

$$m = \frac{0 - 1}{1 - 2} = \frac{-1}{-1} = 1$$

$$y - y_1 = 1(x - x_1)$$

$$y - 0 = 1(x - 1)$$

$$y = x - 1$$

$$x - y = 1 \text{ - test } (0,0)$$

$$0 \leq 1$$

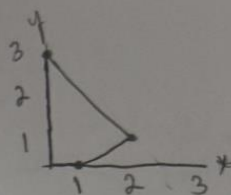
$$x - y \leq 1$$

$$y - y_1 = 0(x - x_1) \quad y \geq 0$$

$$y - 0 = y = 0$$

4 Constraints are:

$$\begin{aligned} x + y &\leq 3 \\ x - y &\leq 1 \\ y &\geq 0 \\ x &\geq 0 \end{aligned}$$



1c.)  $Z = 3x + 2y$  Optimal solutions are the "vertices"

$(2,1)$  - optimal solution

$$Z = 3(2) + 2(1)$$

$$6 + 2$$

$$= 8 \text{ optimal value for } (2,1)$$

2.)  $\max_{x_1, x_2} -20x_1 - 5x_2$  C:  $\begin{aligned} x_1 + 2x_2 &= 2 \\ x_1 - 3x_2 &\leq 3 \\ x_1 &\geq 0 \end{aligned}$

Step 1: Minimization to Maximization

Already a Max problem no change  $\max_{x_1, x_2} -20x_1 - 5x_2$

Step 2: All variables have non-negativity constraint

Replace any variable not restricted to non-negative values, with subtraction of 2 positive

$$x_j \rightarrow x_j' - x_j'' \quad x_j', x_j'' \geq 0$$

$$\max_{x_1, x_2} -20x_1 - 5x_2' + 5x_2''$$

$$\text{Subject to } x_1 + 2x_2' - 2x_2'' = 2$$

$$x_1 - 3x_2' + 3x_2'' \leq 3$$

$$x_1, x_2', x_2'' \geq 0$$

Step 3: Equality to inequality

Replace each equality constraint with a pair of inequalities

$$x_1 + 2x_2' - 2x_2'' = 2$$

$$\downarrow$$

$$x_1 + 2x_2' - 2x_2'' \leq 2$$

$$\nearrow$$

$$x_1 + 2x_2' - 2x_2'' \geq 2$$

Step 4: Greater than equalities to less than

Negate the GE inequalities

$$x_1 + 2x_2' - 2x_2'' \geq 2$$

$$\rightarrow -x_1 - 2x_2' + 2x_2'' \leq -2$$



$$3) \text{ Aggregate} = T(n)/n$$

Multipush = linear with respect to  $A$  ( $k$  objects)

$$\text{Sequence of } n \text{ operations} = n \cdot O(A) = O(nA)$$

$$\text{Aggregate} = \frac{T(n)}{n} = \frac{O(nA)}{n} = O(A) \text{ where } A = k \text{ objects}$$

4.) function FIND-SET( $x$ )

// follow pointer from  $x$  back to its Set object return Head (First)

return  $x.\text{Set}.\text{first}$

function Union( $x, y$ )

// weighted heuristic is Set.Size

Set  $x = x.\text{Set}$

Set  $y = y.\text{Set}$

if Set  $x.\text{Size} \geq \text{Set } y.\text{Size}$

// Set last member of  $x$  to representative of  $y$

Set  $x.\text{last}.\text{next} = \text{Set } y.\text{first}$

// update all of  $y$  members Set properly to point at  $x$

While Set  $y.\text{first} \neq \text{NIL}$

Set  $y.\text{first}.\text{Set} = \text{Set } x$

Set  $y.\text{first} = \text{Set } y.\text{first}.\text{next}$

Set  $x.\text{last} = \text{Set } y.\text{last}$

Set  $x.\text{Size} = \text{Set } x.\text{Size} + \text{Set } y.\text{Size}$

// Destroy Set object for  $y$ 's list. Details for this Do not

Else // handle appending  $x$  to  $y$

Set  $y.\text{last}.\text{next} = \text{Set } x.\text{first}$

While Set  $x.\text{first} \neq \text{NIL}$

Set  $x.\text{first}.\text{Set} = \text{Set } y$

Set  $x.\text{first} = \text{Set } y.\text{first}.\text{next}$

Set  $y.\text{last} = \text{Set } x.\text{last}$

Set  $y.\text{Size} = \text{Set } x.\text{Size} + \text{Set } y.\text{Size}$