

Greatest Common Divisor:

⇒ More formally the positive integer c is said to be the greatest common divisor of a and b if

- ① c is a divisor of a and b
- ② any divisor of a and b is a divisor of c .

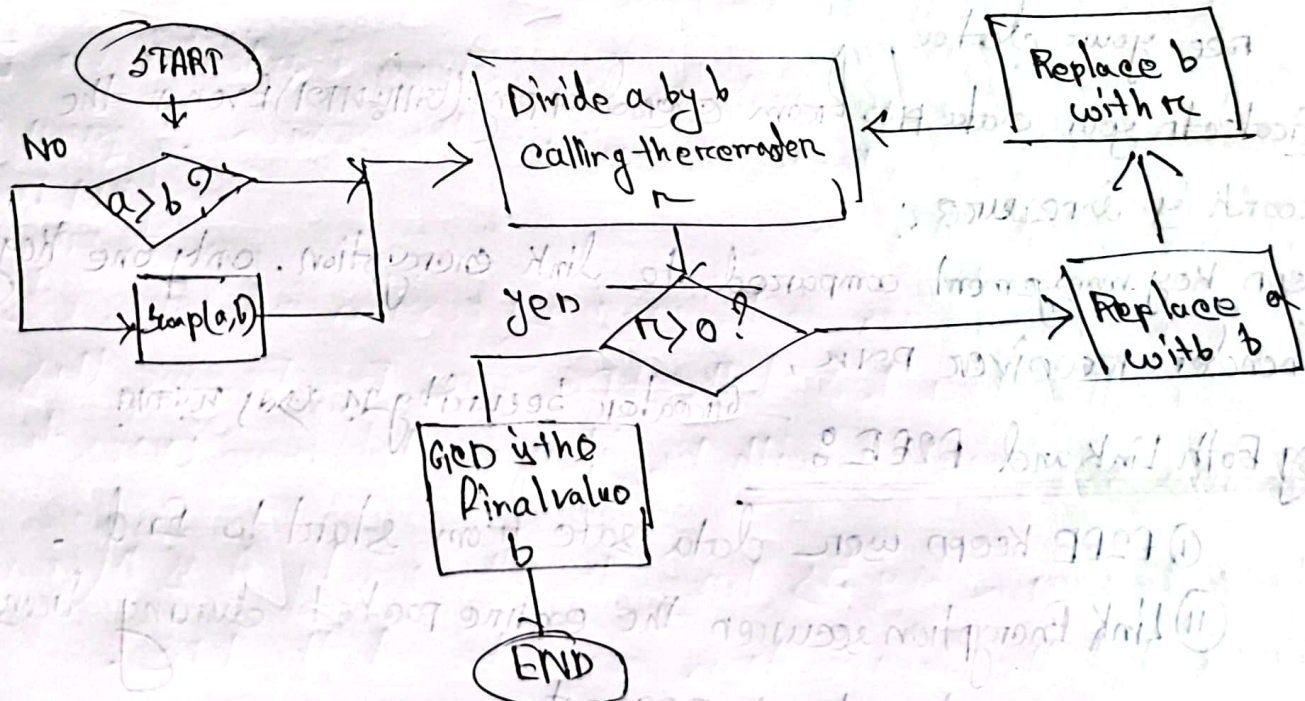
$$\text{gcd}(a, b) = \text{gcd}(|a|, |b|)$$

$$\text{gcd}(60, 24) = \text{gcd}(60, 24) = 12$$

$$\text{gcd}(a, 0) = |a|$$

⇒ Two integers a and b relatively prime if and only if their only common positive factor is 1. This is equivalent to saying that a and b are relatively prime if $\text{gcd}(a, b) = 1$.

$$\text{gcd}(8, 15) = 1 \rightarrow \text{relatively prime.}$$



➤ Link Encryption secures data by encrypting and decrypting information at every node or network switch it passes through, rather than just at the end points.

This ensures that all data including headers and routing information is protected.

Disadvantage: Router switch & Message decrypt & re-encrypt Address & Message Header (since). Message vulnerable to sniffing. Router & Link & Session Key vulnerable to sniffing.

➔ End To End Encryption (E2EE)

- Sender encrypts the message
- Message travels the network without being decrypted.
- Only receiver can decrypt it.

Advantage: ① Router switches and network operation can't see your data.

② Protects your data from eavesdropping (चूँकि) Even if the network is insecure.

③ Less Key management compared to link encryption, only one key per sender-receiver pair.

Using Both Link and E2EE: - Greater security & less management.

① E2EE keeps your data safe from start to end.

② Link Encryption secures the entire packet during transmission across each network segment.

[Router 1] Link Decrypt
Read Header
Link Encrypt Again

[Router 2] Link Decrypt
Read Header
Link Encrypt.

Key Distribution

For symmetric encryption both sender and receiver need to use same secret key to encrypt and decrypt message. Key must be protected from access by others.

The main challenge is sharing the key to parties. A — B

- ① A selects a key and physically delivers it to B.
- ② Third Party selects the key and physically delivers it to A and B.
- ③ If A and B have previously recently used a key, one party can transmit the new key to other, encrypted using old key.
- ④ If A and B each has an encrypted connection to a third party C, C can deliver a key on the encrypted link to A and B.

①, ② practically not possible

Network node level encryption ব্যবহৃত করলে যদি 100 Node থাকে তাহলে তাই বিজ্ঞতা Key distribute করতে হবে।

⇒ KDC (Key Distribution Center) used for key sharing easy and safe way to. two levels of keys are used.

Master Key → প্রত্যেকের Unique. Session Key / temporary key.

⇒ Session key usually used for the duration of a logical connection.

⇒ Session keys are transmitted in encrypted form using a master key that is shared by the key distribution center and an endpoint.

For N entities session key $\frac{N(N-1)}{2}$ master key N.

⇒ Not frequently session key change করে উত্তর দেয়া, কিন্তু বেশি প্রকারে করা হয়। তাহলে Network K slow ২০০ এর Resource খরচ হবে।

Transposition Key distribution scenario.

→ Plai 9) The scenario assumes that each user shares a unique master key with the Key distribution center (KDC).

Plain ② let assume a wishes to establish a logical connection with B and require a one time session key to protect data transmitted over the connection.

A has a master key K_A known only to itself and the KDC.

B n n n n K b n n n n n

The following steps occur:

The following steps occur:

- ① A issues a request to the KDC for a session key to protect a logical connection to B. This message includes A, B and unique identifier N_1 for this transaction. Which we refer to as a nonce.

It should be different ~~nonce~~ nonce \rightarrow timestamp / counter / random number
to guess \leftarrow ~~req~~ req 2, 10, 25

The KDC responds with a message encrypted using K_a .

Message includes \Rightarrow ① One-time session key 'k'.
 \Rightarrow ② The encrypted message.

Message includes \rightarrow ④ The original request: msg including nonrec.

to enable A to match this response ~~as~~ with appropriate request.

➔ In addition the message includes two items intended for B.

→ The one time session key K_{ts} to be used for the session. [$g^{a^2} = K_{ts}$]

→ An identifier of A .

A stores the session key for use in the upcoming session and forwards to B the information generated at the KDC for B. Because this information is encrypted with K_b it is protected from eavesdropping.

Bob now knows the session key K_s , knows the other party A is A (from $G(A)$) and knows that the information originated at the KDC.

At this point a session key has been securely delivered to A and B and they may begin their protected exchange.