CONSERVATION OF THE t-DIGEST SCALE INVARIANT

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ABSTRACT. A t-digest is a compact data structure that allows estimates of quantiles which increased accuracy near q=0 or q=1. This is done by clustering samples from $\mathbb R$ subject to a constraint that the number of points associated with any particular centroid is constrained so that the so-called k-size of the centroid is always ≤ 1 . The k-size is defined using a scale function that maps quantile q to index k. Since the centroids are real numbers, they can be ordered and thus the quantile range of a centroid can be mapped into an interval in k whose size the k-size of that centroid. The accuracy of quantile estimates made using a t-digest depends on the invariance of this constraint even as new data is added or t-digests are merged. This paper provides proofs of this invariance for four practically important scale functions.

1. Introduction

The key property of the t-digest (Dunning and Ertl, 2018) is that centroids either have unit weight or have k-size less than or equal to 1. Adding data to a t-digest or merging t-digests must preserve this constraint. The k-size is defined using a scale function, and there are four scale functions that are important for practical use

$$k_0(q) = \frac{\delta}{2}q$$

$$k_1(q) = \frac{\delta}{2\pi} \sin^{-1}(2q - 1)$$

$$k_2(q) = \frac{\delta}{Z(n)} \log \frac{q}{1 - q}$$

$$k_3(q) = \frac{\delta}{Z(n)} \begin{cases} \log 2q & \text{if } q \le 1/2\\ -\log 2(1 - q) & \text{if } q > 1/2 \end{cases}$$

These scale functions apply progressively more emphasis on the accuracy of quantile estimates near the tails and less emphasis on estimates near q = 1/2.

The size constraint says that any centroid with weight greater than one that spans from quantile q_1 to q_2 has a k-size limited by

$$k(q_2) - k(q_1) \le 1$$

The addition of new data to a t-digest proceeds by the addition of one or more new centroids followed by the merging of consecutive centroids that have combined k-size subject to the size constraint. The merging cannot cause a violation of the constraint, but it is conceivable that the addition of new centroids could indirectly cause a violation due to the shifting of q_1 and q_2 .

In fact, however, no such violation occurs, at least not for the scale functions above. This paper contains proofs of the preservation of this constraint as new data is added to the *t*-digest for each of these four scale functions.

2. Proofs

In all cases, we take n as the total number of sample so far, $q_1 = n_1/n$, and $q_2 = n_2/n$. The centroids to the left of the one we are considering therefore have a total weight of n_1 and our centroid has a weight of $n_2 - n_1$. For each of the scale functions, we need to consider the addition of Δn samples either to the right or left of our centroid. For some functions, addition to the right and left are logically the same due to symmetry and thus only one alternative need be considered.

2.1. For k_0 . The function k_0 is completely symmetric so we need only consider the case of data added to the right of the centroid in question. Subsequent to this addition, the k-size of our centroid will become

$$k_0(q_2') - k_0(q_1') = \frac{\delta}{2(n + \Delta n)} (n_2 - n_1)$$

$$< \frac{\delta}{2n} (n_2 - n_1) = k_0(q_2) - k_0(q_1) \le 1$$

This means that the k-size of a centroid can only shrink with the addition of new data and thus, the constraint on maximum k-size is preserved for k_0 . \square

2.2. For k_1 . The k_1 scaling function is symmetric so we only have to consider the case of new data to the right of the centroid of interest.

Note that

$$\frac{d [k_1(q_2) - k_1(q_1)]}{dn} = \frac{dk_1(q_2)}{dn} - \frac{dk_1(q_1)}{dn}$$
$$= -\frac{1}{n} \left(\sqrt{\frac{n_2}{n - n_2}} - \sqrt{\frac{n_1}{n - n_1}} \right)$$

This is uniformly negative if $n > n_2 > n_1 \ge 0$ which implies that any increase in n will result in a decrease in the k-size of a centroid that spans $q_1 = n_1/n$ to $q_2 = n_2/n$. The constraint is thus preserved for k_1 . \square

2.3. For k_2 . For the scaling function k_2 the same symmetry consideration applies as for k_1 so we only have to consider the addition of data to the right of the centroid of interest. Moreover, $k_2(q)$ increases without bound as $q \to 0$ and $q \to 1$ so the first and the last centroids must have unit weight and need not be considered (so $0 < q_1 < q_2 < 1$). The new k-size of some remaining centroid will be

$$k_{2}(q_{2}') - k_{2}(q_{1}') = \frac{\delta}{Z(n + \Delta n)} \left(\log \frac{n_{2}}{n - n_{2} + \Delta n} - \log \frac{n_{1}}{n - n_{1} + \Delta n} \right)$$
$$= \frac{\delta}{Z(n + \Delta n)} \left(\log \frac{n_{2}}{n_{1}} + \log \frac{n - n_{1} + \Delta n}{n - n_{2} + \Delta n} \right)$$

But Z(n) is non-decreasing and we know that $n > n_2 > n_1 > 0$, so $n - n_2 < n - n_1$ and

$$\frac{n-n_1+\Delta n}{n-n_2+\Delta n}<\frac{n-n_1}{n-n_2}$$

So,

$$k_2(q_2') - k_2(q_1') < \frac{\delta}{Z(n)} \left(\log \frac{n_2}{n_1} + \log \frac{n - n_1}{n - n_2} \right) = k_2(q_2) - k_2(q_1) \le 1$$

Thus

$$k_2(q_2') - k_2(q_1') < k_2(q_2) - k_2(q_1) \le 1$$

and the constraint is preserved for k_2 . \square

2.4. For k_3 . The definition of k_3 has two branches

$$k_3(q) = \frac{\delta}{Z(n)} \begin{cases} \log 2q & \text{if } q \le 1/2\\ -\log 2(1-q) & \text{if } q > 1/2 \end{cases}$$

Except where $q_1 < 1/2 < q_2$, we know by symmetry that we only need to consider the first of these branches, but there are two cases within that branch where the new samples are either to the right of the centroid of interest or to the left.

For the case where data is added to the right, n_1 and n_2 remain unchanged and since the denominator in $q'_i = n_i/(n + \Delta n)$ cancels out, the k-size is unchanged except for the normalizing factor Z(n) which is non-decreasing.

$$k_3(q_2') - k_3(q_1') = \frac{\delta}{Z(n+\Delta)} \left(\log \frac{n_2}{n+\Delta n} - \log \frac{n_1}{n+\Delta n} \right)$$
$$= \frac{\delta}{Z(n+\Delta)} \left(\log \frac{n_2}{n_1} \right)$$
$$\leq \frac{\delta}{Z(n)} \left(\log \frac{n_2}{n_1} \right) = k_3(q_2) - k_3(q_1) \leq 1$$

For the case where data is added to the left, start with the fact $0 < n_1 < n_2$. We know

$$\frac{n_2 + \Delta n}{n_1 + \Delta n} < \frac{n_2}{n_1}$$

From this,

$$k_3(q_2') - k_3(q_1') = \frac{\delta}{Z(n + \Delta n)} \left(\log \frac{n_2 + \Delta n}{n_1 + \Delta n} \right)$$

$$< \frac{\delta}{Z(n + \Delta n)} \left(\log \frac{n_2}{n_1} \right)$$

$$< \frac{\delta}{Z(n)} \left(\log \frac{n_2}{n_1} \right) = k_3(q_2) - k_3(q_1) \le 1$$

This leaves only the case where $q_1 < 1/2 < q_2$. We need only consider the case where the new data is added to the right. Note that

$$k_3(q_2) - k_3(q_1) = \frac{\delta}{Z(n)} \left(-\log\left(1 - \frac{n_2}{n}\right) - \log\frac{n_1}{n} \right)$$
$$= \frac{\delta}{Z(n)} \log\left(\frac{n^2}{n_1(n - n_2)}\right)$$

For the log term, the derivative with respect to n is

$$\frac{\mathrm{d}}{\mathrm{d}n}\log\left(\frac{n^2}{n_1(n-n_2)}\right) = \frac{n(n-2n_2)}{n_1(n-n_2)^2}$$

This is uniformly negative if $n_2 > n/2$. For $n > 2n_2$, however, we revert to the first branch of the k_3 and the k-size no longer changes with increasing n except for the effect of Z(n) which we know cannot make k-size any larger.

This means that for $q_1 < 1/2 < q_2$ we know $k_3(q'_2) - k_3(q'_1) < k_3(q_2) - k_3(q_1)$.

The t-digest size constraint is thus preserved for k_3 in all cases and sub-cases. \square

References

Dunning, T. and O. Ertl (2018). Computing extremely accurate quantiles using t-digests. https://arxiv.org/abs/1902.04023. arXiv:1902.04023 [stat.CO].

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