

Assignment 1

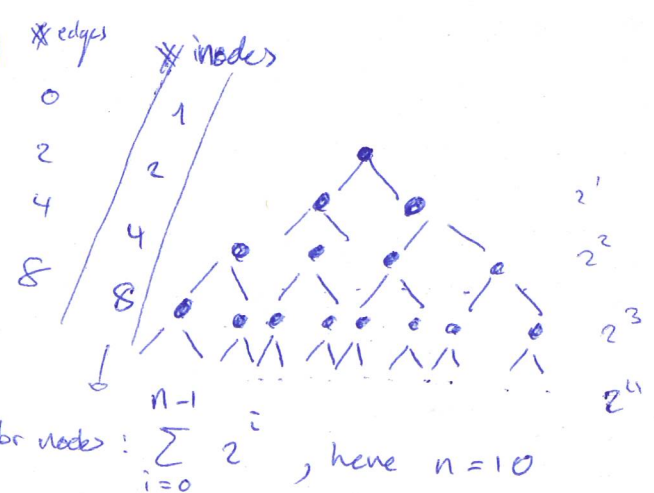
Hierarchical Softmax
&
Huffman Coding \mathcal{D} , document V , vocabulary $p(w) = \frac{\text{\# of times } w \text{ appears in } \mathcal{D}}{|\mathcal{D}|}$, word frequency $|\mathcal{D}|$, \# of words $|V|$, voc size

$$P_w(c) = \frac{\exp\{s(w, c)\}}{\sum_{c \in V} \exp\{s(w, c)\}} \quad , \quad O(|V|) \text{ for each pair } (w, c)$$

approximated by depth of node c in tree $\frac{1}{2}$

$$P_w(c) = \prod_{j=1}^{L(c)-1} P_w((c_j) \rightarrow (c_{j+1}))$$

node of depth j on road to leaf c
(prob of transition) \# edges

3.1 binary tree, $|V| = 2^{10} = 1024$ 

To compute $P_w(c)$, we need to go all the way down to the leaf, the way is composed of 9 edges here.

General case:

$$O(\log |V| - 1) \\ = O(\log |V|)$$

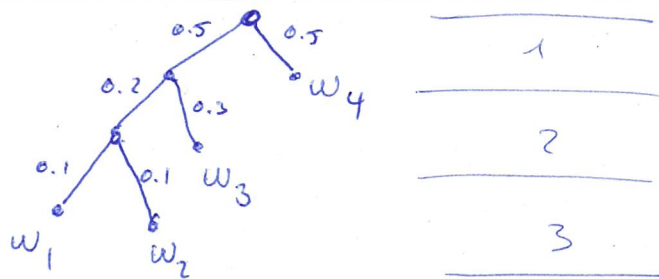
$$\Rightarrow \boxed{\text{\# inner nodes} = 1023}$$

$$\text{Total \# nodes} = \text{inner nodes} + \text{leaf nodes} \\ = 1023 + 1024 = 2047$$

$$\text{\# of edge} = 2046$$

$$3.2 \mid |V| = 4$$

Intuitively, we have:



$$E[L(c)] = 3 \cdot 0.1 + 3 \cdot 0.1 + 2 \cdot 0.3 + 1 \cdot 0.5 = 1.7$$

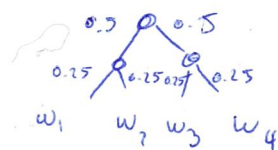
$$\cancel{E[L(c)]} \geq \cancel{E[-\log(p(w_i))]}$$

$$L(w_i) \geq -\log(p(w_i))$$

is true

• worst case: let take the last example.

The worst case is this one:



$$E[L(c)] = 4 \cdot (0.25 \cdot 2) = 2 = \log_2 4$$

general case: $E[L(c)] = O(\log |V|)$