单周期处理器实现插入排序

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一、 实验目的

使用单周期处理器完成插入排序,并将结果显示在七位字码管上。

二、设计方案

采用单周期处理器(同理论课)实现 MIPS 指令集子集,将插入排序 c++ 程序翻译为 MIPS 指令后通过单周期处理器运行,最后将排序后结果 逐个输出到字码管外设地址,使用循环执行无效指令的方式实现 delay 从而达到每个数据维持显示 1s。

三、 算法指令

MIPS 汇编语言实现插入排序的源代码见附录 1。

四、 关键代码清单

ALU.v: 运算单元

ALUControl.v: 运算控制单元

Control.v:CPU 控制单元

CPU.v: 顶层文件

DataMemory.v: 数据存储器

InstructionMemory.v: 指令存储器

RegisterFile.v: 寄存器

 $test_cpu.v:testBench$

xdc_for_singleCycle.xdc: 约束文件insert_sort.asm: 插入排序汇编文件

五、 仿真结果及分析

仿真得到 RAM_data[255:0] 和 LED[31:0] 部分结果如图所示:



图 1: RAM_data[255:0]



图 2: LED[31:0]

根据图 1 中可以看出插入排序的过程,根据图 2 可以看出 LED[8]、LED[9]、LED[10]、LED[11] 依此为 1,其它为零,从而实现字码管动态显示的效果。

六、 综合情况 3

六、 综合情况

硬件资源占用情况和时序报告如图 3 图 4 所示:

Name ^1	Slice LUTs (20800)	Slice Registers (41600)	F7 Muxes (16300)	F8 Muxes (8150)	Slice (8150)	LUT as Logic (20800)	DSPs (90)	Bonded IOB (250)	BUFGCTRL (32)	MMCME2_ADV (5)
∨ N top	4080	9260	1351	512	4096	4080	3	14	2	1
∨ I clk_out_0 (clk_wiz_0)	0	0	0	0	0	0	0	0	2	1
I inst (clk_wiz_0_clk_v	0	0	0	0	0	0	0	0	2	1
∨ I nolabel_line10 (CPU)	4080	9260	1351	512	4096	4080	3	0	0	0
I alu1 (ALU)	27	0	3	0	15	27	3	0	0	0
■ data_memory1 (Data	2261	8236	1090	512	3408	2261	0	0	0	0
	1779	992	258	0	908	1779	0	0	0	0

图 3: 硬件资源占用情况

tup		Hold		Pulse Width					
Worst Negative Slack (WNS):	22.047 ns	Worst Hold Slack (WHS):	0.321 ns	Worst Pulse Width Slack (WPWS):	3.000 ns				
Total Negative Slack (TNS):	0.000 ns	Total Hold Slack (THS):	0.000 ns	Total Pulse Width Negative Slack (TPWS):	0.000 ns				
Number of Failing Endpoints:	0	Number of Failing Endpoints:	0	Number of Failing Endpoints:	0				
Total Number of Endpoints:	18488	Total Number of Endpoints:	18488	Total Number of Endpoints:	9266				

图 4: 时序分析报告

最高工作频率:

$$f_{max} = \frac{1}{T-WNS} = \frac{1}{1/(20\times 10^6) - 22.047\times 10^{-9}} = 35.77MHz$$

下面计算 CPI (仅统计排序部分): MARS 统计插入排序总共共执行指令数如下图,总指令数

$$N = 2417$$

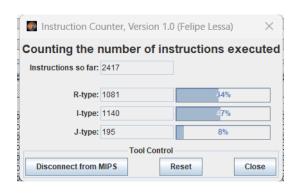


图 5: MIPS 指令数统计

六、 综合情况

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排序部分仿真结果如下图,总用时

$$t=214.7us$$

周期数

$$C = \frac{t}{T} = \frac{214.7 \times 10^{-6}}{0.05 \times 10^{-6}} = 4294$$

										214.700000 us	
lame	Value			100.00	0000 us		150.0000	000 us	200.0000	10 us	250.000000 us
> 😻 [30][31:0]	00000000						0	0000000			
> 🕨 [29][31:0]	00000000						0	0000000			
> 🕨 [28][31:0]	00000000						0	0000000			
> 🕨 [27][31:0]	00000000	-					0	0000000		'	
> 🕨 [26][31:0]	00000000						0	0000000			
> 🕨 [25][31:0]	00000000						0	0000000			
> 🕨 [24][31:0]	00000000						0	0000000			
> 🕨 [23][31:0]	00000000						0	0000000		'	
> 🕨 [22][31:0]	00000000						0	0000000			
> 🕨 [21][31:0]	00000000						0	0000000			
> 😻 [20][31:0]	0000dbe3				000	04955			X		0000dbe3
> 💗 [19][31:0]	0000dac9				0000c39	9			(0000)		0000dac9
> 🕨 [18][31:0]	0000d4b4			0	0004112				000		0000d4h4
> 🕨 [17][31:0]	0000c399			0000391	8			0000dbe3	000		0000c399
> 🕨 [16][31:0]	0000ь783		000	03748			0000dbe3	0000dac9	(000)		0000Ъ783
> 🕨 [15][31:0]	0000acda		00004414			0000	0000dac9	00004454	00001783		0000acda
> 🕨 [14][31:0]	00009899		0000dbe	3		0000	00004414	0000ъ783	0000acda		00009899
> 🕶 [13][31:0]	00008ed9	0000	128d		0000dac9	0000	0000ъ783	0000acda	00009899		00008ed9
W [12][31:0]	00004955	00003c56		0000dac9	0001	0ъ783	0000acda	00009891	00008ed9		00004955
> 🕨 [11][31:0]	000041a8	00009899	00	0000ъ783	000	10acda	00009899	00008ed	9 /	000041	18
> 🕨 [10][31:0]	00004112	0000044e \0000	. (00	0000acda	(00	009899	00008ed9	000041	18	000041	12

图 6: 排序部分仿真结果

因此插入排序部分 CPI 计算如下:

$$CPI = \frac{C}{N} = \frac{4294}{2417} = 1.7766$$

七、 硬件调试情况

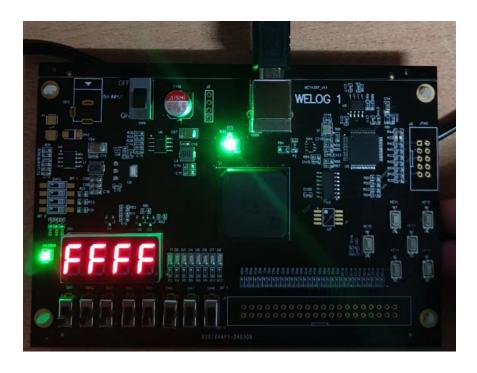


图 7: 实际运行情况

代码可在 FPGA 上正常运行,程序以显示 FFFF 作为结束标志,如图 所示。

八、 附录

(一) MIPS 实现插入排序

```
main:
    addi $s0 $zero 0 #buffer[]
    addi $s2 $zero 0
    lw $s1 0($s0) #($s1)int N = buffer[0]
    addi $v1 $s1 0
insert_sort:? »
    addi $t0 $s0 4 #$t0 = &(buffer[1])
```

```
addi t1 \ s1 \ 0 \ \#t1 = N
    addi t2 \cdot zero 1 \#(t2) int i = 1
insert sort loop:
    beq $t1 $t2 end_insert_sort_loop
search:
    addi $t4 $t2 -1 $\#($t4)$ i = n - 1
    s11 \$t5 \$t2 2 \#\$t5 = 4n
    add $t6 $t0 $t5 \#(\$t6) int tmp = v[n];
    lw $t6 ($t6)
search_loop:
    blt $t4 $zero end_search_loop
    addi $s2 $s2 1
    sll $t5 $t4 2 \#$t5 = 4i
    add t7 t0 t5 #t7 = v[i]
    lw $t7 ($t7)
    ble $t7 $t6 end_search_loop #if (v[i] <= tmp) break;
    addi $t4 $t4 −1 #i—
    j search_loop
end_search_loop:
    addi v0 t4 1 #return i + 1
insert:
    addi $t4 $t2 -1 \#(\$t4) i = n -1
    s11 \$t5 \$t2 2 \#\$t5 = 4n
    add $t6 $t0 $t5 \#(\$t6) int tmp = v[n];
    lw $t6 ($t6)
    addi t7 v0 0 \#t7 = k
insert_loop:
    blt $t4 $t7 end_insert_loop
    sll $t5 $t4 2 #$t5 = 4i
    add $t8 $t0 $t5
    lw \$s4 (\$t8) \#\$s4 = v[i]
    addi $t9 $t8 4 \#$t9 = \&v[i+1]
    sw \$s4 (\$t9) \#v[i+1] = v[i]
```

```
addi $t4 $t4 −1 #i---
    j insert_loop
end_insert_loop:
    sll $t7 $t7 2
    add $t9 $t0 $t7
    sw $t6 \ 0($t9) \ \#v[k] = tmp;
    addi $t2 $t2 1 #i++
    j insert_sort_loop
end_insert_sort_loop:
    sw \$s2 0(\$s0)
    addi $s3 $zero 0 #count
    addi $t0 $zero 4 #&buffer[i]
    addi $k1 $zero 2000 #delay_times
loop:
    addi $t7 $zero 900 #Times
    addi $t8 $zero 0 #OpTimes
decode_loop:
    addi $a0 $zero 16384 #led
    addi $a2 $zero 1 #led_count
    lw $t1 \ 0($t0) \ \#[15:0] in = buffer[i]
    andi v0 $t1 15 \#[3:0] in
    j decode
led1:
    addi $s7 $zero 256 #an[0]
    or $v0 $v0 $s7 #led
    sw $v0 ($a0)
    addi $k0 $zero 0
delay1:
    addi $k0 $k0 1
    beq $k0 $k1 end1
```

```
j delay1
end1:
    lw $t1 0($t0)
    srl $v0 $t1 4 #[7:4] in
    andi $v0 $v0 15
    addi a2 a2 1 \#led\_count += 1
    j decode
led2:
    addi $s7 $zero 512 #an[1]
    or v0 v0 sr \#led
    sw $v0 ($a0)
    addi $k0 $zero 0
delay2:
    addi $k0 $k0 1
    beq k0 \ k1 \ end2
    j delay2
end2:
    lw $t1 0($t0)
    srl $v0 $t1 8 #[11:8]in
    andi $v0 $v0 15
    addi $a2 $a2 1 #led_count += 1
    j decode
led3:
    addi $s7 $zero 1024 #an[2]
    or v0 \ v0 \ s7 \ \#led
    sw $v0 ($a0)
    addi $k0 $zero 0
delay3:
    addi $k0 $k0 1
```

```
beq $k0 $k1 end3
    j delay3
end3:
    lw $t1 0($t0)
    srl $v0 $t1 12 #in[15:12]
    andi v0 v0 15
    addi a2 1 \#led_count +=1
    j decode
led4:
    addi $s7 $zero 2048 #an[3]
    or $v0 $v0 $s7
   sw $v0 ($a0)
    addi $k0 $zero 0
delay4:
    addi $k0 $k0 1
    beq $k0 $k1 end4
    j delay4
end4:
    addi $t8 $t8 1
    bne $t8 $t7 decode_loop #if(OpTimes < Times) decode(in)
    addi $s3 $s3 1 #count++
    addi t0 t0 4 \#i++
    beq s3 v1 end \#if(count = N) break
    j loop
decode:
    addi $a1 $zero 0 #decode_count
    beq $v0 $a1 decode0
```

addi \$a1 \$a1 1 beg \$v0 \$a1 decode1 addi \$a1 \$a1 1 beg \$v0 \$a1 decode2 addi \$a1 \$a1 1 beq \$v0 \$a1 decode3 addi \$a1 \$a1 1 beq \$v0 \$a1 decode4 addi \$a1 \$a1 1 beq \$v0 \$a1 decode5 addi \$a1 \$a1 1 beq \$v0 \$a1 decode6 addi \$a1 \$a1 1 beq \$v0 \$a1 decode7 addi \$a1 \$a1 1 beq \$v0 \$a1 decode8 addi \$a1 \$a1 1 beq \$v0 \$a1 decode9 addi \$a1 \$a1 1 beq \$v0 \$a1 decodeA addi \$a1 \$a1 1 beg \$v0 \$a1 decodeB addi \$a1 \$a1 1 beq \$v0 \$a1 decodeC addi \$a1 \$a1 1 beq \$v0 \$a1 decodeD addi \$a1 \$a1 1 beq \$v0 \$a1 decodeE addi \$a1 \$a1 1 beq \$v0 \$a1 decodeF decode0: addi \$v0 \$zero 63 j decode_jump

```
decode1:
    addi $v0 $zero 6
    j decode_jump
decode2:
    addi $v0 $zero 91
    j decode_jump
decode3:
    addi $v0 $zero 79
    j decode_jump
decode4:
    addi $v0 $zero 102
    j decode_jump
decode5:
    addi $v0 $zero 109
    j decode_jump
decode6:
    addi $v0 $zero 125
    j decode_jump
decode7:
    addi $v0 $zero 7
    j decode_jump
decode8:
    addi $v0 $zero 127
    j decode_jump
decode9:
    addi $v0 $zero 111
    j decode_jump
decodeA:
    addi $v0 $zero 119
    j decode_jump
decodeB:
    addi $v0 $zero 124
    j decode_jump
```

```
decodeC:
    addi $v0 $zero 88
    j decode_jump
decodeD:
    addi $v0 $zero 94
    j decode_jump
decodeE:
    addi $v0 $zero 121
    j decode_jump
decodeF:
    addi $v0 $zero 113
    j decode_jump
decode_jump:
    addi $t9 $zero 1
    beq \$a2 \$t9 \ led1
    addi $t9 $t9 1
    beq $a2 $t9 led2
    addi $t9 $t9 1
    beq $a2 $t9 led3
    addi $t9 $t9 1
    beq \$a2 \$t9 \ led4
end:
    addi $t0 $zero 3953
    sw $t0 ($a0)
    j end
```