

Computer Security means protecting Computer's data from harm, theft or unauthorized access and also preserve Integrity, availability and confidentiality of info system.

Core of Computer Security / Concepts.

- Confidentiality, Integrity, availability.

Confidentiality: It means keeping data secret and not allowing unauthorized access.

Cover 2 concept

Data Confidentiality: only authorized people on system can see or Read private on Confidential Info.

Privacy: Locking personal information. Individuals only

access which is related to him. Only he can read off & clear and stored

Example: GoMail Acc is protected by a password.

Only you can read your email

Techniques used: Encryption, Biometric Login,

Access Control Lists.

Integrity: Data should be correct, complete and not changed by anyone without permission.

Data Integrity: Data is not modified during transfer & rotangles small change can occur big travel.

System Integrity: free from inadvertent unauthorized manipulation of the system.

Eg: If someone file over a Network but someone changes it during transfer its integrity lost.

Techniques used: Hash function & checksum, Digital Signatures.

Availability: Data and system should always be available when needed.

• User should access anytime, anywhere

• This protects against crash or attack that blocks access.

Techniques used: Backups, Redundant system. Protected from (or), attack

Ex (if) hospital server down, doc can't access patient records.

Two other Security Concepts are

Authenticity: Make sure that user is who they say they are.

to read and trusted. [Losing with fingerprint]

Accountability: tracking who did what and when

System Logs show the user actions.

Challengers of CS:

Cyber attack: hackers try to break into system.

(Phising, Dos attack) They need only

one weapon for attack & have

Weak password: It can be guessed one hand.

Insider threat: Attention from people inside the organization.

Cloud computing: Storing data online needs strong protection.

Bouncing security: To much security can make system hard to use.

Changing tech: New threats comes with new tech.

Designer challenge: designer must find and fix.

all weapons for security

Security as an afterthought: Security must be planned.

from design & development done not after makes.

OSI Security Architecture

Security attack, Security Service, Security Mechanisms.

Passive Security Attack: To secretly observe or listen data being transmitted.

Types

• Release of Msg contents.

- reading email, call log files without permission.

Traffic analysis: Even if msg is encrypted, attackers can watch patterns (who talk, to whom, when, abnormal flow).

Active attack: To change, fake, block data being sent.

Masquerade attack: hacks login or someone else.

Replay attack: Captures valid data from transmission and retransmits them to deceive the system.

• Intercept bank transaction and resending it to someone else.

Duplicate transmission.

Modification of Msg: Changing the recipients acce
num in fund transfer Rec.

Denial of Service (DoS): Makes System ~~also unavailable~~
→ flooding a server with fake traffic

Security Services (what we want)

→ It is a communication service that provided by a system to give a specific kind of protection to the system resources. It ensure secure communication across network

Authentication:

Verify the identity of users on system

→ Login with pass & id. and also OTP

Access Control: It controls who can have access to a resource. Grants and denies permission to one system resource

Ex → student can't access teacher data.

Data Confidentiality: protection data from unauthorized users. (use Encryption (AES, RSA))

Ex → what happens we end to end encryption to protect chats & calls.

Data Integrity: Data should be correct or not changed data can be tracked one exactly as send by entity. (use checksum, hashing).

Ex → send file content not be changed.

Non Repudiation: Someone can't deny they performed

Ex → send email with digital signature you can't say I didn't sign it.

Intrusion Detection System (IDS): detect suspicious activities

activities

Ex → Alert admin if unusual login

Security Mechanisms

Specific Security Mechanism:

Encipherment (Encryption)

- Convert data into secret code [Up to end encrypted
End to End encd)

Digital Signature: Confirm identity of the sender
ensure message not changed.

→ a signed Email Confirm identity

Access Control: Access only authorized user

→ Stu Can't access teacher's folder

Data Integrity: Ensure data should be consistent and
not changed

→ download file hash ck me into ck if it changed

File Hashing: (SHA) integrity verification

Hash function
Input: File
Output: Hash value

Positive
Patractive Security Mechanism: (protect Everyone)
help others

Security Recovery: Ensure system can recover from incident
↳ Backup data so that it can recover if
deleted by hacker.

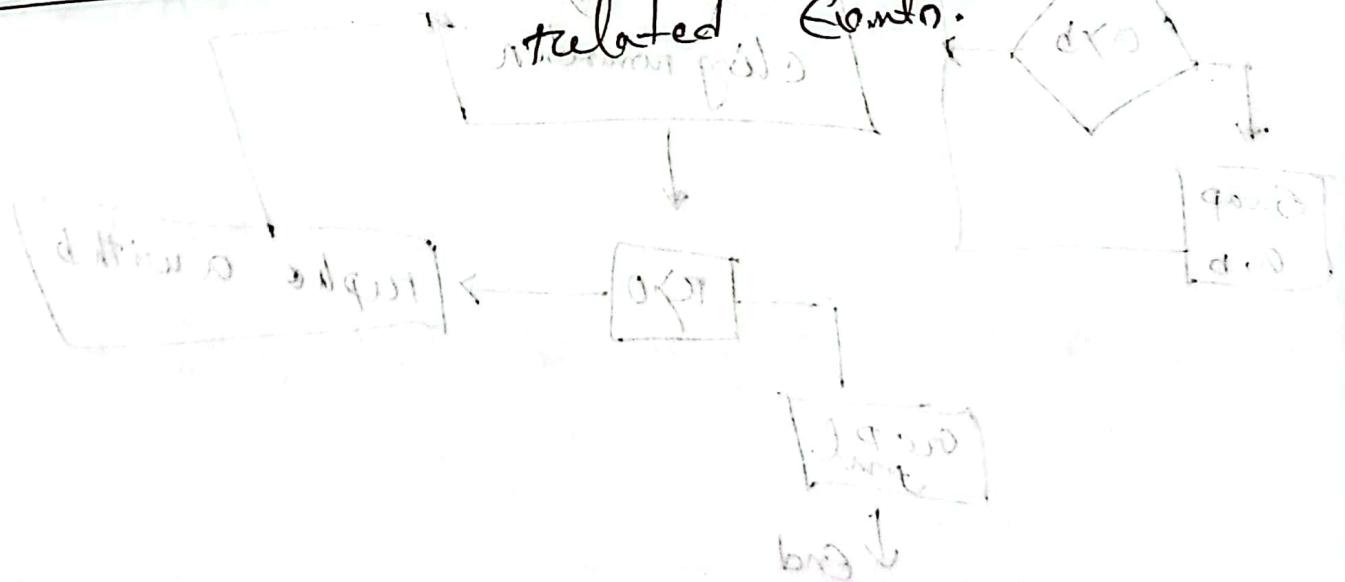
(recovery) warning: ensure system behaves as expected.

Trusted functionality: ensure system behaves as expected.
Security Labels: Page that defines security of all of

data.

Event detection: Identify security related events
(intrusion attempt)

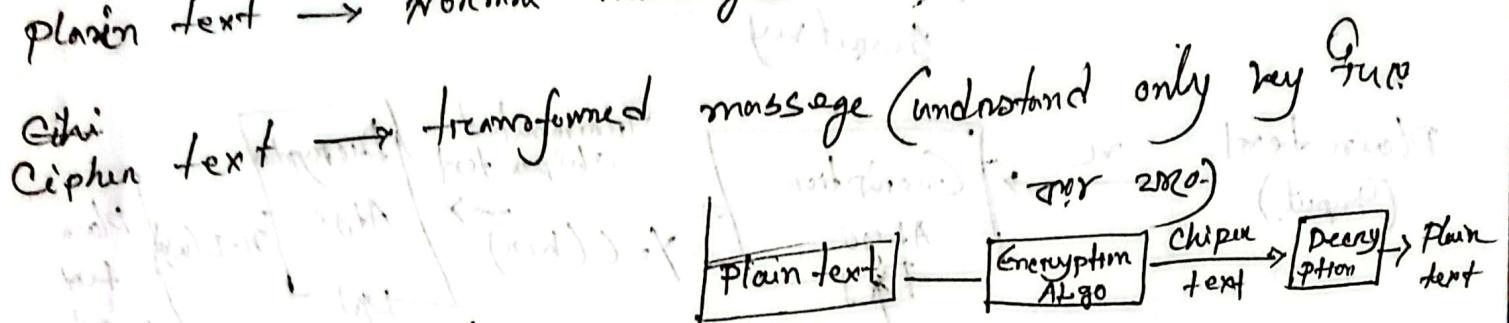
Security audit trails: keeps track of all security related events.



Segment 2

Cryptography is a technique of securing information through the use of codes & secure format. It ensures that only the intended person can read & understand the data.

plain text → Normal message



Feature of Cryptography

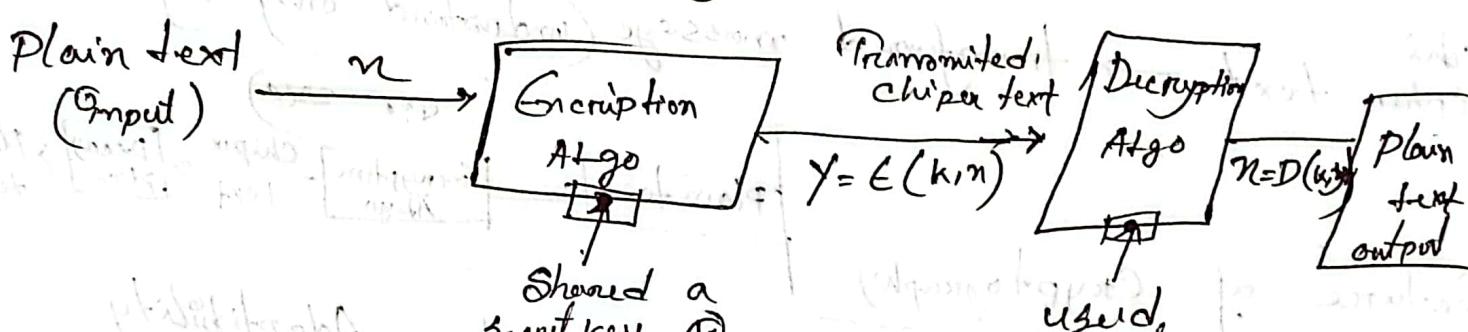
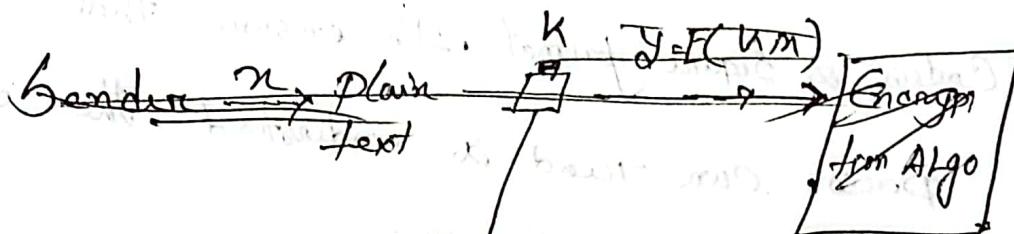
→ Confidentiality, Integrity, Authentication, Adaptability

Symmetric key cryptography / conventional

→ Single key used for encryption and decryption. Exact secret key used for encryption and decryption. Substitutions and transformation performed by algorithm depend on the key of length.

Plaintext: This is original intelligible msg on data that used as algorithm input

Encryption Algorithm: It performs various substitutions and transformations on plain text.

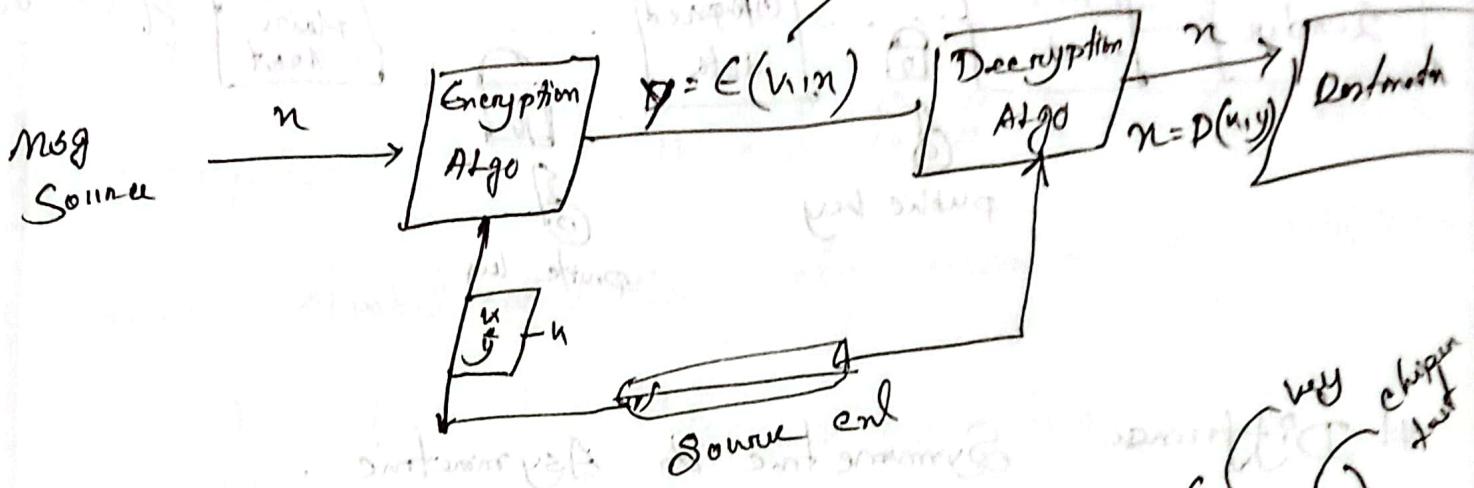


Model of Symmetric encryption

Chiper text: It depends on plain Text and the secret key. It is the coded message (unintelligible)

Decryption Algorithm: Mainly it is a reverse encryption algorithm. It takes Chiper & key and produce original Plain text

Model of Symmetric Cryptosystem



here $y = E(u, n)$

↑
Encryption
Algo

chiper text

plain text

key

$u = D(K, y)$

plain text

Decryption
Algo

Asymmetric key Cryptography : public key

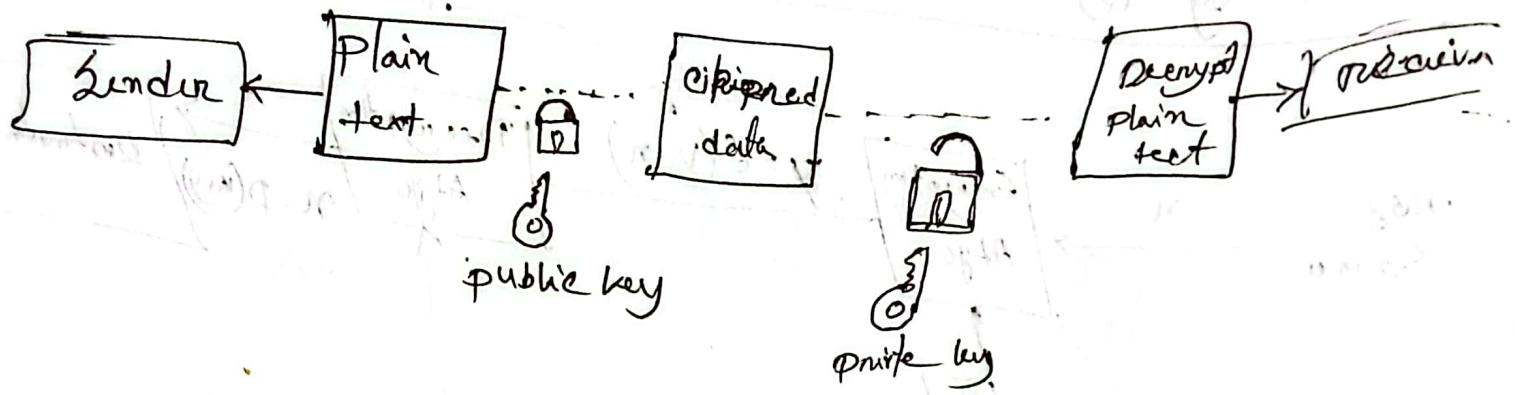
It is a type of encryption where 2 key used.

public - shared with anyone

private - kept secret the owner

WORKS . Sender encrypts the msg using receiver's public key

- only receiver private key can decrypt the msg -



Difference Symmetric & Asymmetric.

Symmetric

- One key same for encryption & decryption
- key must share secretly
- faster & efficient
- if leaked all exposed
- AES, DES, Blowfish
- WiFi, file encryption, Bluetooth
- (secret key used to encrypt and decrypt)

Asymmetric

- Two keys (public & private)
- no need share private key
- slower resource heavy
- more secure for communication
- RSA, ECC, DSA
- Secure email, Secure website
- Online banking

from firm of the bank (user)
transaction history (user)

* One time pad (Vernam cipher)

A B C D E F G H I J K L M
 0 1 2 3 4 5 6 7 8 9 10 11 12

m o p q r s t u v w x y z
 13 14 15 16 17 18 19 20 21 22 23 24 25

Plain Text: HELLO

key: XMEKRL.

1st Convert Alphabets into their value numbers. Start with (0-25)

H E L L O ← plain text

7 4 11 11 24

X M S C K L ← key

23 12 2 10 11

$$\text{cipher} = \underbrace{(\text{key} + \text{plain text})}_{\downarrow} \bmod 26.$$

$$30 \quad 16 \quad 13 \quad 21 \quad 25 \quad (\text{Add})$$

$$4 \quad 16 \quad 13 \quad 21 \quad 25 \quad (\bmod)$$

$$E \quad O \quad N \quad V \quad 2 \quad (\text{Encrypted})$$

Decryption

(key) \times M⁻¹ C^T
23 12 2 10 11

ES NVZ
16 13 21 25 (chipm)

chipm =

$$\text{plaintext} = (\text{Chipm} - \text{key} + 26) \bmod 26$$

4 16 13 21 25 → key ① 21

→ key - 23 - 12 - 2 - 10 - 11 → go down

→ 26 + 21 → 25

+ 26

7 4 11 11 14

↓ H E A L L O

Decrypted msg - HELLO

(background) S V N

Cryptanalysis

It is a process of trying to break a cipher (find original msg or key) without knowing the key.

Attackers use Structure of algorithm, clues about msg and try to find original message (plain text) and find key for decrypt the unreadable cipher text

Brute force attack

Attackers try every possible way to find the key's one by one when correct key found, the msg turns into something readable. almost half of the key tries for find right one

is hard (not much)

(2^8) is hard

is easy ($8 - \text{digit}$)

0 1 2 3 4 5 6 7 8 9
F A E D C B 9 8 7 6 5 4 3 2 1 0
G H I J K L M N O P Q R S T U V W X Y Z

Encryption Sequence

- Cost of break cipher exceeds the value of encrypt.

Info.

- The time required to break the cipher exceeds the lifetime of the info. (Convenience of generation box (key))

Substitution

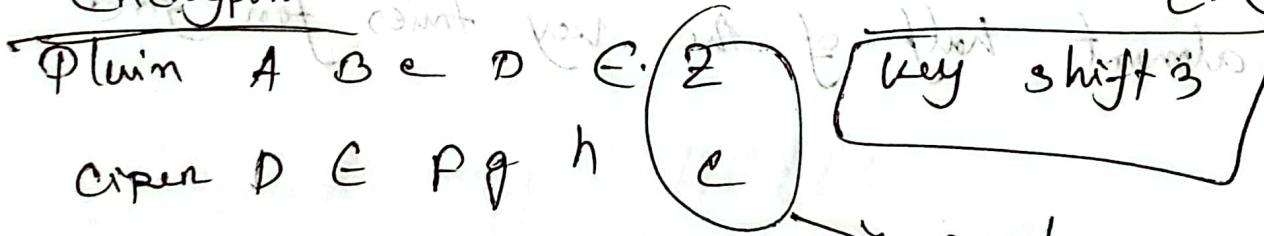
Substitution Cipher (Caesar)

here letter of plaintext are replaced by other

letters only numbered one symbol

Given Tripled 2⁰ 2nd time use 2⁰ or

Encryption



Decryption

$$(\text{Ciper} - 3) \bmod 26$$

D	E	F	G	H	C	$\xrightarrow{3-2}$	$\xrightarrow{=1 \bmod 26}$	$\left[\begin{matrix} 26+1 \\ =25 \end{matrix} \right]$
3	4	5	6	7	2	$\xrightarrow{=1 \bmod 26}$	$\xrightarrow{25 \bmod 26}$	$\left[\begin{matrix} 26+1 \\ =25 \end{matrix} \right]$
A	B	C	D	E	F	$\xrightarrow{25 \bmod 26}$	$\xrightarrow{=25 \bmod 26}$	$\left[\begin{matrix} 26+1 \\ =25 \end{matrix} \right]$

Transposition techniques

change the plain text position so called transposition.

2 way (key term) rail fence (simple)

Column Transposition.

Rowwise

Plain text = Nelson Academy for the best

Row / depth 2

R ₁ -	N	5	A	A	E	y	S	h	E	E	T
R ₂ -	e	o	c	d	m	;	t	s	r	p	u

Chiper text : N S A A e y s h b b h b S, e o c d m i t E C T

Column Transposition

Plain text: Fine Minutes Engineering

Key = 49512 [Cm - 26, 29, 27, 28] 52r key

Cm → ↓ ↓ ↓ ↓ ↓
 4 9 5 1 2
 F I V E M { very 26 under
 T N U T E Cm 29 27 28
 S E N G R T
 N E E R T
 N G R
 4 3 5 1 2
 my key is = 43512

as key Cm put on

Plain text = ~~E T G R Y U N E M E T T P I S N N~~
 4 3 5 1
~~P I S N N~~
 2

E T G R M E T T I N E E G F I S N N
V U N E.

①

Stream Cipher

Stream Cipher

- here Length define.

bit & bytes

- Design is so complex

- Principle is Confusion

• Speed is fasten

• Encryption done by

CFB, OFB

• Decryption done by

XOR

• Vernam cipher (Example)

Block Cipher

• here block size 64 or
128 bits

• design is so simple

• principle is Confusion & Diffusion

Diffusion

• Speed is Slower

• Encryption done by

ECB & CBC

• reverse of Encryption

• DES, AES (Example)

Bigram frequency

A bigram is a sequence of 2 consecutive words from given text.

bigram frequency refers to how many times each pair of words (bigram) appears in the text.

2 bigram pieces are

I Love Coding, I Love mango

→ I Love, — 2 times
Love Coding — 1 time
Love mango → 1 "

frequency 2 times some form
I Love 2 times some form
Love Coding 1 time some form
Love mango 1 time some form
2 times some form

Trigram frequency:

Sequence of 3 consecutive words from a given text.
Trigram frequency refers to how many times each

3 word sequence appears.

\rightarrow I Love Coding, I Love Coding on python

I, Love, Coding. — 2 times

Love Coding on — 1 time

Coding on python — 1 time.

First without program (manual) then

After writing code

Program was Registered and

(Run)

Then it gives output { } with all gathered in
the output

Working on small changes with regards to the

data sets and both major & minor changes

changes made by hand

now

Fig 3.

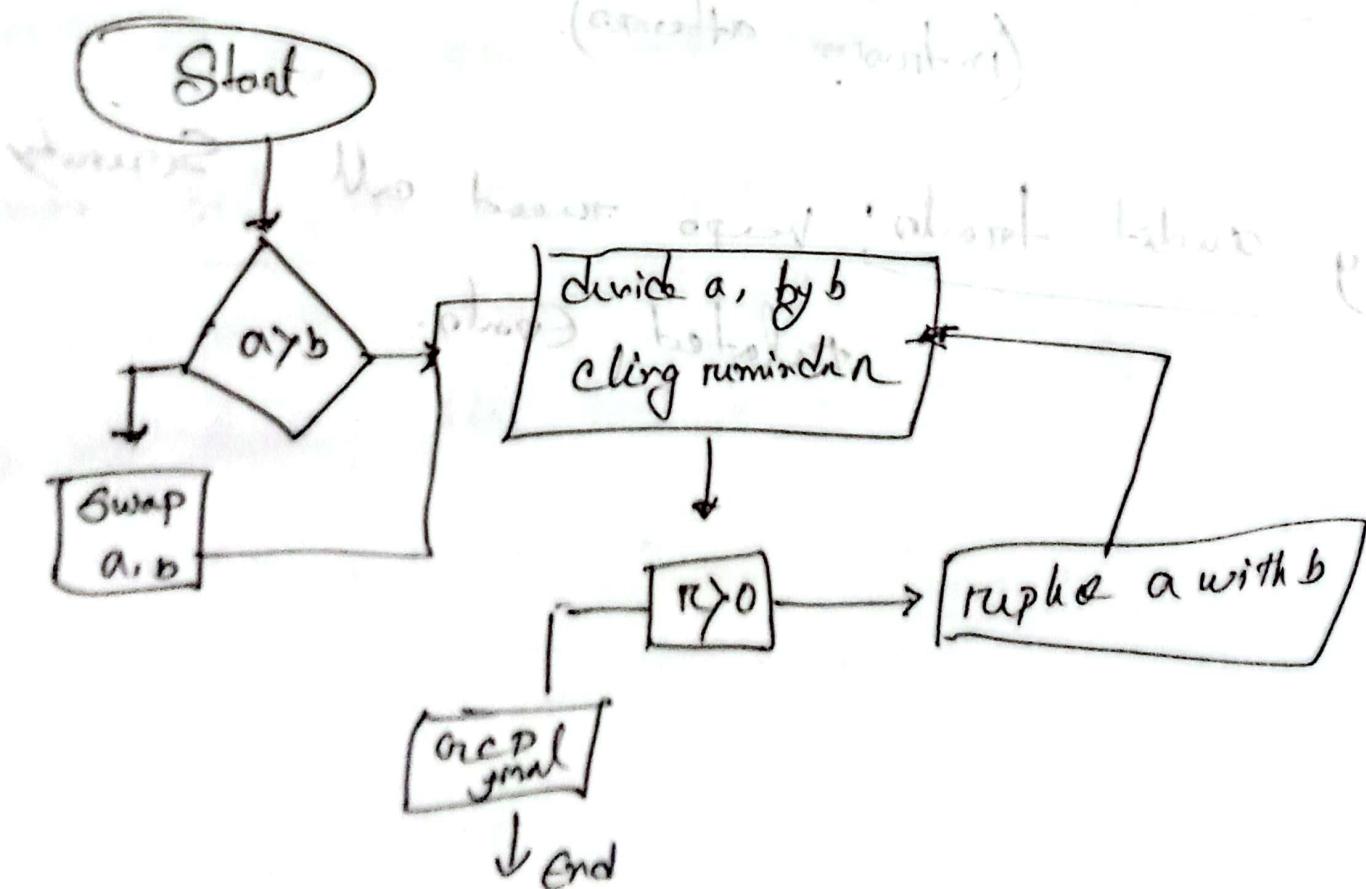
Euclidean Algorithm: Determining the GCD of 2 positive integers.

Two integers are relatively prime (coprime)

If there is no common divisor GCD is 1.

If GCD greater than 1 not coprime.

Euclidean Algorithm to find GCD.



* Jnd GCD of $(450, 120)$? Jnd 2 numbers are co prime.

$$450, 100$$

$$a \mid b$$

a/b

$$450 > 120 \quad (\text{yes})$$

Next divide a/b (calling remainder)

$$\# \quad 450 / 120 = 3 \text{ and remainder } 90$$

$$450 = 3 \times 120 + 90$$

$$\# \quad 120 / 90 = 1 \text{ remainder } 30$$

$$\rightarrow 120 = 1 \times 90 + 30$$

$$\# \quad 90 / 30 = 3 \text{ remainder } 0$$

$$\rightarrow 90 = 3 \times 30 + 0 \text{ hence GCD find}$$

GCD is 30

$$\text{Hence } \text{GCD} = (450, 120) = 30 \text{ not } 1 \quad [\text{Not co prime}]$$

AES

[Advanced Encryption Standard] (symmetric cipher)

AES use to encrypt data (like passwords, files or msg) so that only authorized people can read it.

Diagram AES working Step by Step

Plaintext (Input)

- Original data we want to protect.
Size 16 bytes = 128 bits.

Key (Encrypt the plaintext)

128 bits = 10 rounds, $\frac{192 \text{ bits}}{16 \text{ bits}} = 12 \text{ rounds}$

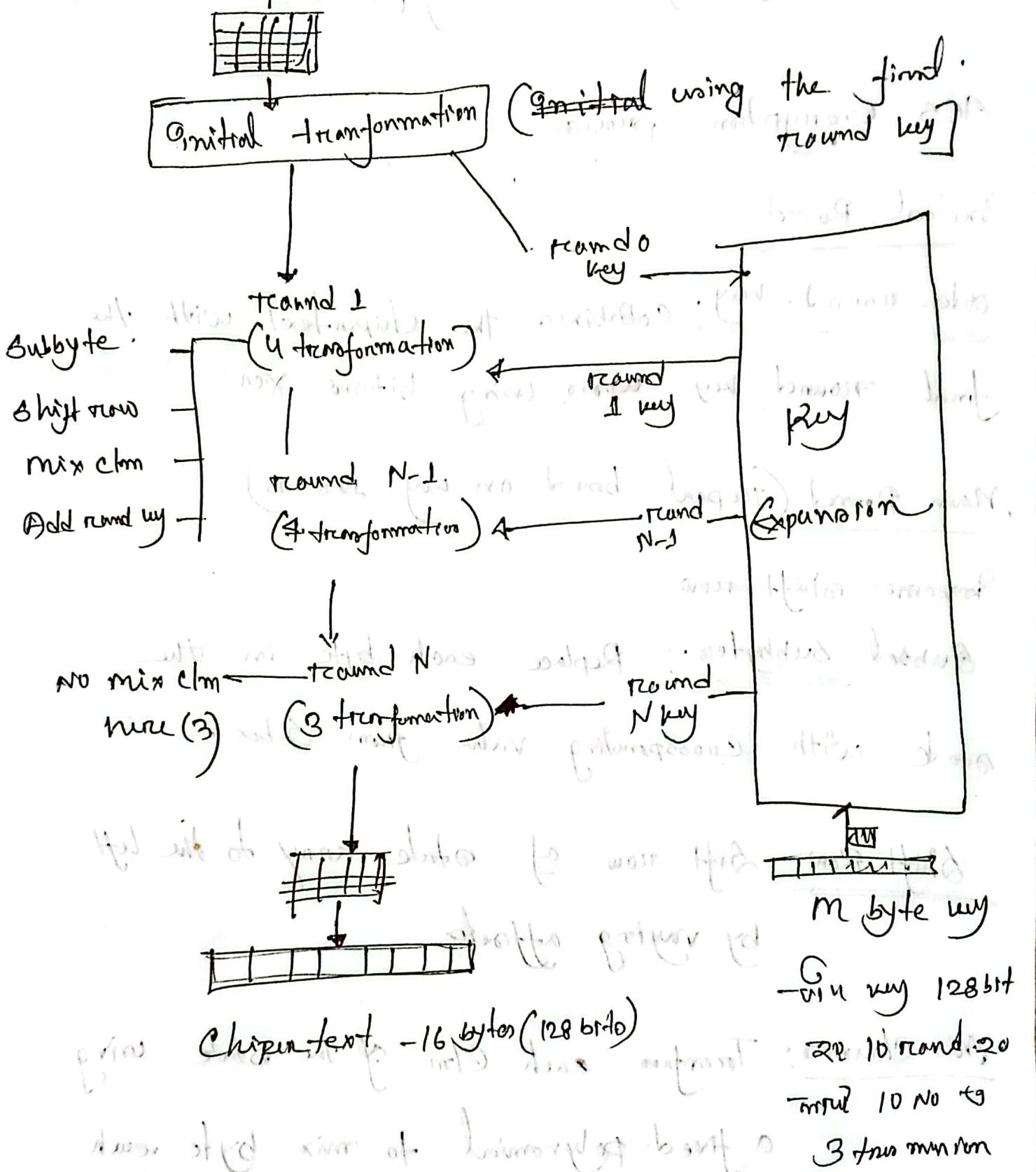
256 bits = 14 rounds

key Expansion:

- Key is expanded into multiple round keys
- One key is used for (round of 12)

transformation.

Plaintext - 16 bytes (128 bits)



AES Important

Security, Efficiency, Versatility (Sensitive Data Secure)

AES Encryption process

Initial Round

Add-round-key: Combines the ciphertext with the final round key using bitwise XOR

Main Round (Repeat based on key $128 = 10$)

Permute right row

Subst subbytes: Replace each byte in the state with corresponding value from S-box

Shift Row: Shift row of state array to the left by varying offsets

fixed (row N+1)

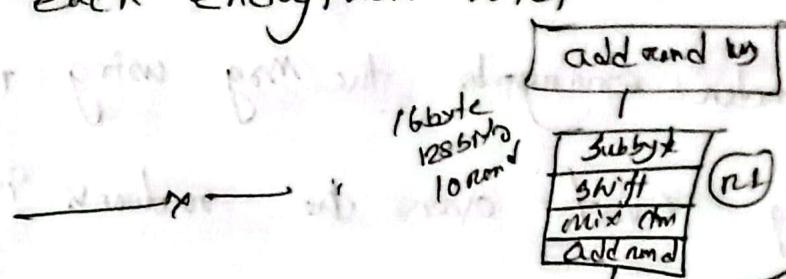
fixed (row N+2)

Mix Columns: Transform each column of the state using a fixed polynomial to mix bytes in each column.

add round key again

final Round: Repeat SubByte, ShiftRow, and AddRoundKey step but omit the Mix Col transformation.

AES Decryption: To retrieve original data, AES performs the inverse of each encryption step in reverse order using Paine key



Link Encryption

Link encryption ~~lectures~~ data by encrypting and decrypting info at every node or network switch.

It passes through ~~path~~ just at end point.

Or ensure all data header & routing info protect.

DisAdv: In every switch msg decryp is mandatory.

for know the address from header so msg can be vulnerable.
use unique key

End to End Encryption

Information has been modified, added and received only here ~~node~~ and host and receiver. Only communication with other hosts is guaranteed that only communicate by private key. It guarantees that only

the intended sender and receiver can access and understand data the data goes through many way.

- Sender encrypts the msg using receiver public key
- Msg sent over the network in encrypted form.
- Receiver decrypts the msg using their private key

Adv - - Router can't see data

- protect data from unauthorized individuals

using Both Link & End-to-end

E2E keeps user data safe from end to end

Link encryption secures the entire packet

during transmission across each network segment

segment

key distribution

It is the process of sharing secret encryption keys securely between parties who want to communicate.

Distributed $A \rightarrow B$. Many ways

- * A can select a key physically delivered A & B
- * third party select key and physically deliver A & B
- * If A & B previously & recently used a key, one party transmit the new key to other encrypting with one
- * If A & B has encrypted connection third party C can deliver a key on encrypted links to A, B

Key Distribution Center (KDC)

It med for key sharing easy and safe way

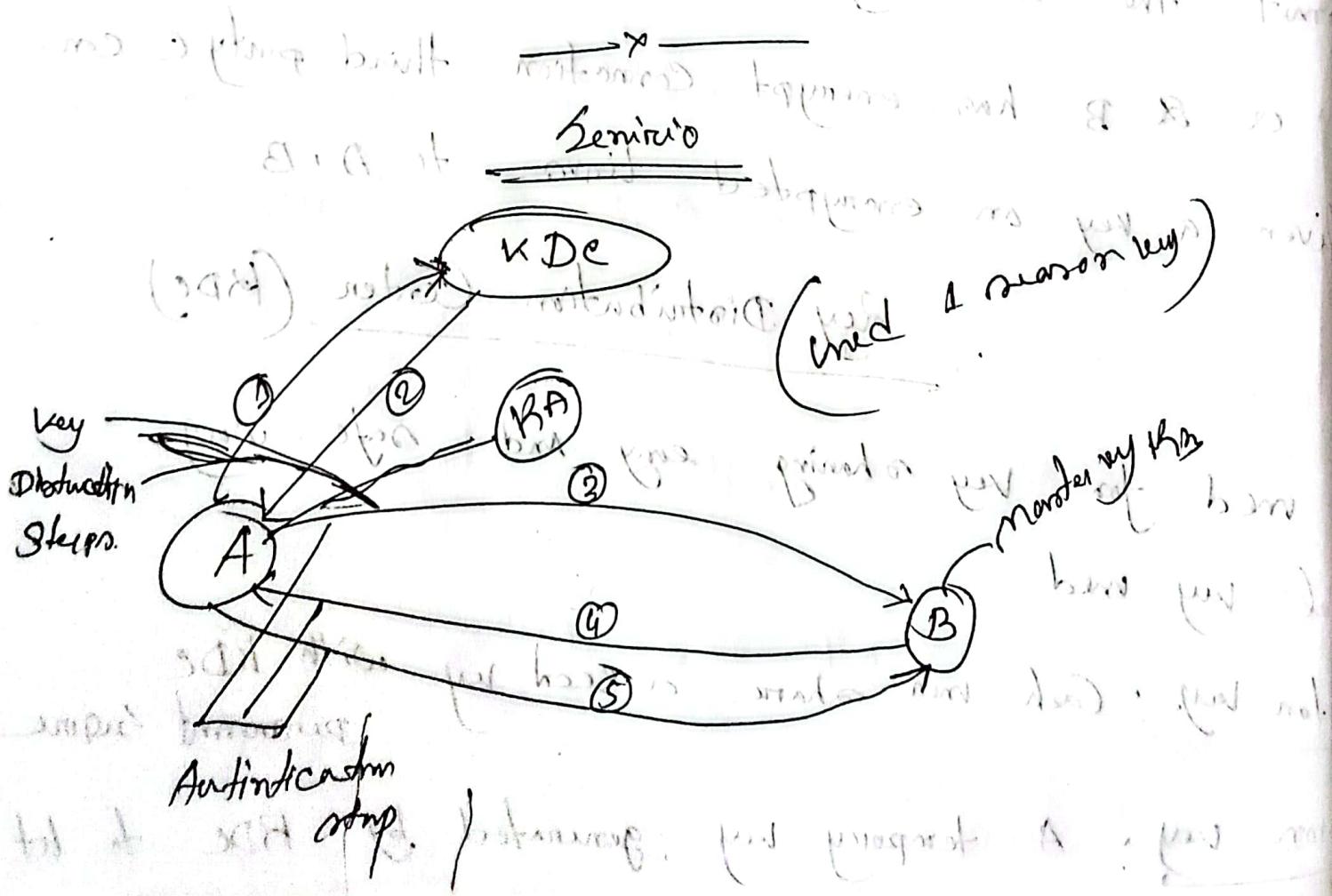
2 level key med

Master key: Each user share a secret key with KDC permanent secure

Session key: A temporary key generated by KDC to let

User communicate securely, need for one session and then discarded.

- Session key ~~distribution~~ transmitted in encrypted form using a master key that is shared by the key distribution center (KDC).
- N entities communicate as many ways $\frac{N(N-1)}{2}$ session keys used at once. N master required. (frequently session key change every 20s)



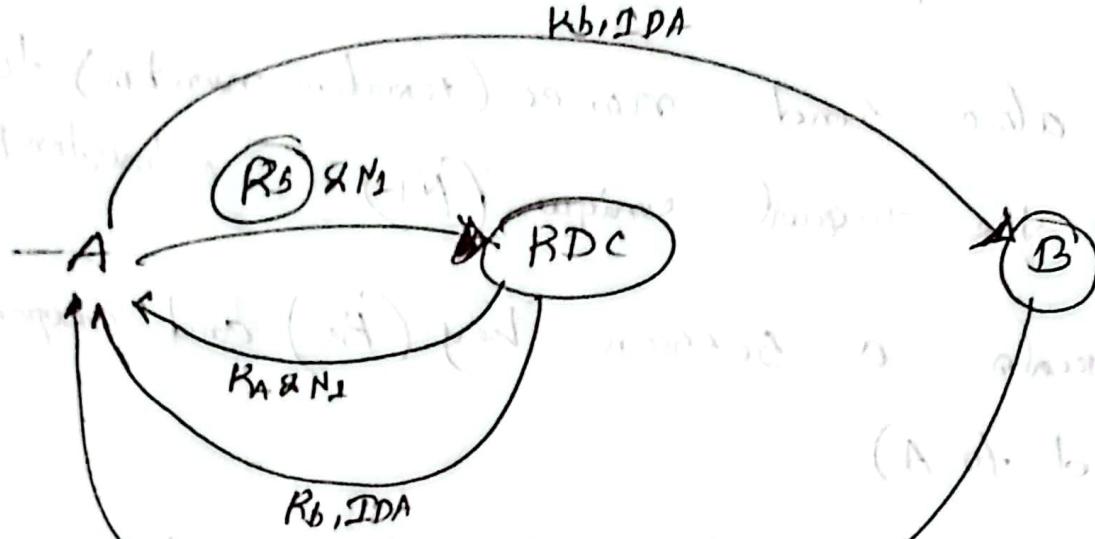
A user A want to send msg B securely

Step using KDC (3rd party)

- ① A sends a tag to KDC for a session key talk B.
- ② ~~The~~ KDC also send nonce (random number) to keep the request unique (N_1) with timestamp.
- ③ The KDC creates a session key (K_S) and prepare reply (both send to A)
 - . one part for A \Rightarrow include a session key and original request(N_1) encrypted with (K_A) master key (A's own key)
 - . another for B \Rightarrow includes session key and $\frac{A's\ ID}{ID_A}$ with B's master key (K_B) [$K_B [K_S] ID_A$]
- ④ A recvs the reply, decrypts 1st part using (K_A) to get session key (K_S), and sends B's part to B (because he can't open 2nd part of B's master key)

⑤ B decrypts it using R_b and gets some session key also known as $K_{b, IDA}$ from A.

Now both A & B securely talk each other.



Session Key Exchange