# University of California, Berkeley College of Engineering Computer Science Division – EECS

Fall 2010

Prof. Michael J. Franklin

#### MIDTERM I

CS 186 Introduction to Database Systems

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This is a **closed book** examination – but you are allowed one 8.5" x 11" sheet of notes (double sided). You should answer as many questions as possible. Partial credit will be given where appropriate. There are 100 points in all. You should read **all** of the questions before starting the exam, as some of the questions are substantially more time-consuming than others.

Write all of your answers directly on this paper. **Be sure to clearly indicate your final answer** for each question. Also, be sure to state any assumptions that you are making in your answers.

#### GOOD LUCK!!!

Problem	Possible	Score
1. Buffer Manager	30	
2. B-Trees	24	
3. Hash Indexes	24	
4. Formal Query Languages	22	
TOTAL	100	

### **Question 1 – Buffer Management** [4 parts, 30 points total]

a) [15 points] Consider a database system with 4 buffer frames (A, B, C, D) and a file of 6 disk pages (1, 2, 3, 4, 5, 6). Assume that you start with an empty buffer pool. A sequence of requests is made to the buffer manager as described in the Request column (below). At certain times a Pin request is immediately followed by an Unpin request (represented as Pin/Unpin), but other times Pin and Unpin requests happen in an interlaced manner with other requests.

Fill in the following table showing the buffer contents after the completion of each operation using the **CLOCK** page replacement policy as described in the book and HW 1. For each page, indicate the pin count (PC) and for unpinned pages, indicate the value of the reference bit (true or false). You should mark unchanged buffer frames with "ditto" ("). As in HW 1 we make the following assumptions in the clock policy:

The pointer doesn't move when filling a free frame in buffer or on a buffer hit.

The pointer advances after replacing a buffer frame

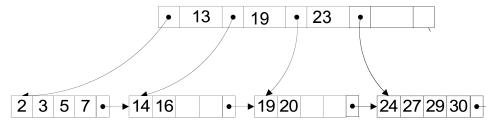
Note: Total points for this problem is 15, (3 each for correct buffer contents at steps T4-T8)

	D	Buffer Frames								
Time	Request	A	В	С	D					
T1	Pin 5	<b>5</b> PC=1	Empty	Empty	Empty					
T2	Pin/Unpin 2	"	2 PC=0 Ref = t	"	"					
Т3	Pin/Unpin 3	"	"	3 PC = 0 Ref =t	"					
T4	Pin 4									
T5	Pin 1									
Т6	Unpin 5									
Т7	Pin/Unpin 2									
Т8	Pin 6									

	SID:								
Question 1 – Buffer Management (continued)  b) [5 points] Consider an access pattern where there are lookups of many different keys in a B+tree. Assume that the B+Tree is 4 levels deep and the buffer pool has only 3 buffer frames. Which replacement policy will provide the best hit rate and why? Circle your answer and be sure to state why. List any assumptions you are making.									
a) LRU	<b>b)</b> MRU	c) Random	d) All are equal						
Why?									
c) [5 points] Consider an access pattern where there are lookups of many different keys in a static hash file (with no directory). Assume that hash file is much bigger than the buffer pool and that the hash function works well on the given data. Which replacement policy will provide the best hit rate and why? Circle your answer and be sure to state why. List any assumptions you are making.									
a) LRU	<b>b)</b> MRU	c) Random	d) All are equal						
d) [5 points] Consider a situation where the access pattern changes suddenly. For example, one set of pages is heavily accessed during the week, but starting at 11:59pm on Friday, a completely different set of pages becomes heavily accessed. Which of the following would you expect to react more quickly to the change in workload (i.e., have a better hit rate once the workload changes). Circle your answer and be sure to state why. List any assumptions you are									
making. a) LRU	P) I EII	c) Not enough	info to tall						
Why?	b) LIO	c) Not chough							

### **Question 2 – B+Trees** [6 parts, 24 points total]

Consider the following B+Tree with order d=2. Recall that the insert algorithm splits full leaf nodes it is inserting to rather than trying to avoid splits by redistributing data entries to neighbors.

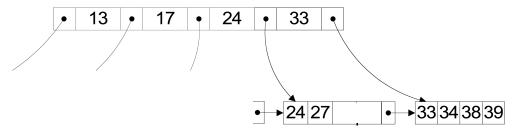


- a) [4 points] Which of the following key values for data entries, <u>if inserted on its own</u>, would cause a LEAF PAGE SPLIT (Circle all that apply):
  - 12 13 18 23
- b) [2 points] For the original tree above, what is the minimum number of keys that you could insert that would cause the ROOT NODE to split? Give an example of a sequence of such inserts.
- c) [2 points] Considering the original B+tree above, which of the following is true? (Circle one only, and briefly explain your answer)
  - A. No deletions have ever been performed on it.
  - B. One or more deletions have been performed on it.
  - C. Not enough information to tell.

Why?

**d)** [4 points] STARTING WITH THE ORIGINAL TREE ABOVE, draw the B+Tree that results from deleting the data entry with the key 14. (you only have to draw the parts of the tree that change).

e) [6 points] Consider the following B+Tree with order d=2 (with some leaf nodes elided)



Draw the B+Tree that results from inserting a data entry with the key 40. (hint: recall that the root node is not required to have d keys in it).

f) [6 points (3 points each)] In the B+Tree algorithm covered in class, we treat the split of a leaf node differently than the split of an index node. Specifically, when splitting a leaf node we make a copy of the new separator key and insert that into the parent node, while when splitting a non-leaf node, we promote a key (i.e., move it) into the parent node. Your project partner argues that it would be simpler to treat all nodes the same – either by copying for all nodes, or promoting for all nodes. For each of these two options, would it work without changing the basic B+Tree search algorithm? If not, why not? If so, is there any disadvantage to doing it compared to the way we did it in class?

Option 1 - Copy for all Nodes: Does it work? YES or NO

If No, say why. If Yes list one disadvantage:

Option 2 – Promote for all Nodes: Does it work? YES or NO

If No, say why. Yes list one disadvantage:

## **Question 3 – Hashing [5 parts, 24 points]:**

a) [10 points] Extendible Hashing Consider the following 4 update operations.

/ L L J		2
Now, operation no.	operation	key value (binary)
with a hash	_insert	25 (011001)
$^{2} 2^{depth}$	insert 1	46 (101110)
shown belgw.	insert 5 7	18 (010010)
Draw the 4 structure and	insert	22 (010110)

consider an extendible hash structur hold up to 4 emtdies,

function  $h(n) = n \mod$ 

an initial state a

extendible hash its contents after the have occurred in the

4 operations have occurred in the order shown. We recommend that you do your scratch work on this page at first. But, this page will not be graded. You MUST put your final answer on the following page!!

	1	
0		
1		

1			
8	10	14	

|--|

Final answer for Question 3(a) - Extendible Hashing:

- ♦ Only this page will be graded for question 3(a).
- ♦ The final structure should have a directory of size 8 so use the template below.
- ♦ Show all buckets and pointers
- Label the directory entries with their corresponding hash value (as on the previous page).
- Make sure to include local depths for all buckets and the global depth of the directory.



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Question 3 – Hashin Now, consider a <i>line</i> Consider the shown in the parts b and c.		x with the following	ng initial state:	sea (bucke (bucke (bucke	t 1)	nserts answer		
	operation no.	operation	key value (bin	nary)				
	1	insert	20 (010100)					
	2	insert	46 (101110)					
	3	insert	18 (010010)					
	4	insert	23 (010111)					
b) [4 points] Assuming the operations are performed in the order specified in the table, which of the four operations will cause the bucket pointed to by the "next" pointer (i.e., "bucket 1") to be split?								
c) [3 points] What will be the value of "Next" after the split (from part b) happens (i.e., which bucket will be scheduled to be split on the subsequent overflow)?								
<b>d)</b> [3 points] For the above structure, how many buckets will there be (ignoring any extra pages used for overflow) at the beginning of round 4 (i.e., when "Level" = 4)?								
e) [4 points] In class we didn't specify if the data entries inside buckets (including overflow chains) should be kept sorted by key value or simply placed in the order they arrive. List one advantage and one disadvantage of keeping the bucket contents sorted. State which operations are affected.								

Advantage (Pro):

Disadvantage (Con):

# **Question 4 – Formal Query Languages [3 parts, 22 points total]**

Consider the following schema that records information about a university (Primary keys are underlined):

Student (sid, sname, address, age)

Course (cid, cname, credits)

Professor(<u>pid</u>, pname, address, age, department)

Grade(sid, cid, pid, grade)

Write queries to answer the following questions. Note that we don't care about efficiency here, but points will be taken off for using unnecessary relations in your queries.

- **a)** [8 points (4 each)] Find the sid and sname of students who have received at least one A in a *four-credit* course.
  - i) Relational Algebra

ii) Relational Calculus

Recall the sch	ema from the previous page (Primary keys are underlined):  Student ( <u>sid</u> , sname, address, age)  Course ( <u>cid</u> , cname, credits)  Professor( <u>pid</u> , pname, address, age, department)  Grade( <u>sid</u> , <u>cid</u> , <u>pid</u> , grade)
	<b>leach)</b> ] Find the sids of students who have taken at least one course and have de of "A" in every course they have taken.
i)	Relational Algebra
ii)	Relational Calculus
	Find the pids of professors who have taught <b>exactly two</b> different classes. Use your tional Algebra or Relational Calculus.

SID:\_\_\_\_\_