

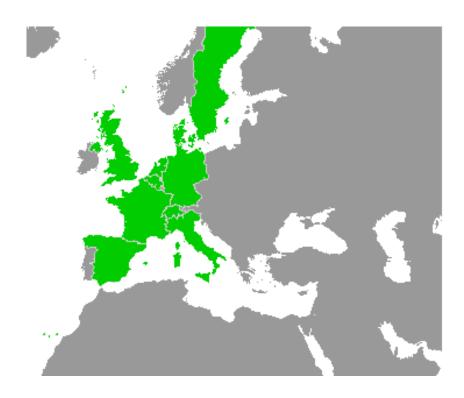
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Work-Package 7: "Toolchain"

# Event-B Model of Subset 026, Section 3.5.3

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# Model Description

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This document describes a formal model of the requirements of section 3.5.3 of the subset 026 of the ETCS specification 3.3.0 [Eur12]. This section describes the establishing of a communication connection between on-board and on-track equipment.

The model is expressed in the formal language Event-B [Abr10] and developed within the Rodin tool [Jas12]. This formalism allows an iterative modeling approach. In general, one starts with a very abstract description of the basic functionality and step-wise adds additional details until the desired level of accuracy of the model is reached. Rodin provides the necessary proof support to ensure the correctness of the refined behavior.

In this document we present an Event-B model of the protocol to initiate a communication session in an ETCS implementation, as implemented from the on-board unit. At first, we describe shortly the background of Event-B, then the overall approach taken to model this section and finally present the model in detail. For each of the iterative modeling steps, we describe the details added by the refinement.

OBU	on board unit
RIU	radio in-fill unit
RBC	radio block centre
SRS	system requirements specification

Table 1. Glossary

#### 1 Short Introduction to Event-B

The formal language Event-B is based on a set-theoretic approach. It is a variant of the B language, with a focus on system level modeling [Abr10]. An Event-B model is separated into a static and a dynamic part.

The dynamic part of an Event-B model describes abstract state machines. The state is represented by a set of state variables. A transition from one state to another is represented by parametrized events which assign new values to the state variables. Event-B allows unbounded state spaces. They are constrained by invariants expressed in first order logic with equality which must be fulfilled in any case. The initial state is created by a special initialization event.

The static part of an Event-B model is represented by contexts. These consist of carrier sets, constants and axioms. The type system of a model is described by means of carrier sets and constraints expressed by axioms.

Event-B is not only comprised of descriptions of abstract state machines and contexts, but also includes a development approach. This approach consists of iterative refinement of the machines until the desired level of detail is reached. In the Rodin tool, proof obligations are automatically created which ensure correct refinement.

Together with the machine invariants, the proof obligations for the refinement are formally proven, creating proof trees. To accomplish this, there are different options: many proof obligations can be discharged by automated provers (e.g., AtelierB, NewPP, Rodin's SMT-plugin), but as the underlying logic is in general undecidable, it is sometimes necessary to use the interactive proof support of Rodin.

Any external actions, e.g., mode changes by the driver or train level changes are modeled via parametrized events. Only events can modify the variables of a machine. An Event-B model is on the system level, events are assumed to be called from a software system into which the

functional model is embedded. The guards of the events assure that any event can only be called when appropriate.

# 2 Modeling Strategy

The section 3.5.3 of the SRS describes how a communication session is established. In its context, the low level EURORADIO network connection (cf. §3.5.1.1) is considered basic functionality and is not part of the modeling.

The basic modeling element are entities which represent one piece of equipment, either on-board, i.e., on the train, or on-track. The model is constructed from the local point of view of an OBU entity. On-track entities are only modeled as possible communication partners.

# 3 Model Overview

Figure 1 shows the structure of the Event-B model. The left column represents the abstract state machines, the right column the contexts. An arrow from one machine to another machine represents a refinement relation, an arrow from a machine to a context represents a sees relation and arrow from one context to another represents an extension relation.

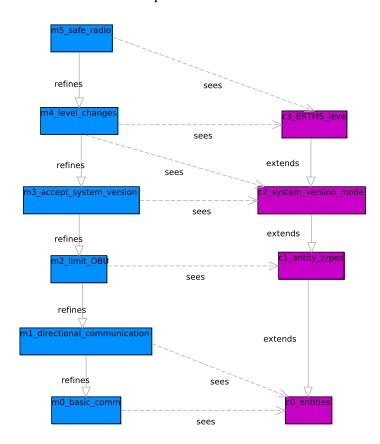


Figure 1. Overview on State Machine and Context Hierarchy

The modeling starts with the very abstract possibility to establish and to terminate a communication session in the machine m0, the set of entities is defined in the context c0. This basic functionality is refined in the succeeding machines to incorporate the different stages of the protocol to establish a session. The contexts further refine the entities to on-track and on-board entities and limit the modeling to the point of view of an OBU.

The machine m1 discerns incoming and outgoing communication sessions, i.e., initiated by the modeled piece of equipment or by an external one. The context c1 introduces the different

types of equipment which are used in m2 to refine the two different protocols for outgoing and incoming sessions and to limit the model to the OBU point of view. c2 introduces the notion of compatible systems. This is used in m3 to identify on-track equipment with a compatible system version. This machine also discerns between accepting and non-accepting RBCs to contact. c3 adds the different ERTMS levels and the relevant train modes to the model. This is used in m4 to model the different situations where a communication session must be established. m5 adds the notion of safe radio connection as low-level prerequisite for a communication session.

The representation of the state machines of the modeled protocols for establishing a communication session is modeled implicitly. The model allows sessions with different partners in parallel (but respects the constraints of the specification like §3.5.3.5.2). The state of the protocol with different partners is tracked by adding / removing these partners from sets representing those different states of the protocol.

#### 4 Model Benefits

The Event-B model in Rodin has some interesting properties which are highlighted here. Some stem from the fact that Rodin is well integrated into the Eclipse platform which renders many useful plugins available, both those explicitly developed for integration with Rodin, but also other without Rodin in mind. Other interesting properties stem from the fact that Rodin and Event-B provide an extensive proof support for properties.

- Refinement The Event-B approach allows iterative development based on refinement. This
  allows starting modeling with a very abstract machine and then step-wise adding more
  detailed behavior. Rodin generates all the necessary proof obligations which are required to
  assure correct refinement.
- Requirements Tracing Rodin provides an extensible EMF model, therefore it is easily possible to trace requirements using the requirements modeling framework of Eclipse (RMF) via the ProR plugin. This allows the usage of requirement documents in the OMG standardized Requirements Interchange Format (ReqIF). Figure 2 shows the requirements tracing using ProR in Rodin.

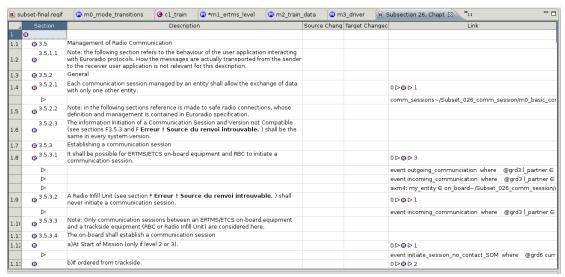


Figure 2. Traceability in ReqIf file using ProR

 Model Animation The Event-B model can be animated via different plugins, e.g., ProB or AnimB. This allows the simulation of the model, by clicking on the activated events and tracking the resulting state of the variables. This technique allows to examine the run-time

behavior of the model, e.g., for testing purposes. There is also ongoing development for a model-based testing plugin in Rodin, which will allow storing and replaying of event sequences.

- **Non-Testable Requirements** The Event-B model supports the specification of invariants which can be formally proven using the proof support of Rodin. This includes for example the non-testable requirement specified in the subset 076 for §3.5.3.2 (see Section 5.5).
- **Safety Properties** Using Rodin's proof support and the formalization as invariants, it is possible to formalize and prove the identified safety property of the case study (see Section 5.7).

# 5 Detailed Model Description

This section describes in more detail the formal model, beginning from the most abstract Event-B machine. For each refinement, in general only the important changes will be shown, the complete model is available as a Rodin project. At each step the additional modeled functionality and its representation will be described. In particular the initialization event is not shown for the refined machines. If not mentioned explicitly, sets are initialized empty, integers with value 0 and Boolean variables with false.

# 5.1 Context 0 - Entities

This context defines the type of entities with whom a communication session can be established. *my\_entity* represents the piece of equipment which is modeled.

```
CONTEXT c0_entities

SETS

entities

CONSTANTS

my_entity

AXIOMS

axm1 : my_entity ∈ entities

END
```

# 5.2 Machine 0 - Basic Communication

This state machine represents the basic functionality. It allows for the creation and the termination of a communication session with another entity. The sessions are represented by the state variable "session" which can contain values of type "entities". The respective events are triggered with a parameter  $l\_partner$  representing the communication partner.

# **Implemented Requirements**

• each session allows for communication between two entities (cf. §3.5.2.1)

```
SEES c0_entities VARIABLES
```

# sessions INVARIANTS inv1: sessions C entities

```
inv1 : sessions \subseteq entities \setminus \{my\_entity\}
EVENTS
Initialisation
   begin
             act1 : sessions := \emptyset
   end
Event establish_communication \widehat{=}
   any
           l\_partner
    where
             grd1: l\_partner \notin sessions
             grd2: l\_partner \neq my\_entity
   then
             act1 : sessions := sessions \cup \{l\_partner\}
   end
Event terminate\_communication =
   any
           l_partner
    where
            grd1: l\_partner \in sessions
   then
            act1 : sessions := sessions \setminus \{l\_partner\}
    end
END
```

# 5.3 Machine 1 - Directional Communication

The first refinement of the machine refines the notion of communication session to incoming sessions, i.e., where another entity initiates a session with  $my\_entity$  and outgoing sessions where  $my\_entity$  initiates the session.

The data refinement is proven by the invariant which states that "sessions" is equal to the disjoint union of "outgoing\_sessions" and "incoming\_sessions". The abstract "establish\_session" event is refined to the two events "incoming\_session" and "outgoing\_session".

```
REFINES m0_basic_comm

SEES c0_entities

VARIABLES

incoming_sessions
outgoing_sessions
INVARIANTS

inv1: partition(sessions, incoming_sessions, outgoing_sessions)

EVENTS
```

```
Event incoming\_communication =
refines establish_communication
   any
          l_partner
   where
            grd1: l\_partner \notin incoming\_sessions \cup outgoing\_sessions
            grd2: l\_partner \neq my\_entity
   then
            act1:incoming\_sessions := incoming\_sessions \cup \{l\_partner\}
   end
Event outgoing_communciation \widehat{=}
refines establish_communication
   any
          l_partner
   where
            grd2: l\_partner \neq my\_entity
            grd1: l\_partner \notin incoming\_sessions \cup outgoing\_sessions
   then
            act1: outgoing\_sessions := outgoing\_sessions \cup \{l\_partner\}
Event terminate communication \widehat{=}
refines terminate_communication
   any
          l_partner
   where
            grd1: l\_partner \in incoming\_sessions \cup outgoing\_sessions
   then
            act1:incoming\_sessions := incoming\_sessions \setminus \{l\_partner\}
            act2 : outgoing\_sessions := outgoing\_sessions \setminus \{l\_partner\}
   end
END
```

# 5.4 Context 1 - Entity Types

The first context extension introduces the different types of entities relevant in this requirement subset, i.e., on-board unit (OBU), radio in-fill unit (RIU) or radio block centre (RBC). Every entity has a unique type. The goal is to model the communication protocol from the point of view of an OBU, therefore the type of *my\_entity* is restricted to OBU.

```
CONTEXT c1_entity_types
EXTENDS c0_entities
CONSTANTS

RBC
RIU
OBU
on_track
```

```
on\_board
AXIOMS
axm1 : partition(entities, RBC, RIU, OBU)
axm2 : on\_track = RIU \cup RBC
axm3 : on\_board = OBU
axm4 : my\_entity \in on\_board
END
```

# 5.5 Machine 2 - On Board Modeling

The next machine refinement adds the notion of being contacted by an on-track entity to establish a communication session. It also adds the first state of the protocol, i.e., entities which should be contacted with the "Initiation of a communication session" message. On-track entities which order "my\_entity" to contact are stored in the "contacted\_by" set, entities to which the first message is sent by "my\_entity" are stored in the set "contacted', representing those which are in the first stage of the protocol.

The invariants prove that  $my\_entity$  will only be in contact with on-track entities and that any entities which are considered for a communication session are on-track entities. Any entity with whom there is already a communication session will not be considered for another session, and finally no radio in-fill unit can initiate a communication session with  $my\_entity$ .

# Implemented Requirements

- It shall be possible for OBU and RBC to initiate communication session (cf. §3.5.3.1)
- RIU cannot initiate a communication session (cf. §3.5.3.2) This invariant is marked as non-testable in Subset-076.

The other invariants ensure that a communication partner is not in different states of the communication protocol at the same time. A session protocol can be started by the order to contact an RBC or directly by the OBU.

```
REFINES m1_directional_communication

SEES c1_entity_types

VARIABLES

contacted
contacted_by
INVARIANTS

inv1 : incoming_sessions ∪ outgoing_sessions ⊆ on_track
inv2 : contacted ⊆ on_track
inv3 : contacted_by ⊆ on_track
inv4 : contacted_by ∩ (incoming_sessions ∪ outgoing_sessions) = ∅
inv5 : contacted ∩ (incoming_sessions ∪ outgoing_sessions) = ∅
inv6 : incoming_sessions ∩ RIU = ∅
inv7 : contacted ∩ contacted_by = ∅
```

```
EVENTS
Event incoming\_communication =
refines incoming_communication
    any
           l_partner
    where
            grd1: l\_partner \notin incoming\_sessions \cup outgoing\_sessions
            grd3: l\_partner \in on\_track \setminus RIU
            grd4: l\_partner \notin contacted
            grd5: l\_partner \notin contacted\_by
   then
            act1:incoming\_sessions := incoming\_sessions \cup \{l\_partner\}
   end
Event receive\_contact\_order =
   any
           l\_partner
    where
            \texttt{grd1} : l\_partner \notin contacted \cup contacted\_by \cup incoming\_sessions \cup outgoing\_sessions
            grd2: l\_partner \in on\_track
   then
            act1: contacted\_by := contacted\_by \cup \{l\_partner\}
   end
Event initiate_session_after_contact \widehat{=}
   any
           l\_partner
    where
            grd2: l\_partner \in contacted\_by
    then
            act1: contacted := contacted \cup \{l\_partner\}
            act2: contacted\_by := contacted\_by \setminus \{l\_partner\}
Event initiate_session_no_contact \widehat{=}
   any
           l_partner
    where
            grd5: l\_partner \notin incoming\_sessions \cup outgoing\_sessions \cup contacted \cup contacted\_by
            grd3: l\_partner \in on\_track
    then
            act2 : contacted := contacted \cup \{l\_partner\}
   end
END
```

# 5.6 Context 2 - System Version Compatibility

The next context extension introduces the notion of compatible system versions. This is modeled as a static property of the on-track equipment wrt.  $my\_entity$ , i.e., the context axiom that "system\_version\_compatible" is a subset of all on-track entities. On this level of abstraction, there is no need for a finer grained notion of compatibility.

```
EXTENDS c1_entity_types

CONSTANTS

system_version_compatible

AXIOMS

axm1 : system_version_compatible ⊆ on_track

END
```

# 5.7 Machine 3 - Accepting RBC and System Version

The next machine refines the contact order events by discerning between the orders to contact an accepting or a non-accepting RBC. The notion of being an accepting RBC is considered to be a dynamic property and therefore modeled as a variable, i.e., the set "accepting".

The "receive\_contact\_order" event is refined by two separate events, one for orders for an accepting RBC and one for a non-accepting RBC. The "outgoing\_communication" event is refined by two events, one for a compatible system version and the other for an incompatible one.

Furthermore, a just established communication session with on-track equipment with an incompatible system version will be terminated immediately after receiving this information. This is modeled by the set "terminating\_session" which holds values of type entities. Only those communication sessions in this set can be closed by the termination event.

# Implemented Requirements

- In case of a non-accepting RBC, all existing communication sessions with other RBCs must be terminated (cf. §3.5.3.5.2)
- After the system version is received by the OBU, the communication session is considered established and (cf. §3.5.3.8)
  - if the system version is compatible, the OBU shall send the session established message to track-side (cf. 3.5.3.8.a)
  - if the system version is incompatible, the OBU shall terminate the session (cf. 3.5.3.8.b)
- Any RBC which is contacted and with whom a communication session is established has a compatible system version (safety requirement from requirements document).

```
REFINES m2_limit_OBU

SEES c2_system_version_mode

VARIABLES

terminating_sessions
accepting
INVARIANTS

inv2 : RBC ∩ outgoing_sessions \ terminating_sessions ⊆ system_version_compatible
inv3 : accepting ⊆ RBC
inv4 : terminating_sessions ⊆ on_track

EVENTS
```

```
Event receive\_information\_compatible =
extends outgoing_communciation
   any
          l_partner
   where
            grd3: l\_partner \in contacted
            grd4: l\_partner \in system\_version\_compatible
   then
            act1: outgoing\_sessions := outgoing\_sessions \cup \{l\_partner\}
            act2: contacted := contacted \setminus \{l\_partner\}
   end
Event receive_information_incompatible \widehat{=}
extends outgoing_communciation
   any
          l_partner
   where
            grd3: l\_partner \in contacted
            grd4: l\_partner \notin system\_version\_compatible
   then
            act1: outgoing\_sessions := outgoing\_sessions \cup \{l\_partner\}
            act2 : contacted := contacted \setminus \{l\_partner\}
            act3: terminating\_sessions := terminating\_sessions \cup \{l\_partner\}
   end
Event receive_contact_order_accept =
refines receive_contact_order
   any
          l_partner
   where
            grd1: l\_partner \notin contacted \cup contacted\_by \cup incoming\_sessions \cup outgoing\_sessions
            grd3: l\_partner \in RIU \cup (RBC \cap accepting)
   then
            act1: contacted\_by := contacted\_by \cup \{l\_partner\}
   end
Event receive_contact_order_non_accept =
refines receive_contact_order
   any
          l_partner
   where
            grd1: l\_partner \notin contacted \cup contacted\_by \cup incoming\_sessions \cup outgoing\_sessions
            grd3: l\_partner \in RIU \cup (RBC \setminus accepting)
   then
            act1 : contacted\_by := contacted\_by \cup \{l\_partner\}
            act2: terminating\_sessions := terminating\_sessions \cup (RBC \cap (incoming\_sessions \cup
          outgoing_sessions))
   end
Event terminate communication \widehat{=}
extends terminate_communication
   any
```

```
l_partner
   where
            grd1: l\_partner \in incoming\_sessions \cup outgoing\_sessions
            grd2: l\_partner \in terminating\_sessions
   then
            act1:incoming\_sessions := incoming\_sessions \setminus \{l\_partner\}
            act2 : outgoing\_sessions := outgoing\_sessions \setminus \{l\_partner\}
            act3: terminating\_sessions := terminating\_sessions \setminus \{l\_partner\}
   end
Event make\_RBC\_accepting =
   any
          l_partner
   where
            grd1: l\_partner \in RBC
   then
            act1: accepting := accepting \cup \{l\_partner\}
   end
Event make\_RBC\_non\_accepting =
   any
          l_partner
   where
            grd1: l\_partner \in accepting
   then
            act1: accepting := accepting \setminus \{l\_partner\}
   end
END
```

# 5.8 Context 3 - ERTMS Levels

The third context introduces the notion of the different ERTMS and the notion of the mission status of a train. The modeled statuses are start of mission (SOM), end of mission (EOM) and the abstract notion of within a mission (MIS), i.e., anything between start and end of the current train mission. At this level of refinement, a more detailed modeling is not necessary.

```
CONTEXT c3_ERTMS_level
SETS

ERTMS_level
train_status
CONSTANTS

NTC
L0
L1
L2
L3
SOM start of mission
```

```
EOM end of mission

MIS while mission

AXIOMS

axm1: partition(ERTMS_level, {NTC, L0, L1, L2, L3})

axm2: partition(train_status, {SOM, EOM, MIS})

END
```

# 5.9 Machine 4 - ERTMS Level Changes

The next refined machine implements the different causes which can trigger the establishing of a communication session. The corresponding events refine the abstract "initiate\_session\_no\_contact" event. For this, the current ERTMS level of the train is tracked, as well as its current mission status.

The indication of a level change, a mission status change, a manual level change and an announced radio hole is modeled by events. These events modify the corresponding indicator variables to signal a change and they modify the corresponding state variables.

This can be illustrated using the manual level change event as example: the Boolean variable "signal\_manual\_level\_change" indicates that the driver manually changed the ERTMS level. It is changed by the "manual\_change\_level" event which is parametrized with the new level and which also modifies the "current\_level" variable which models the current ERTMS level. If the new level is 2 or 3, then the train is required to establish a communication session with trackside. This is realized in the "initiate\_session\_no\_contact\_manual\_change" event which reset the indication variable once the entity has been contacted. Similar events model the initiation because of non-manual level change, mission status change and announced radio holes.

#### Implemented Requirements

- The on-board shall establish a communication session (cf. §3.5.3.4)
  - at start of mission (only if level 2 or 3) (cf. §3.5.3.4.a)
  - if ordered from trackside (cf. §3.5.3.4.b)
  - If a mode change, not considered as an End of Mission, has to be reported to the RBC (only if level 2 or 3) (cf. §3.5.3.4.c)
  - If the driver has manually changed the level to 2 or 3 (cf. §3.5.3.4.d)
  - When the train front reaches the end of an announced radio hole (cf. §3.5.3.4.e)

```
REFINES m3_accept_system_version
SEES c3_ERTMS_level
VARIABLES

current_level
signal_level_change
current_status
signal_status_change
```

signal\_manual\_level\_change

position\_radio\_hole

# signal\_radio\_hole INVARIANTS

```
inv1 : current\_level \in ERTMS\_level
      inv2: signal\_level\_change \in BOOL
      inv3: current\_status \in train\_status
      inv4: signal\_status\_change \in BOOL
      inv5: signal\_manual\_level\_change \in BOOL
      inv6 : position\_radio\_hole \in BOOL
      inv7 : signal\_radio\_hole \in BOOL
EVENTS
Event manual_change_level =
   any
          l_level
   where
           grd1: l\_level \in ERTMS\_level
           grd2 : signal\_manual\_level\_change = FALSE
           grd3 : signal\_level\_change = FALSE
   then
           act1 : signal_manual_level_change := TRUE
           act2 : current_level := l_level
   end
Event change_level \widehat{=}
   any
   where l_level
           grd1: l\_level \in ERTMS\_level
           grd2 : signal\_manual\_level\_change = FALSE
           grd3 : signal\_level\_change = FALSE
   then
           act1 : current_level := l_level
           act2 : signal\_level\_change := TRUE
   end
extends initiate_session_no_contact
   any
          l_partner
   where
           grd5: l\_partner \notin incoming\_sessions \cup outgoing\_sessions \cup contacted \cup contacted\_by
           grd3: l\_partner \in on\_track
           grd6 : current\_status = SOM
           grd7 : current\_level \in \{L2, L3\}
   then
           act2 : contacted := contacted \cup \{l\_partner\}
   end
Event initiate_session_no_contact_manual_change \widehat{=}
extends initiate_session_no_contact
   any
          l_partner
```

```
where

    grd5 : l_partner ∉ incoming_sessions ∪ outgoing_sessions ∪ contacted ∪ contacted_by
    grd3 : l_partner ∈ on_track
    grd6 : current_level ∈ {L2, L3}
    grd7 : signal_manual_level_change = TRUE

then

act2 : contacted := contacted ∪ {l_partner}
    act3 : signal_manual_level_change := FALSE
end

END
```

#### 5.10 Machine 5 - Safe Radio Connection

The next machine refinement specifies handling of the safe radio connection which provides the necessary means to exchange protocol messages on a higher level. The existing established safe radio connections are represented by the set "ER\_connections" which holds values of type entities. Safe radio connections which must be established are modeled by the set "establish\_ER\_connections" while "terminated\_ER\_connections" holds those connections which timed-out. The indication variable "signal\_RBC\_border" signals the crossing of an RBC border which requires to establish a new safe radio connection with a new RBC.

Establishing a communication session then works as follows: if one of the conditions of §3.5.3.4 is fulfilled, the corresponding partner is added to the set "contacted". This is a pre-condition of the events which open a safe radio connection to a communication partner. The initiation message of the protocol is considered to be sent when a communication partner is both "contacted" and "ER\_connections" set. The reception of the system version and the sending of the system established message is modeled via the "receive\_information\_compatible" or "receive\_information\_incompatible" events. Established sessions with incompatible system versions are therefore in the intersection of the "sessions" and "terminating\_sessions" sets.

# Implemented Requirements

- Establish communication session after safe radio connection timeout (cf. §3.5.3.4.f)
- If the communication session is established by an OBU, it shall be preformed according to the following steps (cf. §3.5.3.7)
  - if part of ongoing start of mission procedure (cf. §3.5.3.7.a)
  - if safe radio connection is set up (cf. §3.5.3.7.i)
  - if order to terminate is received (cf. §3.5.3.7.ii)
  - if end of mission is performed (cf. §3.5.3.7.iii)
  - train passes level transition border (cf. §3.5.3.7.iv)
  - order to establish connection with different non-accepting RBC (cf. §3.5.3.7.v)
  - train passes RBC / RBC border (cf. §3.5.3.7.vi)
  - train enters announced radio hole (cf. §3.5.3.7.vii)
  - level 1 is left (RIU only) (cf. §3.5.3.7.viii)
- When the RBC initiates the establishing of a communication session, the OBU shall consider the session established after receiving the initiation message (cf. §3.5.3.10.c)

```
REFINES m4_level_changes
SEES c3 ERTMS level
VARIABLES
     ER_connections
     terminated_ER_connections
     establish_ER_connection
     signal_RBC_border
INVARIANTS
      inv1: terminated\_ER\_connections \subseteq on\_track
      inv2: establish\_ER\_connection \subseteq on\_track
      inv3: (incoming\_sessions \cup outgoing\_sessions) \subseteq ER\_connections
      inv4 : signal\_RBC\_border \in BOOL
EVENTS
Event incoming\_communication =
extends incoming_communication
   any
          l_partner
   where
           grd1: l\_partner \notin incoming\_sessions \cup outgoing\_sessions
           grd3: l\_partner \in on\_track \setminus RIU
           grd4: l\_partner \notin contacted
           grd5: l\_partner \notin contacted\_by
           grd6: l\_partner \in ER\_connections
   then
           act1:incoming\_sessions := incoming\_sessions \cup \{l\_partner\}
   end
Event receive_information_compatible \widehat{=}
extends receive_information_compatible
   any
          l_partner
   where
           grd3: l\_partner \in contacted
           grd4: l\_partner \in system\_version\_compatible
           grd5: l\_partner \in ER\_connections
   then
           act1: outgoing\_sessions := outgoing\_sessions \cup \{l\_partner\}
           act2 : contacted := contacted \setminus \{l\_partner\}
   end
Event receive\_information\_incompatible =
extends receive_information_incompatible
   any
          l_partner
   where
           grd3: l \ partner \in contacted
           grd4: l\_partner \notin system\_version\_compatible
           grd5: l\_partner \in ER\_connections
```

```
then
            act1: outgoing\_sessions := outgoing\_sessions \cup \{l\_partner\}
            act2 : contacted := contacted \setminus \{l\_partner\}
            act3: terminating\_sessions := terminating\_sessions \cup \{l\_partner\}
   end
Event receive_contact_order_accept =
                                              order to contact a RIU or accepting RBC
extends receive_contact_order_accept
   any
          l_partner
   where
            grd1: l\_partner \notin contacted \cup contacted\_by \cup incoming\_sessions \cup outgoing\_sessions
            grd3: l\_partner \in RIU \cup (RBC \cap accepting)
            grd4: l\_partner \notin terminated\_ER\_connections
   then
            act1: contacted\_by := contacted\_by \cup \{l\_partner\}
   end
Event receive_contact_order_non_accept =
extends receive_contact_order_non_accept
   any
          l_partner
   where
            grd1: l\_partner \notin contacted \cup contacted\_by \cup incoming\_sessions \cup outgoing\_sessions
            grd3: l\_partner \in RIU \cup (RBC \setminus accepting)
            grd4: l\_partner \notin terminated\_ER\_connections
   then
            act1: contacted\_by := contacted\_by \cup \{l\_partner\}
            act2: terminating sessions := terminating sessions \cup (RBC \cap (incoming sessions \cup
          outgoing_sessions))
   end
Event initiate_session_after_contact =
extends initiate_session_after_contact
   any
          l_partner
   where
            grd2: l\_partner \in contacted\_by
            grd3: l\_partner \notin terminated\_ER\_connections
   then
            act1: contacted := contacted \cup \{l\_partner\}
            act2: contacted\_by := contacted\_by \setminus \{l\_partner\}
            act3: establish\_ER\_connection := establish\_ER\_connection \cup \{l\_partner\}
   end
Event initiate_session_no_contact_SOM \widehat{=}
extends initiate_session_no_contact_SOM
          l_partner
   where
            grd5: l\_partner \notin incoming\_sessions \cup outgoing\_sessions \cup contacted \cup contacted\_by
```

```
grd3: l\_partner \in on\_track
            grd6: current\_status = SOM
            grd7 : current\_level \in \{L2, L3\}
            grd8: l\_partner \notin terminated\_ER\_connections
   then
            act2 : contacted := contacted \cup \{l\_partner\}
            act3: establish\_ER\_connection := establish\_ER\_connection \cup \{l\_partner\}
   end
Event initiate_session_no_contact_status_change \widehat{=}
extends initiate_session_no_contact_status_change
   any
          l partner
   where
            grd5: l\_partner \notin incoming\_sessions \cup outgoing\_sessions \cup contacted \cup contacted\_by
            grd3: l\_partner \in on\_track
            grd6 : current\_level \in \{L2, L3\}
            grd7 : signal\_status\_change = TRUE
            grd8 : current\_status \neq EOM
            grd9: l\_partner \notin terminated\_ER\_connections
   then
            act2: contacted := contacted \cup \{l\_partner\}
            act3 : signal\_status\_change := FALSE
            act4: establish\_ER\_connection := establish\_ER\_connection \cup \{l\_partner\}
   end
Event initiate_session_no_contact_manual_change \widehat{=}
extends initiate_session_no_contact_manual_change
   anv
          l_partner
   where
            grd5: l\_partner \notin incoming\_sessions \cup outgoing\_sessions \cup contacted \cup contacted\_by
            grd3: l\_partner \in on\_track
            grd6 : current\_level \in \{L2, L3\}
            grd7 : signal_manual_level_change = TRUE
            grd8: l\_partner \notin terminated\_ER\_connections
   then
            act2 : contacted := contacted \cup \{l\_partner\}
            act3 : signal_manual_level_change := FALSE
            act4 : establish\_ER\_connection := establish\_ER\_connection \cup \{l\_partner\}
   end
Event initiate session no contact leave radio hole \widehat{=}
extends initiate_session_no_contact_leave_radio_hole
          l_partner
   where
            grd5: l\_partner \notin incoming\_sessions \cup outgoing\_sessions \cup contacted \cup contacted\_by
            grd3: l\_partner \in on\_track
            grd6: position radio hole = FALSE
            grd7 : signal \ radio \ hole = TRUE
            grd8: l\_partner \notin terminated\_ER\_connections
```

```
then
            act2 : contacted := contacted \cup \{l\_partner\}
            act3 : signal\_radio\_hole := FALSE
            act4: establish\_ER\_connection := establish\_ER\_connection \cup \{l\_partner\}
   end
Event initiate_session_after_timeout \widehat{=}
extends initiate_session_no_contact
   any
          l_partner
   where
            grd5: l\_partner \notin incoming\_sessions \cup outgoing\_sessions \cup contacted \cup contacted\_by
            grd3: l\_partner \in on\_track
            grd6: l\_partner \in terminated\_ER\_connections
   then
            act2: contacted := contacted \cup \{l\_partner\}
            act3: terminated\_ER\_connections := terminated\_ER\_connections \setminus \{l\_partner\}
            act4: establish\_ER\_connection := establish\_ER\_connection \cup \{l\_partner\}
   end
Event establish_ER_connection_SOM \widehat{=}
   any
          l_partner
   where
            grd1: l\_partner \in contacted
            grd2: l\_partner \in establish\_ER\_connection
            grd3 : current\_status = SOM
   then
            act1: establish\_ER\_connection := establish\_ER\_connection \setminus \{l\_partner\}
            act2 : ER\_connections := ER\_connections \cup \{l\_partner\}
   end
Event establish_ER_connection \widehat{=}
   any
          l\_partner
   where
            grd1: l\_partner \in contacted
            grd2: l\_partner \in establish\_ER\_connection
            grd3: current status \neq SOM
   then
            act1: establish\_ER\_connection := establish\_ER\_connection \setminus \{l\_partner\}
            act2 : ER\_connections := ER\_connections \cup \{l\_partner\}
   end
Event est_perform_end_of_mission \widehat{=}
   any
          l_partner
   where
            grd1: l\_partner \in contacted
            grd2: l\_partner \in establish\_ER\_connection
            grd3 : signal\_status\_change = TRUE
            grd4 : current\_status = EOM
```

```
then
            act1: establish\_ER\_connection := establish\_ER\_connection \setminus \{l\_partner\}
   end
Event est\_terminate\_order =
extends drop_contact
   any
          l\_partner
   where
            grd1: l\_partner \in contacted
            grd2: l\_partner \in establish\_ER\_connection
            grd3 : current \ status \neq SOM
   then
            act1: contacted := contacted \setminus \{l\_partner\}
            act2: establish\_ER\_connection := establish\_ER\_connection \setminus \{l\_partner\}
   end
Event est\_pass\_level\_transition =
   any
          l_partner
   where
            grd1: l\_partner \in contacted
            grd2: l \ partner \in establish \ ER \ connection
            grd3 : signal\_level\_change = TRUE
            grd4 : current\_status \neq SOM
            grd5 : current\_level \in \{L0, L1, NTC\}
   then
            act1: establish\_ER\_connection := establish\_ER\_connection \setminus \{l\_partner\}
   end
Event est_pass_radio_hole \widehat{=}
   any
          l\_partner
   where
            grd1: l\_partner \in contacted
            grd2: l\_partner \in establish\_ER\_connection
            grd3 : signal\_radio\_hole = TRUE \land position\_radio\_hole = TRUE
            grd4 : current\_status \neq SOM
   then
            act1: establish\_ER\_connection := establish\_ER\_connection \setminus \{l\_partner\}
   end
Event est_RIU_leave_L1 =
   any
          l_partner
   where
            grd1: l\_partner \in contacted
            grd2: l\_partner \in RIU
            grd3 : signal\_level\_change = TRUE
            grd5 : current\_level \neq L1
            grd4 : current\_status \neq SOM
   then
```

```
act1: establish\_ER\_connection := establish\_ER\_connection \setminus \{l\_partner\}
   end
Event est\_RBC\_border =
   any
          l_partner
   where
            grd1: l\_partner \in contacted
            grd2: l\_partner \in RBC
            grd3 : signal\_RBC\_border = TRUE
            grd4 : current\_status \neq SOM
   then
            act1: establish\_ER\_connection := establish\_ER\_connection \setminus \{l\_partner\}
   end
Event indicate\_RBC\_border =
   any
   where l_flag
            grd1: l\_flag \in BOOL
   then
            act1 : signal\_RBC\_border := l\_flag
   end
Event est_other_RBC_non_accept \widehat{=}
   any
          l_partner
   where
            grd1: l\_partner \in contacted
            grd2: l\_partner \in RBC
            grd3 : RBC \cap accepting \cap contacted\_by \neq \emptyset
            grd4 : current\_status \neq SOM
   then
            act1: establish\_ER\_connection := establish\_ER\_connection \setminus \{l\_partner\}
   end
Event timeout\_ER\_connection =
extends drop_session
   any
          l_partner
   where
            grd1: l\_partner \in incoming\_sessions \cup outgoing\_sessions
            grd3: l\_partner \in ER\_connections
   then
            act1:incoming\_sessions := incoming\_sessions \setminus \{l\_partner\}
            act2 : outgoing\_sessions := outgoing\_sessions \setminus \{l\_partner\}
            act3 : ER\_connections := ER\_connections \setminus \{l\_partner\}
            act4: terminated\_ER\_connections := terminated\_ER\_connections \cup \{l\_partner\}
   end
Event terminate\_communication =
extends terminate communication
   any
```

```
l_partner
   where
            grd1: l\_partner \in incoming\_sessions \cup outgoing\_sessions
            grd2: l\_partner \in terminating\_sessions
            grd3: l\_partner \notin terminated\_ER\_connections
   then
            act1:incoming\_sessions := incoming\_sessions \setminus \{l\_partner\}
            act2 : outgoing\_sessions := outgoing\_sessions \setminus \{l\_partner\}
            act3: terminating\_sessions := terminating\_sessions \setminus \{l\_partner\}
            act4 : ER\_connections := ER\_connections \setminus \{l\_partner\}
   end
Event make\_RBC\_accepting =
extends make_RBC_accepting
   any
          l_partner
   where
            grd1: l\_partner \in RBC
   then
            act1: accepting := accepting \cup \{l\_partner\}
   end
Event make\_RBC\_non\_accepting =
extends make_RBC_non_accepting
   any
          l_partner
   where
            grd1: l\_partner \in accepting
   then
            act1: accepting := accepting \setminus \{l\_partner\}
   end
END
```

# References

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