

CFD General Notation System Advanced Data Format (ADF) User's Guide

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1 Introduction

Advanced Data Format (ADF) is a library of basic database management and I/O subroutines that implements a relatively simple hierarchical database. ADF is written in ANSI C to enhance portability of the software, and the design of the database allows for portability of files from platform to platform. There is also a Fortran interface. The files are self-describing (i.e., it is possible to browse files and determine their contents) and extensible (i.e., many different pieces of software on different platforms may add or modify information).

The routines allow the user to construct a tree structure with their data. (See Figure 1 on p. 4.) This structure is very similar to the directory structures of the UNIX or DOS operating systems. ADF also allows links between nodes within the same file or to different files. This feature works somewhat like "soft links" in the UNIX operating system. The major difference between the aforementioned directory structures and ADF is that the nodes not only contain information about their children (next lower-level nodes) but may also contain data.

The installation package includes the source code to ADF, the Fortran interface, sample files, and a simple file browser.

1.1 History of Project

ADF was developed as part of the CFD General Notation System(CGNS) project. The purpose of the CGNS project is to define the data models for Navier-Stokes based Computational Fluid Dynamics (CFD) technology, develop a set of standard interface data structures for those data models, and develop the software that will allow implementation of those data structures in existing and future CFD analysis tools. The CGNS system consists of a collection of conventions, and software conforming to these conventions, for the storage and retrieval of CFD data. Adherence to these conventions is intended to facilitate the exchange of CFD data between sites, between application codes such as solvers and grid generators, and across computing platforms.

Once the data models (called the Standard Interface Data Structures, or SIDS) were defined, a project was started to write software that could faithfully reproduce that information on disk. ADF is the result of that project. While ADF was developed specifically for the CFD process, it is quite general and has no built-in knowledge of CFD; therefore, it should be applicable to storing any type of data that lends itself to a hierarchical definition.

1.2 Other Software Interfaces and ADF

Two other software packages were investigated as part of this project. The first database interface investigated was the Hierarchical Data Format (HDF) developed at the National Center for Supercomputing Applications at the University of Illinois. This database system has a large user base and support and has been in existence more than 6 years with utilities and graphical routines written with both a C and a Fortran interface. The limitation of HDF was that it was not truly hierarchical despite its name. Any hierarchy has to be built using naming conventions. Since the CGNS data models indicated a natural hierarchical structure, it seemed appropriate to develop database software that worked in that mode by design. The ADF design considerations are summarized in Appendix E.

The second was the Common File Format (CFF) developed by McDonnell Douglas Aerospace. CFF is a second-generation database management system that provides a unifying file structure for CFD data. The purpose of the Common File is to insulate the user from the myriad of different

computer types that make up computer systems so that the user or application programmer may process a file from another machine without performing explicit conversions. CFF was written in Fortran; however, it was felt that portability and extensibility could be enhanced using C. Much of the experience gained from the McDonnell Douglas group is incorporated into ADF due to the cooperative efforts of personnel from the CFD group at McDonnell Douglas Aerospace.

1.3 Organization of Manual

The main section of this manual explains the basics of the hierarchical structure of an ADF database. In addition, the concept of a node as the basic building block of the hierarchy is developed in detail. The remaining sections of this manual are extensive. They provide a glossary of terms and conventions, as well as information related to the ADF version releases, version control numbering, and architectures that are supported by ADF. There are two examples, one in Fortran and one in C, in Appendix J and Appendix K respectively, that implement the structure illustrated in Figure 1 on p. 4. These examples should help familiarize the new user with ADF. Design considerations, file optimization, and portability issues are also discussed. The individual ADF core routines are described in detail in Appendix I. They are categorized into database-level routines, data structure and management subroutines, data query subroutines, data I/O subroutines, and some miscellaneous utility subroutines. Numerous examples are included to clarify the use of each subroutine. Appendix G provides a summary of ADF error messages, and Appendix H lists default values for various parameters and limits on dimensions of arrays.

2 The ADF Software Library

The ADF library is a hierarchical database system that is built around the concept of a "node". Each node contains information about itself and its ancestors and possibly data (e.g., arrays, vectors, character strings, etc.). Each of these nodes, in turn, can be connected to an arbitrary number of children, each of which is itself a node. In this system, an ADF node contains user-accessible information related to identification, name, type, and amount of data associated with it, and pointers to child nodes. Basic nodal information includes:

- a unique ID (essentially a file location pointer)
- a name (character string) used to describe the node and its data
- a label (character string) an additional field used to describe the node and its data. It is analogous to, but not exactly the same as, the name.
- information describing the type and amount of data
- data
- IDs of child nodes

There are no restrictions on the number of child nodes that a node can have associated with it in the ADF database. This structure allows the construction of a hierarchical database as shown in Figure 1 on p. 4. As illustrated in the figure, it is possible to reference nodes in a second file (ADF_File_Two) from the original file (ADF_File_One) . This is the concept of "linking."

A node knows about itself and its children, but it does not know anything about its parent. This means that it is possible to traverse "down" the tree by making queries about what lies below the current node, but it is not possible to traverse "up" the tree by making queries about nodes above a given node. If it is desired to move back up the tree, the user must keep track of that information.

All ADF files start with the root node, named "ADF MotherNode". This node is created automatically when a new file is opened. There is only one root node in an ADF file.

2.1 ADF Node Attributes

There is a single building block called a node (see Figure 1) used to construct the hierarchical structure. Each node may have zero to many subnodes that are associated with it, as well as its own data. The following are a list of attributes accessible by the user for a node in the ADF hierarchical database system.

Child Table A table of file pointers and names for each of the node's children.

Data The data associated with a node.

Data Type A 32-byte character field, blank filled, case sensitive. Specifies the type

of data (e.g., real, integer, complex) associated with this node. The

supported data types are listed in Table 1 on p. 5.

Dimensions An integer vector containing the number of elements within each di-

mension. For example, if the array A was declared (using Fortran) as A(10,20), the Dimension vector would contain two entries (10,20).

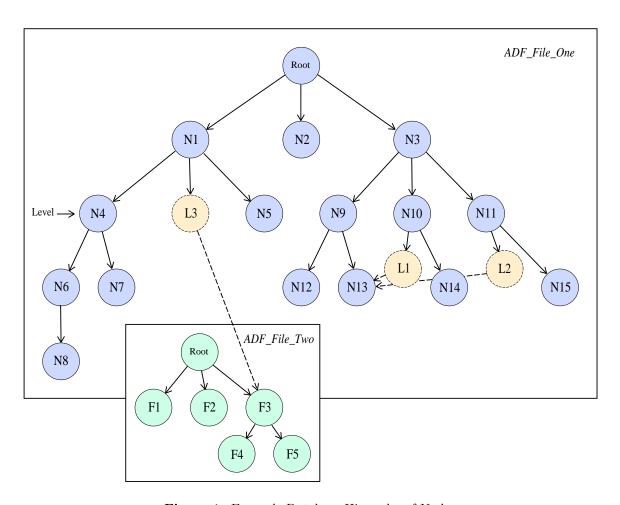


Figure 1: Example Database Hierarchy of Nodes

ID

A unique identifier to access a given node within a file. This field contains sufficient information for ADF to locate the node within a file. For any given node, the ID is generated only after the file it resides in has been opened by a program and the user requests information about the node. The ID is valid only within the program that opened the file and while that file is open. If the file is closed and reopened, the ID for any given node will be different. Within different programs, the node ID for the same node will be different. The ID is never actually written into a file.

Label

A 32-byte character field. The rules for Labels are identical to those for Names. Unlike names, Labels do not have to be unique. The Label field was introduced to allow "data typing" similar to the "typedef" concept in C. Using the Label field in this way allows programs to know some additional information about the use of the node itself or its child nodes and to call specific subroutines to read the data or react in specific ways upon detection of the type.

Name

A 32-byte character field. The names of child nodes directly attached

Table 1: Data Types

Data Type	Notation
No Data	MT (i.e., eMpTy)
Integer 32	14
Integer 64	18
Unsigned Integer 32	U4
Unsigned Integer 64	U8
Real 32	R4
Real 64	R8
Complex 64	X4
Complex 128	X8
Character (unsigned byte)	C1
Byte (unsigned byte)	B1
Link	LK

to a parent node must be unique. For example, in Figure 1, all nodes directly attached to N3 must have unique names. When a request to create a new node is made, ADF checks the requested name against the other names of the child nodes of the specified parent. If the requested name is not unique, ADF returns an error.

Legal characteristics of a name are a A-Z, a-z, 0-9, and special characters (ASCII values from 32 to 126, except for the forward slash "/" (ASCII number 47)). Names will be blank filled to 32 bytes; they are case sensitive. Leading blanks are discarded and trailing blanks are ignored, whereas internal blanks are significant.

Note: Names passed to ADF from C must have the null "\0" character appended to them. Names returned from ADF through the C interface will have the null character appended to them. Therefore, C programs should allocate 33 bytes for any Name in order to accommodate the null character.

Fortran programs can allocate 32 characters for Names. The Fortran interface takes care of adding or removing the null character as required.

Names of Subnodes

A list of names of the subnodes (children) of a node. (This is the information contained in the child table.)

Number of Dimensions

The dimensionality of the data. ADF views all data as an array and can handle from zero (i.e., no data) to 12 dimensions. A "0" is used if the data type is empty. Thus, a scalar is viewed as a vector with one dimension and length 1.

Number of Subnodes

The number of child nodes directly attached to any given node. Each node can have zero or more child nodes directly associated with it.

Pointer

An address, from the point of view of a programming language. Pointers are like jumps, leading from one part of the data structure to another.

2.1.1 Data Type Definitions

Structure

It is possible to define a "structure" in ADF similar to the way that "struct" is defined in the C programming language. ADF will treat each instance of the structure as a single instance of data. A structure is described by using a string. For example: "I4,I4,R8,R4".

Notes:

- Structures can be only as complex as can be described in 32 characters.
- This construct is not very portable; therefore, its use is highly discouraged.

Character

A character type is intended for ASCII-type character information. There may be different system architectures that use different representations for translating data. Byte data type should be used for pure byte or bit data usage.

Precision

The R4 and R8 are single and double precision, respectively, on a 32-bit system architecture. On the Cray, single precision is R8 and double precision may be R8 or R16, depending on the compile settings. ADF tracks the number of bits to guarantee precision.

Link

A link is denoted by "LK" in Table 1 on p. 5 and defines the linkage between nodes and subnodes. A link provides a mechanism for referring to a node that physically resides in a different part of the hierarchy or a different file. A link within ADF parallels a soft link in the UNIX operating system in that it does not guarantee that the referenced node exists. ADF will "resolve" the link only when information is requested about the linked node.

2.2 Acquiring the Software and Documentation

The ADF software is distributed as part of the CGNS Library, available from SourceForge, at http://sourceforge.net/projects/cgns. This manual, as well as the other CGNS documentation, is available in both HTML and PDF format from the CGNS documentation web site, at http://www.grc.nasa.gov/www/cgns/.

A ADF Glossary of Terms

ADF The initialism (acronym) for Advanced Data Format.

Child One of the subnodes of a Parent. A child node does not have knowledge of its

parent node. The user must keep track of this relationship.

Database A hierarchy of ADF nodes. By use of links, it may physically span multiple files.

File An ADF file, which a single root node and its underlying structure.

ID A unique identifier to access a given node within a file. This 8-byte field contains

sufficient information for ADF to locate the node within a file. For any given node, the ID is generated only after the file it resides in has been opened by a program and the user requests information about the node. The ID is valid only within the program that opened the file and while that file is open. If the file is closed and reopened, the ID for any given node will be different. Within different programs, the node-ID for the same node will be different. The ID is

never actually written into a file.

Link-Node A special type of node. Links are created using the ADF_Link subroutine. The data type of this node is LK, and its data is a one-dimensional array containing the

name of the file (if other than the current file) containing the node to be linked and the full path name in that file from the root node to the desired node.

Links provide a mechanism for referring to a node that physically resides in a different part of the hierarchy. The node pointed to by a link may or may not reside in the same file as the link itself. A link within ADF is very similar to a "soft" link in the UNIX operating system in that it does not guarantee that the referenced node exists. ADF will "resolve" the link only when information is requested about the node. If the ID of a link-node is used in an ADF call, the effect of the call is the same as if the ID of the linked-to node was used. Note that a link node does not have children itself. In Figure 1 on p. 4, the children seen for L3 are F4 and F5. If a child is "added" to L3, then in reality, the child is added to F3. There are specialized subroutines provided to create link nodes and

The single component used to construct an ADF database.

Node name A node has a 32-character name. Every child node directly under a given parent

must have a unique name. Legal characteristics in a name are A–Z, a–z, 0–9, and special characters (ASCII values from 32 to 126, omitting the forward slash "/", ASCII number 47). Names will be blank filled to 32 bytes; they are case sensitive. Leading blanks are discarded and trailing blanks are ignored, whereas

internal blanks are significant.

extract the link details.

Node

Parent A node that has subnodes directly associated with it.

Pathname Within a database, nodes can be referenced using the name of a node along with its parent ID, or by using a "pathname" whose syntax is roughly the same as a path name in the UNIX environment. A pathname that begins with a leading slash "/" is assumed to begin at the root node of the file. If no leading slash is given,

the name is assumed to begin at the node specified by the parent ID. Although there is a 32-character limitation on the node Name, there is no restriction on

the length of the path name. For example, equivalent ways to refer to node ${\tt N8}$ in Figure 1 are:

- $\bullet\,$ Node-ID for N6 and name = "N8"
- \bullet Node-ID for N4 and name = "N6/N8"
- Node-ID for N1 and name = "N4/N6/N8"
- \bullet Node-ID for the Root_Node and name = "/N1/N4/N6/N8"

B History of ADF Version Releases

This section contains a list of all the versions of ADF Software Libraries that have been released since the first release in December 1995. Along with each release, we include a brief synopsis of the improvements made in this version. In the future, newer versions may be implemented and released and they will also be accessible on the World Wide Web.

ADF A01 (December 1995)	This was the first release of the software. It includes the majority of the subroutines that are summarized in this document, and it was internally distributed for use at Boeing and to members of the working group (NASA and McDonnell Douglas).
ADF A02, A03 (July 1996)	This release includes improvements and some bug fixes for the ${\rm A}01$ version.
ADF B01 (September 1996)	This release includes the first release of the ADF library and with the following subroutines added: Delete a Node, Add a Name (Change) on a Node, Change the Parent Node (move a Child).
ADF C00 (October 1996)	This release includes the performance enhancements and support for more platforms (Cray T90, SGI 6.2, DEC Alpha, Intel Paragon).
ADF AXX (Future Release)	It is anticipated that there will be future releases of the ADF core libraries, and they will include at a minimum the following additional routines (plus others that will provide enhancements as they are needed): Delete an Existing Database and Garbage Collection. Garbage Collection will redistribute the ADF file to use free-space that is not located at the end of the file.

C ADF File System Architectures

The following platform architectures have been tested and used both for functionality testing of the ADF core software libraries and for testing and running the prototype.

Table 2: Platform Architectures

Release	Machine	OS Version	Native Format
A01	Cray	Unicos 8.0	Unicos 8.0
A01	SGI/IRIS	4.0.5	IEEE Big Endian ¹
B01	HP	9.05	IEEE Big Endian
B01	SGI/IRIS	5.03	IEEE Little Endian
C00	Intel Paragon	_	IEEE Little Endian
C00	Dec Alpha	_	IEEE Little Endian
C00	SGI/IRIS	6.2	IEEE Big Endian
C00	Cray T90	Unicos 9.02	Cray Format

¹In the table, "Endian" refers to the ordering of bytes in a multi-byte number. Big endian is a computer architecture in which, within a given multi-byte numeric representation, the most significant byte has the lowest address (the word is stored "big-end-first"). Little endian is a computer architecture in which, within a given 16- or 32-bit word, bytes at lower addresses have lower significance (the word is stored "little-end-first").

D ADF File Version Control Numbering

The ADF file version control number scheme is described below. The format for the version number is a field of six digits or characters:

AXXxxx

where:

A Major revision number. Major internal structure changes. This number is not expected to change very often, if at all, because the backward compatibility is only available by explicit policy decision.

One alphabetic character.

Range of values: A-Z, a-z

In the unlikely event of reaching **z**, then use any other unused printable ASCII character, except the blank or symbols used by the "what" command

@()#~>\

XX Minor revision number. New features, minor changes, and bug fixes. Backward but not forward compatible.

Two-digit hexadecimal number (uppercase letters).

Range of values: 00-FF

Reset to 00 with changes in major revision number.

Incremental number. Incremented with every new version of the library (even if no changes are made to the file format). Files are forward and backward compatible.

Three-digit hexadecimal number (lowercase letters).

Range of values: 000-fff

Does not reset.

The following definitions are used:

Forward compatible Older versions of libraries can read and write to files created by newer

versions of libraries.

Backward compatible Newer versions of libraries can read and write to files created by older

versions of libraries.

E ADF Design Considerations

This section provides a summary of the design considerations that were used in the construction of the ADF software library.

E.1 ADF File Header Information

Every ADF file has a header section that contains information about the file itself. The following information from this header is available to the user:

- ADF version number of the library that created the database.
- File creation date and time.
- File modification date and time.
- Data format used in database (IEEE big endian, IEEE little endian, etc.)

[Under data format, "endian" refers to the ordering of bytes in a multi-byte number. Big endian is a computer architecture in which, within a given multi-byte numeric representation, the most significant byte has the lowest address (the word is stored "big-end-first"). Little endian is a computer architecture in which, within a given 16- or 32-bit word, bytes at lower addresses have lower significance (the word is stored "little-end-first").]

E.2 ADF File Optimizations

To optimize the performance of ADF, the following techniques have been incorporated into the ADF software to enhance performance of the file:

- Use block-based, unbuffered (raw mode) I/O where available, with block sizes of 4096 bytes.
- Align medium to large data chunks on block boundaries.
- If the data size is equal to or greater than 2048, then
 - align data to the next block and
 - add extra to free (garbage) lists
- Avoid, where possible, small- to medium-sized chunks of data that span block boundaries.
- Align, where possible, to the next block and add extra to free (garbage) lists.

Allow data space to grow by linking data chunks. It is possible to increase the last dimension; doing so will extend the data. However, internal dimensional changes will corrupt the existing data.

The pointer table for child nodes will also contain the child node names.

E.3 ADF File Portability

To address code portability and future needs, the following design decisions were made:

- Use larger than 32-bit file pointers to allow for files larger than 4 Gigabytes. (C routine lseek may not handle this, but ADF files should allow it.)
- Use 48-bit pointers within a block of data. 32-bit pointers to 4096-byte blocks and a 16-bit pointer to a position within the block. This allows for files with 4 gigabyte blocks, in other words $2^{32} \times 4096 = 17.5922 + 12$ bytes or 17.59 Tera bytes.
- The ID pointers will be 64-bit coded IDs. Users may use double data type (real*8). This is parceled as follows: 32 bits for block number, 12 bits for block offset, and 20 bits for file identifier.
- Encode the integer information, other than data, in ASCII, Hex-based notation.

E.4 ADF File Error Checking

Error checking has been implemented for the ADF file and is summarized here. Each item in the ADF file will have item-specific boundary tags surrounding it to provide file-based corruption checking. For variable-sized items, the associated boundary tags will include file-based size information. Information will be written to the disk in a sequence that will not allow corrupt files. For example, when adding a new child node, the complete child information will be written before the parent's child-table is updated. An ADF-core subroutine for downloading data to disk will be provided.

E.5 ADF Source Code Considerations

The ADF library of source code will incorporate the use of UNIX "what" strings for the ADF version number and also RCS versioning information in the source code and in the object code. The source code is written in portable ANSI C, using POSIX-defined system calls.

E.6 ADF Node Header Information

The following information is contained in an ADF node header:

- node boundary tag
- name (32 characters)
- label (32 characters)
 - number of subnodes (num_sub_nodes) is an integer (ASCII, Hex encoded in file).
 - address pointer of the subnodes indicated by the variable.
- sub_nodes, which are pointers (ASCII, Hex encoded in file) to a table (the "child table") of file pointers and names for each of the node's children. Note that the child name information is redundant and included for performance.
- data type is specified in data_type (32 characters allowed).

- dimensionality of the data, called num_dimensions. It is ASCII, Hex encoded in the file.
- dimension values are listed in the integer array of 12, dimension_values. The dimension values are Hex encoded.
- integer value for the number of data chunks is found in num_data_chunks (ASCII, Hex encoded).
- data address, which is as follows:
 - if $num_data_chunks = 1$, then a file pointer to data
 - if num_data_chunks > 1, then a file pointer to a table of num_data_chunks file pointers and associated data sizes
- ending node boundary tag.

E.7 Fortran Character Array Portability Concerns

Fortran character arrays are different from any other array type because they inherently include declared length information. Abstractly, they are a compound type: an array and an integer. The ANSI standards do not specify the implementation mechanism for handling this data type and so it is left to the vendor. As one might expect, vendors have devised different policies. This is particularly evident in how argument lists are created and used. The matter is further complicated when writing functions in other languages that are to be called from Fortran.

To keep the interfaces simple and to keep the Fortran and C data I/O calls similar (as opposed to having separate data I/O functions for character data), ADF suggests abiding by the following rules (these are required for Cray T90-mode users):

- Do not pass character arrays as the actual data arguments to any ADF read or write function unless that node has been defined with a data type of C1.
- If a node has been defined as data type C1, then pass character arrays only as the actual data arguments to all ADF read and write function.

Note to Cray T90-mode users: The above rules must be followed. In addition, the given node must be available and have its data type correctly defined. Error handling is not possible otherwise, and ADF will abort or fail regardless of how the error state flag is set.

E.8 Integer 64 Data Type Portability Concerns

For portability reasons, it is suggested that the use of the I8 data type be restricted to a 64-bit environment.

E.9 Compound Data Types Portability Concerns

For the transportability reasons discussed below, use of compound data types is not recommended.

When using compound data types (e.g., with structures in C), it is important to be aware of data alignment issues. If one is not careful, the actual size of the structure in memory may be larger

than the sum of the individual members. The total size depends on the order and word boundary alignment requirements of the specific data types. This is platform and compiler dependent and not handled by ADF. In order to provide the greatest portability (at least up to 64-bit environments), it is recommended that

- 8-byte data types (I8, R8) be aligned on 8 byte boundaries, and
- data types smaller than the word size be padded to a size equal to the word size.

So a 4-byte data type (e.g., I4[1]) needs "padding" (e.g., I4[2]) if it is to be followed by an 8-byte data type. And assuming a word size of 4 bytes, all C1-data-type elements need dimension values of multiples of 4 bytes (e.g., C1[4], C1[8], etc.). To be even more careful, size everything in multiples of 8 bytes, for example, "C1[8], I4[2], C1[16], R4[6], R8[5], I8[1]".

For a given architecture and compiler, and taking into consideration the restrictions given above, compound data types should work. It is more portable and highly recommended that users write out the individual components of a structure into separate nodes. It would probably be best to copy the individual components of a list of structures into an appropriate array type and write the temporary array out using the write-all or write-strided routines.

F ADF Conventions and Implementations

С

All input strings are to be null terminated. All returned strings will have the trailing blanks removed and will be null terminated. Variables declared to hold Names, Labels, and Data-Types should be at least 33 characters long. *ADF.h* has a number of variables defined. An example declaration would be:

char name[ADF_NAME_LENGTH+1];

Fortran

Strings will be determined using inherited length. Returned strings will be blank filled to the specified length. All returned names will be left justified and blank filled on the right. There will be no null character. An example declaration would be:

PARAMETER ADF_NAME_LENGTH=32 CHARACTER*(ADF_NAME_LENGTH) NAME

ID

A unique identifier to access a given node within a file. This 8-byte field contains sufficient information for ADF to locate the node within a file. For any given node, the ID is generated only after the file it resides in has been opened by a program and the user requests information about the node. The ID is valid only within the program that opened the file and while that file is open. If the file is closed and reopened, the ID for any given node will be different. Within different programs, the node ID for the same node will be different. The ID is not ever actually written into a file.

The declaration for variables that will hold node IDs should be for an 8-byte real number.

The ID is actually a 64-bit combination of a system-generated file index along with the block and offset of the location of the node on the disk. In general, users do not need to know the internal coding of this information.

error_return

The error code for the ADF routines is the following:

-1 No error.

 $n \ (> 0)$ ADF error code. The routine ADF_Error_Message is used to get a textual description of the error.

0 ADF should never return the value zero.

Indexing

All indexing is Fortran-like in that the starting index is 1 and the last is \mathbb{N} for N items in an index or array dimension. The array structure is assumed to be the same as in Fortran with the first array dimension varying the fastest and the last dimension varying the slowest.

The index starting at one is used in ADF_Read_Data, ADF_Write_Data, and ADF_Children_Names.

The user should be aware of the differences in array indexing between Fortran and C. The subroutines ADF_Read_All_Data and ADF_Write_All_Data merely take a pointer to the beginning of the data, compute how much data is to be read/written, and process as many bytes as have been requested. Thus, these routines effectively make a copy of memory onto disk or vice versa. Given this convention, it is possible for a C program to use standard C conventions for array indexing and use ADF_Write_All_Data to store the array on disk. Then a Fortran program might use ADF_Read_All_Data to read the data set. Unless

the user is aware of the structure of the data, it is possible for the array to be transposed relative to what is expected.

The implications of the assumed array structure convention can be quite subtle. The subroutines ADF_Write_Data and ADF_Read_Data assume the Fortran array structure in order to index the data. Again, unless the user is aware of the implications of this, it is possible to write an array on disk and later try to change a portion of the data and not change the correct numbers.

As long as users are aware of how their data structure maps onto ADF, there will not be any problems.

G ADF Error Messages

Table 3: ADF Error Messages

Number	Error Message
-1	No Error
1	Integer number is less than a given minimum value
2	Integer value is greater than given maximum value
3	String length of zero of blank string detected
4	String length longer than maximum allowable length
5	String length is not an ASCII-Hex string
6	Too many ADF files opened
7	ADF file status was not recognized
8	ADF file open error
9	ADF file not currently opened
10	ADF file index out of legal range
11	Block/offset out of legal range
12	A string pointer is null
13	FSEEK error
14	FWRITE error
15	FREAD error
16	Internal error: Memory boundary tag bad
17	Internal error: Disk boundary tag bad
18	File Open Error: NEW - File already exists ²
19	ADF file format was not recognized
20	Attempt to free the RootNode disk information
21	Attempt to free the FreeChunkTable disk information
22	File Open Error: OLD - File does not $exist^3$
23	Entered area of unimplemented code
24	Subnode entries are bad
25	Memory allocation failed
26	Duplicate child name under a parent node
27	Node has no dimensions
28	Node's number of dimensions is not in legal range
29	Specified child is not a child of the specified parent
30	Data-Type is too long
31	Invalid Data-Type
32	A pointer is null

 $Continued\ on\ next\ page$

 $^{^{2}}$ The user is trying to create a new file and give it a name. The system has responded that the name has already been used.

 $^{^{3}}$ The user wants to open an existing file that supposedly has the given name. The system has responded that no file by that name exists.

Table 3: ADF Error Messages (Continued)

Number	Error Message
33	Node had no data associated with it
34	Error zeroing out of memory
35	Requested data exceeds actual data available
36	Bad end value
37	Bad stride values
38	Minimum value is greater than maximum value
39	The format of this machine does not match a known signature ⁴
40	Cannot convert to or from an unknown native format
41	The two conversion formats are equal; no conversion done
42	The data format is not supported on a particular machine
43	File close error
44	Numeric overflow/underflow in data conversion
45	Bad start value
46	A value of zero is not allowable
47	Bad dimension value
48	Error state must be either a 0 (zero) or a 1 (one)
49	Dimensional specifications for disk and memory are unequal
50	Too many link levels are used; may be caused by a recursive link
51	The node is not a link. It was expected to be a link.
52	The linked-to node does not exist
53	The ADF file of a linked node is not accessible
54	A node ID of 0.0 is not valid
55	Incomplete data when reading multiple data blocks
56	Node name contains invalid characters
57	ADF file version incompatible with this library version
58	Nodes are not from the same file
59	Priority stack error
60	Machine format and file format are incomplete
61	Flush error

⁴When ADF wakes up, it tries to figure out what data format the machine it is running on uses (e.g., IEEE big endian, IEEE little endian, Cray). If it doesn't recognize the format, it can't convert files created on other platforms to the current one, so it issues this error message and punts.

H Default Values and Sizes and Limits of Dimensions and Arrays

The following default values, sizes, and limits are defined in the header file ADF.h.

Table 4: Default Values and Sizes

Attribute	Limit, Size, or Value
Data type length	32-byte character field
Date length	32-byte character field
File name length	1024-byte character field
Format length	20-byte character field
Label length	32-byte character field
Maximum link depth	100
Maximum dimension	12
Maximum length of error string	80 characters
Maximum link data size	4096
Name length	32-byte character field
Length of status	32-byte character field
Version length	32-byte character field
Maximum children	None
Pointer size	48-bit pointer is used for a block size of 4096
	32-bit pointer is used for blocks less than 4096
	16-bit pointer is used for blocks within blocks

I ADF Library of Subroutines

This appendix contains a listing of all the ADF routines for release 2.0. Each is described in detail for both C and Fortran use. There are examples, hints, and diagnostics included in some of the synopses and discussions of the routines.

I.1 Major Functional Groupings of Subroutines

The routines are summarized under the following major functional groupings:

Database-Level Routines

С	Fortran	Description	Page
ADF_Database_Close	ADFDCLO	Close an opened database	31
ADF_Database_Delete	ADFDDEL	Delete an existing database ⁵	32
ADF_Database_Get_Format	ADFDGF	Get the data format in existing database	33
ADF_Database_Open	ADFDOPN	Open a database	29
ADF_Database_Set_Format	ADFDSF	Set the data format in an existing database	34

Data Structure and Management Routines

С	Fortran	Description	Page
ADF_Create	ADFCRE	Create a node	35
ADF_Delete	ADFDEL	Delete a node	39
ADF_Children_Names	ADFCNAM	Get the child names of a node	43
ADF_Number_of_Children	ADFNCLD	Get the number of children of a node	47
ADF_Get_Node_ID	ADFGNID	Get a unique identifier of a node	48
ADF_Get_Name	ADFGNAM	Get the name of a node	51
ADF_Put_Name	ADFPNAM	Put (change) the name of a node	52
ADF_Move_Child	ADFMOVE	Change a parent (move a child)	55
ADF_Link	ADFLINK	Create a link	59
ADF_Is_Link	ADFISLK	Test if a node is a link	66
ADF_Get_Link_Path	ADFGLKP	Get the path information from a link	68
ADF_Get_Root_ID	ADFGRID	Get the root ID of the ADF file	71

 $^{^5\}mathrm{Not}$ implemented in the current release

Data Query Routines

С	Fortran	Description	Page
ADF_Get_Label	ADFGLB	Get a label	73
ADF_Set_Label	ADFSLB	Set a label	75
ADF_Get_Data_Type	ADFGDT	Get the data type	76
ADF_Get_Number_of_Dimensions	ADFGND	Get the number of dimensions	79
ADF_Get_Dimension_Values	ADFGDV	Get the dimension values of a node	80
ADF_Put_Dimension_Information	ADFPDIM	Set the data type and dimension information	81

Data I/O Routines

С	Fortran	Description	Page
ADF_Read_Data	ADFREAD	Read the data from a node (with partial capabilities)	83
ADF_Read_All_Data	ADFRALL	Read all the data into a contiguous memory space	90
ADF_Read_Block_Data	ADFRBLK	Read a contiguous block of data from a node	91
ADF_Write_Data	ADFWRIT	Write the data to a node (with partial capabilities)	92
ADF_Write_All_Data	ADFWALL	Write all the data from a contiguous memory space	100
ADF_Write_Block_Data	ADFWBLK	Write a contiguous block of data to a node	101

Miscellaneous Routines

C	Fortran	Description	Page
ADF_Flush_to_Disk	ADFFTD	Flush the data to a disk	102
ADF_Database_Garbage_Collection	ADFDGC	Garbage collection ⁶	103
ADF_Error_Message	ADFERR	Return an error message	104
ADF_Set_Error_State	ADFSES	Set the error state	105
ADF_Get_Error_State	ADFGES	Get the error state	107
ADF_Database_Version	ADFDVER	Get the ADF file version ID	109
ADF_Library_Version	ADFLVER	Get the ADF library version ID	111

⁶Not implemented in the current release

Alphabetical Listing of Subroutines

All the routines are listed alphabetically below, according to their C name, along with the page number for each routine.

C	Fortran	Description	Page
ADF_Children_Names	ADFCNAM	Get the child names of a node	43
ADF_Create	ADFCRE	Create a node	35
ADF_Database_Close	ADFDCLO	Close an opened database	31
ADF_Database_Delete	ADFDDEL	Delete an existing database ⁷	32
ADF_Database_Garbage_Collection	ADFDGC	Garbage collection ⁸	103
ADF_Database_Get_Format	ADFDGF	Get the data format in existing database	33
ADF_Database_Open	ADFDOPN	Open a database	29
ADF_Database_Set_Format	ADFDSF	Set the data format in an existing database	34
ADF_Database_Version	ADFDVER	Get the ADF file version ID	109
ADF_Delete	ADFDEL	Delete a node	39
ADF_Error_Message	ADFERR	Return an error message	104
ADF_Flush_to_Disk	ADFFTD	Flush the data to a disk	102
ADF_Get_Data_Type	ADFGDT	Get the data type	76
ADF_Get_Dimension_Values	ADFGDV	Get the dimension values of a node	80
ADF_Get_Error_State	ADFGES	Get the error state	107
ADF_Get_Label	ADFGLB	Get a label	73
ADF_Get_Link_Path	ADFGLKP	Get the path information from a link	68
ADF_Get_Name	ADFGNAM	Get the name of a node	51
ADF_Get_Node_ID	ADFGNID	Get a unique identifier of a node	48
ADF_Get_Number_of_Dimensions	ADFGND	Get the number of dimensions	79
ADF_Get_Root_ID	ADFGRID	Get the root ID of the ADF file	71
ADF_Is_Link	ADFISLK	Test if a node is a link	66
ADF_Library_Version	ADFLVER	Get the ADF library version ID	111
ADF_Link	ADFLINK	Create a link	59
ADF_Move_Child	ADFMOVE	Change a parent (move a child)	55
ADF_Number_of_Children	ADFNCLD	Get the number of children of a node	47
ADF_Put_Dimension_Information	ADFPDIM	Set the data type and dimension information	81
ADF_Put_Name	ADFPNAM	Put (change) the name of a node	52
ADF_Read_All_Data	ADFRALL	Read all the data into a contiguous memory space	90
ADF_Read_Block_Data	ADFRBLK	Read a contiguous block of data from a node	91

Continued on next page

⁷Not implemented in the current release ⁸Not implemented in the current release

С	Fortran	Description	Page
ADF_Read_Data	ADFREAD	Read the data from a node (with partial capabilities)	83
ADF_Set_Error_State	ADFSES	Set the error state	105
ADF_Set_Label	ADFSLB	Set a label	75
ADF_Write_All_Data	ADFWALL	Write all the data from a contiguous memory space	100
ADF_Write_Block_Data	ADFWBLK	Write a contiguous block of data to a node	101
ADF_Write_Data	ADFWRIT	Write the data to a node (with partial capabilities)	92

I.3 Database-Level Routines

ADF_Database_Open — Open a Database

ADF_Database_Open (filename, status, format, root_ID, error_return)			
Language	C	Fortran	
Routine Name	ADF_Database_Open	ADFDOPN	
Input	const char *filename	character*(*) filename	
	const char *status character*(*) status		
	const char *format	character*(*) format	
Output	double *root_ID	real*8 root_ID	
	int *error_return	integer error_return	

filenameA legal file name that may include a relative or absolute path where it is directly

usable by the C fopen() system routine (no environment expansion is done).

statusSimilar to a Fortran OPEN() status. Input is required: there is no default. Allowable values are:

> READ ONLY File must exist; writing is not allowed.

OLD File must exist; reading and writing are allowed.

NEW File must not exist.

SCRATCH Temporary new file is created with a system name, and *filename*

is ignored. The temporary file is deleted when the program exits

or the file is closed.

UNKNOWN OLD if file exists or else NEW is used.

Specifies the numeric format for the file. This field is used only when a file is created and is ignored when status = OLD. Allowable values are:

NATIVE Use the native numeric format of the computer that creates the file. This is the default for a new file if the input string for format is null. Note that if the native numeric format is not one of the supported formats listed here, then the file

cannot be read on machines using any other format.

Use the IEEE big endian format. IEEE_BIG Use the IEEE little endian format. IEEE_LITTLE CRAY Use the native CRAY format.

 $root_ID$ The root identity of the database.

format

Error return code. (See Appendix G.) $error_return$

This routine, ADF_Database_Open, opens a new or existing database. If links to other ADF files exist in the current file, they will be opened only as required. Using this routine is similar to opening a file in Fortran with the corresponding clarifiers, such as whether it is READ_ONLY, OLD, NEW, or named as SCRATCH file.

The format of the file, which is ignored when the status of the file is OLD, is used when the file is first created. Big endian is a binary format in which the most significant byte or bit comes first,

whereas in little endian, the most significant byte or bit comes last. To specify the format more explicitly, you can use the following formats:

```
IEEE_BIG_32
IEEE_BIG_64
IEEE_LITTLE_32
IEEE_LITTLE_64
```

for IEEE big or little endian formats.

Example

This example opens a new database using the native format of the host computer. Note that the default format is specified by using the empty string. In the C programming language, a null string could have been used.

```
PROGRAM TEST

CHARACTER*(80) MSG

REAL*8 RID

INTEGER IERR

CALL ADFDOPN('db.adf','NEW',' ',RID,IERR)

IF (IERR .GT. 0) THEN

CALL ADFERR(IERR,MSG)

PRINT *,MSG

STOP

ENDIF

STOP

END
```

${\tt ADF_Database_Close} - Close \ a \ Database$

ADF_Database_Close (root_ID,error_return))			
Language	C	Fortran	
Routine Name	ADF_Database_Close ADFDCLO		
Input	Input const double *root_ID real*8 root_ID		
Output	int *error_return	integer error_return	

root_ID The root identification of the database. This can be a valid node ID for this database.

error_return Error return code. (See Appendix G.)

This routine, ADF_Database_Close, closes an existing database, as well as the other ADF files that may be attached through links. For example, if there is another ADF file that is opened and linked to this database, only the file stream associated with this database will be closed. This routine is similar to the close of a file in Fortran.

Example

This example closes a database. Note that while the root ID is used in the call to ADFDCLO, any valid node ID for this file will work. Also, in general, it is not necessary to close open ADF files when the program exits normally. The standard shutdown procedures will flush all buffers and bring files up to date. The primary use of ADFDCLO is to clean up file tables or to release unused files.

```
PROGRAM TEST

CHARACTER*(80) MSG

REAL*8 RID

INTEGER IERR

CALL ADFDOPN('db.adf','NEW',' ',RID,IERR)

IF (IERR .GT. 0) THEN

CALL ADFERR(IERR,MSG)

PRINT *,MSG

STOP

ENDIF

.
...do useful stuff (hopefully)

.
CALL ADFDCLO(RID,IERR)

STOP

END
```

${\tt ADF_Database_Delete} - Delete \ a \ File$

ADF_Database_Delete (filename,error_return))		
Language	C	Fortran
Routine Name	ADF_Database_Delete	ADFDDEL
Input	const char *filename	character*(*) filename
Output	int *error_return	integer error_return

filename A legal file name of an existing ADF database. The filename may include a

relative or absolute path where it is directly usable by the C fopen() system

routine (no environment expansion is done).

This routine, ADF_Database_Delete, deletes an existing database file. It does not delete links referenced in the database. This routine is similar to the deletion of a file in Fortran.

Note: This routine will be implemented in a future release.

 ${\tt ADF_Database_Get_Format} \ -- \ \mathit{Get the Data Format}$

ADF_Database_Get_Format (root_ID,format,error_return))		
Language	age C Fortran	
Routine Name	ADF_Database_Get_Format	ADFDGF
Input	const double *root_ID	real*8 root_ID
Output	char *format	character*(*) format
	int *error_return	integer error_return

root_ID Any valid node ID for the given ADF file.

format The format for the file.

error_return Error return code. (See Appendix G.)

This routine, ADF_Database_Get_Format, gets the data format for an existing database.

Example

This example opens an existing ADF database, creates a new node, and then uses the node ID for the new node to ask what the file type is. Note that the file *format* is ignored because the database already exists.

```
PROGRAM TEST
   CHARACTER*(80) MSG
   CHARACTER*(32) FORM
   REAL*8 RID, CID
   INTEGER IERR
   CALL ADFDOPN('db.adf','OLD',' ',RID,IERR)
   IF (IERR .GT. 0) THEN
      CALL ADFERR (IERR, MSG)
      PRINT *,MSG
      STOP
   ENDIF
   CALL ADFCRE(RID, 'junk_node', CID, IERR)
   CALL ADFDGF(CID, FORM, IERR)
   PRINT *,'FILE FORMAT = ',FORM
   STOP
   END
```

${\tt ADF_Database_Set_Format} \ -- \ Set \ the \ Data \ Format$

ADF_Database_Set_Format (root_ID,format,error_return))		
Language	ge C Fortran	
Routine Name	ADF_Database_Set_Format	ADFDSF
Input	const double *root_ID	real*8 root_ID
Output	char *format	character*(*) format
	int *error_return	integer error_return

 ${\it root_ID}$ The root identity of the database.

format The numeric format for the file.

This routine, ADF_Database_Set_Format, sets the data format in an existing database.

Note: Use with extreme caution. This routine is needed only for the data conversion utilities and not intended for the general user.

I.4 Data Structure and Management Routines

ADF_Create — Create a Node

ADF_Create (PID, name, ID, error_return)		
Language	C	Fortran
Routine Name	ADF_Create	ADFCRE
Input	const double PID	real*8 PID
	const char *name	character*(*) name
Output	double *ID	real*8 ID
	int *error_return	integer error_return

PID ID of the parent node of the created child node.

name Name of the parent node.

ID The ID of the newly created child node.

This routine, ADF_Create, creates a new node (not a link) as a child of a given parent node.

Default node header values in this new node are:

- label = blank
- number of subnodes = 0
- datatype = MT
- number of dimensions = 0
- data = NULL

Example

This example opens a database and creates a node under the root node. Note that the default values for a newly created node are label = ' ', datatype = MT, dimension = null, data = none. These are reset as required using the routines ADFSLB, ADFPDIM and ADFWALL/ADFWRIT. Note also that the root node is named "ADF MotherNode". This name is generated when the database is first opened. If desired, it could be reset using ADFPNAM.

```
PROGRAM TEST
C
PARAMETER (MAXCHR=32)
C
CHARACTER*(MAXCHR) NODNAM
C
C *** NODE IDS
C
REAL*8 RID,PID,CID
INTEGER I,J,IERR,NUMCLD
```

```
С
C *** OPEN DATABASE
C
      CALL ADFDOPN('db.adf','NEW',' ',RID,IERR)
С
C *** CREATE NODES AT FIRST LEVEL
\mathsf{C}
      DO 150 I = 1,3
         WRITE(NODNAM, '(A7, I1)')'PARENT.', I
         CALL ADFCRE(RID, NODNAM, PID, IERR)
С
C ***** CREATE NODES AT SECOND LEVEL
С
         NUMCLD = I*I
         DO 110 J = 1, NUMCLD
            WRITE(NODNAM, '(A6,I1,A1,I1)')'CHILD.',I,'x',J
            CALL ADFCRE(PID, NODNAM, CID, IERR)
       CONTINUE
  110
C ***** PRINT NODE NAMES JUST CREATED
         CALL PRTCLD(PID)
  150 CONTINUE
С
C *** PRINT NAMES OF NODES ATTACHED TO ROOT NODE
С
      CALL PRTCLD(RID)
\mathsf{C}
      STOP
      END
C
C ******** SUBROUTINES **********
C
      SUBROUTINE ERRCHK(IERR)
С
C *** CHECK ERROR CONDITION
С
      CHARACTER*80 MESS
      IF (IERR .GT. 0) THEN
         CALL ADFERR(IERR, MESS)
         PRINT *, MESS
         CALL ABORT('ADF ERROR')
      ENDIF
      RETURN
      END
      SUBROUTINE PRTCLD(PID)
C
C *** PRINT TABLE OF CHILDREN GIVEN A PARENT NODE-ID
C
      PARAMETER (MAXCLD=10)
```

```
PARAMETER (MAXCHR=32)
         REAL*8 PID
         CHARACTER*(MAXCHR) NODNAM, NDNMS(MAXCLD)
         CALL ADFGNAM(PID, NODNAM, IERR)
         CALL ERRCHK(IERR)
         CALL ADFNCLD(PID, NUMC, IERR)
         CALL ERRCHK(IERR)
         WRITE(*,120)NODNAM,NUMC
     120 FORMAT(/, ' PARENT NODE NAME = ',A,/,
        Х
                      NUMBER OF CHILDREN = ',12,/,
              ,
        X
                      CHILDREN NAMES: ')
        NLEFT = NUMC
         ISTART = 1
         --- TOP OF DO-WHILE LOOP
     130 CONTINUE
            CALL ADFCNAM(PID, ISTART, MAXCLD, LEN(NDNMS),
        Х
                         NUMRET, NDNMS, IERR)
           CALL ERRCHK(IERR)
            WRITE(*,140)(NDNMS(K),K=1,NUMRET)
     140
            FORMAT(2(8X,A))
            NLEFT = NLEFT - MAXCLD
            ISTART = ISTART + MAXCLD
         IF (NLEFT .GT. 0) GO TO 130
         RETURN
         END
The resulting output is:
  PARENT NODE NAME = PARENT.1
      NUMBER OF CHILDREN = 1
       CHILDREN NAMES:
          CHILD.1.1
  PARENT NODE NAME = PARENT.2
      NUMBER OF CHILDREN = 4
       CHILDREN NAMES:
          CHILD.2.1
                                                   CHILD.2.2
          CHILD.2.3
                                                   CHILD.2.4
  PARENT NODE NAME = PARENT.3
      NUMBER OF CHILDREN = 9
       CHILDREN NAMES:
          CHILD.3.1
                                                   CHILD.3.2
          CHILD.3.3
                                                   CHILD.3.4
          CHILD.3.5
                                                   CHILD.3.6
          CHILD.3.7
                                                   CHILD.3.8
          CHILD.3.9
  PARENT NODE NAME = ADF MotherNode
      NUMBER OF CHILDREN = 3
       CHILDREN NAMES:
```

PARENT.1 PARENT.2 PARENT.3

ADF_Delete — Delete a Node

ADF_Delete (PID,ID,error_return)		
Language	C	Fortran
Routine Name	ADF_Delete	ADFDEL
Input	const double PID	real*8 PID
	const double ID	real*8 ID
Output	int *error_return	integer error_return

PID The ID of the node's parent.
ID The ID of the node to use.
error_return Error return code. (See Appendix G.)

In general, this routine, ADF_Delete, deletes a node and all of its children. Given the starting node, a recursive search is done down the hierarchy, deleting all nodes. If a "link" is encountered during the deletion (i.e., the specified node or any of its children), then the link is deleted, and the downward search stops. That is, the link information is deleted, but not the actual node it refers to.

To understand the deletion of a node that is a link, it must be remembered that a link is merely a reference to another node. Therefore, the deletion of a node that is a link is the deletion of that reference, not the referred node itself. The reason for this is that a link may actually point to data in another file that may be owned by another user. Therefore, it would not be proper for ADF to try to delete that node. Therefore, ADF stops at the link.

When a node is deleted, any links that reference it are left "dangling." In other words, the existing links to the node still reference the node, but if ADF is asked to resolve that reference, it will determine that the referred to node doesn't exist and will return an error flag.

Note that the parent ID of the node to be deleted is required. This is due to the fact that child nodes do not know the ID of their parent node. Thus when a node is deleted, in order for the child table of the parent to be updated properly, the parent ID must be supplied as an input.

Example

This example opens a database and creates three nodes attached to the root node. It also generates nodes to each of these base nodes. Then one of the base nodes is deleted. Not only is the node "PARENT.2" deleted; all of its children are deleted at the same time.

```
PROGRAM TEST
C
PARAMETER (MAXCHR=32)
C
CHARACTER*(MAXCHR) NODNAM
C
C *** NODE IDS
C
REAL*8 RID,PID,CID
```

```
INTEGER I, J, IERR, NUMCLD
С
C *** OPEN DATABASE
C
      CALL ADFDOPN('db.adf','NEW',' ',RID,IERR)
C
C *** CREATE NODES AT FIRST LEVEL
      DO 150 I = 1,3
         WRITE(NODNAM, '(A7,I1)')'PARENT.',I
         CALL ADFCRE(RID, NODNAM, PID, IERR)
C ***** CREATE NODES AT SECOND LEVEL
C
         NUMCLD = I*I
         DO 110 J = 1, NUMCLD
            WRITE(NODNAM, '(A6, I1, A1, I1)')'CHILD.', I, '.', J
            CALL ADFCRE(PID, NODNAM, CID, IERR)
         CONTINUE
  150 CONTINUE
C *** PRINT NAMES OF NODES ATTACHED TO ROOT NODE
C
      CALL PRTCLD(RID)
C *** PRINT NAMES OF CHILDREN UNDER PARENT.2
C
      CALL ADFGNID(RID, 'PARENT.2', PID, IERR)
      CALL PRTCLD(PID)
С
C *** NOW DELETE PARENT.2
      CALL ADFDEL(RID, PID, IERR)
      CALL PRTCLD(RID)
C
C *** JUST FOR GRINS, LOOK FOR CHILDREN UNDER ORIGINAL ID
С
      CALL ADFGNID(RID, '/PARENT.2/CHILD.2.1', CID, IERR)
      CALL ERRCHK(IERR)
C
      STOP
      END
С
C ******* SUBROUTINES **********
C
      SUBROUTINE ERRCHK(IERR)
C
C *** CHECK ERROR CONDITION
С
      CHARACTER*80 MESS
      IF (IERR .GT. 0) THEN
```

```
CALL ADFERR(IERR, MESS)
            PRINT *, MESS
            CALL ABORT('ADF ERROR')
         ENDIF
         RETURN
         END
         SUBROUTINE PRTCLD(PID)
  С
  C *** PRINT TABLE OF CHILDREN GIVEN A PARENT NODE-ID
  С
         PARAMETER (MAXCLD=10)
         PARAMETER (MAXCHR=32)
         REAL*8 PID
         CHARACTER*(MAXCHR) NODNAM, NDNMS(MAXCLD)
         CALL ADFGNAM(PID, NODNAM, IERR)
         CALL ERRCHK(IERR)
         CALL ADFNCLD(PID, NUMC, IERR)
         CALL ERRCHK(IERR)
         WRITE(*,120)NODNAM,NUMC
     120 FORMAT(/,' PARENT NODE NAME = ',A,/,
                    NUMBER OF CHILDREN = ',12,/,
        X
               ,
        X
                      CHILDREN NAMES: ')
        NLEFT = NUMC
         ISTART = 1
         --- TOP OF DO-WHILE LOOP
     130 CONTINUE
            CALL ADFCNAM(PID, ISTART, MAXCLD, LEN(NDNMS),
        Х
                         NUMRET, NDNMS, IERR)
            CALL ERRCHK(IERR)
            WRITE(*,140)(NDNMS(K),K=1,NUMRET)
     140
           FORMAT(2(8X,A))
            NLEFT = NLEFT - MAXCLD
            ISTART = ISTART + MAXCLD
         IF (NLEFT .GT. 0) GO TO 130
         RETURN
         END
The resulting output is:
  PARENT NODE NAME = ADF MotherNode
       NUMBER OF CHILDREN = 3
       CHILDREN NAMES:
          PARENT.1
                                                   PARENT.2
          PARENT.3
  PARENT NODE NAME = PARENT.2
      NUMBER OF CHILDREN = 4
       CHILDREN NAMES:
          CHILD.2.1
                                                   CHILD.2.2
          CHILD.2.3
                                                   CHILD.2.4
```

PARENT NODE NAME = ADF MotherNode
NUMBER OF CHILDREN = 2
CHILDREN NAMES:
PARENT.1

PARENT.3

ADF 29: Specified child is NOT a child of the specified parent. IOT Trap $\mbox{\sc Abort}$ - core dumped

ADF_Children_Names — Get the Names of the Child Nodes

ADF_Children_Names (PID,istart,imax_num,imax_name_length,inum_ret,names, error_return)		
Language	C Fortran	
Routine Name	ADF_Children_Names	ADFCNAM
Input	const double PID	real*8 PID
	const int istart	integer istart
	const int imax_num	integer imax_num
	const int imax_name_length	integer imax_name_length
Output	int *inum_ret	integer inum_ret
	char *names	character*(*) names
	int *error_return	integer error_return

PID The ID of the parent node to use.

istart The nth child's name (to start with the first, use istart = 1).

imax_num The maximum number of names to return.

imax_name_length The number of characters allocated to hold the name of each child node.

inum_ret The number of names returned.

names The names of the children.

This routine, ADF_Children_Names, returns the child names directly associated with a parent node. The names of the children are not guaranteed to be returned in any particular order. For example, if four child nodes were created in the order: node1, node2, node3, node4, when the call to ADFCNAM is made, there is no guarantee that the order of the node names in the character array names will be the same.

The reason for not guaranteeing node ordering has to do with efficient use of disk space. Although the concept of "linked lists" works fine in central memory, it is not particularly efficient on disk. Therefore, static tables are used to maintain parent/child lists. The order in which children names are returned is the order found in the static table. If a child node is deleted, an empty slot is created and will be used by the next child node created under that parent node.

The indexing of a child is Fortran-like and begins with 1, but as noted above, this does not imply a notion of node ordering. To start with a child node listed as the first index, use an *istart* value of 1.

C Programming Notes

• Node names can be up to 32 characters. ADF appends the null (\0) character to the end of each node name. Therefore, at least 33 characters should be allocated for each node name to be returned.

• imax_name_length is used to "stride" through the character array passed to ADF. For example, if 50 characters per name were allocated, the first character of the first name would be in the first allocated byte (position [0]), and the second name would start in the 51st byte (position [50]) and so on. The length of the individual names can be determined using the strlen function.

Fortran Programming Notes

• Node names can be up to 32 characters. The returned names are left justified and blank filled within the array. No null (\0) character is appended to the name; therefore, an appropriate declaration would be CHARACTER*(32).

Example

This example creates three nodes attached to the root node. It then creates a varying number of child nodes under each of the base nodes. Lastly, it queries the database to find out how many children were created and then gets the child names. The thing to notice in this example is that the character array NDNMS is not large enough to hold all the names under PARENT.3. The array NDNMS will hold only five names at a time; therefore, a loop was set up to read a subset of the name list during each pass. Note also that the call to ADFCNAM may request more names than are present. If this occurs, the routine will return all that is available in the output array names and return that number in the <code>inum_ret</code> variable.

```
PROGRAM TEST
C
      PARAMETER (MAXCLD=5)
      PARAMETER (NDATA=10)
      PARAMETER (MAXCHR=32)
C
      CHARACTER*(MAXCHR) NODNAM, NDNMS(MAXCLD)
С
      RID - ROOT ID
C
      CID - CHILD ID
С
      PID - PARENT ID
C
      REAL*8 RID, CID, PID
      INTEGER I,J,K,IERR,NUMCLD,NLEFT,ISTART
C *** OPEN DATABASE
      CALL ADFDOPN('DB.ADF', 'NEW', '', RID, IERR)
      CALL ERRCHK(IERR)
C
C *** CREATE NODES AT FIRST LEVEL
      DO 150 I = 1,3
         WRITE(NODNAM, '(A7, I1)')'PARENT.', I
         CALL ADFCRE(RID, NODNAM, PID, IERR)
         CALL ERRCHK(IERR)
C ***** CREATE NODES AT SECOND LEVEL
```

```
С
            NUMCLD = I*I
            DO 110 J = 1, NUMCLD
               WRITE(NODNAM, '(A6, I1, A1, I1)')'CHILD.', I, '.', J
               CALL ADFCRE(PID, NODNAM, CID, IERR)
               CALL ERRCHK(IERR)
     110
            CONTINUE
   С
   C ***** GET NUMBER AND NAMES OF CHILDREN JUST CREATED
   С
            AND PRINT THEM OUT
   С
            CALL ADFGNAM(PID, NODNAM, IERR)
            CALL ERRCHK(IERR)
            CALL ADFNCLD(PID, NUMC, IERR)
            CALL ERRCHK(IERR)
            WRITE(*,120)I,NODNAM,NUMC
     120
            FORMAT(' LEVEL = ',12,' PARENT NODE NAME = ',A,/,
                      NUMBER OF CHILDREN = ',12,/,
        X
                   ,
                         CHILDREN NAMES: ')
            NLEFT = NUMC
            ISTART = 1
            --- TOP OF DO-WHILE LOOP
   C
     130
            CONTINUE
               CALL ADFCNAM(PID, ISTART, MAXCLD, LEN(NDNMS),
        X
                            NUMRET, NDNMS, IERR)
               CALL ERRCHK(IERR)
               PRINT *,' FETCHED: ',NUMRET,' NAMES'
               WRITE(*,140)(NDNMS(K),K=1,NUMRET)
     140
               FORMAT(8X,A)
               NLEFT = NLEFT - MAXCLD
               ISTART = ISTART + MAXCLD
            IF (NLEFT .GT. 0) GO TO 130
     150 CONTINUE
         STOP
         END
   C
   C ******* SUBROUTINES **********
         SUBROUTINE ERRCHK(IERR)
         CHARACTER*80 MESS
         IF (IERR .GT. 0) THEN
            CALL ADFERR(IERR, MESS)
            PRINT *, MESS
            CALL ABORT ('ADF ERROR')
         ENDIF
         RETURN
         END
The resulting output is:
```

PARENT = 1 PARENT NODE NAME = PARENT.1

```
NUMBER OF CHILDREN = 1
   CHILDREN NAMES:
    FETCHED
                    1 NAMES
      CHILD.1.1
PARENT = 2 PARENT NODE NAME = PARENT.2
   NUMBER OF CHILDREN = 4
   CHILDREN NAMES:
             4 NAMES
    FETCHED
      CHILD.2.1
      CHILD.2.2
      CHILD.2.3
      CHILD.2.4
PARENT = 3 PARENT NODE NAME = PARENT.3
   NUMBER OF CHILDREN = 9
   CHILDREN NAMES:
    FETCHED
             5 NAMES
      CHILD.3.1
      CHILD.3.2
      CHILD.3.3
      CHILD.3.4
      CHILD.3.5
    FETCHED
                4 NAMES
      CHILD.3.6
      CHILD.3.7
      CHILD.3.8
      CHILD.3.9
```

ADF_Number_of_Children — Get the Number of Children Nodes

ADF_Number_of_Children (PID,num_children,error_return)		
Language	nguage C Fortran	
Routine Name	ADF_Number_of_Children	ADFNCLD
Input	const double PID	real*8 PID
Output	int *num_children	integer num_children
	int *error_return	integer error_return

PID The ID of the parent node to use.

 $num_children$ The number of children directly associated with this node.

This routine, ADF_Number_of_Children, returns the number of child names directly associated with a parent node.

Example

See the example for ADF_Children_Names.

ADF_Get_Node_ID — Get the ID of a Child Node

ADF_Get_Node_ID (PID,name,ID,error_return)		
Language	C	Fortran
Routine Name	ADF_Get_Node_ID	ADFGNID
Input	const double PID	real*8 PID
	const char *name	character*(*) name
Output	double *ID	real*8 ID
	int *error_return	integer error_return

PID The ID of the parent node.
name The name of the node.
ID The ID of the named node.
error_return Error return code. (See Appendix G.)

This routine, ADF_Get_Node_ID, returns the ID of a child node, given the parent node ID and the name of the child node. To return the ID of the root node in an ADF file, use any known ID in the ADF file and a name of "/". The syntax for *name* is essentially the same as the "path name" within the UNIX operating system. It might look like /level.1/level.2/node. The name of the node may be one of two forms. If the name begins with a "/", then the name is relative to the root node for the associated database. If the name does not begin with a "/", then the name is relative to the parent node associated with the given PID.

To return the ID of the root node an ADF file, use any known ID in the ADF file and a name of "/".

Example

This example illustrates the various ways to access a node. Note that when a full path is specified (i.e., a leading slash "/" is specified), all that is required for the ID is any valid ID for the associated database. It will probably be clearer to others if the root ID is used in that situation.

```
PROGRAM TEST
C
      PARAMETER (MAXCHR=32)
C
      CHARACTER*(MAXCHR) NODNAM
С
C
      RID - ROOT ID
C
      AL1ID - LEVEL 1 ID
C
      AL2ID - LEVEL 2 ID
C
      AL3ID - LEVEL 3 ID
С
      CID - CHILD ID
      REAL*8 RID, AL1ID, AL2ID, AL3ID, CID
```

```
С
C *** OPEN DATABASE
C
      CALL ADFDOPN('DB.ADF', 'NEW', '', RID, IERR)
      CALL ERRCHK(IERR)
С
C *** CREATE NODE AT FIRST LEVEL
      CALL ADFCRE(RID, 'LEVEL.1', AL1ID, IERR)
      CALL ERRCHK(IERR)
С
C *** CREATE NODE AT SECOND LEVEL
С
      CALL ADFCRE(AL1ID, 'LEVEL.2', AL2ID, IERR)
      CALL ERRCHK(IERR)
C
C *** CREATE NODE AT THIRD LEVEL
С
      CALL ADFCRE(AL2ID, 'LEVEL.3', AL3ID, IERR)
      CALL ERRCHK(IERR)
C
C *** EQUIVALENT WAYS TO GET THE LOWER LEVEL NODE ID
C ***** FULL PATH NAME
С
      CALL ADFGNID(RID, '/LEVEL.1/LEVEL.2/LEVEL.3', CID, IERR)
      CALL ERRCHK(IERR)
      CALL ADFGNAM(CID, NODNAM, IERR)
      PRINT *,' '
      PRINT *, ' FULL PATH EXAMPLE: ROOT NODE ID: NODE NAME = ', NODNAM
C
C ***** FULL PATH NAME - GIVEN ANY VALID NODE ID FOR FILE
C
      CALL ADFGNID(AL3ID, '/LEVEL.1/LEVEL.2/LEVEL.3', CID, IERR)
      CALL ERRCHK(IERR)
      CALL ADFGNAM(CID, NODNAM, IERR)
      PRINT *,' '
      PRINT *,' FULL PATH EXAMPLE - VALID NODE ID: NODE NAME = ', NODNAM
C
C ***** PARTIAL PATH NAME
С
      CALL ADFGNID(AL1ID, 'LEVEL.2/LEVEL.3', CID, IERR)
      CALL ERRCHK(IERR)
      CALL ADFGNAM(CID, NODNAM, IERR)
      PRINT *,' '
      PRINT *, ' PARTIAL PATH EXAMPLE: NODE NAME = ', NODNAM
C ***** DIRECT USE OF PARENT ID
С
      CALL ADFGNID(AL2ID, 'LEVEL.3', CID, IERR)
      CALL ERRCHK(IERR)
```

```
CALL ADFGNAM(CID, NODNAM, IERR)
         PRINT *,' '
         PRINT *, ' GIVEN PARENT NAME EXAMPLE: NODE NAME = ', NODNAM
         END
   С
   C ******** SUBROUTINES ***********
         SUBROUTINE ERRCHK(IERR)
         CHARACTER*80 MESS
         IF (IERR .GT. 0) THEN
            CALL ADFERR(IERR, MESS)
            PRINT *, MESS
           CALL ABORT('ADF ERROR')
         ENDIF
         RETURN
         END
The resulting output is:
   FULL PATH EXAMPLE: ROOT NODE ID: NODE NAME = LEVEL.3
   FULL PATH EXAMPLE - VALID NODE ID: NODE NAME = LEVEL.3
   PARTIAL PATH EXAMPLE: NODE NAME = LEVEL.3
   GIVEN PARENT NAME EXAMPLE: NODE NAME = LEVEL.3
```

ADF_Get_Name — Get the Name of a Node

ADF_Get_Name (ID,name,error_return)		
Language	\mathbf{C}	Fortran
Routine Name	ADF_Get_Name	ADFGNAM
Input	const double ID	real*8 ID
Output	char *name	character*(*) name
	int *error_return	integer error_return

ID The ID of the node to use.

name The simple name of the node.

This routine, ADF_Get_Name, returns the 32-character name of a node, given the node's ID. In C, the name will be null terminated after the last nonblank character; therefore, 33 characters should be used (32 for the name, plus 1 for the null). In Fortran, the null character is not returned; therefore, the character variable declaration for name should be for 32 characters (e.g., CHARACTER*(32) NAME).

Example

See the example for ADF_Get_Node_ID.

 ${\tt ADF_Put_Name} \ -- \ Put \ a \ Name \ on \ a \ Node$

ADF_Put_Name (PID,ID,name,error_return)			
Language	e C Fortran		
Routine Name	ADF_Put_Name	ADFPNAM	
Input	const double PID	real*8 PID	
	const double ID	real*8 ID	
	const char *name	character*(*) name	
Output	int *error_return	integer error_return	

PID The ID of the node's parent.
ID The ID of the node to use.
name The new name of the node.
error_return Error return code. (See Appendix G.)

This routine, ADF_Put_Name, changes the name of a node.

Note that the parent ID of the node to be deleted is required. This is due to the fact that child nodes do not know the ID of their parent node. Thus, when a node is deleted, in order for the child table of the parent to be updated properly, the parent ID must be supplied as an input.

Warning: If the node is pointed to by a link node, changing the node's name will break the link. Example

This example illustrates the creation of a node with an initial name. Later, the name of the node is changed. The routine ADFCNAM is used to get the new name from the parent's information table.

```
PROGRAM TEST
C
      PARAMETER (MAXCHR=32)
C
      CHARACTER*(MAXCHR) NODNAM, CLDNAM
C
С
      RID - ROOT ID
С
      CID - CHILD ID
C
      REAL*8 RID, AL1ID, AL2ID, AL3ID, CID
C
C *** OPEN DATABASE
      CALL ADFDOPN('db.adf','NEW',' ',RID,IERR)-*9+
C
      CALL ERRCHK(IERR)
C *** CREATE NODE
```

```
С
         CALL ADFCRE(RID, 'LEVEL.1', CID, IERR)
         CALL ERRCHK(IERR)
   C
   C *** GET NODE NAME AND CHECK PARENTS TABLE
         CALL ADFGNAM(CID, NODNAM, IERR)
         CALL ERRCHK(IERR)
         PRINT *,' '
         PRINT *, ' NODE NAME = ', NODNAM
         CALL ADFCNAM(RID, 1, 1, LEN(CLDNAM),
                      NUMRET, CLDNAM, IERR)
         CALL ERRCHK(IERR)
         PRINT *, ' NODE NAME IN PARENTS TABLE = ', CLDNAM
   C
   C *** CHANGE THE NODE NAME
         CALL ADFPNAM(RID, CID, 'NEW_NAME', IERR)
         CALL ERRCHK(IERR)
   С
   C *** GET NEW NODE NAME AND CHECK PARENTS TABLE
   С
         CALL ADFGNAM(CID, NODNAM, IERR)
         CALL ERRCHK(IERR)
         PRINT *,' '
         PRINT *, ' NEW NODE NAME = ', NODNAM
         CALL ADFCNAM(RID, 1, 1, LEN(CLDNAM),
                      NUMRET, CLDNAM, IERR)
         CALL ERRCHK(IERR)
         PRINT *, ' NODE NAME IN PARENTS TABLE = ', CLDNAM
         STOP
         END
   C
   C ******* SUBROUTINES **********
   С
         SUBROUTINE ERRCHK(IERR)
         CHARACTER*80 MESS
         IF (IERR .GT. 0) THEN
            CALL ADFERR(IERR, MESS)
            PRINT *, MESS
            CALL ABORT('ADF ERROR')
         ENDIF
         RETURN
         END
The resulting output is:
   NODE NAME
                 = LEVEL.1
   NODE NAME IN PARENTS TABLE = LEVEL.1
   NEW NODE NAME = NEW_NAME
```

NODE NAME IN PARENTS TABLE = NEW_NAME

ADF_Move_Child — Move a Child Node to a Different Parent

ADF_Move_Child (PID,ID,NPID,error_return)		
Language	C Fortran	
Routine Name	ADF_Move_Child	ADFMOVE
Input	const double PID	real*8 PID
	const double ID	real*8 ID
	double NPID	real*8 NPID
Output	int *error_return	integer error_return

PID ID of the node's current parent.

ID of the node to use.

NPID ID of the node's new parent.

This routine, ADF_Move_Child, deletes the given child node from the current parent's child table and adds it to the new parent's child table. ADF_Move_Child is restricted to moves within the same physical file. If the node is pointed to by a link-node, moving the node's name will break the link.

Example

This example creates a simple hierarchy. It then picks up a node from its original parent and moves it to a new parent node.

```
PROGRAM TEST
С
      PARAMETER (MAXCHR=32)
С
      CHARACTER*(MAXCHR) NODNAM
C
C *** NODE IDS
      REAL*8 RID,PID,CID,PID1,PID3
      INTEGER I, J, IERR
С
C *** OPEN DATABASE
С
      CALL ADFDOPN('db.adf','NEW',' ',RID,IERR)
      CALL ERRCHK(IERR)
C *** CREATE NODES AT FIRST LEVEL
      WRITE(*,100)
  100 FORMAT(/,' *** ORIGINAL DATABASE SETUP ***')
      DO 150 I = 1,3
```

```
WRITE(NODNAM,'(A7,I1)')'PARENT.',I
         CALL ADFCRE(RID, NODNAM, PID, IERR)
         CALL ERRCHK(IERR)
C ***** CREATE NODES AT SECOND LEVEL
         NUMCLD = I*I
         DO 110 J = 1, NUMCLD
            WRITE(NODNAM, '(A6, I1, A1, I1)')'CHILD.', I, '.', J
            CALL ADFCRE(PID, NODNAM, CID, IERR)
            CALL ERRCHK(IERR)
  110
         CONTINUE
С
C ***** GET NUMBER AND NAMES OF CHILDREN JUST CREATED
         AND PRINT THEM OUT
C
         CALL PRTCLD(PID)
  150 CONTINUE
C *** PICK UP NODE /PARENT.3/CHILD.3.4 AND MOVE IT
С
      TO /PARENT.1
С
      CALL ADFGNID (RID, 'PARENT.3', PID3, IERR)
      CALL ERRCHK(IERR)
C
      CALL ADFGNID (PID3, 'CHILD.3.4', CID, IERR)
      CALL ERRCHK(IERR)
C
      CALL ADFGNID(RID, 'PARENT.1', PID1, IERR)
      CALL ERRCHK(IERR)
С
      CALL ADFMOVE(PID3,CID,PID1,IERR)
      CALL ERRCHK(IERR)
C *** CHECK TO MAKE SURE THE NODE WAS ACTUALLY MOVED
      WRITE(*,160)
  160 FORMAT(/, '*** PARENT.1 AND PARENT.3 AFTER MOVE ***')
      CALL PRTCLD(PID1)
      CALL PRTCLD(PID3)
C
      STOP
      END
C
C ******** SUBROUTINES **********
C
      SUBROUTINE ERRCHK(IERR)
C
C *** CHECK ERROR CONDITION
      CHARACTER*80 MESS
```

```
IF (IERR .GT. 0) THEN
            CALL ADFERR(IERR, MESS)
            PRINT *, MESS
            CALL ABORT('ADF ERROR')
         ENDIF
         RETURN
         END
         SUBROUTINE PRTCLD(PID)
  C
  C *** PRINT TABLE OF CHILDREN GIVEN A PARENT NODE-ID
  С
         PARAMETER (MAXCLD=10)
         PARAMETER (MAXCHR=32)
         REAL*8 PID
         CHARACTER*(MAXCHR) NODNAM, NDNMS(MAXCLD)
         CALL ADFGNAM(PID, NODNAM, IERR)
         CALL ERRCHK(IERR)
         CALL ADFNCLD(PID, NUMC, IERR)
         CALL ERRCHK(IERR)
         WRITE(*,120)NODNAM, NUMC
     120 FORMAT(/, ' PARENT NODE NAME = ',A,/,
        χ,
                     NUMBER OF CHILDREN = ',12,/,
        X
                      CHILDREN NAMES: ')
         NLEFT = NUMC
         ISTART = 1
        --- TOP OF DO-WHILE LOOP
     130 CONTINUE
           CALL ADFCNAM(PID, ISTART, MAXCLD, LEN(NDNMS),
        Х
                         NUMRET, NDNMS, IERR)
           CALL ERRCHK(IERR)
            WRITE(*,140)(NDNMS(K),K=1,NUMRET)
     140
         FORMAT(2(8X,A))
           NLEFT = NLEFT - MAXCLD
            ISTART = ISTART + MAXCLD
         IF (NLEFT .GT. 0) GO TO 130
         RETURN
         END
The resulting output is:
   *** ORIGINAL DATABASE SETUP ***
  PARENT NODE NAME = PARENT.1
       NUMBER OF CHILDREN = 1
       CHILDREN NAMES:
          CHILD.1.1
  PARENT NODE NAME = PARENT.2
      NUMBER OF CHILDREN = 4
      CHILDREN NAMES:
```

CHILD.2.1	CHILD.2.2
CHILD.2.3	CHILD.2.4
PARENT NODE NAME = PARENT.3 NUMBER OF CHILDREN = 9 CHILDREN NAMES:	
CHILD.3.1	CHILD.3.2
CHILD.3.3	CHILD.3.4
CHILD.3.5	CHILD.3.6
CHILD.3.7	CHILD.3.8
CHILD.3.9	
*** PARENT.1 AND PARENT.3 AFTER MOVE	***
PARENT NODE NAME = PARENT.1 NUMBER OF CHILDREN = 2 CHILDREN NAMES:	
NUMBER OF CHILDREN = 2	CHILD.3.4
NUMBER OF CHILDREN = 2 CHILDREN NAMES:	CHILD.3.4
NUMBER OF CHILDREN = 2 CHILDREN NAMES: CHILD.1.1 PARENT NODE NAME = PARENT.3 NUMBER OF CHILDREN = 8	CHILD.3.4
NUMBER OF CHILDREN = 2 CHILDREN NAMES: CHILD.1.1 PARENT NODE NAME = PARENT.3 NUMBER OF CHILDREN = 8 CHILDREN NAMES:	
NUMBER OF CHILDREN = 2 CHILDREN NAMES: CHILD.1.1 PARENT NODE NAME = PARENT.3 NUMBER OF CHILDREN = 8 CHILDREN NAMES: CHILD.3.1	CHILD.3.2

ADF_Link — Create a Link to a Node

ADF_Link (PID,name,file,name_in_file,ID,error_return)			
Language	Language C Fortran		
Routine Name	ADF_Link	ADFLINK	
Input	const double PID	real*8 PID	
	const char *name	character*(*) name	
	const char *file	character*(*) file	
	const char *name_in_file	character*(*) name_in_file	
Output	double ID	real*8 ID	
	int *error_return	integer error_return	

PID	The ID of the node's parent.	
name	The name of the link node.	
file	The file name to use for the link directly usable by a C open() routine. If blank (null), the link is assumed to be within the same file as the parent (PID).	
$name_in_file$	The name of the node that the link will point to. This can be a simple or a compound node.	
ID	The ID of the created node.	
$error_return$	Error return code. (See Appendix G.)	

This routine, ADF_Link, will create a link (reference) to a node somewhere within the same ADF database file or another ADF database file. The node that the newly created link node refers to (points to) does not have to exist when the link is created. The ADF library does not check to make sure that the referenced node actually exists at the time ADF_Link is called. However, when information from the referenced node is requested by routines such as ADF_Get_Label or ADF_Read_Data, the referenced node is then accessed. If the referenced node is not in existence at that time, an error will occur. ADF_Link behaves like a "soft link" in the UNIX operating system.

Example 1

This example creates a link to another node that exists in the same physical file. Note that the length of <code>name_in_file</code> is not limited to 32 characters but can be any length required to fully specify the desired node.

```
PROGRAM TEST
C
PARAMETER (MAXCHR=32)
C
CHARACTER*(MAXCHR) NODNAM, NODLBL, TSTLBL
CHARACTER*(72) FN, PATH
C
C *** NODE IDS
```

```
С
      REAL*8 RID, PID, CID, PID1, PID3
      INTEGER I, J, IERR
C
C *** OPEN DATABASE
C
      CALL ADFDOPN('db.adf','NEW',' ',RID,IERR)
      CALL ERRCHK(IERR)
С
C *** CREATE NODES AT FIRST LEVEL
      WRITE(*,100)
  100 FORMAT(/,' *** ORIGINAL DATABASE SETUP ***')
      DO 150 I = 1,3
         WRITE(NODNAM, '(A7, I1)')'PARENT.', I
         CALL ADFCRE(RID, NODNAM, PID, IERR)
         CALL ERRCHK(IERR)
С
C ***** CREATE NODES AT SECOND LEVEL
С
         NUMCLD = I*I
         DO 110 J = 1, NUMCLD
            WRITE(NODNAM, '(A6, I1, A1, I1)')'CHILD.', I, '.', J
            CALL ADFCRE(PID, NODNAM, CID, IERR)
            CALL ERRCHK(IERR)
            WRITE(NODLBL, 105) I, J
  105
            FORMAT('LABEL STRING IN CHILD.', I1,'.', I1)
            CALL ADFSLB(CID, NODLBL, IERR)
            CALL ERRCHK(IERR)
  110
         CONTINUE
C
C ***** GET NUMBER AND NAMES OF CHILDREN JUST CREATED
C
         AND PRINT THEM OUT
C
         CALL PRTCLD(PID)
  150 CONTINUE
C
C *** LINK NODE /PARENT.3/CHILD.3.4 TO /PARENT.1
C
      CALL ADFGNID(RID, 'PARENT.1', PID1, IERR)
      CALL ERRCHK(IERR)
С
      CALL ADFLINK(PID1, 'LINKED_NODE', '',
                   '/PARENT.3/CHILD.3.4',CID,IERR)
      CALL ERRCHK(IERR)
C
C *** CHECK TO MAKE SURE THE NODE WAS ACTUALLY LINKED
      WRITE(*,160)
  160 FORMAT(/,'*** PARENT.1 AFTER LINK ***')
      CALL PRTCLD(PID1)
```

```
С
C *** FOR FINAL CONFIRMATION, GET ORIGINAL LABEL
С
     GOING THROUGH NEW LINK
С
      CALL ADFGLB(CID, TSTLBL, IERR)
      CALL ERRCHK(IERR)
      WRITE(*,170)TSTLBL
  170 FORMAT(/, 'LINKED_NODE LABEL = ',A)
С
      STOP
      END
C
C ******* SUBROUTINES **********
C
      SUBROUTINE ERRCHK(IERR)
C
C *** CHECK ERROR CONDITION
С
      CHARACTER*80 MESS
      IF (IERR .GT. 0) THEN
         CALL ADFERR (IERR, MESS)
         PRINT *, MESS
         CALL ABORT('ADF ERROR')
      ENDIF
      RETURN
      END
      SUBROUTINE PRTCLD(PID)
С
C *** PRINT TABLE OF CHILDREN GIVEN A PARENT NODE-ID
      PARAMETER (MAXCLD=10)
     PARAMETER (MAXCHR=32)
     REAL*8 PID
      CHARACTER*(MAXCHR) NODNAM, NDNMS(MAXCLD)
      CALL ADFGNAM(PID, NODNAM, IERR)
      CALL ERRCHK(IERR)
      CALL ADFNCLD(PID, NUMC, IERR)
      CALL ERRCHK(IERR)
      WRITE(*,120)NODNAM,NUMC
  120 FORMAT(/, ' PARENT NODE NAME = ',A,/,
     X ' NUMBER OF CHILDREN = ',12,/,
     X
                  CHILDREN NAMES: ')
     NLEFT = NUMC
     ISTART = 1
    --- TOP OF DO-WHILE LOOP
  130 CONTINUE
         CALL ADFCNAM(PID, ISTART, MAXCLD, LEN(NDNMS),
     X
                      NUMRET, NDNMS, IERR)
         CALL ERRCHK(IERR)
         WRITE(*,140)(NDNMS(K),K=1,NUMRET)
```

```
140
           FORMAT(2(8X,A))
            NLEFT = NLEFT - MAXCLD
            ISTART = ISTART + MAXCLD
         IF (NLEFT .GT. 0) GO TO 130
         RETURN
         END
The resulting output is:
   *** ORIGINAL DATABASE SETUP ***
  PARENT NODE NAME = PARENT.1
      NUMBER OF CHILDREN = 1
       CHILDREN NAMES:
          CHILD.1.1
  PARENT NODE NAME = PARENT.2
       NUMBER OF CHILDREN = 4
       CHILDREN NAMES:
          CHILD.2.1
                                           CHILD.2.2
          CHILD.2.3
                                           CHILD.2.4
  PARENT NODE NAME = PARENT.3
      NUMBER OF CHILDREN = 9
       CHILDREN NAMES:
          CHILD.3.1
                                          CHILD.3.2
          CHILD.3.3
                                           CHILD.3.4
          CHILD.3.5
                                          CHILD.3.6
          CHILD.3.7
                                          CHILD.3.8
          CHILD.3.9
   *** PARENT.1 AFTER LINK ***
  PARENT NODE NAME = PARENT.1
       NUMBER OF CHILDREN = 2
       CHILDREN NAMES:
          CHILD.1.1
                                           LINKED_NODE
  LINKED_NODE LABEL = LABEL STRING IN CHILD.3.4
Example\ 2
  This example illustrates the linking of nodes across files.
  PROGRAM TEST
  С
         PARAMETER (MAXCHR=32)
         CHARACTER*(MAXCHR) TSTLBL
```

```
C *** NODE IDS
C
      REAL*8 RID, PID, CID
      INTEGER IERR
С
C *** 1.) OPEN 1ST DATABASE
С
      2.) CREATE TWO NODES
С
      3.) PUT LABEL ON 2ND NODE
С
      4.) CLOSE DATABASE
      CALL ADFDOPN('db1.adf','NEW',' ',RID,IERR)
      CALL ERRCHK(IERR)
      CALL ADFCRE(RID, 'DB1_NODE1', PID, IERR)
      CALL ERRCHK(IERR)
      CALL ADFCRE(PID, 'DB1_NODE2', CID, IERR)
      CALL ERRCHK(IERR)
      CALL ADFSLB(CID, 'LABEL IN FILE.1: NODE2', IERR)
      CALL ERRCHK(IERR)
      CALL ADFDCLO(RID, IERR)
      CALL ERRCHK(IERR)
C
C *** 1.) OPEN 2ND DATABASE
      CALL ADFDOPN('db2.adf','NEW',' ',RID,IERR)
      CALL ERRCHK(IERR)
      CALL ADFCRE(RID, 'DB2_NODE1', PID, IERR)
      CALL ERRCHK(IERR)
      CALL ADFCRE(PID, 'DB2_NODE2', CID, IERR)
      CALL ERRCHK(IERR)
С
C *** LINK NODE /DB1_NODE1/DB1_NODE2 TO /DB2_NODE1
C
      CALL ADFLINK(PID, 'LINKED_NODE', 'db1.adf',
                    '/DB1_NODE1/DB1_NODE2',CID,IERR)
      CALL ERRCHK(IERR)
C
C *** CHECK TO MAKE SURE THE NODE WAS ACTUALLY LINKED
      WRITE(*,160)
  160 FORMAT(/,'*** PARENT AFTER LINK ***')
      CALL PRTCLD(PID)
С
C *** FOR FINAL CONFIRMATION, GET ORIGINAL LABEL
C
      GOING THROUGH NEW LINK
С
      CALL ADFGLB(CID, TSTLBL, IERR)
      CALL ERRCHK(IERR)
      WRITE(*,170)TSTLBL
  170 FORMAT(/, 'LINKED_NODE LABEL = ', A)
C
      STOP
```

```
END
C
C ******* SUBROUTINES **********
      SUBROUTINE ERRCHK(IERR)
С
C *** CHECK ERROR CONDITION
      CHARACTER*80 MESS
      IF (IERR .GT. 0) THEN
         CALL ADFERR (IERR, MESS)
         PRINT *, MESS
         CALL ABORT('ADF ERROR')
      ENDIF
      RETURN
      END
      SUBROUTINE PRTCLD(PID)
C *** PRINT TABLE OF CHILDREN GIVEN A PARENT NODE-ID
С
      PARAMETER (MAXCLD=10)
      PARAMETER (MAXCHR=32)
      REAL*8 PID
      CHARACTER*(MAXCHR) NODNAM, NDNMS(MAXCLD)
      CALL ADFGNAM(PID, NODNAM, IERR)
      CALL ERRCHK(IERR)
      CALL ADFNCLD(PID, NUMC, IERR)
      CALL ERRCHK(IERR)
      WRITE(*,120)NODNAM,NUMC
  120 FORMAT(/,' PARENT NODE NAME = ',A,/,
        ,
                NUMBER OF CHILDREN = ',12,/,
                   CHILDREN NAMES: ')
     NLEFT = NUMC
     ISTART = 1
     --- TOP OF DO-WHILE LOOP
  130 CONTINUE
         CALL ADFCNAM(PID, ISTART, MAXCLD, LEN(NDNMS),
    X
                      NUMRET, NDNMS, IERR)
         CALL ERRCHK(IERR)
         WRITE(*,140)(NDNMS(K),K=1,NUMRET)
  140
        FORMAT(2(8X,A))
         NLEFT = NLEFT - MAXCLD
         ISTART = ISTART + MAXCLD
      IF (NLEFT .GT. 0) GO TO 130
      RETURN
      END
```

The resulting output is:

```
*** PARENT AFTER LINK ***
```

I ADF Library of Subroutines

PARENT NODE NAME = DB2_NODE1
NUMBER OF CHILDREN = 2
CHILDREN NAMES:
DB2_NODE2

LINKED_NODE

LINKED_NODE LABEL = LABEL IN FILE.1: NODE2

ADF_Is_Link — See If the Node Is a Link

ADF_Is_Link (ID,link_path_length,error_return)				
Language	C	Fortran		
Routine Name	ADF_Is_Link	ADFISLK		
Input	const double ID	real*8 ID		
Output	int *link_path_length	integer link_path_length		
	int *error_return	integer error_return		

ID The ID of the node to use.

link_path_length This returned value is zero if the node is not a link. If the node referenced by ID is a link within the same file, the number of characters in the path is returned. If the node referenced by ID is a link in another file, the sum of the number of characters in the referenced file name and the number of characters in the path +1 is returned.

error_return Error return code. (See Appendix G.)

This routine, ADF_Is_Link, tests to see if the node is a link. If the actual data type of the node is LK (created with ADF_Link), the routine returns the link path length; otherwise it returns 0.

Example

This example creates a link into a second file, and then calls ADF_Is_Link to determine whether the requested node is indeed a link or a normal node.

```
PROGRAM TEST
C
      PARAMETER (MAXCHR=32)
C
      CHARACTER*(MAXCHR) TSTLBL
C
C *** NODE IDS
      REAL*8 RID, PID, CID
      INTEGER IERR
C
C *** 1.) OPEN 1ST DATABASE
C
      2.) CREATE TWO NODES
С
      3.) PUT LABEL ON 2ND NODE
\mathsf{C}
      4.) CLOSE DATABASE
      CALL ADFDOPN('db1.adf','NEW',' ',RID,IERR)
      CALL ERRCHK(IERR)
      CALL ADFCRE(RID, 'DB1_NODE1', PID, IERR)
      CALL ERRCHK(IERR)
      CALL ADFCRE(PID, 'DB1_NODE2', CID, IERR)
```

```
CALL ERRCHK(IERR)
      CALL ADFDCLO(RID, IERR)
      CALL ERRCHK(IERR)
C
C *** 1.) OPEN 2ND DATABASE
      CALL ADFDOPN('db2.adf','NEW',' ',RID,IERR)
      CALL ERRCHK(IERR)
      CALL ADFCRE(RID, 'DB2_NODE1', PID, IERR)
      CALL ERRCHK(IERR)
      CALL ADFCRE(PID, 'DB2_NODE2', CID, IERR)
      CALL ERRCHK(IERR)
С
C *** LINK NODE FILE 1:/DB1_NODE1/DB1_NODE2 TO /DB2_NODE1
      CALL ADFLINK(PID, 'LINKED_NODE', 'db1.adf',
                   '/DB1_NODE1/DB1_NODE2',CID,IERR)
      CALL ERRCHK(IERR)
C *** CHECK TO MAKE SURE THE NODE WAS ACTUALLY LINKED
      CALL ADFISLK(CID, LEN, IERR)
      CALL ERRCHK(IERR)
      PRINT *, ' PATH LENGTH FROM LINK IS: ',LEN
С
      STOP
      END
С
C ******* SUBROUTINES **********
С
      SUBROUTINE ERRCHK(IERR)
C
C *** CHECK ERROR CONDITION
      CHARACTER*80 MESS
      IF (IERR .GT. 0) THEN
         CALL ADFERR(IERR, MESS)
         PRINT *, MESS
         CALL ABORT('ADF ERROR')
      ENDIF
      RETURN
      END
```

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The resulting output is:

PATH LENGTH FROM LINK IS:

ADF_Get_Link_Path — Get the Path Information From a Link

ADF_Get_Link_Path (ID,file,name_in_file,error_return)				
Language	C	Fortran		
Routine Name	ADF_Get_Link_Path	ADFGLKP		
Input	const double ID	real*8 ID		
Output	char *file	character*(*) file		
	char *name_in_file	<pre>character*(*) name_in_file</pre>		
	int *error_return	integer error_return		

ID The ID of the node to use.

file The file name to use for the link. It is directly usable by the C open() routine.

If blank (null), the link is within the same file.

 $name_in_file$ The name of the node that the link points to.

This routine, ADF_Get_Link_Path, gets the path information from a link. If the node is a link node, the routine returns the path information; otherwise it returns an error.

Example

This example opens two ADF files. A link is created in the second file (db2.adf) that references a node in the first file (db1.adf). ADFGLKP is used to extract the file name and path that the link actually points to. It is not anticipated that this information will be useful to the normal user. For most applications, it is best to allow ADF to resolve the link internally, transparent to the user.

```
PROGRAM TEST
С
      PARAMETER (MAXCHR=32)
C
      CHARACTER*(MAXCHR) TSTLBL
      CHARACTER*(40) FILENM, PATH
С
C *** NODE IDS
C
      REAL*8 RID, PID, CID
      INTEGER IERR
C *** 1.) OPEN 1ST DATABASE
      2.) CREATE TWO NODES
С
      3.) PUT LABEL ON 2ND NODE
С
      4.) CLOSE DATABASE
C
      CALL ADFDOPN('db1.adf','NEW',' ',RID,IERR)
      CALL ADFCRE(RID, 'DB1_NODE1', PID, IERR)
```

```
CALL ADFCRE(PID,'DB1_NODE2',CID,IERR)
      CALL ADFDCLO(RID, IERR)
C
C *** 1.) OPEN 2ND DATABASE
С
      CALL ADFDOPN('db2.adf','NEW',' ',RID,IERR)
      CALL ADFCRE(RID, 'DB2_NODE1', PID, IERR)
      CALL ADFCRE(PID, 'DB2_NODE2', CID, IERR)
С
C *** LINK NODE FILE 1:/DB1_NODE1/DB1_NODE2 TO /DB2_NODE1
С
      CALL ADFLINK(PID, 'LINKED_NODE', 'db1.adf',
                   '/DB1_NODE1/DB1_NODE2',CID,IERR)
С
C *** CHECK TO MAKE SURE THE NODE WAS ACTUALLY LINKED
C
      CALL ADFGLKP (CID, FILENM, PATH, IERR)
      PRINT *, ' INFORMATION FROM LINK: '
      PRINT *,' FILE: ',FILENM
      PRINT *,' PATH: ',PATH
С
      STOP
      END
С
C ******* SUBROUTINES **********
С
      SUBROUTINE ERRCHK(IERR)
С
C *** CHECK ERROR CONDITION
С
      CHARACTER*80 MESS
      IF (IERR .GT. 0) THEN
         CALL ADFERR(IERR, MESS)
         PRINT *, MESS
         CALL ABORT('ADF ERROR')
      ENDIF
      RETURN
      END
      SUBROUTINE PRTCLD(PID)
С
C *** PRINT TABLE OF CHILDREN GIVEN A PARENT NODE-ID
С
      PARAMETER (MAXCLD=10)
      PARAMETER (MAXCHR=32)
      REAL*8 PID
      CHARACTER*(MAXCHR) NODNAM, NDNMS(MAXCLD)
      CALL ADFGNAM(PID, NODNAM, IERR)
      CALL ERRCHK(IERR)
      CALL ADFNCLD(PID, NUMC, IERR)
      CALL ERRCHK(IERR)
```

```
WRITE(*,120)NODNAM,NUMC
120 FORMAT(/, ' PARENT NODE NAME = ',A,/,
  X ' NUMBER OF CHILDREN = ',12,/,
X ' CHILDREN NAMES:')
   NLEFT = NUMC
    ISTART = 1
   --- TOP OF DO-WHILE LOOP
130 CONTINUE
      CALL ADFCNAM(PID, ISTART, MAXCLD, LEN(NDNMS),
                    NUMRET, NDNMS, IERR)
       CALL ERRCHK(IERR)
       WRITE(*,140)(NDNMS(K),K=1,NUMRET)
140
       FORMAT(2(8X,A))
       NLEFT = NLEFT - MAXCLD
       ISTART = ISTART + MAXCLD
    IF (NLEFT .GT. 0) GO TO 130
    RETURN
    END
```

The resulting output is:

INFORMATION FROM LINK:

FILE: db1.adf

PATH: /DB1_NODE1/DB1_NODE2

ADF_Get_Root_ID — Get the Root ID for the ADF System

ADF_Get_Root_ID (ID,root_ID,error_return)		
Language	C Fortran	
Routine Name	ADF_Get_Root_ID	ADFGRID
Input	const double ID real*8 ID	
Output	double root_ID	real*8 root_ID
	int *error_return	integer error_return

ID Any valid node ID for the given ADF database.root_ID The ID of the root nodeerror_return Error return code. (See Appendix G.)

This routine, ADF_Get_Root_ID, returns the root ID for an ADF file when given any valid node ID for that file.

Example

This example illustrates that the root node ID can be obtained at any time by using a currently valid node ID for that file.

```
PROGRAM TEST
С
      PARAMETER (MAXCHR=32)
С
      CHARACTER*(MAXCHR) NODNAM, ROOTNM
С
C *** NODE IDS
C
      REAL*8 RID, PID, CID, TESTID
      INTEGER I, J, IERR, NUMCLD
С
C *** OPEN DATABASE
С
      CALL ADFDOPN('db.adf', 'NEW', '', RID, IERR)
      CALL ADFGNAM(RID, ROOTNM, IERR)
      PRINT *, ' AFTER OPENING FILE, ROOT NAME = ', ROOTNM
С
C *** CREATE NODES AT FIRST LEVEL
      DO 150 I = 1,2
         WRITE(NODNAM, '(A7, I1)')'PARENT.', I
         CALL ADFCRE(RID, NODNAM, PID, IERR)
         TESTID = 0.0
         ROOTNM =''
         CALL ADFGRID (PID, TESTID, IERR)
```

```
CALL ADFGNAM (TESTID, ROOTNM, IERR)
         WRITE(*,100)NODNAM,ROOTNM
         FORMAT('USING NODE ID FROM: ',A,
  100
                ' ROOT NAME = ',A)
C
C ***** CREATE NODES AT SECOND LEVEL
C
         NUMCLD = 2*I
         DO 110 J = 1, NUMCLD
            WRITE(NODNAM, '(A6, I1, A1, I1)')'CHILD.', I, '.', J
            CALL ADFCRE(PID, NODNAM, CID, IERR)
            TESTID = 0.0
            ROOTNM =''
            CALL ADFGRID (PID, TESTID, IERR)
            CALL ADFGNAM(TESTID, ROOTNM, IERR)
            WRITE(*,100)NODNAM,ROOTNM
  110
         CONTINUE
  150 CONTINUE
      STOP
      END
```

The resulting output is:

```
AFTER OPENING FILE, ROOT NAME = ADF MotherNode

USING NODE ID FROM: PARENT.1 ROOT NAME = ADF MotherNode

USING NODE ID FROM: CHILD.1.1 ROOT NAME = ADF MotherNode

USING NODE ID FROM: CHILD.1.2 ROOT NAME = ADF MotherNode

USING NODE ID FROM: PARENT.2 ROOT NAME = ADF MotherNode

USING NODE ID FROM: CHILD.2.1 ROOT NAME = ADF MotherNode

USING NODE ID FROM: CHILD.2.2 ROOT NAME = ADF MotherNode

USING NODE ID FROM: CHILD.2.3 ROOT NAME = ADF MotherNode

USING NODE ID FROM: CHILD.2.4 ROOT NAME = ADF MotherNode
```

I.5 Data Query Routines

ADF_Get_Label — Get the String in a Node's Label Field

ADF_Get_Label (ID, label, error_return)		
Language	C	Fortran
Routine Name	ADF_Get_Label	ADFGLB
Input	const double ID	real*8 ID
Output	char *label	character*(*) label
	int *error_return	integer error_return

ID The ID of the node to use.
 label The 32-character label of the node.
 error_return Error return code. (See Appendix G.)

This routine returns the 32-character string stored in the node's label field. In C, the label will be null terminated after the last nonblank character. Therefore, in general, 33 characters should be allocated (32 for the label, plus 1 for the null). In Fortran, the label is left justified and blank filled to the right. The null character is not returned in Fortran, therefore, the variable declaration can be for 32 characters (e.g., CHARACTER*(32)).

Example

```
PROGRAM TEST
С
      PARAMETER (MAXCHR=32)
С
      CHARACTER*(MAXCHR) NODNAM, LABL
C
C *** NODE IDS
С
      REAL*8 RID, CID
С
C *** OPEN DATABASE
      CALL ADFDOPN('db.adf','NEW',' ',RID,IERR)
      CALL ADFCRE(RID, 'NODE 1', CID, IERR)
      CALL ADFSLB(CID, 'THIS IS A NODE LABEL', IERR)
С
      CALL ADFGNAM(CID, NODNAM, IERR)
      CALL ADFGLB(CID, LABL, IERR)
С
      PRINT *,' NODE NAME = ',NODNAM
      PRINT *, LABEL
                          = ',LABL
С
```

STOP END

The resulting output is:

NODE NAME = NODE 1 LABEL = THIS IS A NODE LABEL

ADF_Set_Label — Set the String in a Node's Label Field

ADF_Set_Label (ID, label, error_return)			
Language	C Fortran		
Routine Name	ADF_Set_Label	ADFSLB	
Input	const double ID	real*8 ID	
	const char *label	character*(*) label	
Output	int *error_return	integer error_return	

ID The ID of the node to use.

 $\begin{array}{ll} \it label & The \ 32\mbox{-character label for the node.} \\ \it error_return & Error \ return \ code. \ (See \ Appendix \ G.) \\ \end{array}$

This routine, ADF_Set_Label, sets the 32-character string in the node's label field.

Example

See the example for ADF_Get_Label.

ADF_Get_Data_Type — Get the String in a Node's Data-Type Field

ADF_Get_Data_Type (ID,data_type,error_return)			
Language	C Fortran		
Routine Name	ADF_Get_Data_Type	ADFGDT	
Input	const double ID	real*8 ID	
Output	char *data_type	character*(*) data_type	
	int *error_return	integer error_return	

ID The ID of the node to use.

data_type The 32-character data type field stored in the node information header.

error_return Error return code. (See Appendix G.)

This routine, ADF_Get_Data_Type, returns the 32-character string in the node's data-type field. In C, the label will be null terminated after the last nonblank; therefore, at least 33 characters (32 for the label, plus 1 for the null) should be allocated for the data_type string. In Fortran, the null character is not returned; therefore, the declaration for the string data_type can be for 32 characters.

Example

This example illustrates the process of creating a node, storing data in it, querying the node for the information in the header, and reading the data back out.

```
PROGRAM TEST
C
      PARAMETER (MAXCHR=32)
      PARAMETER (MAXROW=2)
      PARAMETER (MAXCOL=10)
С
      CHARACTER*(MAXCHR) NODNAM, LABL
      CHARACTER*(MAXCHR) DTYPE
      REAL R4DATI (MAXROW, MAXCOL), R4DATO (MAXROW, MAXCOL)
      INTEGER IDIMI(2),IDIMO(2)
C
C *** NODE IDS
C
      REAL*8 RID, CID
C
C *** OPEN DATABASE
      CALL ADFDOPN('db.adf', 'NEW', '', RID, IERR)
C *** GENERATE SOME DATA
C
      IDIMI(1) = MAXROW
      IDIMI(2) = MAXCOL
```

```
DO 200 ICOL = 1, MAXCOL
           DO 100 IROW = 1, MAXROW
              R4DATI(IROW, ICOL) = 2.0*ICOL*IROW
           CONTINUE
    100
    200 CONTINUE
  С
  C *** GENERATE A NODE AND PUT DATA IN IT
        CALL ADFCRE(RID, 'NODE 1', CID, IERR)
        CALL ADFSLB(CID, 'LABEL FOR NODE 1', IERR)
        CALL ADFPDIM(CID, 'R4', 2, IDIMI, IERR)
        CALL ADFWALL(CID, R4DATI, IERR)
  С
  C *** GET INFORMATION FROM NODE
  C
        CALL ADFGNAM(CID, NODNAM, IERR)
        CALL ADFGLB(CID, LABL, IERR)
        CALL ADFGDT(CID, DTYPE, IERR)
        CALL ADFGND(CID, NDIM, IERR)
        CALL ADFGDV(CID, IDIMO, IERR)
        CALL ADFRALL(CID, R4DATO, IERR)
  С
        PRINT *, NODE NAME
                                   = ',NODNAM
                                     = ',LABL
        PRINT *, LABEL
        PRINT *, DATA TYPE = ',DTYPE
        PRINT *, ' NUMBER OF DIMENSIONS = ', NDIM
        PRINT *, DIMENSIONS
                               = ',IDIMO
        PRINT *,' DATA:'
        WRITE(*,300)((R4DATO(I,J),I=1,MAXROW),J=1,MAXCOL)
    300 FORMAT(2(5X,F10.2))
  С
        STOP
        END
The resulting output is:
  NODE NAME
                     = NODE 1
  LABEL
                      = LABEL FOR NODE 1
  DATA TYPE = R4
  NUMBER OF DIMENSIONS =
  DIMENSIONS
                                               10
  DATA:
           2.00
                        4.00
           4.00
                         8.00
           6.00
                        12.00
           8.00
                        16.00
          10.00
                       20.00
          12.00
                       24.00
          14.00
                       28.00
          16.00
                        32.00
                 36.00
          18.00
```

20.00 40.00

ADF_Get_Number_of_Dimensions — Get the Number of Node Dimensions

ADF_Get_Number_of_Dimensions (ID,num_dims,error_return)		
Language	C Fortran	
Routine Name	ADF_Get_Number_of_Dimensions	ADFGND
Input	const double ID	real*8 ID
Output	int *num_dims	integer num_dims
	int *error_return	integer error_return

ID The ID of the node to use.

num_dims The integer dimension value.

This routine, ADF_Get_Number_of_Dimensions, returns the number of data dimensions used in the node. Values will be returned only for the number of dimensions defined in the node. If the number of dimensions for the node is zero, an error is returned.

Example

See the example for ADF_Get_Data_Type.

ADF_Get_Dimension_Values — Get the Values of the Node Dimensions

ADF_Get_Dimension_Values (ID,dim_vals[],error_return)		
Language	C Fortran	
Routine Name	ADF_Get_Dimension_Values	ADFGDV
Input	const double ID	real*8 ID
Output	int *dim_vals[]	integer dim_vals()
	int *error_return	integer error_return

ID The ID of the node to use.

dim_vals The array (list) of dimension values.error_return Error return code. (See Appendix G.)

This routine, ADF_Get_Dimension_Values, returns the array (list) of dimension values for a node. Values will be returned only for the number of dimensions defined in the node. If the number of dimensions for the node is zero, an error is returned.

Example

See the example for ADF_Get_Data_Type.

 ${\tt ADF_Put_Dimension_Information} \ -- \ Set \ or \ Change \ the \ Data \ Type \ and \ Dimensions \ of \ a \ Node$

ADF_Put_Dimension_Information (ID,data_type,dims,dim_vals[],error_return)			
Language	C	Fortran	
Routine Name	ADF_Put_Dimension_Information	ADFPDIM	
Input	const double ID	real*8 ID	
	const char *data_type	character*(*) data_type	
const int dims		integer dims	
int dim_vals[] integer dim_vals()		integer dim_vals()	
Output	int *error_return	integer error_return	

ID The ID of the node to use.

data_type The 32-character data type of the node. The valid user-definable data types are:

<u>Notation</u>
MT
14
18
U4
U8
R4
R8
X4
Х8
C1
B1

Compound data types can be defined as a combination of types ("I4,I4,R8"), an array ("I4[25]"), or a combination of types and arrays ("I4,C1[20],R8[3]"). They can contain up to 32 characters. This style of data type definition is not very portable across platforms; therefore, it is not recommended.

dims The number of dimensions of this node. dims can be a number from 0 to 12. "0" means no data. The dimension of an array can range from 1 (vector) to 12.

The array (list) of dimension values for this node. dim_vals is a vector of integers that define the size of the array in each dimension as defined by dims. If the dims is zero, the dims_vals are not used. The valid range of dim_vals is from 1 to 2,147,483,648. The total data size in bytes, calculated by the data_type size times the dimension values, cannot exceed 2,147,483,648 for any one node.

 $dim_{-}vals$

This routine, ADF_Put_Dimension_Information, sets or changes the data type and dimension information for a node.

Note: When this routine is used to change the data type or number of dimensions of an existing node, any data currently associated with the node are lost. The dimension values can be changed and the data space will be extended as needed. Be very careful changing the dimension values. The layout of the data is assumed to be Fortran-like. If the left-most dimension values are changed, the data are not shifted around on disk to account for this; only the amount of data is changed. Therefore, the indexing into the data will be changed. In general, it is safe to change the right-most dimension value. For example, if the array of dimension values was (10,11,20,50), then changing the 10, 11, or 20 is very risky, whereas changing the 50 should be safe.

Note: See the discussion of Fortran character array portability in Appendix E.7.

$\underline{Example}$

See the example for ADF_Get_Data_Type.

I.6 Data I/O Routines

Note: For all data I/O routines, the system is based on indexing starting from 1 and not 0. (That is, the first element in an array is indexed as 1 and not zero.)

 ${\tt ADF_Read_Data} \ -- \ Read \ the \ Data \ From \ a \ Node \ Having \ Stride \ Capabilities$

ADF_Read_Data (ID,s_start[],s_end[],s_stride[],m_num_dims,m_dims[],m_start[], m_end[],m_stride[],data,error_return)

Language	C	Fortran	
Routine Name	ADF_Read_Data	ADFREAD	
Input	const double ID	real*8 ID	
	const int s_start[]	integer s_start()	
	const int s_end[]	integer s_end()	
const int s_stride[]		integer s_stride()	
	const int m_num_dims	integer m_num_dims	
	const int m_dims[]	integer m_dims()	
	const int m_start[]	integer m_start()	
	const int m_end[]	integer m_end()	
	<pre>const int m_stride[]</pre>	integer m_stride()	
Output	char *data	character*(*) data	
	int *error_return	integer error_return	

ID	The ID of the node to use.
$s_start[]$	The starting index to use for each dimension of the array within the database node (1D array; i.e., list of indices). The maximum number of dimensions an array is allowed in ADF is 12.
$s_end[]$	The ending index to use for each dimension of the array within the database node $(1D \text{ array}; \text{ i.e., list of indices}).$
$s_stride[]$	The stride value to use for each dimension of the array within the database node (1D array; i.e., list of indices).
m_num_dims	The number of dimensions to use in memory.
$m_dims[]$	The dimension values to use for the array in memory (1D array; i.e., list of indices).
$m_start[]$	The starting index to use for each dimension of the array in memory (1D array; i.e., list of indices).
$m_end[]$	The ending index to use for each dimension of the array in memory (1D array; i.e., list of indices).
$m_stride[]$	The stride value to use for each dimension of the array in memory (1D array; i.e., list of indices).

```
data The starting address of the data in memory.

error_return

Error return code. (See Appendix G.)
```

This routine, ADF_Read_Data, provides general purpose read capabilities. It allows for a general specification of the starting location within the data well as fixed step lengths (strides) through the data from the initial position. This capability works for both the data on disk and the data being stored in memory. One set of integer vectors (s_start, etc.) is used to describe the mapping of the data within the node, and a second set of integer vectors (m_start) is used to describe the mapping of the desired data within memory.

There can be a significant performance penalty for using ADF_Read_Data when compared with ADF_Read_All_Data. If performance is a major consideration, it is best to organize data to take advantage of the speed of ADF_Read_All_Data.

The data are stored in both memory and on disk in "Fortran ordering." That is, the first index varies the fastest.

ADF_Read_Data will not accept "negative" indexing. That is, it is not possible to reverse the ordering of the data from the node into memory.

Be careful when writing data using ADF_Write_All_Data and then using ADF_Read_Data to randomly access the data. ADF_Write_All_Data takes a starting address in memory and writes N words to disk, making no assumption as to the order of the data. ADF_Read_Data assumes that the data have Fortran-like ordering to navigate through the data in memory and on disk. It assumes that the first dimension varies the fastest. It would be easy for a C program to use the default array ordering (last dimension varying fastest) and write the data out using ADF_Write_All_Data. Then another program might use ADF_Read_Data to access a subsection of the data, and the routine would not return what was expected.

Note: If all the data type of the node is a compound data type, such as ("I4[3],R8"), the partial capabilities will access one or more of these 20-byte data entities. You cannot access a subset of an occurrence of the data type.

Note: See the discussion of Fortran character array portability in Appendix E.7.

Example 1

This example shows ADFREAD being used to emulate the same capabilities as those in ADFRALL.

```
PROGRAM TEST
C

PARAMETER (MAXROW=10)
PARAMETER (MAXCOL=3)
C

REAL R4ARRI(MAXROW, MAXCOL)
REAL R4ARRO(MAXROW, MAXCOL)
INTEGER IDIMI(2), IDIMO(2)
INTEGER IDBEG(2), IDEND(2), IDINCR(2)
INTEGER IMBEG(2), IMEND(2), IMINCR(2)
C
C *** NODE IDS
C
REAL*8 RID,CID
C
```

```
C *** OPEN DATABASE
      CALL ADFDOPN('db.adf','NEW',' ',RID,IERR)
С
C *** GENERATE SOME DATA
      IDIMI(1) = MAXROW
      IDIMI(2) = MAXCOL
      DO 200 ICOL = 1, MAXCOL
         DO 100 IROW = 1, MAXROW
            R4ARRI(IROW, ICOL) = 2.0*ICOL*IROW
  100
         CONTINUE
  200 CONTINUE
      PRINT *, 'ORIGINAL ARRAY:'
      WRITE(*,300)((R4ARRI(I,J),J=1,MAXCOL),I=1,MAXROW)
  300 FORMAT(3(5X,F10.2))
C *** GENERATE A NODE AND PUT DATA IN IT
      CALL ADFCRE(RID, 'NODE 1',CID, IERR)
      CALL ADFSLB(CID, 'LABEL FOR NODE 1', IERR)
      CALL ADFPDIM(CID, 'R4', 2, IDIMI, IERR)
      CALL ADFWALL(CID, R4ARRI, IERR)
С
C *** GET INFORMATION FROM NODE
C
C *** GET DATA FROM NODE (EXACTLY EQUIVALENT TO ADFRALL)
С
      IDBEG(1) = 1
      IDEND(1) = MAXROW
      IDINCR(1) = 1
C
      IDBEG(2) = 1
      IDEND(2) = MAXCOL
      IDINCR(2) = 1
C
      IDIMO(1) = MAXROW
      IDIMO(2) = MAXCOL
C
      IMBEG(1) = 1
      IMEND(1) = MAXROW
      IMINCR(1) = 1
С
      IMBEG(2) = 1
      IMEND(2) = MAXCOL
      IMINCR(2) = 1
      CALL ADFREAD (CID, IDBEG, IDEND, IDINCR,
     X
                   2, IDIMO, IMBEG, IMEND, IMINCR,
                   R4ARRO, IERR)
      CALL ERRCHK(IERR)
C
```

```
PRINT *,' ARRAY PULLED FROM DISK USING ADFREAD:'
      WRITE(*,300)((R4ARRO(I,J),J=1,MAXCOL),I=1,MAXROW)
С
      STOP
      END
С
C ******** SUBROUTINES **********
С
      SUBROUTINE ERRCHK(IERR)
С
C *** CHECK ERROR CONDITION
      CHARACTER*80 MESS
      IF (IERR .GT. 0) THEN
        CALL ADFERR(IERR, MESS)
        PRINT *, MESS
         CALL ABORT('ADF ERROR')
      ENDIF
      RETURN
      END
```

The resulting output is:

ORIGINAL ARRAY:			
2.00	4.00	6.00	
4.00	8.00	12.00	
6.00	12.00	18.00	
8.00	16.00	24.00	
10.00	20.00	30.00	
12.00	24.00	36.00	
14.00	28.00	42.00	
16.00	32.00	48.00	
18.00	36.00	54.00	
20.00	40.00	60.00	
ARRAY PULLED FROM	DISK USING ADFREA	AD:	
ARRAY PULLED FROM 2.00	DISK USING ADFREA	AD: 6.00	
2.00	4.00	6.00	
2.00 4.00	4.00 8.00	6.00 12.00	
2.00 4.00 6.00	4.00 8.00 12.00	6.00 12.00 18.00	
2.00 4.00 6.00 8.00	4.00 8.00 12.00 16.00	6.00 12.00 18.00 24.00	
2.00 4.00 6.00 8.00 10.00	4.00 8.00 12.00 16.00 20.00	6.00 12.00 18.00 24.00 30.00	
2.00 4.00 6.00 8.00 10.00 12.00	4.00 8.00 12.00 16.00 20.00 24.00	6.00 12.00 18.00 24.00 30.00 36.00	
2.00 4.00 6.00 8.00 10.00 12.00 14.00	4.00 8.00 12.00 16.00 20.00 24.00 28.00	6.00 12.00 18.00 24.00 30.00 36.00 42.00	

$\underline{Example~2}$

This example illustrates some of the flexibility available with ADF_Read_Data. An array is created

and written to disk using ADFWALL. Then every other entry in the second column is read back into every other element of a vector.

```
PROGRAM TEST
С
      PARAMETER (MAXROW=10)
      PARAMETER (MAXCOL=3)
С
      REAL R4ARRI (MAXROW, MAXCOL), R4VECO (MAXROW)
      INTEGER IDIMD(2)
      INTEGER IDBEG(2), IDEND(2), IDINCR(2)
С
C *** NODE IDS
C
      REAL*8 RID, CID
C
C *** OPEN DATABASE
C
      CALL ADFDOPN('db.adf','NEW',' ',RID,IERR)
C
C *** GENERATE SOME DATA
C
      IDIMD(1) = MAXROW
      IDIMD(2) = MAXCOL
      DO 200 ICOL = 1, MAXCOL
         DO 100 IROW = 1, MAXROW
            R4ARRI(IROW, ICOL) = 2.0*ICOL*IROW
  100
         CONTINUE
  200 CONTINUE
С
      DO 250 I = 1,MAXROW
         R4VECO(I) = 0.0
  250 CONTINUE
С
      PRINT *, 'ORIGINAL ARRAY:'
      WRITE(*,300)((R4ARRI(I,J),J=1,MAXCOL),I=1,MAXROW)
  300 FORMAT(3(5X,F10.2))
C *** GENERATE A NODE AND PUT DATA IN IT
      CALL ADFCRE(RID,'NODE 1',CID,IERR)
      CALL ADFSLB(CID, 'LABEL FOR NODE 1', IERR)
      CALL ADFPDIM(CID, 'R4', 2, IDIMD, IERR)
      CALL ADFWALL(CID, R4ARRI, IERR)
С
C *** GET DATA FROM NODE USING STRIDED READ
C
C ***** TAKE EVERY OTHER NUMBER FROM THE 2ND COLUMN OF THE ARRAY
С
         AND PUT IT IN SEQUENTIALLY IN A VECTOR IN MEMORY
C *** DATABASE STRIDE INFORMATION
```

```
С
        IDBEG(1) = 1
        IDEND(1) = MAXROW
        IDINCR(1) = 2
  С
        IDBEG(2) = 2
        IDEND(2) = 2
        IDINCR(2) = 1
  С
  C *** MEMORY STRIDE INFORMATION
  С
        NDIMM = 1
        IDIMM = MAXROW
        IMBEG = 1
        IMEND = MAXROW
        IMINCR = 2
  C
        CALL ADFREAD(CID, IDBEG, IDEND, IDINCR,
                     NDIMM, IDIMM, IMBEG, IMEND, IMINCR,
       Χ
                     R4VECO, IERR)
        CALL ERRCHK(IERR)
  С
        PRINT *,' VECTOR WITH DATA EXTRACTED FROM ARRAY'
        WRITE(*,400)(R4VECO(J), J=1,MAXROW)
    400 FORMAT(3(5X,F10.2))
  С
        STOP
        END
  С
  C ******** SUBROUTINES **********
  C
        SUBROUTINE ERRCHK(IERR)
  С
  C *** CHECK ERROR CONDITION
        CHARACTER*80 MESS
        IF (IERR .GT. 0) THEN
           CALL ADFERR(IERR, MESS)
           PRINT *, MESS
           CALL ABORT('ADF ERROR')
        ENDIF
        RETURN
        END
The resulting output is:
  ORIGINAL ARRAY:
            2.00
                         4.00
                                         6.00
            4.00
                          8.00
                                       12.00
                                    18.00
                      12.00
            6.00
```

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8.00	16.00	24.00
10.00	20.00	30.00
12.00	24.00	36.00
14.00	28.00	42.00
16.00	32.00	48.00
18.00	36.00	54.00
20.00	40.00	60.00
VECTOR WITH DATA	EXTRACTED FROM ARRAY	<i>!</i>
4.00	0.00	12.00
0.00	20.00	0.00
28.00	0.00	36.00
0.00		

 ${\tt ADF_Read_All_Data} \ -- \ {\it Read\ All\ the\ Data\ From\ a\ Node}$

ADF_Read_All_Data (ID,data,error_return)				
Language	C Fortran			
Routine Name	ADF_Read_All_Data	ADFRALL		
Input	const double ID	real*8 ID		
Output	char *data	character*(*) data		
	int *error_return	integer error_return		

ID The ID of the node to use.

data The starting address of the data in memory.

This routine, ADF_Read_All_Data, reads all data from a node. It reads all the node's data and returns them into a contiguous memory space.

The disk performance of ADF_Read_All_Data is very good. The routine issues a single read command to the system for the entire data set; therefore, it is as fast as the system can return the data.

Note: See the discussion of Fortran character array portability in Appendix E.7.

Example

See the example for ADF_Get_Data_Type.

ADF_Read_Block_Data — Read a Contiguous Block of Data From a Node

ADF_Read_Block_Data (ID,b_start,b_end,data,error_return)		
Language	C	Fortran
Routine Name	ADF_Read_Block_Data	ADFRBLK
Input	const double ID	real*8 ID
	const long b_start	integer b_start
	const long b_end	integer b_end
Output	char *data	character*(*) data
	int *error_return	integer error_return

ID The ID of the node to use. b_start The starting point of the block in token space. b_end The ending point of the block in token space. data The starting address of the data in memory. $error_return$ Error return code. (See Appendix G.)

This routine, ADF_Read_Block_Data, reads a block of data from a node and returns it into a contiguous memory space.

Note: See the discussion of Fortran character array portability in Appendix E.7.

 ${\tt ADF_Write_Data} - Write\ the\ Data\ to\ a\ Node\ Having\ Stride\ Capabilities$

ADF_Write_Data (ID,s_start[],s_end[],s_stride[],m_num_dims,m_dims[],m_start[],		
Language	C	Fortran
Routine Name	ADF_Write_Data	ADFWRIT
Input	const double ID	real*8 ID
	const int s_start[]	integer s_start()
	const int s_end[]	integer s_end()
	const int s_stride[]	integer s_stride()
	const int m_num_dims	integer m_num_dims
	const int m_dims[]	integer m_dims()
	const int m_start[]	integer m_start()
	const int m_end[]	integer m_end()
	const int m_stride[]	integer m_stride()
	char *data	character*(*) data
Output	int *error return	integer error return

ID	The ID of the node to use.
$s_start[]$	The starting index to use for each dimension of the array within the database node (1D array; i.e., list of indices). The maximum number of dimensions an array is allowed in ADF is 12.
$s_end[]$	The ending index to use for each dimension of the array within the database node (1D array; i.e., list of indices).
$s_stride[]$	The stride value to use for each dimension of the array within the database node (1D array; i.e., list of indices).
m_num_dims	The number of dimensions to use in memory.
$m_dims[]$	The dimension values to use for the array in memory (1D array; i.e., list of indices). $$
$m_start[]$	The starting index to use for each dimension of the array in memory (1D array; i.e., list of indices).
$m_end[]$	The ending index to use for each dimension of the array in memory (1D array; i.e., list of indices).
$m_stride[]$	The stride value to use for each dimension of the array in memory (1D array; i.e., list of indices).
data	The starting address of the data in memory.
$error_return$	Error return code. (See Appendix G.)

This routine, ADF_Write_Data, provides general purpose write capabilities. It allows offsets and strides within both the data in memory and the node on disk. One set of integer vectors (s_start, etc.) is used to describe the mapping of the data within the node, and a second set of integer vectors (m_start, etc.) is used to describe the mapping of the desired data within memory.

There can be a significant performance penalty for using ADF_Write_Data when compared with ADF_Write_All_Data. If performance is a major consideration, it is best to organize data to take advantage of the speed of ADF_Write_All_Data.

The data are stored in both memory and on disk in "Fortran ordering." That is, the first index varies the fastest.

ADF_Write_Data will not accept "negative" indexing. That is, it is not possible to reverse the ordering of the data from the node into memory.

Be careful when using ADF_Read_All_Data to randomly access data that has been written using ADF_Write_Data. ADF_Read_All_Data takes a starting address in memory and takes N contiguous words from disk, making no assumption as to the order of the data. ADF_Write_Data assumes that the data have Fortran-like ordering to navigate through the data on disk and in memory. It assumes that the first dimension varies the fastest.

Note: See the discussion of Fortran character array portability in Appendix E.7.

Example 1

This example uses ADF_Write_Data to perform exactly the same task as ADF_Write_All_Data. ADF_Write_All_Data should be used whenever possible for performance reasons.

```
PROGRAM TEST
С
      PARAMETER (MAXCHR=32)
      PARAMETER (MAXROW=10)
      PARAMETER (MAXCOL=3)
С
      CHARACTER*(MAXCHR) NODNAM, LABL
      CHARACTER*(MAXCHR) DTYPE
      REAL R4ARRI (MAXROW, MAXCOL)
      REAL R4ARRO (MAXROW, MAXCOL)
      INTEGER IDIMI(2),IDIMO(2)
      INTEGER IDBEG(2), IDEND(2), IDINCR(2)
      INTEGER IMBEG(2), IMEND(2), IMINCR(2)
С
C *** NODE IDS
C
      REAL*8 RID, CID
C
C *** OPEN DATABASE
C
      CALL ADFDOPN('db.adf', 'NEW', '', RID, IERR)
C *** GENERATE SOME DATA
      IDIMI(1) = MAXROW
      IDIMI(2) = MAXCOL
```

```
DO 200 ICOL = 1, MAXCOL
         DO 100 IROW = 1, MAXROW
            R4ARRI(IROW, ICOL) = 2.0*ICOL*IROW
  100
         CONTINUE
  200 CONTINUE
      PRINT *, 'ORIGINAL ARRAY:'
      WRITE(*,300)((R4ARRI(I,J),J=1,MAXCOL),I=1,MAXROW)
  300 FORMAT(3(5X,F10.2))
С
C *** GENERATE A NODE AND PUT DATA IN IT
C
      THIS IS EXACTLY EQUIVALENT TO USING ADFWALL
С
      CALL ADFCRE(RID, 'NODE 1', CID, IERR)
      CALL ADFSLB(CID, 'LABEL FOR NODE 1', IERR)
      CALL ADFPDIM(CID, 'R4', 2, IDIMI, IERR)
C
      IDBEG(1) = 1
      IDEND(1) = MAXROW
      IDINCR(1) = 1
С
      IDBEG(2) = 1
      IDEND(2) = MAXCOL
      IDINCR(2) = 1
С
      IDIMO(1) = MAXROW
      IDIMO(2) = MAXCOL
C
      IMBEG(1) = 1
      IMEND(1) = MAXROW
      IMINCR(1) = 1
С
      IMBEG(2) = 1
      IMEND(2) = MAXCOL
      IMINCR(2) = 1
С
      CALL ADFWRIT(CID, IDBEG, IDEND, IDINCR, 2, IDIMO, IMBEG,
                    IMEND, IMINCR, R4ARRI, IERR)
      CALL ERRCHK(IERR)
С
C *** GET INFORMATION FROM NODE
      CALL ADFGNAM(CID, NODNAM, IERR)
      CALL ADFGLB(CID, LABL, IERR)
      CALL ADFGDT(CID, DTYPE, IERR)
      CALL ADFGND (CID, NDIM, IERR)
      CALL ADFGDV(CID, IDIMO, IERR)
      CALL ADFRALL(CID, R4ARRO, IERR)
      CALL ERRCHK(IERR)
С
      PRINT *,' '
      PRINT *,' NODE NAME
                                     = ',NODNAM
```

```
= ',LABL
       PRINT *, LABEL
       PRINT *, LABEL = ',LABL PRINT *, DATA TYPE = ',DTYPE
       PRINT *, ' NUMBER OF DIMENSIONS = ', NDIM
       PRINT *, DIMENSIONS
                           = ',IDIMO
       PRINT *,' ADFRALL DATA:'
       WRITE(*,300)((R4ARRO(I,J),J=1,MAXCOL),I=1,MAXROW)
  С
       STOP
       END
  С
  C ******* SUBROUTINES **********
  С
       SUBROUTINE ERRCHK(IERR)
  C
  C *** CHECK ERROR CONDITION
  C
       CHARACTER*80 MESS
       IF (IERR .GT. 0) THEN
          CALL ADFERR(IERR, MESS)
          PRINT *, MESS
          CALL ABORT('ADF ERROR')
       ENDIF
       RETURN
       END
The resulting output is:
  ORIGINAL ARRAY:
         6.00
                                 6.00
12.00
                                  18.00
                                   24.00
                                   30.00
                                   36.00
                                   42.00
                                   48.00
                                   54.00
                                   60.00
                   = NODE 1
  NODE NAME
  LABEL
                   = LABEL FOR NODE 1
             = R4
  DATA TYPE
  NUMBER OF DIMENSIONS =
                               2
              = 10
                                     3
  DIMENSIONS
  ADFRALL DATA:
          2.00 4.00
                                    6.00
           4.00
                      8.00
                                   12.00
          6.00 12.00
8.00 16.00
10.00 20.00
                                   18.00
                                   24.00
                                   30.00
```

12.00	24.00	36.00
14.00	28.00	42.00
16.00	32.00	48.00
18.00	36.00	54.00
20.00	40.00	60.00

$Example\ 2$

This example illustrates the capability to write a full matrix to an ADF file and then use ADF_Write_Data to rewrite selected portions of the matrix with new data from a much smaller data structure.

```
PROGRAM TEST
C
      PARAMETER (MAXCHR=32)
      PARAMETER (MAXROW=10)
      PARAMETER (MAXCOL=3)
С
      CHARACTER*(MAXCHR) NODNAM, LABL
      CHARACTER*(MAXCHR) DTYPE
      REAL R4ARRI (MAXROW, MAXCOL), R4VEC (MAXCOL)
      REAL R4ARRO(MAXROW, MAXCOL)
      INTEGER IDIMI(2), IDIMO(2), IDIMM(2)
      INTEGER IDBEG(2), IDEND(2), IDINCR(2)
      INTEGER IMBEG(2), IMEND(2), IMINCR(2)
C
C *** NODE IDS
C
      REAL*8 RID, CID
C
C *** OPEN DATABASE
C
      CALL ADFDOPN('db.adf','NEW',' ',RID,IERR)
C
C *** GENERATE SOME DATA
      IDIMI(1) = MAXROW
      IDIMI(2) = MAXCOL
      DO 200 ICOL = 1, MAXCOL
         DO 100 IROW = 1, MAXROW
            R4ARRI(IROW, ICOL) = 2.0*ICOL*IROW
  100
         CONTINUE
         R4VEC(ICOL) = 2.2*ICOL
  200 CONTINUE
      PRINT *, 'ORIGINAL ARRAY:'
      WRITE(*,300)((R4ARRI(I,J),J=1,MAXCOL),I=1,MAXROW)
  300 FORMAT(3(5X,F10.2))
C
C *** GENERATE A NODE AND WRITE THE ARRAY IN IT
      CALL ADFCRE(RID, 'NODE 1', CID, IERR)
```

```
CALL ADFSLB(CID, 'LABEL FOR NODE 1', IERR)
      CALL ADFPDIM(CID, 'R4', 2, IDIMI, IERR)
      CALL ADFWALL(CID, R4ARRI, IERR)
      CALL ERRCHK(IERR)
С
C *** GET INFORMATION FROM NODE (JUST TO PROVE ITS RIGHT)
С
      CALL ADFGNAM(CID, NODNAM, IERR)
      CALL ADFGLB(CID, LABL, IERR)
      CALL ADFGDT(CID, DTYPE, IERR)
      CALL ADFGND(CID, NDIM, IERR)
      CALL ADFGDV(CID, IDIMO, IERR)
      CALL ADFRALL(CID, R4ARRO, IERR)
      CALL ERRCHK(IERR)
C
      PRINT *,' '
      PRINT *,' NODE NAME
                                    = ', NODNAM
      PRINT *, LABEL
                                     = ',LABL
      PRINT *,' DATA TYPE
                                    = ',DTYPE
      PRINT *, ' NUMBER OF DIMENSIONS = ', NDIM
      PRINT *, DIMENSIONS
                            = ',IDIMO
      PRINT *, 'ORIGINAL DATA ON DISK:'
      WRITE(*,300)((R4ARRO(I,J),J=1,MAXCOL),I=1,MAXROW)
С
C *** NOW, USING A VECTOR WITH NEW DATA IN IT, SCATTER
C
      IT INTO THE DATABASE (THIS MODIFIES THE 5TH ROW
С
      OF THE MATRIX)
C
      IDBEG(1) = 5
      IDEND(1) = 5
      IDINCR(1) = 1
C
      IDBEG(2) = 1
      IDEND(2) = MAXCOL
      IDINCR(2) = 1
C
      NMDIM = 1
      IDIMM(1) = MAXCOL
      IMBEG(1) = 1
      IMEND(1) = MAXCOL
      IMINCR(1) = 1
С
      CALL ADFWRIT (CID, IDBEG, IDEND, IDINCR,
     Х
                   NMDIM, IDIMM, IMBEG, IMEND, IMINCR,
     X
                   R4VEC, IERR)
      CALL ERRCHK(IERR)
C *** NOW PULL THE REVISED ARRAY OFF DISK AND PRINT IT
С
      CALL ADFRALL(CID, R4ARRO, IERR)
      CALL ERRCHK(IERR)
```

```
С
       PRINT *,' '
       PRINT *,' AFTER SCATTER:'
       WRITE(*,300)((R4ARRO(I,J),J=1,MAXCOL),I=1,MAXROW)
  С
  STOP
  END
  С
  C ********* SUBROUTINES **********
  С
       SUBROUTINE ERRCHK(IERR)
  С
  C *** CHECK ERROR CONDITION
       CHARACTER*80 MESS
       IF (IERR .GT. 0) THEN
          CALL ADFERR(IERR, MESS)
          PRINT *, MESS
          CALL ABORT('ADF ERROR')
       ENDIF
       RETURN
       END
The resulting output is:
  ORIGINAL ARRAY:
          2.00
                      4.00
                                    6.00
          4.00
                      8.00
                                  12.00
                     12.00
          6.00
                                  18.00
                    16.00
20.00
          8.00
                                  24.00
          10.00
                                  30.00
          12.00
                     24.00
                                  36.00
                     28.00
                                  42.00
          14.00
          16.00
                     32.00
                                  48.00
          18.00
                     36.00
                                  54.00
          20.00
                     40.00
                                  60.00
  NODE NAME
                  = NODE 1
  LABEL
                   = LABEL FOR NODE 1
  DATA TYPE
                  = R4
  NUMBER OF DIMENSIONS =
                               2
  DIMENSIONS
                             10
                                          3
  ORIGINAL DATA ON DISK:
          2.00
                      4.00
                                    6.00
          4.00
                      8.00
                                   12.00
                   12.00
          6.00
                                  18.00
          8.00
                     16.00
                                  24.00
                     20.00
          10.00
                                  30.00
          12.00
                     24.00
                                  36.00
          14.00
                     28.00
                                  42.00
```

I ADF Library of Subroutines

16.00	32.00	48.00
18.00	36.00	54.00
20.00	40.00	60.00
AFTER SCATTER:		
2.00	4.00	6.00
4.00	8.00	12.00
6.00	12.00	18.00
8.00	16.00	24.00
2.20	4.40	6.60
12.00	24.00	36.00
14.00	28.00	42.00
16.00	32.00	48.00
18.00	36.00	54.00
20.00	40.00	60.00

${\tt ADF_Write_All_Data} \ -- \ \mathit{Write} \ \mathit{All} \ \mathit{the} \ \mathit{Data} \ \mathit{to} \ \mathit{a} \ \mathit{Node}$

ADF_Write_All_Data (ID,data,error_return)		
Language	C	Fortran
Routine Name	ADF_Write_All_Data	ADFWALL
Input	const double ID	real*8 ID
	const char *data	character*(*) data
Output	int *error_return	integer error_return

ID The ID of the node to use.

data The starting address of the data in memory.

This routine, ADF_Write_All_Data, writes all data to a node. It copies all the node's data from a contiguous memory space into a contiguous disk space.

The disk performance of ADF_Write_All_Data is very good. The routine issues a single write command to the system for the entire data set; therefore, it is as fast as the system can put the data on disk.

Note: See the discussion of Fortran character array portability in Appendix E.7.

Example

See the example for ADF_Get_Data_Type.

ADF_Write_Block_Data — Write a Contiguous Block of Data To a Node

ADF_Write_Block_Data (ID,b_start,b_end,data,error_return)		
Language	C	Fortran
Routine Name	ADF_Write_Block_Data	ADFWBLK
Input	const double ID	real*8 ID
	const long b_start	integer b_start
	const long b_end	integer b_end
	char *data	character*(*) data
Output	int *error_return	integer error_return

ID The ID of the node to use. b_start The starting point of the block in token space. b_end The ending point of the block in token space. data The starting address of the data in memory. $error_return$ Error return code. (See Appendix G.)

This routine, ADF_Write_Block_Data, writes a contiguous block of data from memory to a node. *Note:* See the discussion of Fortran character array portability in Appendix E.7.

I.7 Miscellaneous Routines

ADF_Flush_to_Disk — Flush the Data to the Disk

ADF_Flush_to_Disk (ID,error_return)		
Language	\mathbf{C}	Fortran
Routine Name	ADF_Flush_to_Disk	ADFFTD
Input	const double ID	real*8 ID
Output	int *error_return	integer error_return

ID The ID of the node to use.

This routine, ADF_Flush_to_Disk, flushes data to disk; it is used to force any modified information to be flushed to the physical disk. This ensures that data will not be lost if a program aborts. The control of when to flush all data to disk is provided to the user rather than flushing the data every time it is modified, which would result in reduced performance.

ADF_Database_Garbage_Collection — Flush the Data to the Disk

ADF_Database_Garbage_Collection (ID,error_return)		
Language	\mathbf{C}	Fortran
Routine Name	ADF_Database_Garbage_Collection ADFDGC	
Input	const double ID	real*8 ID
Output	int *error_return	integer error_return

ID The ID of the node to use.

This routine, ADF_Database_Garbage_Collection, redistributes the data within a file to use free space within the file. This free space is not located at the end of the file, and it may have been created during node deletions or other file operations. Neighboring free spaces will be merged.

Note: For better file compaction, a utility could be written to copy an ADF file into a newly created ADF file without wasted space.

Note: This routine is currently not implemented.

ADF_Error_Message — Get a Description of the Error

ADF_Error_Message (error_code,error_string)		
Language	C	Fortran
Routine Name	ADF_Error_Message	ADFERR
Input	const int error_code	integer error_code
Output	char *error_string	character*(*) error_string

error_code The error return code from any ADF routine. (See Appendix G.)

error_string An 80-byte description for the specified error. If ADF_Error_Message cannot find a message corresponding to the input error_code, the error string "Unknown

error#nnn" will be returned.

This routine, ADF_Error_Message, returns a textual error message when given the error return code from any ADF routine.

Example

See the example for ADF_Set_Error_State.

ADF_Set_Error_State — Set the Error State Flag

ADF_Set_Error_State (error_state,error_return)			
Language	C Fortran		
Routine Name	ADF_Set_Error_State ADFSES		
Input	const int error_state	integer error_state	
Output	int *error_return integer error_return		

error_state Flag specifying the action to take when an error occurs; 0 to return the error

code and continue, 1 to abort. The default is 0.

This routine, ADF_Set_Error_State, sets the error state flag. This flag controls the error handling convention for all ADF routines: either return the error codes or print an error message and abort the program on error. The flag covers all open ADF files associated with the current program, and it is not done on a file-by-file basis. The default state for the ADF interface is to return error codes and not abort.

Example

```
PROGRAM TEST
      PARAMETER (MAXCHR=32)
С
      CHARACTER*(MAXCHR) NODNAM
      CHARACTER*(80) MESS
С
C *** NODE IDS
С
      REAL*8 RID, CID1, CID2
С
C *** OPEN DATABASE
C
      CALL ADFDOPN('db.adf','NEW',' ',RID,IERR)
С
C *** CREATE 2 NODES
      CALL ADFCRE(RID, 'NODE 1', CID1, IERR)
      CALL ADFCRE(RID, 'NODE 2', CID2, IERR)
С
      CALL ADFGES (IDEFS, IERR)
      PRINT *,' DEFAULT ERROR STATE = ',IDEFS
C *** REQUEST NODE NAME FOR A NODE THAT DOES NOT EXIST
      PRINT *, ' *** ON ERROR CONTINUE'
```

```
CALL ADFGNAM(CID3, NODNAM, IERR)
         CALL ADFERR(IERR,MESS)
         PRINT *,'
                      ADF ERROR OCURRED, MESSAGE: ', MESS
         PRINT *,''
   С
  C *** SET ABORT ON ERROR FLAG
   С
         INEWS = 1
         CALL ADFSES (INEWS, IERR)
         PRINT *,' *** ABORT ON ERROR SET'
   С
   C *** REQUEST NODE NAME FOR A NODE THAT DOES NOT EXIST
   С
         CALL ADFGNAM(CID3, NODNAM, IERR)
         PRINT *,' HELLO WORLD'
   С
         STOP
         END
The resulting output is:
  DEFAULT ERROR STATE =
                                     0
     *** ON ERROR CONTINUE
         ADF ERROR OCURRED, MESSAGE:
   ADF 10: ADF file index out of legal range.
     *** ABORT ON ERROR SET
   ADF 10: ADF file index out of legal range.
   ADF Aborted: Exiting
```

${\tt ADF_Get_Error_State} \ -- \ \mathit{Get the Error State}$

ADF_Get_Error_State (error_state,error_return)			
Language	C Fortran		
Routine Name	ADF_Get_Error_State	ADFGES	
Input			
Output	int *error_state	integer error_state	
	int *error_return	integer error_return	

error_state Flag specifying the action to take when an error occurs; 0 to return the error code and continue, 1 to abort. The default is 0.

This routine, ADF_Get_Error_State, returns the currently set error state.

Example

```
PROGRAM TEST
C *** NODE IDS
С
      REAL*8 RID
\mathsf{C}
C *** OPEN DATABASE
С
      CALL ADFDOPN('db.adf','NEW',' ',RID,IERR)
С
C *** CHECK THE DEFAULT ERROR STATE
С
      CALL ADFGES (IDEFS, IERR)
С
C *** SET THE ERROR STATE TO ABORT ON ERROR
С
      CALL ADFSES(1, IERR)
С
C *** MAKE SURE STATE WAS SET AS DESIRED
С
      CALL ADFGES (NDEFS, IERR)
С
C *** PRINT OUT RESULTS
      PRINT *, ' DEFAULT ERROR STATE = ', IDEFS
      PRINT *, ' RESET ERROR STATE = ', NDEFS
С
      STOP
      END
```

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The resulting output is:

DEFAUL	T ERRO	OR STA	TE =	C
RESET	ERROR	STATE	: =	1

 $\begin{tabular}{ll} ADF_Database_Version $--$ Get the Version Number of the ADF Library That Created the ADF Database \end{tabular}$

ADF_Database_Version (root_ID, version, creation_date, modification_date, error_return)			
Language	C	Fortran	
Routine Name	ADF_Database_Version	ADFDVER	
Input	constant double root_ID	real*8 root_ID	
Output	char *version	character*(*) version	
	char *creation_date	character*(*) creation_date	
	char *modification_date	character*(*) modifica-	
		tion_date	
	int *error_return	integer error_return	

root_ID The ID of the root node in the ADF file.

version A 32-byte character string containing the version ID.

creation_date A 32-byte character string containing the creation date of the file.

modification_date A 32-byte character string containing the last modification date of the file.

error_return Error return code. (See Appendix G.)

This routine, ADF_Database_Version, returns the version number of the ADF library that created an ADF database, the file creation date and time, and the last modification date and time. A modified ADF database will take on the version ID of the current ADF library version if it is higher than the version indicated in the file. On return, "version" contains a six-character string. The meaning of the characters are described in detail in Appendix D.

Example

```
PROGRAM TEST
C
PARAMETER (MAXCHR=32)
C
CHARACTER*(MAXCHR) NODNAM
CHARACTER*(MAXCHR) CVER,LVER,CDATE,MDATE
CHARACTER*(80) MESS
C
C *** NODE IDS
C
REAL*8 RID,CID1,CID2
C
C *** OPEN DATABASE
C
CALL ADFDOPN('db.adf','NEW',','RID,IERR)
```

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```
C *** CREATE 2 NODES
C CALL ADFCRE(RID,'NODE 1',CID1,IERR)
CALL ADFCRE(RID,'NODE 2',CID2,IERR)
C CALL ADFDVER(RID,CVER,CDATE,MDATE,IERR)
CALL ADFLVER(LVER,IERR)
PRINT *,' VERSION INFORMATION:'
PRINT *,' ADF LIBRARY USED FOR CREATION: ',CVER
PRINT *,' CREATION DATE : ',CDATE
PRINT *,' MODIFICATION DATE : ',MDATE
PRINT *,' ADF LIBRARY BEING USED : ',LVER
C STOP
END
```

The resulting output is:

VERSION INFORMATION:

ADF LIBRARY USED FOR CREATION: ADF Database Version A01007 CREATION DATE : Thu Apr 24 15:41:55 1997 MODIFICATION DATE : Thu Apr 24 15:41:55 1997 ADF LIBRARY BEING USED : ADF Library Version C01

ADF_Library_Version — Get the Version Number of the ADF Library That the Application Program is Currently Using

ADF_Library_Version (version, error_return)			
Language	\mathbf{C}	Fortran	
Routine Name	ADF_Library_Version ADFLVER		
Input			
Output	char *version	character*(*) version	
	int *error_return integer error_return		

version A 32-byte character string containing the ADF library version ID information.

This routine, ADF_Library_Version, gets the ADF library version number. This is the version number of the ADF library that the application program is currently using. For this routine, the format of the version ID is "ADF Library Version AXXXxxx".

Example

See the example for ADF_Database_Version.

J Sample Fortran Program

The following Fortran program builds the ADF file shown in Figure 1 on p. 4.

```
PROGRAM TEST
С
С
      SAMPLE ADF TEST PROGRAM TO BUILD ADF FILES ILLUSTRATED
C
      IN THE EXAMPLE DATABASE FIGURE
C
      PARAMETER (MAXCHR=32)
С
      CHARACTER*(MAXCHR) TSTLBL,DTYPE
      CHARACTER*(MAXCHR) FNAM, PATH
C
      REAL*8 RID, PID, CID, TMPID, RIDF2
      REAL A(4,3),B(4,3)
      INTEGER IC(6), ID(6)
      INTEGER IERR
      INTEGER IDIM(2),IDIMA(2),IDIMC,IDIMD
С
      DATA A /1.1,2.1,3.1,4.1,
     X
              1.2,2.2,3.2,4.2,
              1.3,2.3,3.3,4.3/
      DATA IDIMA /4,3/
C
      DATA IC /1,2,3,4,5,6/
      DATA IDIMC /6/
С
C
      SET ERROR FLAG TO ABORT ON ERROR
С
      CALL ADFSES(1, IERR)
С
C *** 1.) OPEN 1ST DATABASE (ADF_FILE_TWO.ADF)
      2.) CREATE THREE NODES AT FIRST LEVEL
С
      3.) PUT LABEL ON NODE F3
      4.) PUT DATA IN F3
С
      5.) CREATE TWO NODES BELOW F3
      6.) CLOSE DATABASE
С
      CALL ADFDOPN('adf_file_two.adf','NEW',' ',RID,IERR)
      RIDF2 = RID
      CALL ADFCRE(RID, 'F1', TMPID, IERR)
      CALL ADFCRE(RID, 'F2', TMPID, IERR)
      CALL ADFCRE(RID, 'F3', PID, IERR)
      CALL ADFSLB(PID, 'LABEL ON NODE F3', IERR)
      CALL ADFPDIM(PID, 'R4', 2, IDIMA, IERR)
      CALL ADFWALL (PID, A, IERR)
C
      CALL ADFCRE(PID, 'F4', CID, IERR)
С
```

```
CALL ADFCRE(PID, 'F5', CID, IERR)
C
      CALL ADFDCLO(RID, IERR)
C
C *** 1.) OPEN 2ND DATABASE
С
      2.) CREATE NODES
\mathsf{C}
      3.) PUT DATA IN N13
C
      CALL ADFDOPN('adf_file_one.adf','NEW',' ',RID,IERR)
C
C
      THREE NODES UNDER ROOT
С
      CALL ADFCRE(RID, 'N1', TMPID, IERR)
      CALL ADFCRE(RID, 'N2', TMPID, IERR)
      CALL ADFCRE(RID, 'N3', TMPID, IERR)
C
С
      THREE NODES UNDER N1 (TWO REGULAR AND ONE LINK)
C
      CALL ADFGNID(RID,'N1',PID,IERR)
      CALL ADFCRE(PID, 'N4', TMPID, IERR)
      CALL ADFLINK(PID, 'L3', 'adf_file_two.adf', '/F3', TMPID, IERR)
      CALL ADFCRE(PID,'N5',TMPID,IERR)
С
С
      TWO NODES UNDER N4
      CALL ADFGNID (PID, 'N4', CID, IERR)
      CALL ADFCRE(CID,'N6',TMPID,IERR)
      CALL ADFCRE(CID, 'N7', TMPID, IERR)
C
С
      ONE NODE UNDER N6
С
      CALL ADFGNID (RID, '/N1/N4/N6', PID, IERR)
      CALL ADFCRE(PID, 'N8', TMPID, IERR)
С
С
      THREE NODES UNDER N3
С
      CALL ADFGNID(RID, 'N3', PID, IERR)
      CALL ADFCRE(PID, 'N9', TMPID, IERR)
      CALL ADFCRE(PID, 'N10', TMPID, IERR)
      CALL ADFCRE(PID, 'N11', TMPID, IERR)
С
С
      TWO NODES UNDER N9
С
      CALL ADFGNID(PID, 'N9', CID, IERR)
      CALL ADFCRE(CID, 'N12', TMPID, IERR)
      CALL ADFCRE(CID, 'N13', TMPID, IERR)
C
C
      PUT LABEL AND DATA IN N13
С
      CALL ADFSLB(TMPID, 'LABEL ON NODE N13', IERR)
      CALL ADFPDIM(TMPID, '14', 1, IDIMC, IERR)
```

```
CALL ADFWALL (TMPID, IC, IERR)
С
С
      TWO NODES UNDER N10
C
      CALL ADFGNID(RID, '/N3/N10', PID, IERR)
      CALL ADFLINK(PID, 'L1', '', '/N3/N9/N13', TMPID, IERR)
      CALL ADFCRE(PID, 'N14', TMPID, IERR)
С
C
      TWO NODES UNDER N11
C
      CALL ADFGNID(RID, '/N3/N11', PID, IERR)
      CALL ADFLINK(PID, 'L2', '', '/N3/N9/N13', TMPID, IERR)
      CALL ADFCRE(PID, 'N15', TMPID, IERR)
C
C *** READ AND PRINT DATA FROM NODES
С
      1.) NODE F5 THROUGH LINK L3
C
      CALL ADFGNID(RID, '/N1/L3', PID, IERR)
      CALL ADFGLB(PID, TSTLBL, IERR)
      CALL ADFGDT(PID, DTYPE, IERR)
      CALL ADFGND (PID, NUMDIM, IERR)
      CALL ADFGDV(PID, IDIM, IERR)
      CALL ADFRALL (PID, B, IERR)
      PRINT *,' NODE F3 THROUGH LINK L3:'
      PRINT *,' LABEL = ',TSTLBL
      PRINT *,' DATA TYPE = ',DTYPE
      PRINT *,' NUM OF DIMS = ', NUMDIM
      PRINT *,' DIM VALS
                            = ',IDIM
      PRINT *,' DATA:'
      WRITE(*,100)((B(J,I),I=1,3),J=1,4)
  100 FORMAT(5X,3F10.2)
C
C
      2.) N13
C
      CALL ADFGNID(RID, 'N3/N9/N13', PID, IERR)
      CALL ADFGLB(PID, TSTLBL, IERR)
      CALL ADFGDT(PID, DTYPE, IERR)
      CALL ADFGND (PID, NUMDIM, IERR)
      CALL ADFGDV(PID, IDIMD, IERR)
      CALL ADFRALL(PID,ID,IERR)
      PRINT *,' '
      PRINT *, ' NODE N13:'
      PRINT *,' LABEL
                              = ',TSTLBL
      PRINT *,' DATA TYPE = ',DTYPE
      PRINT *,' NUM OF DIMS = ', NUMDIM
      PRINT *,'
                DIM VALS
                             = ',IDIMD
      PRINT *,'
                DATA:'
      WRITE(*,200)(ID(I),I=1,6)
  200 FORMAT(5X,616)
C
      3.) N13 THROUGH L1
```

```
С
      CALL ADFGNID(RID, 'N3/N10/L1', TMPID, IERR)
      CALL ADFGLB (TMPID, TSTLBL, IERR)
      CALL ADFRALL (TMPID, ID, IERR)
      PRINT *,' '
      PRINT *,' NODE N13 THROUGH LINK L1:'
      PRINT *,' LABEL
                        = ',TSTLBL
      PRINT *,' DATA:'
      WRITE(*,200)(ID(I),I=1,6)
C
C
      4.) N13 THROUTH L2
C
      CALL ADFGNID(RID, 'N3/N11/L2', CID, IERR)
      CALL ADFGLB(CID,TSTLBL,IERR)
      CALL ADFRALL(CID, ID, IERR)
      PRINT *,' '
      PRINT *, ' NODE N13 THROUGH LINK L2:'
      PRINT *,' LABEL = ',TSTLBL
      PRINT *,' DATA:'
      WRITE(*,200)(ID(I),I=1,6)
C
С
     PRINT LIST OF CHILDREN UNDER ROOT NODE
С
      CALL PRTCLD(RID)
С
С
      PRINT LIST OF CHILDREN UNDER N3
C
      CALL ADFGNID (RID, 'N3', PID, IERR)
      CALL PRTCLD(PID)
С
C
     REOPEN ADF_FILE_TWO AND GET NEW ROOT ID
С
      CALL ADFDOPN('adf_file_two.adf','OLD',' ',RID,IERR)
      PRINT *,' '
      PRINT *, ' COMPARISON OF ROOT ID: '
      PRINT *, ' ADF_FILE_TWO.ADF ORIGINAL ROOT ID = ',RIDF2
     PRINT *, 'ADF_FILE_TWO.ADF NEW ROOT ID = ',RID
С
      STOP
      END
С
C ******* SUBROUTINES **********
C
      SUBROUTINE PRTCLD(PID)
С
C *** PRINT TABLE OF CHILDREN GIVEN A PARENT NODE-ID
     PARAMETER (MAXCLD=10)
     PARAMETER (MAXCHR=32)
     REAL*8 PID
      CHARACTER*(MAXCHR) NODNAM, NDNMS(MAXCLD)
```

```
CALL ADFGNAM(PID, NODNAM, IERR)
        CALL ADFNCLD(PID, NUMC, IERR)
        WRITE(*,120)NODNAM,NUMC
    120 FORMAT(/,' PARENT NODE NAME = ',A,/,
                    NUMBER OF CHILDREN = ',12,/,
       X
       X
                     CHILDREN NAMES: ')
        NLEFT = NUMC
        ISTART = 1
       --- TOP OF DO-WHILE LOOP
    130 CONTINUE
           CALL ADFCNAM(PID, ISTART, MAXCLD, LEN(NDNMS),
       Х
                        NUMRET, NDNMS, IERR)
           WRITE(*,140)(NDNMS(K),K=1,NUMRET)
     140
           FORMAT(8X,A)
           NLEFT = NLEFT - MAXCLD
           ISTART = ISTART + MAXCLD
        IF (NLEFT .GT. 0) GO TO 130
        RETURN
        END
The resulting output is:
  NODE F3 THROUGH LINK L3:
    LABEL
                = LABEL ON NODE F3
    DATA TYPE = R4
    NUM OF DIMS = 2
    DIM VALS
              = 4 3
    DATA:
            1.10
                    1.20
                               1.30
            2.10
                                2.30
                      2.20
            3.10
                      3.20
                                3.30
            4.10
                      4.20
                                4.30
  NODE N13:
                = LABEL ON NODE N13
    LABEL
    DATA TYPE = 14
    NUM OF DIMS = 1
    DIM VALS
              = 6
    DATA:
                  2
                         3 4 5
                                             6
           1
  NODE N13 THROUGH LINK L1:
    LABEL
                = LABEL ON NODE N13
    DATA:
                                4
                  2
                         3
                                      5
                                             6
           1
  NODE N13 THROUGH LINK L2:
           = LABEL ON NODE N13
    DATA:
           1
                2 3 4 5
                                             6
```

Advanced Data Format (ADF) User's Guide

```
PARENT NODE NAME = ADF MotherNode

NUMBER OF CHILDREN = 3

CHILDREN NAMES:

N1

N2

N3

PARENT NODE NAME = N3

NUMBER OF CHILDREN = 3

CHILDREN NAMES:

N9

N10

N11

COMPARISON OF ROOT ID:

ADF_FILE_TWO.ADF ORIGINAL ROOT ID = 1.2653021189994324-320

ADF_FILE_TWO.ADF NEW ROOT ID = 4.7783097267391979-299
```

K Sample C Program

The following C program builds the ADF file shown in Figure 1 on p. 4. /* Sample ADF test program to build adf files illustrated in example database figure. #include <stdio.h> #include <ctype.h> #include <string.h> #include "../include/ADF.h" void print_child_list(double node_id); main () /* --- Node header character strings */ char label[ADF_LABEL_LENGTH+1]; char data_type[ADF_DATA_TYPE_LENGTH+1]; char file_name[ADF_FILENAME_LENGTH+1]; char path[ADF_MAX_LINK_DATA_SIZE+1]; /* --- Node id variables */ double root_id,parent_id,child_id,tmp_id,root_id_file2; /* --- Data to be stored in database */ float $a[3][4] = {$ 1.1,2.1,3.1,4.1, 1.2,2.2,3.2,4.2, 1.3,2.3,3.3,4.3 int a_dimensions[2] = {4,3}; int $c[6] = \{1,2,3,4,5,6\};$ int c_dimension = 6; /* --- miscellaneous variables */ int error_flag, i, j; int error_state = 1; int num_dims, dim_d, d[6], dims_b[2]; float b[3][4]; /* ----- begin source code ----- */ /* --- set database error flag to abort on error */

ADF_Set_Error_State(error_state,&error_flag);

```
/* ----- build file: adf_file_two.adf ----- */
/* --- 1.) open database
       2.) create three nodes at first level
       3.) put label on node f3
       4.) put some data in node f3
       5.) create two nodes below f3
       6.) close database */
 ADF_Database_Open("adf_file_two.adf", "new", " ", &root_id, &error_flag);
 root_id_file2 = root_id;
 ADF_Create(root_id, "f1", &tmp_id, &error_flag);
 ADF_Create(root_id, "f2", &tmp_id, &error_flag);
 ADF_Create(root_id, "f3", &parent_id, &error_flag);
 ADF_Set_Label(parent_id, "label on node f3", &error_flag);
 ADF_Put_Dimension_Information(parent_id, "R4", 2, a_dimensions, &error_flag);
 ADF_Write_All_Data(parent_id,(char *)(a),&error_flag);
 ADF_Create(parent_id, "f4", &child_id, &error_flag);
 ADF_Create(parent_id, "f5", &child_id, &error_flag);
ADF_Database_Close(root_id,&error_flag);
/* ----- build file: adf_file_one.adf ----- */
/* open database and create three nodes at first level */
 ADF_Database_Open("adf_file_one.adf", "new", " ", &root_id, &error_flag);
 ADF_Create(root_id, "n1", &tmp_id, &error_flag);
 ADF_Create(root_id, "n2", &tmp_id, &error_flag);
 ADF_Create(root_id, "n3", &tmp_id, &error_flag);
/* put three nodes under n1 (two regular and one link) */
 ADF_Get_Node_ID(root_id, "n1", &parent_id, &error_flag);
 ADF_Create(parent_id, "n4", &tmp_id, &error_flag);
 ADF_Link(parent_id,"13","adf_file_two.adf","/f3",&tmp_id,&error_flag);
 ADF_Create(parent_id, "n5", &tmp_id, &error_flag);
/* put two nodes under n4 */
 ADF_Get_Node_ID(parent_id, "n4", &child_id, &error_flag);
 ADF_Create(child_id, "n6", &tmp_id, &error_flag);
ADF_Create(child_id,"n7",&tmp_id,&error_flag);
/* put one nodes under n6 */
 ADF_Get_Node_ID(root_id,"/n1/n4/n6",&parent_id,&error_flag);
ADF_Create(parent_id, "n8", &tmp_id, &error_flag);
/* put three nodes under n3 */
 ADF_Get_Node_ID(root_id, "n3", &parent_id, &error_flag);
 ADF_Create(parent_id, "n9", &tmp_id, &error_flag);
 ADF_Create(parent_id, "n10", &tmp_id, &error_flag);
 ADF_Create(parent_id, "n11", &tmp_id, &error_flag);
/* put two nodes under n9 */
```

```
ADF_Get_Node_ID(parent_id, "n9", &child_id, &error_flag);
ADF_Create(child_id, "n12", &tmp_id, &error_flag);
ADF_Create(child_id, "n13", &tmp_id, &error_flag);
/* put label and data in n13 */
ADF_Set_Label(tmp_id, "Label on Node n13", &error_flag);
ADF_Put_Dimension_Information(tmp_id,"i4",1,&c_dimension,&error_flag);
ADF_Write_All_Data(tmp_id,(char *)(c),&error_flag);
/* put two nodes under n10 (one normal, one link) */
ADF_Get_Node_ID(root_id, "/n3/n10", &parent_id, &error_flag);
ADF_Link(parent_id,"11"," ","/n3/n9/n13",&tmp_id,&error_flag);
ADF_Create(parent_id,"n14",&tmp_id,&error_flag);
/* put two nodes under n11 (one normal, one link) */
ADF_Get_Node_ID(root_id,"/n3/n11",&parent_id,&error_flag);
ADF_Link(parent_id,"12"," ","/n3/n9/n13",&tmp_id,&error_flag);
ADF_Create(parent_id, "n15", &tmp_id, &error_flag);
/* ----- finished building adf_file_one.adf ----- */
/* ----- access and print data ----- */
/* access data in node f3 (adf_file_two.adf) through link 13 */
ADF_Get_Node_ID(root_id,"/n1/13",&tmp_id,&error_flag);
ADF_Get_Label(tmp_id,label,&error_flag);
ADF_Get_Data_Type(tmp_id,data_type,&error_flag);
ADF_Get_Number_of_Dimensions(tmp_id,&num_dims,&error_flag);
ADF_Get_Dimension_Values(tmp_id,dims_b,&error_flag);
ADF_Read_All_Data(tmp_id,(char *)(b),&error_flag);
printf (" node f3 through link l3:\n");
printf ("
            label
                        = %s\n",label);
printf ("
            data_type = %s\n",data_type);
printf (" num of dims = %5d\n",num_dims);
                       = %5d %5d\n",dims_b[0],dims_b[1]);
printf ("
            dim vals
printf (" data:\n");
for (i=0; i<=3; i++)
    for (j=0; j<=2; j++)
        printf("
                   %10.2f",b[j][i]);
    printf("\n");
  };
/* access data in node n13 */
ADF_Get_Node_ID(root_id, "/n3/n9/n13", &tmp_id, &error_flag);
ADF_Get_Label(tmp_id,label,&error_flag);
ADF_Get_Data_Type(tmp_id,data_type,&error_flag);
ADF_Get_Number_of_Dimensions(tmp_id,&num_dims,&error_flag);
ADF_Get_Dimension_Values(tmp_id,&dim_d,&error_flag);
```

```
ADF_Read_All_Data(tmp_id,(char *)(d),&error_flag);
printf (" node n13:\n");
                         = %s\n", label);
printf ("
            label
printf ("
            data_type = %s\n",data_type);
            num of dims = %5d\n",num_dims);
printf ("
printf ("
                        = %5d\n",dim_d);
            dim val
printf ("
            data:\n");
for (i=0; i<=5; i++)
    printf("
                 %-4d",d[i]);
  };
printf("\n\n");
/* access data in node n13 through l1 */
ADF_Get_Node_ID(root_id, "/n3/n10/l1", &tmp_id, &error_flag);
ADF_Get_Label(tmp_id,label,&error_flag);
 ADF_Read_All_Data(tmp_id,(char *)(d),&error_flag);
printf (" node n13 through l1:\n");
printf ("
            label
                        = %s\n", label);
printf (" data:\n");
 for (i=0; i<=5; i++)
  {
    printf("
                 %-4d",d[i]);
printf("\n\n");
/* access data in node n13 through 12 */
ADF_Get_Node_ID(root_id,"/n3/n11/12",&tmp_id,&error_flag);
 ADF_Get_Label(tmp_id,label,&error_flag);
ADF_Read_All_Data(tmp_id,(char *)(d),&error_flag);
printf (" node n13 through 12:\n");
printf ("
            label
                         = %s\n'', label);
printf (" data:\n");
for (i=0; i<=5; i++)
    printf("
                %-4d",d[i]);
  };
printf("\n\n");
/* print list of children under root node */
print_child_list(root_id);
/* print list of children under n3 */
 ADF_Get_Node_ID(root_id,"/n3",&tmp_id,&error_flag);
print_child_list(tmp_id);
/* re-open adf_file_two and get new root id */
 ADF_Database_Open("adf_file_two.adf","old"," ",&root_id,&error_flag);
printf (" Comparison of root id:\n");
printf (" adf_file_two.adf original root id = %x\n",root_id_file2);
printf (" adf_file_two.adf new root id = %x\n",root_id);
```

```
}
   void print_child_list(double node_id)
      print table of children given a parent node-id
      char node_name[ADF_NAME_LENGTH+1];
      int i, num_children, num_ret, error_return;
      ADF_Get_Name(node_id,node_name,&error_return);
      ADF_Number_of_Children(node_id,&num_children,&error_return);
      printf ("Parent Node Name = %s\n",node_name);
      printf (" Number of Children = %2d\n",num_children);
      printf (" Children Names:\n");
      for (i=1; i<=num_children; i++)</pre>
          ADF_Children_Names(node_id,i,1,ADF_NAME_LENGTH+1,
              &num_ret,node_name,&error_return);
          printf ("
                        %s\n",node_name);
                                                                          };
       printf ("\n");
  }
The resulting output is:
  node f3 through link 13:
     label
                 = label on node f3
     data_type = R4
     num of dims =
     dim vals
                             3
     data:
             1.10
                            1.20
                                           1.30
             2.10
                            2.20
                                           2.30
                            3.20
             3.10
                                           3.30
             4.10
                            4.20
                                           4.30
  node n13:
     label
                 = Label on Node n13
               = i4
     data_type
     num of dims = 1
     dim val
                 = 6
     data:
       1
                2
                         3
                                  4
                                           5
                                                     6
  node n13 through 11:
                = Label on Node n13
     label
     data:
       1
                2
                         3
                                  4
                                           5
                                                     6
  node n13 through 12:
     label
             = Label on Node n13
```

```
data:
                    3
                             4 5 6
   1
Parent Node Name = ADF MotherNode
 Number of Children = 3
 Children Names:
    n1
    n2
    n3
Parent Node Name = n3
 Number of Children = 3
  Children Names:
    n9
    n10
    n11
Comparison of root id:
 adf_file_two.adf original root id = 18
  adf_file_two.adf new root id = 2a
```