

4.1

Domain: All values of type order

for values t of type tree
and m of type int

Property $P(t)$: If $\max t = \text{Some } m$, then $\text{Member}(m, t)$ \wedge

Order: $R(t_1, t_2) : t_1$ is a proper substructure of t_2 .

Partition: $\{\text{Leaf}\} \{\text{Branch}(l, k, r)\}$ | l and r are values of type int tree and k is a value of type int.

Case 1: $t = \text{Leaf}$

Property holds vacuously because $\max \{\text{Leaf}\} = \text{None}$

Case 2: $t = \text{Branch}(l, k, r)$

By assumption, we know $\max t = \text{Some } m$

2nd Induction

Domain: All values of type order for int values m and k

Property $Q(o)$: if $\text{compare_int } m \ k = 0$, then $P(t)$

Order: $R(o_1, o_2)$ false (empty relation)

Partition: $\{\text{less}\}, \{\text{Equal}\}, \{\text{Greater}\}$

Case a: $o = \text{Less}$

This situation is not dealt with in \max , so holds

Case b: $o = \text{Equal}$

If $o = \text{Equal}$, by definition of \max

$\max k = \text{Some } m$, by $P(t)$

$\text{Member}(m, k)$, and because k is a proper

substructure of t , $\text{Member}(m, t)$

Case c: $o = \text{Greater}$

If $o = \text{greater}$, by definition of \max

$\max r = \text{Some } m$, by $P(t)$

$\text{Member}(m, r)$, because r is a proper substructure of t

$\text{Member}(m, t)$

° Because either $o = \text{Equal}$ or $o = \text{Greater}$ will be true
° $P(t)$.

4.2

Domain: all values of type order

Property $P(t)$: If $\text{max } t = \text{Some } m$, then $\text{AllLess}(m, \text{delete } m \ t)$
and If $\text{max } t = \text{None}$, then $t = \text{Leaf}$, for values of t of type tree and m ,
order: $R(t_1, t_2)$: t_1 is a proper substructure of t_2 . + type int

Partition: $\{\text{Some } m \mid m \text{ is a value of type int}\} \ \{\text{None}\}$

Case 1: $\text{max } t = \text{None}$

by the definition of max, $\text{max } t$ returns None
when t is matched with leaf, therefore
 $t = \text{Leaf}$.

Case 2: $\text{max } t = \text{Some } m$

- by the definition of max on $\text{BST}(t)$, m is the greatest int value in the tree t .

- $\text{AllLess}(m, \text{delete } m \ t)$ is true if m is greater than every value in the tree $\text{delete } m \ t$.

- We know m is the greatest value in t , so if delete correctly removes m from t , $P(t)$ holds.

- By lemma 1, If $\text{max } t = \text{Some } m$, $\text{AllLess}(m, \text{delete } m \ t)$

Lemma 1: for values i of type int and tr of type tree,
if $\text{Member}(i, tr)$ delete i tr removes i from tr .

Domain: all values of type order

Property $P(i)$: If $\text{Member}(i, tr)$, delete i tr removes i from tr

Order: $R(t_1, t_2)$: t_1 is a proper substructure of t_2

Partition: $\{\text{Member}(i, tr) = \text{true}\} \ \{\text{Member}(i, tr) = \text{false}\}$

Case 1: $\text{Member}(i, tr) = \text{false}$

by thm 4.1 this is an impossible case.

Case 2: $\text{Member}(i, tr) = \text{true}$

~ by def of delete with the understanding that $\text{Member}(i, tr)$
delete will reach the 'Equal' case in the 2nd matching

~ From here there are 2 cases the right side of the tree matches with Leaf or —.

~ If r matches with leaf, this means there is nothing in the right side, so the only remaining part is the left side, so delete returns the left side of br .

~ If r matches with $—$, delete will then check if l is a leaf or not, it matches $max\ l$ with some m and None

* in the case of None, this means l is empty and delete will return r .

* in the case of some m , this means l is not empty, so delete returns a tree with both r and l .