

A

In order to prove if the conclusion is derivable from the hypotheses I am going to use the Deduction in first-order logic theorem.

Let U_1, U_2, \dots, U_m be first order formulas, and V the conclusion. V is deducible from U_1, U_2, \dots, U_m if there exists a sequence of formulas (f_1, f_2, \dots, f_m) such that $f_m = V$ and we have:

$$\forall i \in \{1, \dots, m\}$$

a) $f_i \in \mathcal{A}_{PL}$ (axiom of predicate logic)

b) $(\exists i \in \{1, \dots, m\}) f_i$ is one of the hypotheses

c) We can derive ~~another formula~~ f_i from true formulas which ~~already exist~~ f_i from true formulas which exist in the sequence using modes rules.

d) We can obtain f_i using the universal generalization on an existing formula in the sequence.

Blades de la Academia
9.11. A

H_1 : Anyone who plays tennis or soccer is healthy.

H_2 : Anyone who plays chess is smart

H_3 : All CS students play chess

H_4 : Samuel is a CS student and plays tennis

C : Samuel is smart and healthy

D - is the domain: the universe of people

Samuel - constant $\in D$

Predicate symbols:

unary:

$CS: D \rightarrow \{T, F\}$, $CS(x) = T$ if x is a CS student

$chess: D \rightarrow \{T, F\}$, $chess(x) = T$ if x plays chess

$tennis: D \rightarrow \{T, F\}$, $tennis(x) = T$ if x plays tennis

$soccer: D \rightarrow \{T, F\}$, $soccer(x) = T$ if x plays soccer

$healthy: D \rightarrow \{T, F\}$, $healthy(x) = T$ if x is healthy

$smart: D \rightarrow \{T, F\}$, $smart(x) = T$ if x is smart

$$H_1: (\forall x) ((\text{tennis}(x) \vee \text{soccer}(x)) \rightarrow \text{healthy}(x))$$

$$H_2: (\forall x) (\text{chess}(x) \rightarrow \text{smart}(x))$$

$$H_3: (\forall x) (\text{CS}(x) \rightarrow \text{chess}(x))$$

$$H_4: \text{CS}(\text{Samuel}) \wedge \text{tennis}(\text{Samuel})$$

$$C: \text{Smart}(\text{Samuel}) \wedge \text{healthy}(\text{Samuel})$$

$$H_4 \vdash_{\text{simplification}} \text{CS}(\text{Samuel}) = \gamma_5$$

$$U \wedge V \vdash_{\text{simplif}} U$$

$$U \wedge V \vdash_{\text{simplif}} V$$

$$H_3 \vdash_{\text{univ. inst}} \text{CS}(\text{Samuel}) \rightarrow \text{chess}(\text{Samuel}) = \gamma_6$$

universal
instantiation

$$(\forall x) U(x) \vdash_{\text{univ. inst}} U(t)$$

t - term (variable or constant of the domain)

$$\gamma_5, \gamma_6 \vdash_{\text{mp}} \text{chess}(\text{Samuel}) = \gamma_4$$

modus
ponens

$$V, V \rightarrow U \vdash_{\text{mp}} U$$

$$H_2 \vdash_{\text{univ-inst}} \text{clen}(\text{Samuel}) \rightarrow \text{smart}(\text{Samuel}) = \mathcal{I}_8$$

$$\mathcal{I}_4, \mathcal{I}_8 \vdash_{\text{mp}} \text{smart}(\text{Samuel}) = \mathcal{I}_9$$

$$H_4 \vdash_{\text{simplification}} \text{tennis}(\text{Samuel}) = \mathcal{I}_{10}$$

$$H_1 \vdash_{\text{univ-inst}} (\text{tennis}(\text{Samuel}) \vee \text{soccer}(\text{Samuel})) \rightarrow \text{healthy}(\text{Samuel}) = \mathcal{I}_{11}$$

$$\mathcal{I}_{10}, \mathcal{I}_{11} \vdash_{\text{mp}} \text{healthy}(\text{Samuel}) = \mathcal{I}_{12}$$

$$\mathcal{I}_9, \mathcal{I}_{12} \vdash_{\text{conjunction}} \text{smart}(\text{Samuel}) \wedge \text{healthy}(\text{Samuel}) = C$$

$(H_1, H_2, H_3, H_4, \mathcal{I}_5, \mathcal{I}_6, \mathcal{I}_7, \mathcal{I}_8, \mathcal{I}_9, \mathcal{I}_{10}, \mathcal{I}_{11}, \mathcal{I}_{12})$ is the proof of C . Therefore Samuel is smart and healthy.