Documentation

ADT Map – implementation on a hash table, collision resolution by coalesced chaining.

Problem statement:

A student implemented an agenda to keep track of the number of hours spent for studying per day during a month(30 days). The application offers the possibility to add, delete and check the amount of time spent for learning. In addition, it can compute the average of worked hours per week.

Justification:

A map is a container where the elements are pairs <key, value>. Similarly the chosen problem uses <day, hours> pairs. Keys have to be unique in a map and so are the days in a month(integers from 1 to 30). Moreover, each key has one single associated value (amount of hours). The problem is suitable for this implementation because the hash table provides a virtually direct access to objects(number of hours spent for studying) based on a key (day of the month --unique Integer). In addition, each key is associated with a value and elements can be found, inserted, and removed using the integer index as an array index

❖ Domain of the ADT Map:

 $M = \{m \mid m \text{ is a map with elements } e = (k, v), \text{ where } k \in TKey \text{ and } v \in TValue\}$

Interface of the ADT:

subalgorithm init(m)

descr: creates a new empty map

pre: true

post: $m \in M$, m is an empty map.

subalgorithm destroy(m)

descr: destroys a map

pre: $m \in M$

post: m was destroyed

subalgorithm add(m, k, v)

descr: add a new key-value pair to the map

pre: $m \in M$, $k \in TKey$, $v \in TValue$ post: $m' \in M$, $m' = m \cup \langle k, v \rangle$

subalgorithm remove(m, k, v)

descr: removes a pair with a given key from the map

pre: $m \in M$, $k \in TKey$

post: v ∈ TValue, where

$$v \leftarrow \begin{cases} v', \text{ if } \exists < k, \, v' > \in m \text{ and } m' \in M, \, m' = m \backslash < k, \, v' > \\ \\ 0 \, \text{TValue, otherwise} \end{cases}$$

function search(m, k, v)

descr: searches for the value associated with a given key in the map

pre: $m \in M$, $k \in TKey$

post: v ∈ TValue, where

$$v \leftarrow \begin{cases} v', & \text{if } \exists < k, v' > \in m \\ 0 & \text{TValue, otherwise} \end{cases}$$

function getIterator(m, it)

descr: returns an iterator for a map

pre: $m \in M$

post: it \in I, it is an iterator over m.

function getSize(m)

descr: returns the number of pairs from the map

pre: $m \in M$

post: size ← the number of pairs from m

function getKeys(m, s)

descr: returns the set of keys from the map

pre: $m \in M$

post: $s \in S$, s is the set of all keys from m

o function getValues(m, b)



descr: returns a bag with all the values from the map

pre: $m \in M$

post: $b \in B$, b is the bag of all values from m

function getPairs(m, s)

descr: returns the set of pairs from the map

pre: $m \in M$

post: $s \in S$, s is the set of all pairs from m

❖ Interface Iterator:

o subalgorithm init(it, ht):

descr: creates an iterator over the give sh table

pre: list ∈ HT

post: it ϵ I, it is an iterator over the given hash table

subalgorithm destroy(it):

descr: destroys the given iterator

pre: it ∈ I

post: it was destroyed

o function getCurrent(it):

descr: returns the current element in the hash table

pre: it ∈ I

post: getCurrent <- the current noac(pair<key,value>) in the table

function isValid(it):

descr: checks if the iterator is valid

pre: it ∈ I

post: isValid <- true if the iterator is still valid, false otherwise

o subalgorithm next(it):

descr: moves the iterator to the next node in the hash table

pre: it ∈ I

post: it' ϵ I, it' is positioned on the next node in the hash table

ADT Representation:

∘ Node: □

key: TKey value: TVal

o Hashī e:

T: Node[]

next: Integer[] m: Integer

firstFree: Integer h: TFunction

o Iterator:

list: Hash Table currentPos: Integer