LP Modeling Using Julia/ JuMP

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Our Goal

- In OR-related research, chances are you'll need to solve a linear or mixed integer program at some point.
- In this class, we're going to learn how to model LPs and MIPs in Julia, solve them, and do interesting things with them.
- We're going to use a new programming language called Julia and a package for Julia called JuMP (Julia for Mathematical Programming).

Background of Julia and JuMP

Julia:

- "high-level, high-performance, open-source dynamic language for technical computing"
- Dynamic nature allows rapid development, but Julia is also fast
 - Within factor of 2 of C/C++
- See Lubin and Dunning (2013) (http://arxiv.org/abs/1312.1431) for comprehensive experiments

JuMP:

 Algebraic Modeling Language (AML) for Julia developed by Miles and Iain

Outline

- Julia basics
- JuMP basics
- Facility location exercises
- Network revenue management exercises

References

- For further information on Julia:
 - Speed tutorial: http://learnxinyminutes.com/docs/julia/
 - Doc: http://docs.julialang.org/en/latest/manual/
- For further information on JuMP:
 - Doc: https://jump.readthedocs.org/en/release-0.2/jump.html
 - Issues: https://github.com/JuliaOpt/JuMP.jl/issues
- You may find it useful to have these open as we go through the class.

Getting started with Julia

- Open up a terminal window, and type in
 - julia

Basic stuff

- Simple arithmetic:
 - 3 + 7
 - 4.8 / 9.3
 - 11^5
- Comparisons:
 - 12 <= 5
 - !(5 == 8)
 - 5!=8
 - · (2014 % 4) == 0
- Strings:
 - "Velibor and Vishal are office mates"
 - "Velibor has \$(25 % 2) cat"

Variables

- Variables are easy:
 - x = 5
 - y = 12
 - z = x y
 - Z
 - W
- Everything has a type:
 - typeof(z) #Int64
 - typeof(z + 1.5) #Float64
 - typeof($z \ge 0$) #Bool

Arrays

- Arrays are for sequences of values, indexed from 1 to n:
 - $\cdot v = [1, 2, 3]$
 - v(1) # you'll get an error different from MATLAB
 - v[1] # access elements using []; starts from 1
 - v[0] # gives an error!
 - v[end] # last element
- 2D arrays are made the same way as in MATLAB:
 - A = [1 1 1; -1 -1 -2; 0 0 0]
 - B = [0.1 0.5 0.4; 0.2 0.3 0.5; 0.8 0.1 0.1]
 - B[2,3]
 - A + B
 - B^100
 - B[1,:]

Arrays continued

- Modifying arrays:
 - a = [1, 2, 3]
 - a[end+1] = 4 # won't work!
- Push/pop commands:
 - push!(a, 4) # append 4 to end of a
 - pop!(a) # remove last element of a; 4
 - a # a is back to being [1, 2, 3]
 - z = []
 - push!(z, 5) # what happens?
 - z = Int64[]
 - push!(z,5) # how about now?

Exercise 1a

- Create the arrays A, x and b as follows:
 - A = rand(10,3); x = rand(3,1); b = rand(10,1)
- Write an expression that evaluates to
 - true if the inner product of the fourth row of A and x is greater than the fourth element of b; and
 - false otherwise.
- What do you find?

Exercise 1a's lesson

- Julia is much stricter with arrays than MATLAB!
- For example:
 - A = rand(10,3); x = rand(3,1); b = rand(10,1)
 - A[4,:] * x b[4]
 - A[4,:] * x b[4] < 0 # what happens?
 - (A[4,:] * x b[4])[1] < 0 # how about now?
- Yet another trap!
 - z = [1:1:3] # what is it?
 - z[:] # what do you get?
 - z[:,:] # how about now?

Initializing arrays

- List comprehensions:
 - oddNumbers = [2*i 1 for i in 1:10]
 - evenNumbers = [j + 1 for j in oddNumbers]
- Can build matrices this way:
 - A = [i + j for i in 1:5, j in 1:5]

Tuples

- Tuples are "fixed" arrays:
 - \cdot t = (1, 2, 3)
- Can access them like arrays, but can't change them:
 - t[1] # gives back 1
 - t[1] = 5 # error!
- Tuples can be initialized in other ways:
 - x = tuple(1, 2, 3)
 - y = [1:5]
 - z = tuple(y...)

Dictionaries

- Dictionaries are like arrays, except they can be indexed by arbitrary objects:
 - mydict = Dict() # makes an empty dictionary.
 - mydict["Velibor"] = "is alright" # add key-value pairs...
 - mydict["Vishal"] = "is awesome!"
 - myotherdict = ["King's Landing"=> 1, "Winterfell"=> 2, "Qarth"=>3]
 - myotherdict["Pyke"] # gives an error
- Keys of dictionary:
 - keys(mydict)
- Values of dictionary:
 - values(mydict)

Conditionals and loops

```
• If-then-else:
   • x = 2
   • if x > 6
           x = x + 20
     elseif (x < 5)
           x = factorial(x)
     else
           println("Not changing x")
     end
for loops:
   • for k = 1:5
           println(k<sup>2</sup>)
     end
while loops:
   while x < 2000</li>
           print("*")
           \dot{x} = 2 \times \dot{x}
```

end

Functions

Functions are defined using the keyword function:

```
    function convexcomb(x,y, theta)
    (1 – theta) * x + theta * y
    end
```

 Functions by default return the last statement, but you can also use return:

```
    function convexcomb2(x, y, theta)
        if (theta < 0 || theta > 1)
            println("Bad theta!!!")
        else
            return (1 – theta) * x + theta * y
        end
        end
```

Exercise 1b

- Write a function that takes a matrix A, vectors x and b, and returns a vector containing the row indices for which Ax >= b does not hold.
- There are many ways to do this!
- Hints:
 - .< performs element-wise comparisons
 - "find" function returns indices of non-zero elements for Int64 or Float64 arrays, true elements for Bit arrays

Scripts in Julia

- Julia scripts are plain text files with the extension ".jl".
- In a terminal window, you can run one (say script.jl) by passing it as a command line argument to Julia – e.g.,
 - julia script.jl

Scripts continued

- Within a session, you can use the keyword 'using' to load everything in a script.
- E.g. suppose test.jl was:
 - function VeliborsFunction(x, y)
 return gamma(x)*gamma(y)/gamma(x+y)
 end
 - println("Velibor's function loaded!")
- In Julia session:
 - using test.jl
 - VeliborsFunction(1,2)
 - VeliborsFunction(0.5,0.5)

Time to JuMP!

Why an AML?

 Algorithms like simplex work with an optimization problem that is expressed in the form

minimize
$$c^T x$$

subject to $Ax = b$
 $x \ge 0$

- Functions like linprog take A, b, c explicitly as inputs, solve this problem and return an optimal x.
- While it is possible to represent the data A, b, c explicitly for any given model, doing the conversion – from a mathematical description, to solver-ready form – can be quite laborious.

Why an AML? (continued)

- An AML allows you to specify an optimization problem in a compact way that you, as a human being, can understand.
 - Optimization variables are represented by programming variables.
 - Constraints are built with expressions resemble.
 - Throughout, we can use all of the tools available to us from programming languages – conditionals, loops, functions, etc. – to build our model.
- Once the model is built in this way, the AML takes care of converting your description to solver-ready data (A, b, c).

Let's model a simple LP

- Run the following:
 - using JuMP
 - m = Model()
 - @defVar(m, x >= 0)
 - @defVar(m, 0 <= y <= 3)
 - @addConstraint(m, x + y <= 3)
 - @addConstraint(m, 5x + 3y <= 11)
 - @setObjective(m, Max, 1x + 4y)
 - print(m)
- What did we just model?

Now solve it!

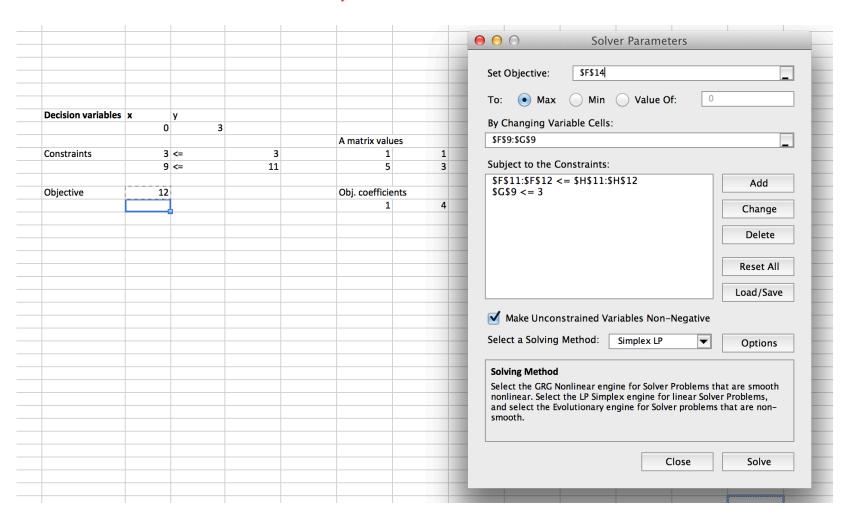
- Run the following:
 - status = solve(m)
 - println("Z = ", getObjectiveValue(m))
 - println("x = ", getValue(x))
 - println("y = ", getValue(y))

Without an AML, in MATLAB

In MATLAB, with linprog:

```
c = [-1; -4];
A = [1 1; 5 3];
b = [3; 11];
x = linprog(c, A, b, [], [], [0;0], [inf; 3])
```

Without an AML, in Excel



Building more general models

- More general ways to build things:
 - m2 = Model()
 - @defVar(m2, x[1:10, 1:2] >= 0)
 - @addConstraint(m2, sum{ x[i,j], i =1:10, j=1:2} <= 1)
 - primes = [2, 3, 5, 7]
 - @addConstraint(m2, sum{x[i,j], i in primes, j = 1:2} >=0.2)
 - @addConstraint(m2, sum{x[i,j], i =1:10, j =1:2; i + j <= 5} >= 0.05)
 - myTest(i,j) = i + j <= 5 # what did we do here?
 - @addConstraint(m2, sum{x[i,j], i = 1:10, j=1:2; myTest(i,j)} >= 0.05)

Solver options

- The Model() constructor accepts a specification of a solver – e.g.,
 - using JuMP, Gurobi
 - m = Model(solver = GurobiSolver())
- The solver constructor (GurobiSolver()) accepts parameters. Parameter names/values follow the same naming/meaning as within the solver – e.g.,
 - m = Model(solver = GurobiSolver(Method = 2, Crossover = 0))
 - m = Model(solver = GurobiSolver(Presolve = 0))

Facility Location

minimize $\sum \sum d_{ij} x_{ij}$ $i = 1 \ i = 1$ subject to $x_{ij} \leq y_j$, $\forall i = 1, ..., M, j = 1, ..., N$, $\sum x_{ij} = 1, \quad \forall i = 1, \dots, M,$ j=1 $\sum y_j \le K,$ i=1 $0 \le x_{ij} \le 1, \quad \forall i = 1, \dots, M, \ j = 1, \dots, N,$

 $y_i \in \{0, 1\}, \quad \forall j = 1, \dots, N.$

Without an AML...

- In MATLAB, using linprog?
 - Think about how you would need to keep track of indices of many rows and columns of A matrix will have double indices (rows because of x_ij <= y_j, columns because of x_ij).
 - This can be nasty to debug (trust me).
- In Excel, using Solver?
 - Could potentially handle it have large plots of spreadsheet space where rows and columns correspond to i and j.
 - But what if you have new data? What if the dimensions of the problem change?

Exercise 2

- Suppose that we live in a 1D world, where customers and facilities are located on a line as follows:
 - customerLocs = [3, 7, 9, 10, 12, 15, 18, 20]
 - facilityLocs = [1, 5, 10, 12, 24]
 - d_ij = absolute value distance between customer i and facility j
- Model and solve this problem in Julia using JuMP!
 (Set K = 3.)
- HINT: to make a variable binary, use "Bin" in @defVar:
 - @defVar(m, z[1:5], Bin)

Exercise 3

- In MIPs, it is sometimes useful to know not only the best solution, but also the second best, third best and so on.
- Modify your code so that, after the first solve, the model is solved another three times.
- Each time, add a constraint that cuts off the current facility decision, re-solve the model and output the objective value.

Exercise 3 Hint

Suppose our current solution is

$$\{\bar{y}_j\}_{j=1}^N$$

Suppose

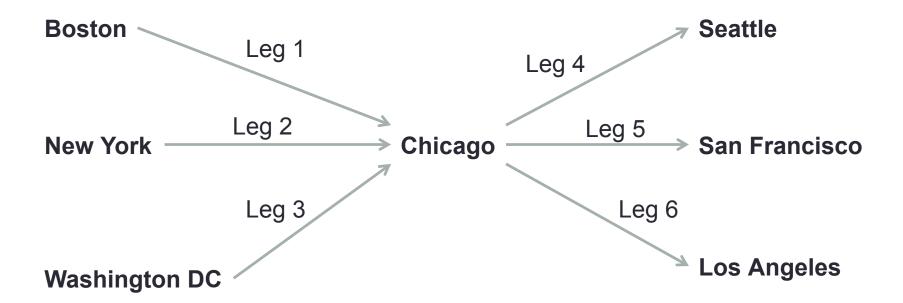
$$J_1 = \{j \mid \bar{y}_j = 1\}$$
 $J_0 = \{j \mid \bar{y}_j = 0\}$

Then the following constraint cuts off the current solution:

$$\sum_{j \in J_1} y_j + \sum_{j \in J_0} (1 - y_j) \le N - 1$$

Network revenue management

 Suppose we are running an airline that operates the following network:



Network revenue management - II

- We sell fares that may use capacity on one or two of the links in the network.
 - E.g.: we may sell Boston -> Chicago, Chicago -> LA, or Boston -> LA via Chicago.
- For now, we'd like to decide how many requests of each fare to accept, so as to maximize our revenue, subject to:
 - How much capacity is on each flight.
 - How many requests we expect of each fare, given that we are selling over T days.

LO Formulation

```
A_{\ell,f} = \begin{cases} 1 & \text{if fare } f \text{ uses link/leg } \ell, \\ 0 & \text{otherwise.} \end{cases}
R_f = \text{revenue from fare } f
p_f = \begin{cases} \text{probability of daily request} \\ \text{being of fare } f \end{cases}
T = \text{length of horizon}
```

LO Formulation

maximize
$$\sum_{f \in \mathcal{F}} r_f x_f$$
 subject to
$$\sum_{f \in \mathcal{F}} A_{\ell,f} x_f \le c_{\ell}, \quad \forall \ell \in \mathcal{L},$$

$$0 \le x_f \le T \cdot p_f, \quad \forall f \in \mathcal{F}.$$

Data

- We have the following data:
 - legs = [1, 2, 3, 4, 5, 6]
 - fareLegs = [(1), (2), (3), (4), (5), (6), (1,4), (1,5), (1,6), (2,4), (2,5), (2,6), (3,4), (3,5), (3,6)]
 - fareProbabilities = [0.06, 0.096, 0.046, 0.073, 0.159, 0.067, 0.043, 0.019, 0.112, 0.075, 0.031, 0.044, 0.012, 0.0210, 0.1130]
 - fareRevenues = [40, 30, 30, 10, 40, 10, 190, 80, 90, 70, 60, 190, 60, 50, 10]
 - legCapacities = [20, 20, 20, 20, 20, 20]
 - T = 100

- Formulate the LO problem in Julia and JuMP, and solve it.
 What is the optimal revenue?
- Hint: what happens when you try this:
 - z = (1, 2, 3) # what is the type of z?
 - 2 in z # what happens?
 - 4 in z # what happens?

- Create a function to solve the problem, that takes
 - legs, fareLegs, fareProbabilities, fareRevenues, legCapacities, T
- as inputs.
- You may use Ex5_NRM_Function_test.jl to test it.

Constraint references

- Like variables, constraints can have a "name" in JuMP, this is a constraint reference.
- Constraint references are created using the @defConstrRef macro.
- The output of @addConstraint can then be assigned to a constraint reference.
- Example:
 - m = Model(:Max)
 - @defVar(m, x[1:4] >= 0)
 - @defConstrRef myCons[1:4]
 - for i =1:4 myCons[i] = @addConstraint(m, x[i] <= 10)

Dual information

- JuMP also provides a function, getDual, for obtaining dual information.
- If cRef is a constraint reference, then getDual(cRef) returns the shadow price of the constraint.
- If var is a variable, then getDual(var) returns the reduced cost of the variable.

- Modify the function from Exercise 5 to return the shadow prices of the leg constraints in an array, as well as the optimal objective.
- You may use Ex6_NRM_DualFunction_test.jl to test it.

Online NRM

- In reality, the problem is an online problem.
- On each day, there is a request for a fare. We must decide whether to accept it or reject it.
 - Accepting a fare now gives us revenue, but uses capacity that could be used on more lucrative fares later.
 - Rejecting a fare keeps our capacity intact, but we do not earn any revenue.
- This is a notoriously difficult stochastic control problem.

Bid-price control

Let

dual price of leg
$$\ell$$
's capacity constraint, $\lambda_{\ell}(\vec{c},t) = \text{when the remaining capacity is } \vec{c}$ and t days remain.

 In BPC, the action to take (when there is sufficient capacity) is given by comparing the dual prices to the revenue of the fare:

$$\pi_t(f) = \begin{cases} \text{Accept} & \text{if } \sum_{\ell \in \mathcal{L}} A_{\ell,f} \lambda_{\ell}(\vec{c}, t) \leq r_f, \\ \text{Reject} & \text{otherwise.} \end{cases}$$

- Ex7_NRM_SimulateControl.jl contains two simulation functions – one for BPC and one for a greedy policy. The BPC function is missing two lines of code that you need to fill in.
- To test it, use Ex7_NRM_SimulateControl_test.jl.
- How does BPC compare to the greedy policy?

Thank you for listening!

- For further information on Julia:
 - Speed tutorial: http://learnxinyminutes.com/docs/julia/
 - Doc: http://docs.julialang.org/en/latest/manual/
- For further information on JuMP:
 - Doc: https://jump.readthedocs.org/en/release-0.2/jump.html
 - Issues: https://github.com/JuliaOpt/JuMP.jl/issues
 - Contact Miles and Jain
- If you find JuMP useful, please cite Miles and lain's paper!
 - M. Lubin and I. Dunning, "Computing in Operations Research using Julia", under review.