

TABLE 2.2 Basic asymptotic efficiency classes

Class	Name	Comments
1	<i>constant</i>	Short of best-case efficiencies, very few reasonable examples can be given since an algorithm's running time typically goes to infinity when its input size grows infinitely large.
$\log n$	<i>logarithmic</i>	Typically, a result of cutting a problem's size by a constant factor on each iteration of the algorithm (see Section 4.4). Note that a logarithmic algorithm cannot take into account all its input or even a fixed fraction of it: any algorithm that does so will have at least linear running time.
n	<i>linear</i>	Algorithms that scan a list of size n (e.g., sequential search) belong to this class.
$n \log n$	<i>linearithmic</i>	Many divide-and-conquer algorithms (see Chapter 5), including mergesort and quicksort in the average case, fall into this category.
n^2	<i>quadratic</i>	Typically, characterizes efficiency of algorithms with two embedded loops (see the next section). Elementary sorting algorithms and certain operations on $n \times n$ matrices are standard examples.
n^3	<i>cubic</i>	Typically, characterizes efficiency of algorithms with three embedded loops (see the next section). Several nontrivial algorithms from linear algebra fall into this class.
2^n	<i>exponential</i>	Typical for algorithms that generate all subsets of an n -element set. Often, the term "exponential" is used in a broader sense to include this and larger orders of growth as well.
$n!$	<i>factorial</i>	Typical for algorithms that generate all permutations of an n -element set.

- a.** $2n(n-1)/2 \in O(n^3)$ **b.** $2n(n-1)/2 \in O(n^2)$
c. $2n(n-1)/2 \in \Theta(n^3)$ **d.** $2n(n-1)/2 \in \Omega(n)$

3. For each of the following functions, indicate the class $\Theta(g(n))$ the function belongs to. (Use the simplest $g(n)$ possible in your answers.) Prove your assertions.

- a.** $(n^3 + 1)^6$ **b.** $\sqrt{10n^4 + 7n^2 + 3n}$
c. $2n \lg(2n+2)^3 + (n^2 + 2)^2 \lg n$ **d.** $3^{n+1} + 3^{n-1}$
e. $2 \log_2 n$