

CME 213, ME 339—Spring 2021

Eric Darve, ICME, Stanford



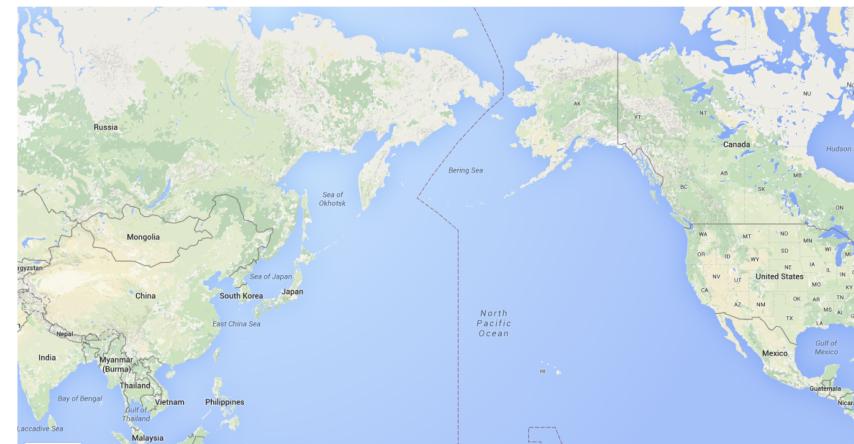
“Software is like entropy: It is difficult to grasp, weighs nothing, and obeys the Second Law of Thermodynamics; i.e., it always increases.” (Norman Augustine)

Concurrency and latency

Imagine you are a pencil manufacturer.

You outsource your manufacturing plants to China but your market is in the US.

How do you organize the logistics of the transport?

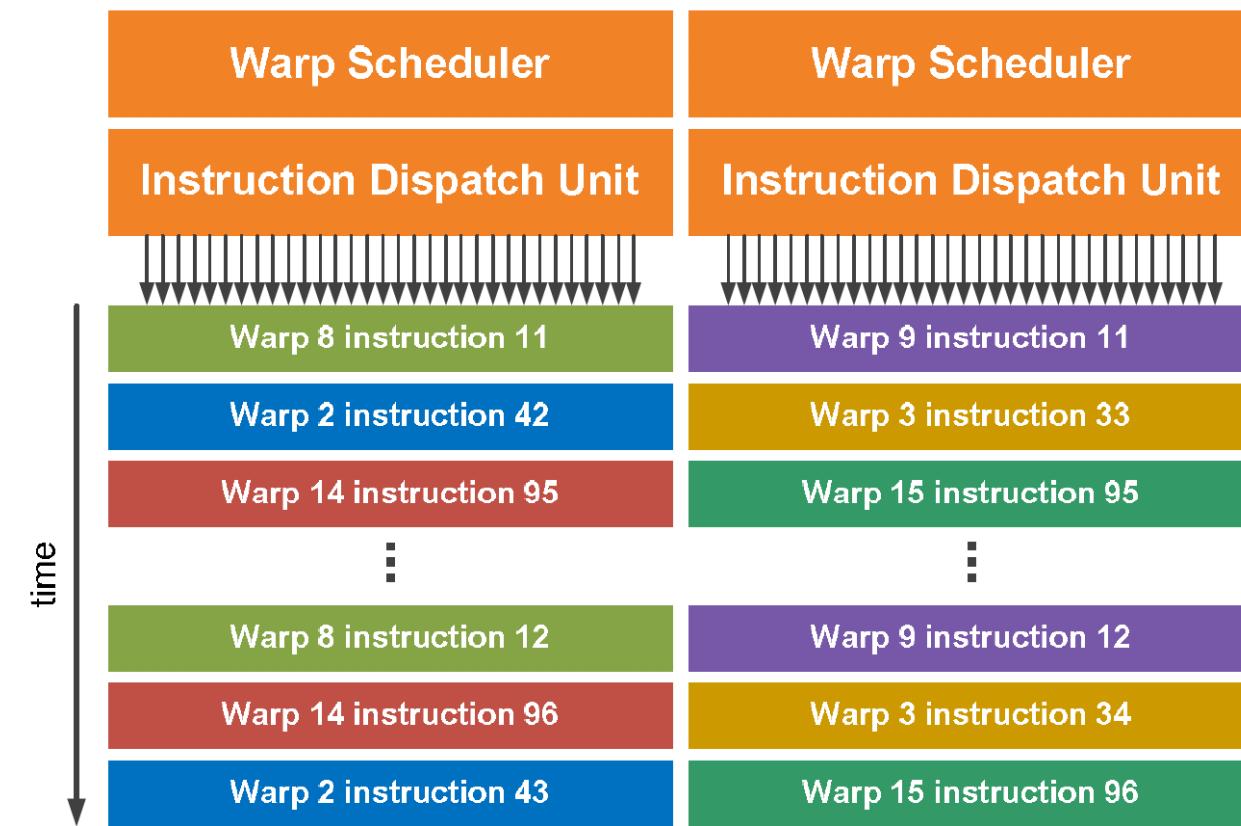


Concurrency is used to hide long latencies:

- memory access
- floating point units
- any long sequence of operations

Processors are optimized in the same way

Hide latency through concurrency



How to maximize concurrency?

- Have as many live threads as possible
- Instruction-level parallelism

Hardware limits

- Max dim. of grid: y/z 65,535 (x is large, $2^{31} - 1$)
- Max dim. of block: x/y 1,024; z 64
- Max # of threads per block: 1,024
- Max blocks per SM: 16
- Max resident warps: 32
- Max threads per SM: 1,024
- # of 4-byte registers per SM: 64 K
- Max shared mem per SM: 64 KB

[Spreadsheet](#)

How can we make sense of this?

CUDA API

Achieve best potential occupancy; recommended parameter selections

`cudaOccupancyMaxActiveBlocksPerMultiprocessor`

Reports an occupancy prediction based on the block size and shared memory usage of a kernel

`cudaOccupancyMaxPotentialBlockSize`

`cudaOccupancyMaxPotentialBlockSizeVariableSMem`

Returns the minimum grid size and recommended block size to achieve maximum occupancy

Occupancy spreadsheet!

[CUDA Occupancy Calculator](#)

CUDA Occupancy Calculator

Just follow steps 1, 2, and 3 below! (or [click here for help](#))

1.) Select Compute Capability (click): **7.5** [\(Help\)](#)

2.) Enter your resource usage:
 Threads Per Block **256** [\(Help\)](#)
 Registers Per Thread **30** [\(Help\)](#)
 User Shared Memory Per Block (bytes) **4224** [\(Help\)](#)

(Don't edit anything below this line)

3.) GPU Occupancy Data is displayed here and in the graphs:
 Active Threads per Multiprocessor **1024** [\(Help\)](#)
 Active Warps per Multiprocessor **32**
 Active Thread Blocks per Multiprocessor **4**
 Occupancy of each Multiprocessor **100%**

Physical Limits for GPU Compute Capability: **7.5**

Threads per Warp	32
Max Warps per Multiprocessor	32
Max Thread Blocks per Multiprocessor	16
Max Threads per Multiprocessor	1024
Maximum Thread Block Size	1024
Registers per Multiprocessor	65536
Max Registers per Thread Block	65536
Max Registers per Thread	255
Shared Memory per Multiprocessor (bytes)	65536
Max Shared Memory per Block	65536
Register allocation unit size	256
Register allocation granularity	warp
Shared Memory allocation unit size	256
Warp allocation granularity	4
Shared Memory Per Block (bytes) (CUDA runtime use)	0

Allocated Resources

	Per Block	Limit Per SM	= Allocatable Blocks Per SM
Warp (Threads Per Block / Threads Per Warp)	8	32	4
Registers (Warp limit per SM due to per-warp reg count)	8	64	8
Shared Memory (Bytes)	4352	65536	15

Note: SM is an abbreviation for (Streaming) Multiprocessor

Maximum Thread Blocks Per Multiprocessor

Blocks/SM	* Warps/Block	= Warps/SM
4	8	32

Limited by Max Warps or Max Blocks per Multiprocessor **4** **8** **32**

Limited by Registers per Multiprocessor **8**

Limited by Shared Memory per Multiprocessor **15**

Note: Occupancy limiter is shown in orange

Physical Max Warps/SM = 32
 Occupancy = 32 / 32 = 100%

CUDA Occupancy Calculator

Version: **11.1**

[Copyright and License](#)

Click Here for detailed instructions on how to use this occupancy calculator.
[For more information on NVIDIA CUDA, visit http://developer.nvidia.com/cuda](#)

Your chosen resource usage is indicated by the red triangle on the graphs. The other data points represent the range of possible block sizes, register counts, and shared memory allocation.

Impact of Varying Block Size

My Block Size, 256

Multiprocessor/Warp Occupancy

Threads Per Block

Impact of Varying Shared Memory Usage Per Block

Shared Memory, 4352

Multiprocessor/Warp Occupancy

Shared Memory Per Block

Impact of Varying Register Count Per Thread

My Register Count, 30

Multiprocessor/Warp Occupancy

Registers Per Thread

12 / 61

Command line

```
$ nvcc --ptxas-options=-v -o transpose transpose.cu
```

--ptxas-options=-v

Output

```
ptxas info    : Compiling entry function '...fastTranspose...' for 'sm_75'  
ptxas info    : Function properties for ...fastTranspose...  
    0 bytes stack frame, 0 bytes spill stores, 0 bytes spill loads  
ptxas info    : Used 30 registers, 4224 bytes smem, 384 bytes cmem[0]
```

Occupancy calculator demo

CUDA Occupancy Calculator

Just follow steps 1, 2, and 3 below! (or click here for help)

1.) Select Compute Capability (click):

7.5

(Help)

1.b) Select Shared Memory Size Config (bytes)

65536

2.) Enter your resource usage:

Threads Per Block

256

(Help)

Registers Per Thread

30

User Shared Memory Per Block (bytes)

4224

(Don't edit anything below this line)

3.) GPU Occupancy Data is displayed here and in the graphs:

Active Threads per Multiprocessor

1024

(Help)

Active Warps per Multiprocessor

32

Active Thread Blocks per Multiprocessor

4

Occupancy of each Multiprocessor

100%

Physical Limits for GPU Compute Capability:

7.5

Threads per Warp

32

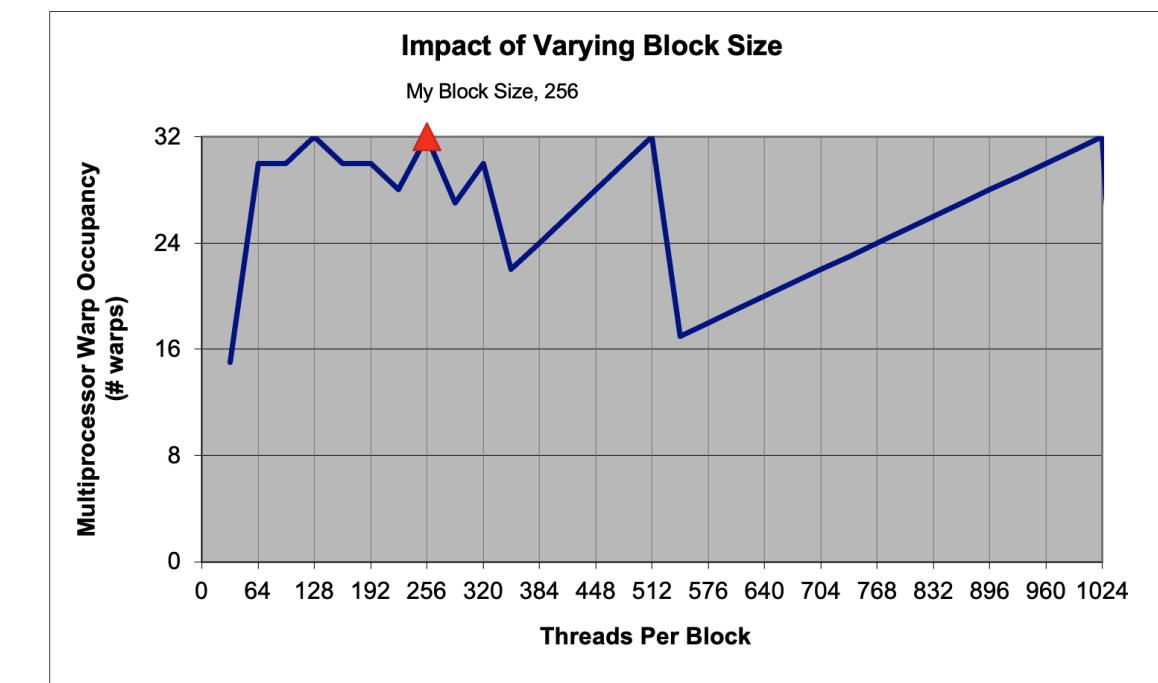
Max Warps per Multiprocessor

32

[Click Here for detailed instructions on how to use this occupancy calculator.](#)

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Branching and divergent execution path

32 threads = 1 warp

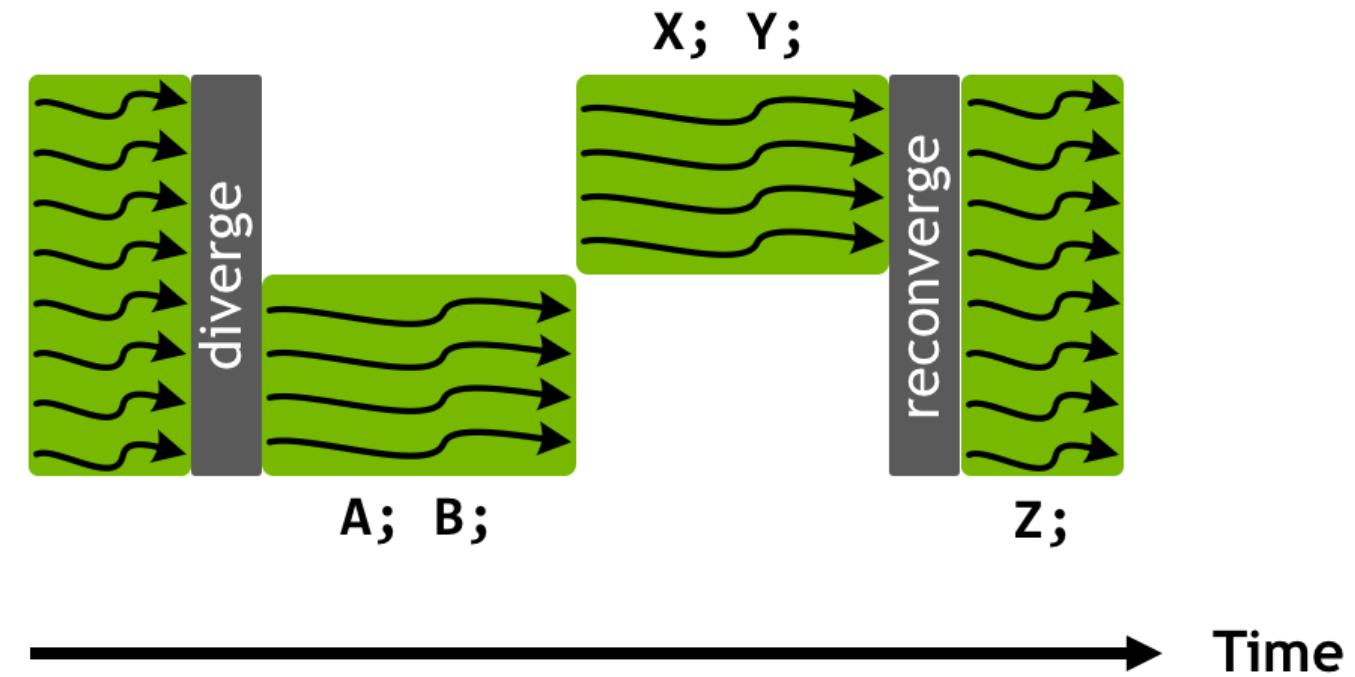
SIMT: Single Instruction, Multiple Thread

Up to Pascal

Single program counter shared amongst all 32 threads

What happens with branches?

```
if (threadIdx.x < 4) {  
    A;  
    B;  
} else {  
    X;  
    Y;  
}  
Z;
```



Active mask

Specifies which threads of the warp are active at any given time

Execution becomes serialized

```
__global__ void branch_thread(float* out){  
    int tid = threadIdx.x;  
    if (tid%2 == 0) {  
        ...;  
    } else {  
        ...;  
    }  
}
```



```
--global__ void branch_warp(float* out){  
    int wid = threadIdx.x/32;  
    if (wid%2 == 0) {  
        ...;  
    } else {  
        ...;  
    }  
}
```



Volta!

Mitigates this problem a bit

But fundamentally performance hit is still significant

Pre-Volta

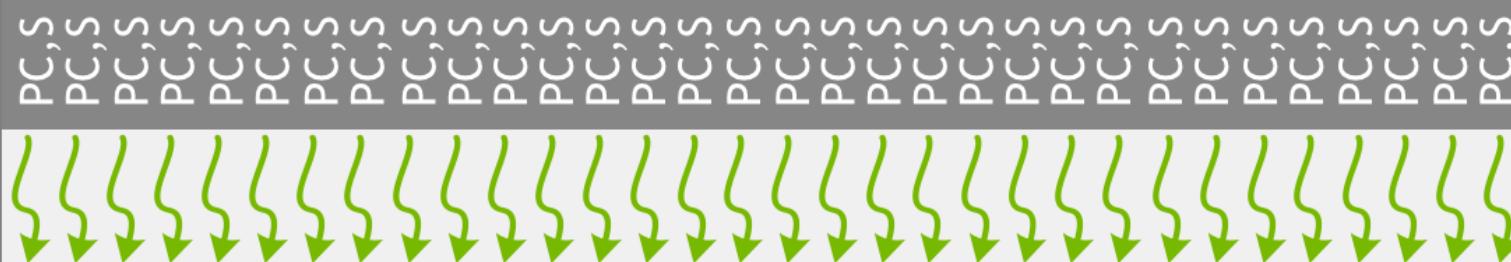
Program
Counter (PC)
and Stack (S)



32 thread warp

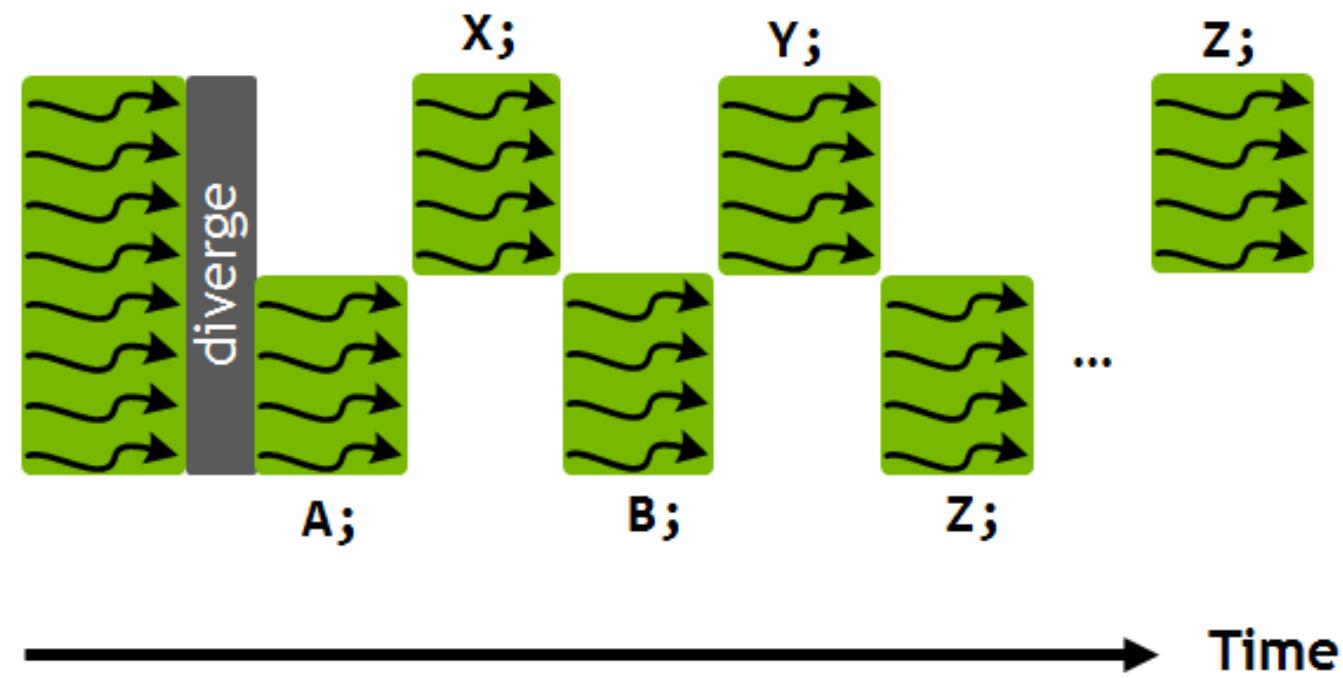
Volta

Convergence
Optimizer



32 thread warp with independent scheduling

```
if (threadIdx.x < 4) {  
    A;  
    B;  
} else {  
    X;  
    Y;  
}  
Z;
```



Gives more flexibility in programming and performance.

Avoids certain deadlock bugs due to divergence.

Volta allows the correct execution of many concurrent algorithms.

Definition: **starvation-free algorithm**

Algorithms that are guaranteed to execute correctly so long as the system ensures that all threads have adequate access to a contended resource.

```
__device__ void insert_after(Node *a, Node *b)
{
    Node *c;
    lock(a); lock(a->next);
    c = a->next;

    a->next = b;
    b->prev = a;

    b->next = c;
    c->prev = b;

    unlock(c); unlock(a);
}
```

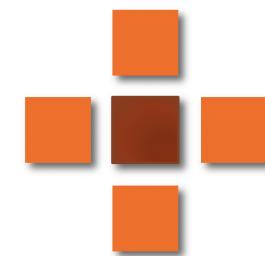
Volta ensures that the thread holding the lock is able to make progress.

Homework 4

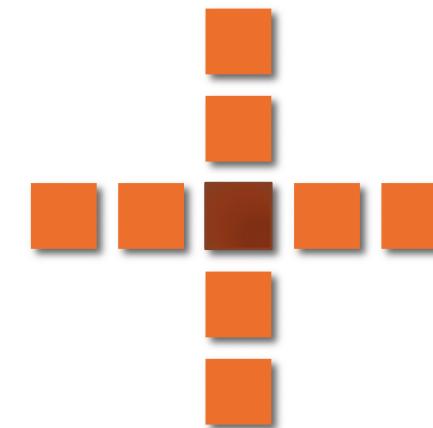
$$\frac{\partial T}{\partial t} = \frac{\partial^2 T}{\partial x^2} + \frac{\partial^2 T}{\partial y^2}$$

Finite-difference

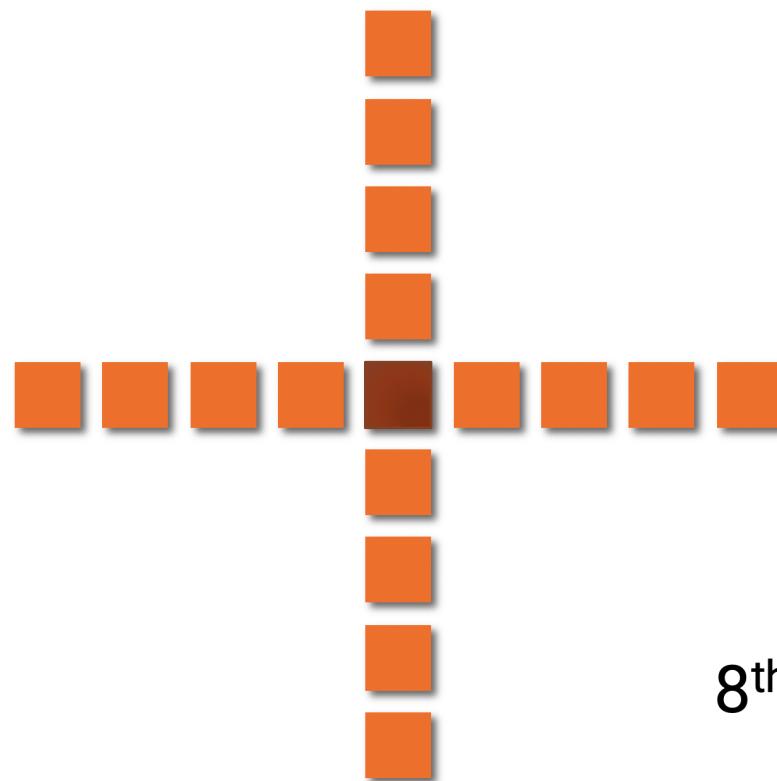
$$T^{n+1} = AT^n$$



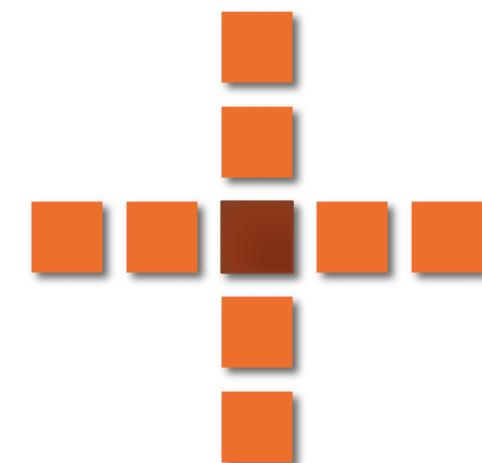
2nd order stencil



4th order stencil

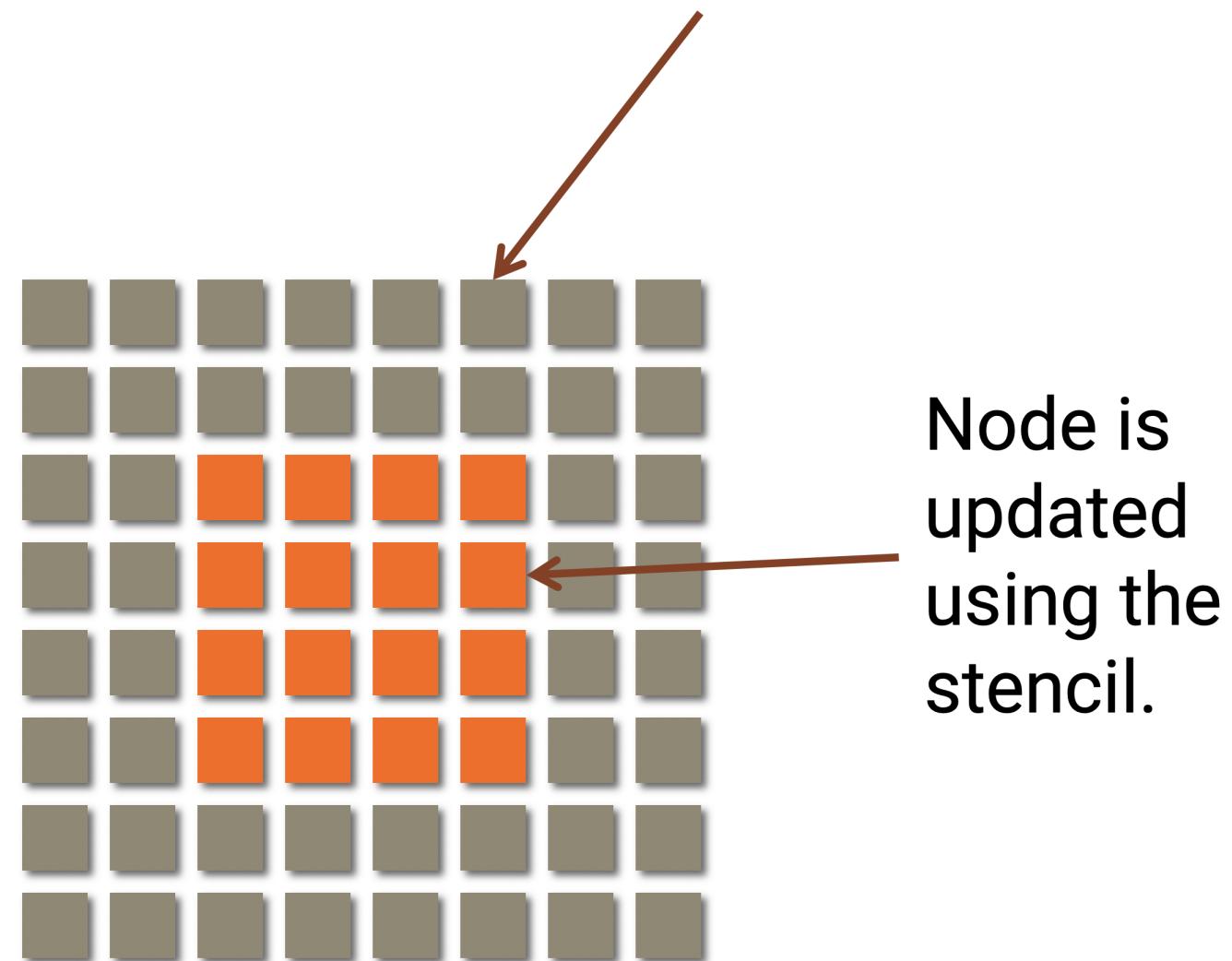


8th order stencil



4th order stencil

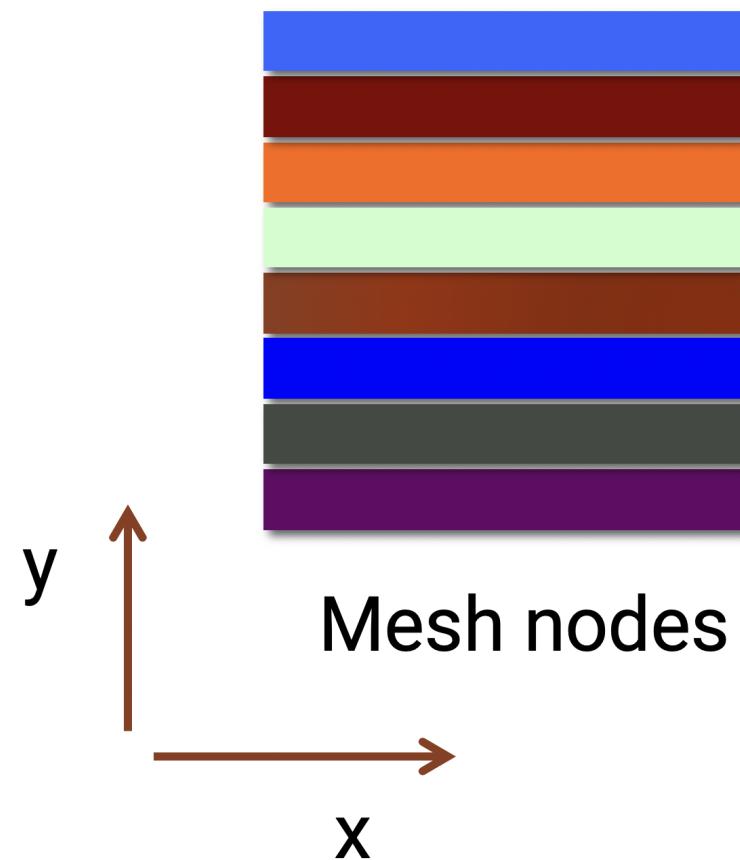
Updated in a separate routine;
boundary conditions are used



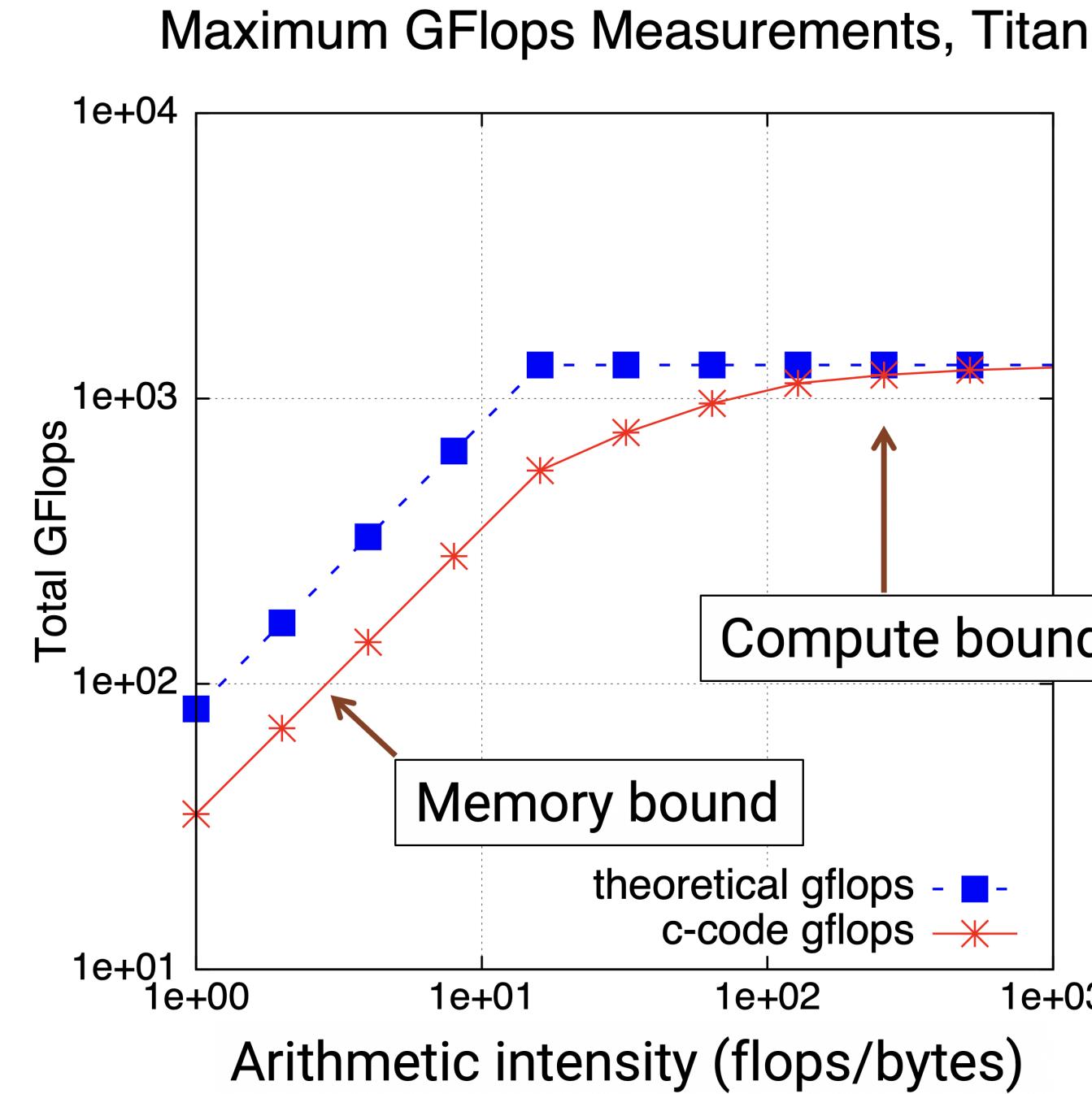
Node is
updated
using the
stencil.

Goal of the homework

Implement a CUDA routine to update nodes inside the domain using a finite-difference centered stencil

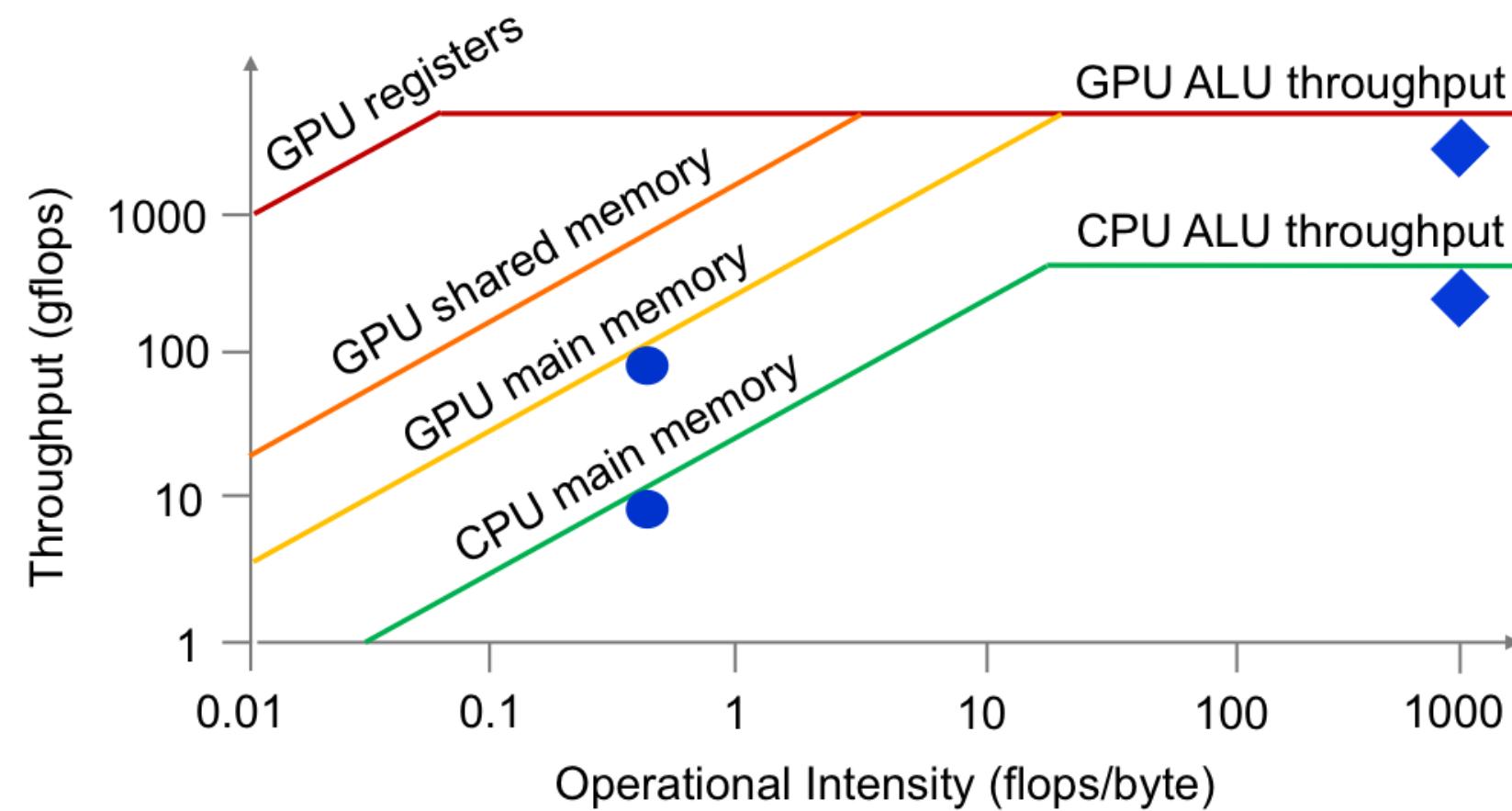


NVIDIA
K20x GPU

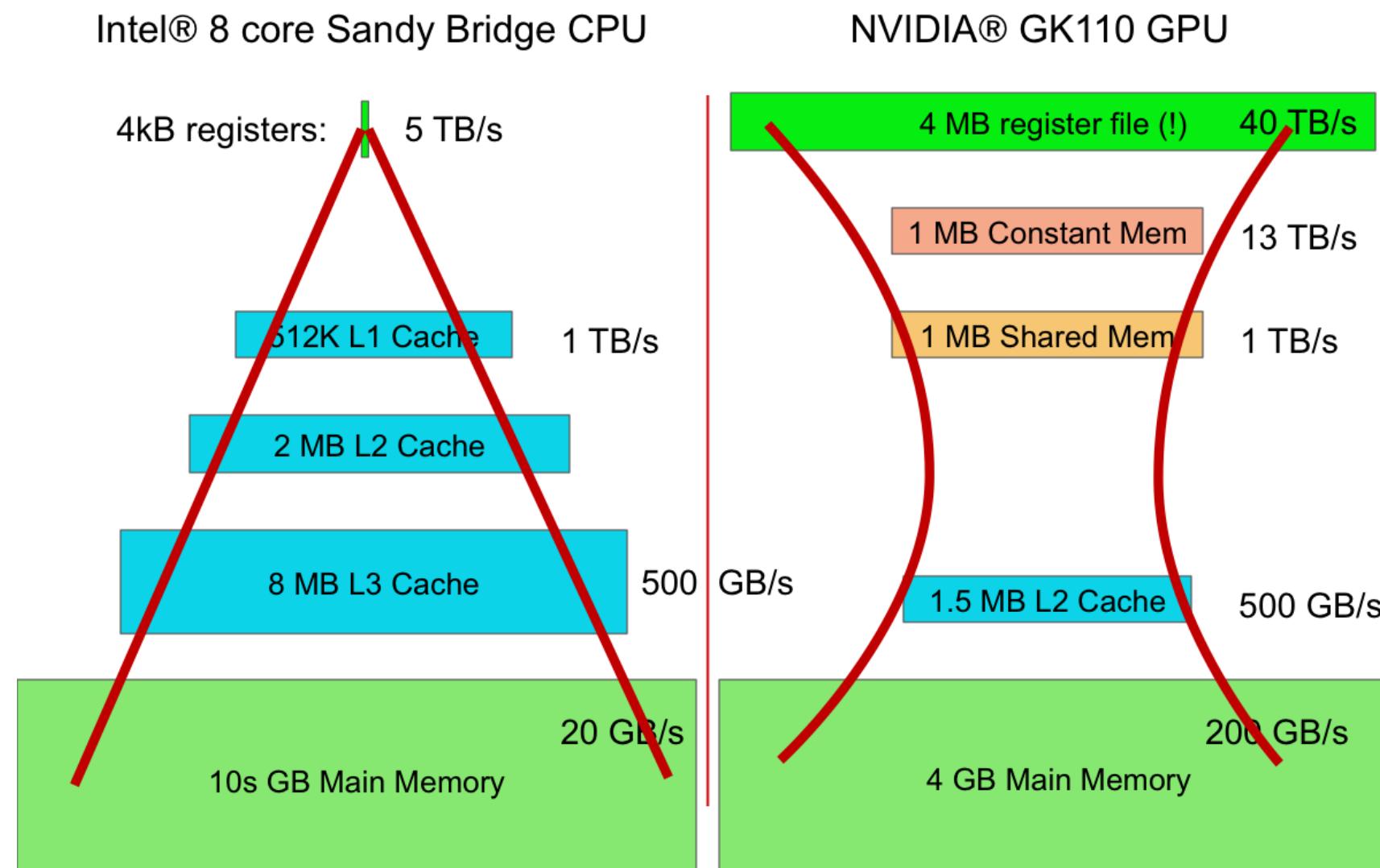


Roofline Design – Matrix kernels

- Dense matrix multiply 
- Sparse matrix multiply 



Where is my Memory?



Where are we in the roofline plot?

Order 2 stencil

How many flops?

How many words?

```
case 2:  
    return curr[0] + xcfl * (curr[-1] + curr[1] - 2.f * curr[0]) +  
        ycfl * (curr[width] + curr[-width] - 2.f * curr[0]);
```

How many flops?

10 additions / multiplications

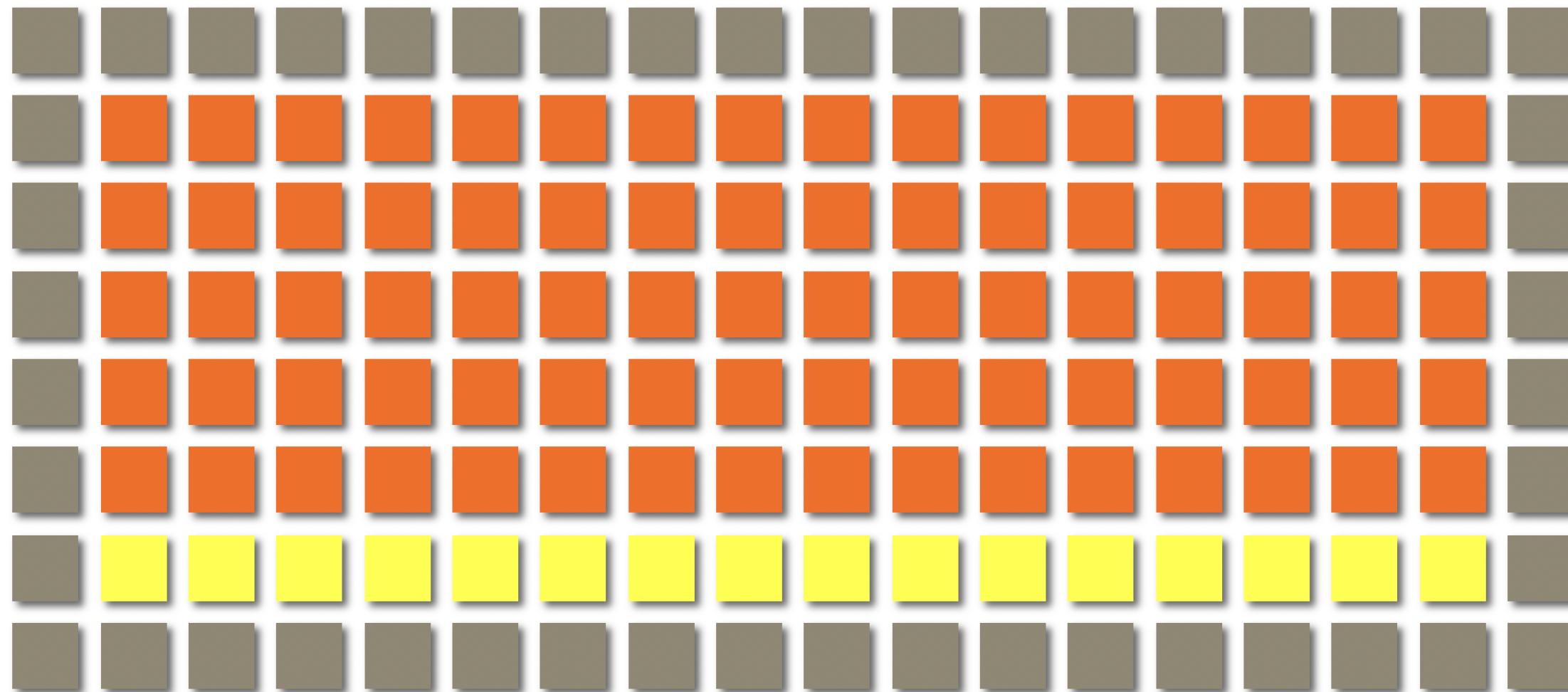
How many words?

- Read: 5
- Write: 1
- Total: 6

Two main ideas

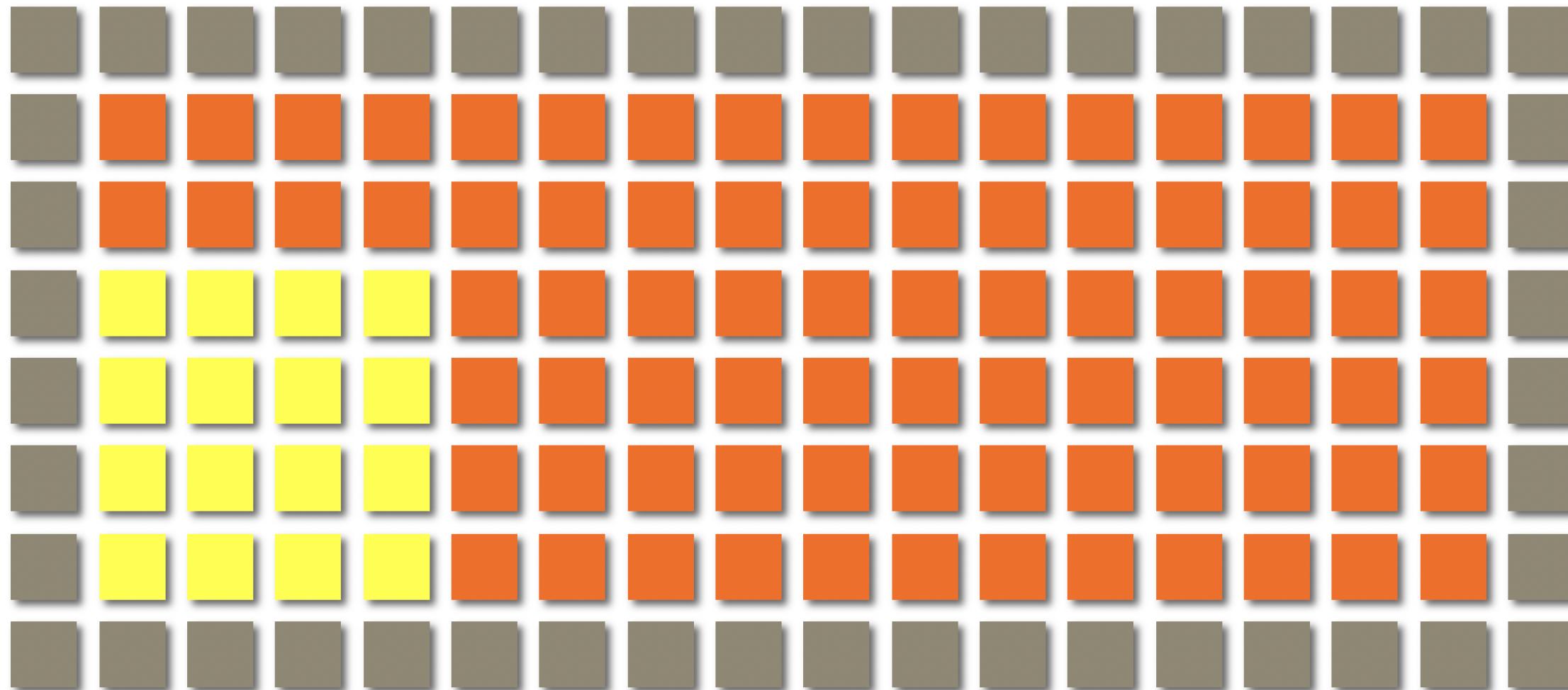
1. Use cache or shared memory
2. Memory accesses should be coalesced

Idea 1



- Thread-block = 16 threads.
- Reads: $16*3+2$. Writes: 16. Total = 66
- Flops: $10*16 = 160$
- Ratio: flops/word = 2.4

Idea 2



Improvement

Ratio: flops/word

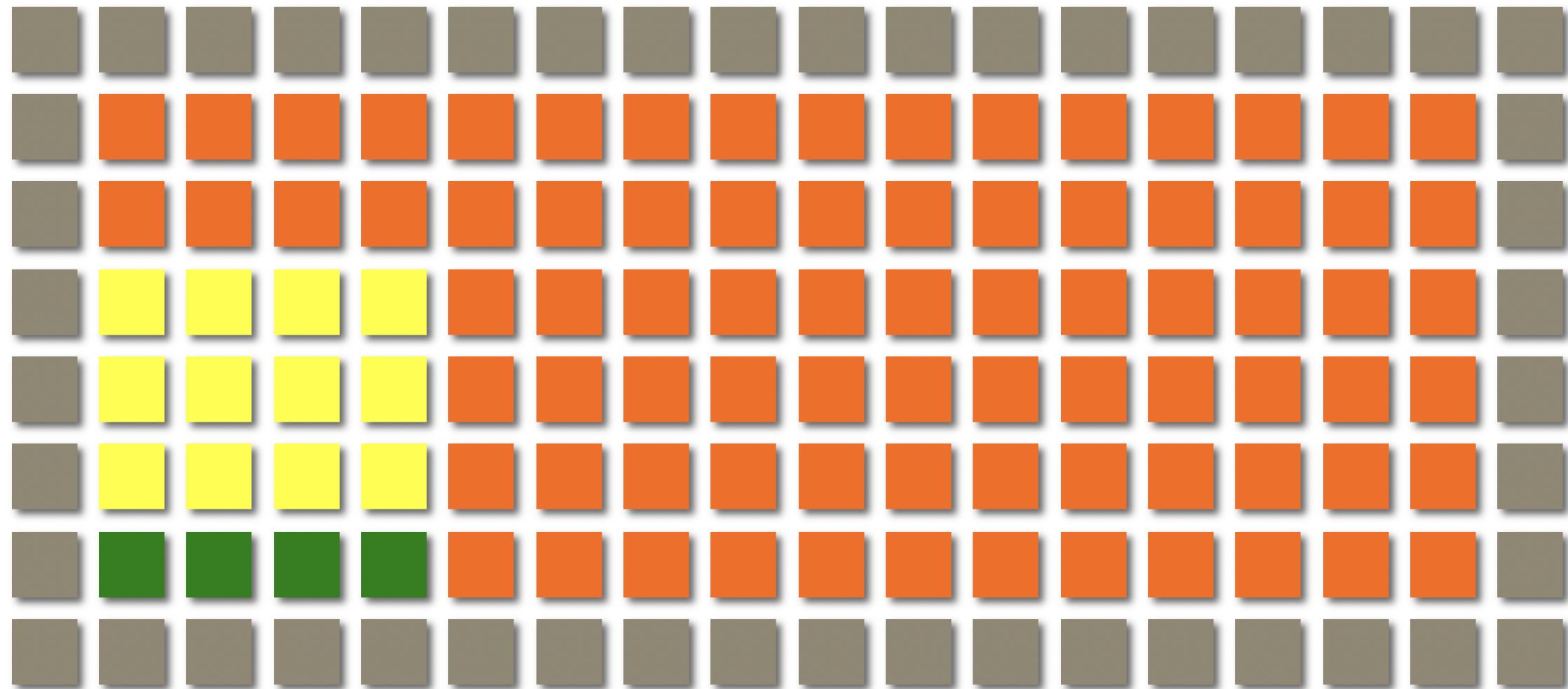
$2.4 \rightarrow 3.3$

Peak

For an $n \times n$ block:

- Memory traffic: $2n^2 + 4n$
- Flops: $10n^2$

Maximum intensity: 5 flops/words

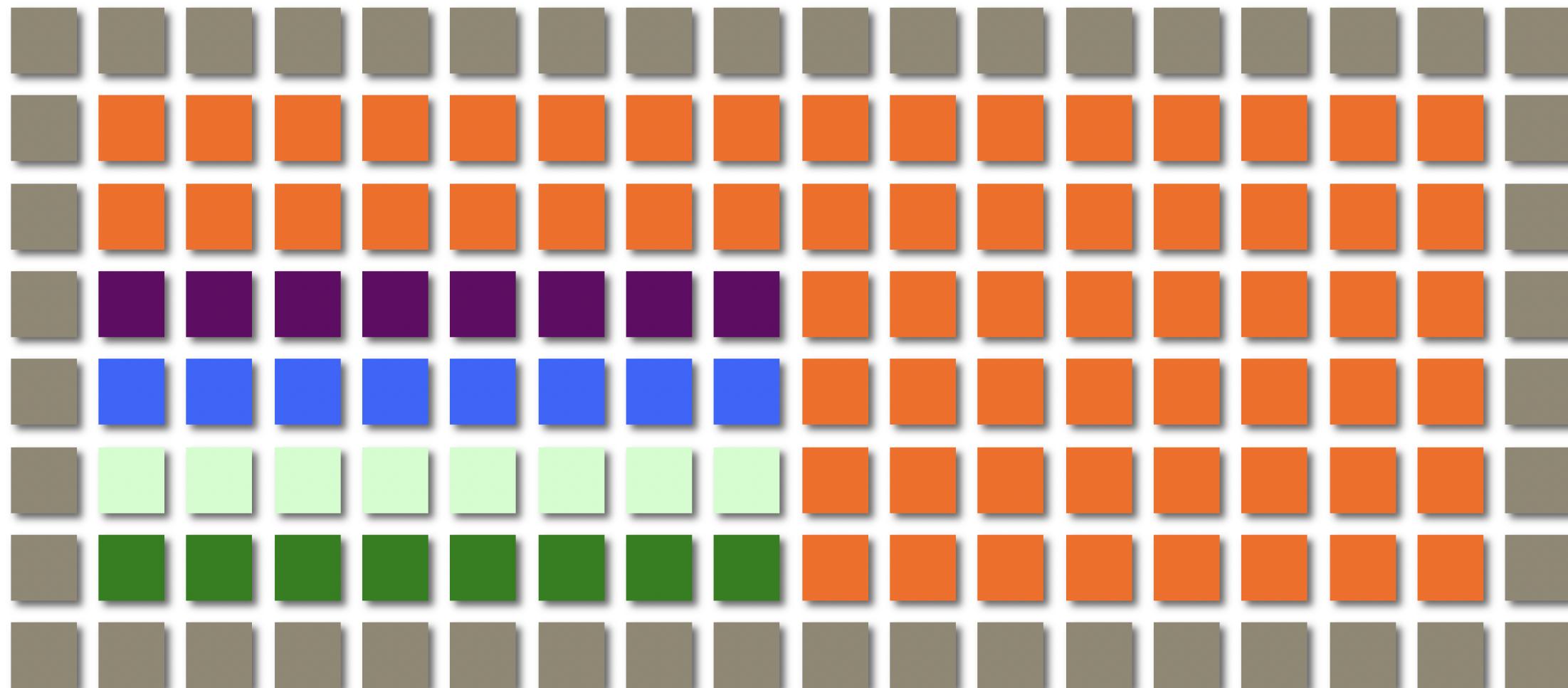


Only mesh nodes along x are contiguous.

This size must be a multiple of 32.

A warp must work on a chunk aligned along x .

Warp memory access



Thread block should be rectangular

Along x: dimension 64

Along y: determines number of threads in block

Example: 512 → dimension y = 8

Check your bounds



Shared memory algorithm

Step 1: all threads load data in shared memory



Load in shared memory

Step 2: threads inside the domain apply the stencil and write to the output array



Update inside nodes

Sample output

```
$ ./main -gsb
Order: 8, 4096x4096, 100 iterations
          time (ms)      GBytes/sec
    CPU        2681.71        45.0444
  Global       48.7767        2476.51
    Block       37.4602        3224.65
  Shared       33.9051        3562.77

          L2Ref          L1Inf          L2Err
  Global       0.447065  3.57628e-07  4.38207e-08
  Block       0.447065  3.57628e-07  4.38207e-08
  Shared       0.447065  3.57628e-07  4.38207e-08
```

If you compile with the **-G** option:

```
$ ./main -gsb
Order: 8, 4096x4096, 100 iterations
      time (ms)      GBytes/sec
    CPU        2683.12        45.0206
  Global       208.855       578.373
    Block      111.014      1088.12
  Shared      136.626      884.138

      L2Ref      L1Inf      L2Err
  Global      0.447065        0        0
    Block      0.447065        0        0
  Shared      0.447065        0        0
```

Order 2

```
$ ./main -gsb
Order: 2, 4096x4096, 100 iterations
          time (ms)      GBytes/sec
      CPU          908.71        44.3104
  Global         37.6773       1068.69
    Block         35.3023       1140.59
  Shared         36.0975       1115.46

          L2Ref          LInf          L2Err
  Global         0.418194  2.98023e-07  4.87715e-08
    Block         0.418194  2.98023e-07  4.87715e-08
  Shared         0.418194  2.98023e-07  4.87715e-08
```

What is the best blocking strategy?

Go as wide as you can along x while satisfying memory and concurrency constraints

Then loop along y reusing loaded data

